

Due: 04.02.17

Instructor: Dr. Pawan Kumar

Maximum Marks: 45

INSTRUCTIONS:

The codes can be written in either C, C++, MATLAB, Octave, FORTRAN, or Python. However, it is also recommended to write these codes in C/C++. Usually, it is a good practice to write two versions of the same algorithm, first in MATLAB or Python, and then the same in C/C++. Write useful comments, and do proper indentation.

For testing, write a code that generates random matrix of desired size. In MATLAB, this can be done with `A = rand(m,n)`.

The questions prefixed with ******* are “somewhat” challenging, they may not be considered for grading. A grace period of 1 week may be provided after due date.

1. (Matrix Multiplication) Given two matrices A and B . Write a code for matrix-matrix multiplication when [2+2+3+3+2]
 1. A and B are **dense** matrices.
 2. A and B are **banded** matrices. Use the banded storage scheme discussed in class.
 3. (Write this in only C/C++) A and B are **sparse** matrices (a matrix that has significantly more number of zeros compared to number of non-zeros). For this case, consider the following storage schemes.
 - (a) **Coordinate Storage Format (COO)**: In this format, the matrix entries are stored in three arrays, namely, `row_indices`, `col_indices`, and `val` :
 - i. The array `row_indices` contains the row indices of non-zero entries. It is of length `nz`. Here `nz` refers to the total number of non-zeros.
 - ii. The array `col_indices` contains the column indices. It is of length `nz`.
 - iii. Array `val` contains the matrix entries at the corresponding row and column. It is also of length `nz`.

For example, for the following matrix

$$\begin{bmatrix} 1 & 0 & 0 & 2 & 0 \\ 3 & 4 & 0 & 5 & 0 \\ 6 & 0 & 7 & 8 & 9 \\ 0 & 0 & 10 & 11 & 0 \\ 0 & 0 & 0 & 0 & 12 \end{bmatrix}, \quad (1)$$

the arrays `row_indices`, `col_indices`, and `val` are given as follows

```
row_indices = [5 3 3 2 1 1 4 2 3 2 3 4],
col_indices = [5 5 3 4 1 4 4 1 1 2 4 3],
val = [12 9 7 5 1 2 11 3 6 4 8 10].
```

Note that the entries of the array `val` are not written in ordered way, for example, the first entry is 12, which is the (5,5)*th* entry of the matrix.

(b) **Compressed Sparse Row (CSR) Format:** Here again, the matrix data is stored in three arrays, namely, `val`, `col_indices`, and `row_pointers`:

- i. All the matrix entries are stored in arrays `val` row by row. Along each row, they are stored from smallest column number to the largest. The length of `val` is `nz`.
- ii. An integer array `col_indices` contains the column indices of the elements stored in the array `val` above.
- iii. An integer array `row_pointers` contains the pointers to the beginning of each row in the arrays `val` and `col_indices`. Thus, the content of `row_pointers(i)` is the position in arrays `val` and `col_indices` where the *i*-th row starts.

For the matrix (1) above, the CSR format is given as follows

```
val = [1 2 3 4 5 6 7 8 9 10 11 12],
col_indices = [1 4 1 2 4 1 3 4 5 3 4 5],
row_pointers = [1 3 6 10 12 13].
```

4. Determine the storage complexity and flops for all three items above.

2. **(LU Factorization)** Given a matrix $A \in \mathbb{R}^{n \times n}$.

[3+2+2+8]

1. Write a code to compute the LU factorization of A . The factors L and U must be overwritten in A . Name this function `mylu.*`. Write a code that implements forward substitution (`foreward.*`) with a lower triangular matrix L . Similarly, write a code that implements backward substitution (`back.*`) with an upper triangular matrix U .
2. Use algorithms above to write a code, for example, `lu_solve.*` that takes the matrix A and a right hand side vector b as inputs, and it outputs the solution to the equation $Ax = b$ by first doing forward substitution, i.e., by solving $Lt = b$, then performing the backward substitution, i.e., by solving $Ux = t$. Check your solution by computing $\|b - Ax\|_2$, it should be close to zero (in double precision arithmetic). [Note that LU factors are stored in A , so keep another copy of A to check your solution.]
3. Adapt your code to optimally compute the LU factorization of a Hessenberg matrix.
4. ***** (Write this in C/C++)** Write a LU factorization routine when the matrix A is stored in CSR format. Name this function `lu_sparse.*`. The L and U factors must also be stored in CSR format. Then write a routine `forward_sparse.*` to do forward substitution for a sparse lower triangular matrix stored in CSR format, similarly, write a backward substitution routine `backward_sparse.*` for a sparse upper triangular matrix stored in CSR format. As before, to solve $Ax = b$, you need to first solve $Lt = b$, which is a forward substitution, and then solve $Ux = t$, which is a backward substitution to obtain the solution x to the given linear system $Ax = b$.

3. (QR Factorization) Given $A \in \mathbb{R}^{m \times n}$.

[2+2+2+3+3+6]

1. Given a vector $x \in \mathbb{R}^m$, write a code that computes $v \in \mathbb{R}^m$ with $v(1) = 1$, and $\beta \in \mathbb{R}$ such that the Householder matrix $P = I_m - \beta vv^T$ is orthogonal and $Px = \|x\|_2 e_1$. Name this function `house.*`.
2. Write a code `myqr.*` that computes Householder QR. The Q and R factors must be overwritten in A . Then use this factorization to solve the linear system $Ax = b$ by using the fact that

$$Ax = b \implies QRx = b \implies Rx = Q^T b.$$

Since R is upper triangular, back substitution easily gives the desired solution x .

3. Adapt your code to optimally compute the LU factorization of a Hessenberg matrix.
4. Write a code `recursive_qr.*` for recursive implementation of QR algorithm as discussed in class.
5. Let $m > n$, then the linear system $Ax = b$ may not have a solution. In this case, we may seek a minimum norm solution, for example, a least squares solution. Use the QR algorithm above to write a code that solves the following least squares problem

$$\min_{x \in \mathbb{R}^n} \|b - Ax\|_2.$$

6. ***** (Write this in C/C++)** Repeat item 2 above, when the given matrix A is stored in CSR format. Determine the flop count in this case.