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chapter: 17

(I) inserting n elements using

(a) Aggregate method

→ The table doubles in size when it runs out of space.

→ So, if the original size is 1, after insertion it doubles to size 2, after 2 more insertions, it doubles to size 4 etc.

→ In general, after k doublings, the size is 2^k .

* Pseudo code:

→ initialize table with capacity = 1.

for $i = 1$ to n ,

if if table is Full:

new table = create new table with size $2 * \text{current size}$ copy element from old table to new table

table = new table

insert element i into table

let $k = \log(n+1) - 1$

total cost = $O(n) * k$
 $= O(n \log n)$

Amortize cost per insertion
 $= O(\log n)$

Runtime per insertion is
 $O(\log n)$ //

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Total time is $O(n) * \log(n+1)$,

(b) Accounting method

* charge 2 units for each insertion.

→ when the table doubles in size from m to $2m$, credit m units.

→ The credit exactly pay for the copy cost of $O(m)$

→ Total credit is $m + 2m + 4m + \dots + \frac{n}{2}m = O(n)$

* Pseudo code:

→ Initialize table with capacity = 1
For $i = 0$ to n :

if table is Full:

newtable = create new table
with size 2 * current size
copy elements from old table to new
table, table = newtable

insert element i into table

initialize charges = 0

initialize credits = 0

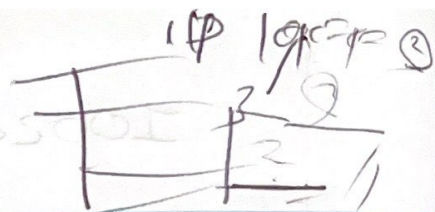
for $i = 1$ to n :

charges += 2

if table doubled in size from
 m to $2m$:

credit += m

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$$\text{Total charges} = 2 \times n = O(n)$$

$$\text{Total credits} = n + 2n + \dots + \frac{n}{2} = O(n)$$

Amortized cost per insertion

$$= \text{Total} / n = O(n) / n$$

$$= 1$$

Runtime per insertion $O(1)$

$$\text{Total time} = O(n)$$