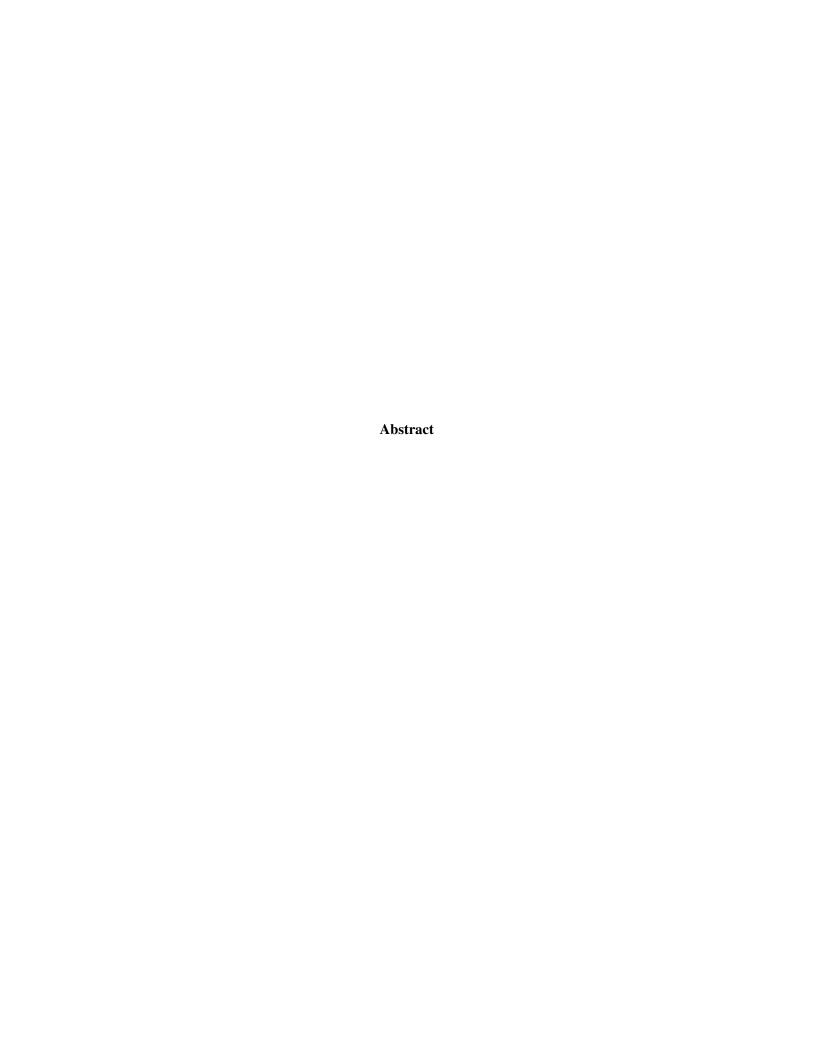
Predator Prey Model

Names Here

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Introduction

In this task we worked as a group to implement a 2D, sequential predator-prey algorithm with spatial diffusion in Java. We tried to design the program in such a way as to take advantage of Java's object-orientated, modular nature. Due to the large size of our group we also decided to develop a GUI for the program, although the code was designed in such a way as to be usable from the command line as well.

—is it?—

There is always some compromise between readability and performance in coding, and we tried to keep this balance in mind when designing our code. We made best use of class structure whilst still keeping our implementation of the given algorithm as efficient as possible.

The predator prey algorithm implemented in this project crudely models the interaction between population densities of different animals, specifically Hares and Pumas. Each population has a self-interaction coefficient (birth for Hares, death for Pumas) and a coefficient describing how it interacts with other species' populations.

These populations exist on a 'world' consisting of land and water, and are also able to diffuse across land squares with a rate determined by a diffusion coefficient. Thus, each population has N+1 coefficients which determine it's behaviour, where N is the number of animals.

