

A Pseudocode for the second question of the project of COMP 6651

Aria Adibi

Algorithm Pseudocode for the Question 2 of the project

```
1  Sort the intervals based on their end time; then “rename” them as  $I_1, \dots, I_n$ .
2  Add  $I_0 = (-\infty, end_0)$  as sentinel interval such that  $end_0$  is strictly smaller than
   start point of all intervals. //The purpose of the sentinel is to not check whether all
   intervals are considered or not in the following loops. Also note that in the graph this
   interval is a isolated vertex, therefore it will not change the answer.
3  Define  $d[i][j]$  where  $0 \leq i \leq n$  and  $0 \leq j \leq 1$  as an integer array and
   initialize it to all 0. //Note that integer overflow issue is ignored here; that was the
   reason for the modulo.

4  //Initialization
5   $d[0][0] = 0$ ;       $d[0][1] = 1$ 
6  for  $i = 1$  to  $n$ 
7       $d[i][0]$ 
8      //Partition the possibilities with respect to the latest possible interval that is in
   the Dominating Set
9       $j = i - 1$ ;      latestStartSoFar =  $I_i.start$ 
10     while  $I_j.end \geq latestStartSoFar$ 
11          $d[i][0] = d[i][0] + d[j][1]$ 
12         latestStartSoFar =  $\max\{latestStartSoFar, I_j.start\}$ 
13          $j = j - 1$ 
14      $d[i][1]$ 
15     //First, consider the cases in which at least one of the neighbors of  $I_i$  which is
   more restrictive than  $I_i$  is a member of Dominating Set, i.e.  $I_{neighbor}.start < I_i.start$ 
16     //Note that the covered intervals by  $I_i$ , i.e.  $I_{neighbor}.start \geq I_i.start$ , add no
   additional restrictions, therefore it is as if they are not there and just the final answer
   is  $\times 2$  for each, member of the Dominating Set or not.
17      $j = i - 1$ ;      numCoveredIntervals = 0
18     while  $I_j.end \geq I_i.start$ 
19         if  $I_j.start \geq I_i.start$ 
20             numCoveredIntervals = numCoveredIntervals + 1
21         else
22              $d[i][1] = d[i][1] + 2^{numCoveredIntervals} \times d[j][1]$ 
23          $j = j - 1$ 
24     //Second and final case is when no restrictive neighbor of  $I_i$  is in the Dominating
   Set. Note that now  $j$  is pointing to the first interval that is not a neighbor of  $I_i$ 
25      $d[i][1] = d[i][1] + 2^{numCoveredIntervals} \times (d[j][0] + d[j][1])$ 
26 print  $d[n][0] + d[n][1]$ 
```