Model

$$H_{ij}^{new} = H_{ij}^{old} + \Delta t (C_1 H_{ij}^{old} + C_2 H_{ij}^{old} P_{ij}^{old} + l (H_{i+1j}^{old} + H_{i-1j}^{old} + H_{ij+1}^{old} + H_{ij-1}^{old} - H_{ij}^{old}))$$

```
for k=1:NumberOfAnimals if (k=ThisAnimal) then N_k^{new} = N_k^{old} + \Delta t C_k(k) N_k^{old} else N_k^{new} = N_k^{old} + \Delta t C_k(k) N_k^{old} N_{ThisAnimal}^{old} end if end  \qquad \qquad \text{where}  C_{Hare} = \left( \begin{array}{cc} H & PH \end{array} \right)
```

Design & Implementation

- **3.1** Code
- 3.1.1 Structure
- 3.1.2 Algorithm
- 3.1.3 Input/Output
- 3.1.4 GUI
- 3.2 Tools
- 3.2.1 SVN
- 3.2.2 Makefile
- 3.2.3 Unit Testing

Every class of the project form a functional unit according to the design elaborated at the beginning of the

has been tested using the Java JUnit framework

The libraries necessary to the tests are thus accessed from the test cases using Java import statements and extending the TestCase superclass in the class declaration:

Listing 3.1: Test case headers

```
1 import org.junit.Before;
2 import org.junit.Test;
3 import junit.framework.TestCase;
```

```
5 public class TestAnimal extends TestCase { 6 \ldots
```

Then the test cases, for each class, follow the normal directives in use for Unit testing with JUnit. The most relevant methods were submitted to thorough tests, where their properties, return values and correct functioning were evaluated by comparing the actual results they provide with the expected ones. These procedures were implemented by test methods in each test case using the format required by the JUnit standards.

Every test case included a setup() method that builds the initial object and its different parameters. Directives like @Before and @Test for every method declaration were thus inserted in the code were appropriate. The functioning of a method is then assessed comparing the expected results with the actual ones with JUnit testing methods such as AssertEquals(), AssertNotNull(), etc.

Listing 3.2: Use of JUnit directives in test cases

```
1 @Before
2 public void setUp() {
3
      testAnimal = new Animal(numbAnimals);
4
5
      testAnimal.setName("Puma");
      testAnimal.setDiffCo(diffCoIn);
6
7
8 @Test
9 public void testAnimal() {
     assertNotNull(testAnimal);
11
      assertNotNull(testAnimal2);
12 }
13 ...
```

Performance Analysis

4.1 Testing

Tests were preformed in order to profile the code's performance and scalability. The unix *t*ime command was used to time the code runs, with 'user' and 'system' time being summed to find the total computation time.

The relationship between total computation time per cell and the number of cells was measured in order to quantify the overhead needed by the code (see Figure 4.1). We found that, below a grid size of $\sim \! 100$ by 100, overhead became important, whereas at larger grid sizes the computation time scaled linearly with the total number of cells. Any scaling worse than linear (such as total computation time being proportional to the NumberOfCells^{1.2} would be a huge (and unecessary) inefficieny, especially with larger grids.

Some overhead is inevitable in any code, but the amount of effort put into simplifying a problem (such as creating arrays of neighbours) should depend on the expected size of the problem. Our code could probably have had less overhead, allowing it to run faster with smaller gridsizes, but this would likely led to it having a worse (linear) scaling with cell number.

4.2 Analysis

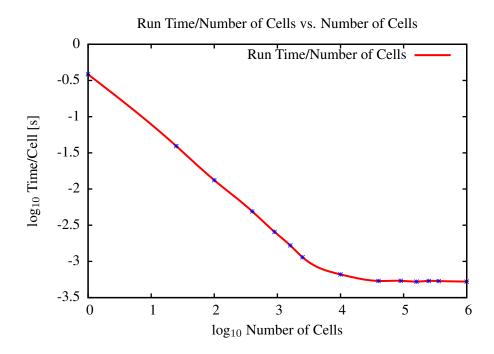


Figure 4.1: This is a graph of total time spent on each cell, which shows that overheads such as array initialization and output (i.e. things not directly involved in the simulation) take up a significant fraction of computation time with grids smaller than ~ 100 by 100.

Conclusions

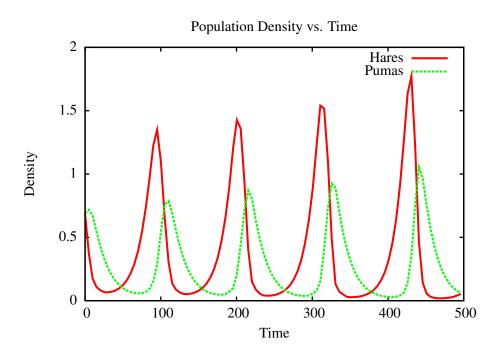


Figure 5.1: Population density vs. time for Hares and Pumas. This periodic behaviour is typical of predator-prey interactions.

Group Evaluation