

Assignment 5

Q.1. Construct LL Parsing table for the given grammars.

$$S \rightarrow iEtS \mid iEtSes \mid a$$

$$E \rightarrow b$$

After removing the left recursion,

$$S \rightarrow iEtSS' \mid a$$

$$S' \rightarrow eS \mid \epsilon$$

$$E \rightarrow b$$

First

Follow

$$S \rightarrow \{i, a\}$$

$$S \rightarrow \{e, \$\}$$

$$S' \rightarrow \{e, \epsilon\}$$

$$S' \rightarrow \{e, \$\}$$

$$E \rightarrow \{b\}$$

$$E \rightarrow \{t\}$$

Predictive Parsing table

	i	t	a	e	b	\$
S	$S \rightarrow iEtSS'$		$S \rightarrow a$			
S'				$S' \rightarrow eS$		$S' \rightarrow \epsilon$
E					$E \rightarrow b$	

Q.2. Perform shift reduce parsing for the given input string

$$S \rightarrow S + S$$

$$S \rightarrow S * S$$

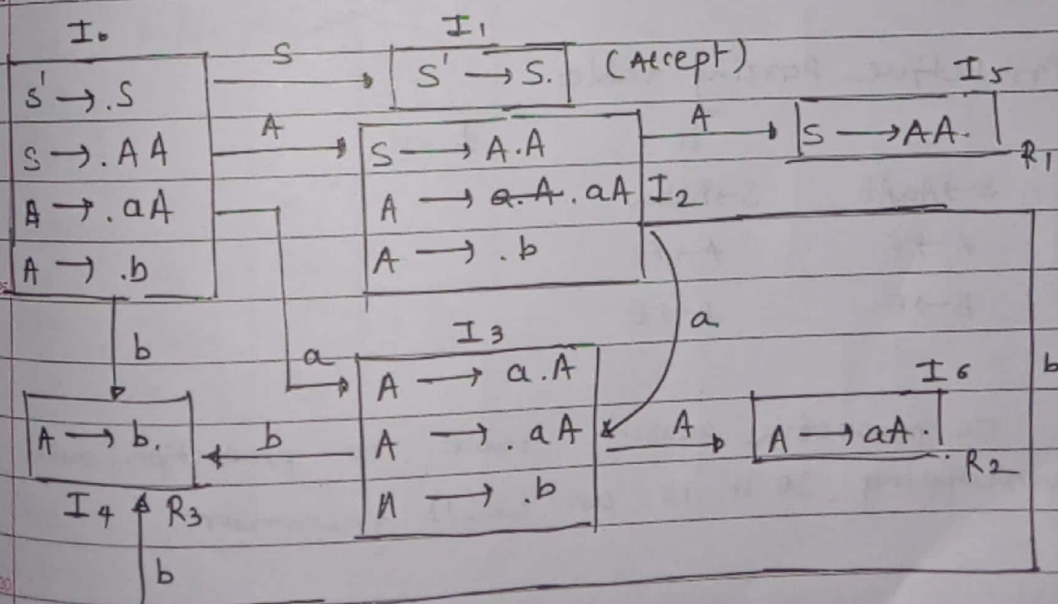
$$S \rightarrow id$$

input string : $id + id + id$

Input: id + id + id \$	Action: shift	Stack: \$
+ id + id \$	Reduce { $s \rightarrow id$ }	\$ id
+ id + id \$	shift	\$ s
id + id \$	shift	\$ s +
+ id \$	Reduce { $s \rightarrow id$ }	\$ s + id
+ id \$	shift	\$ s + s
+ id \$	Reduce { $s \rightarrow s + s$ }	\$ s
id \$	shift	\$ s +
\$	Reduce shift	\$ s + id
\$	Reduce { $s \rightarrow id$ }	\$ s + s
\$	Reduce { $s \rightarrow s + s$ }	\$ s
\$	Accept	\$ s

Q3 Construct LR parsing table for the given grammars.

- $s \rightarrow AA$
 $A \rightarrow aA \mid b$
- 1) $s' \rightarrow .s$
 - 2) $s \rightarrow .AA$
 - 3) $A \rightarrow .aA$
 - 4) $A \rightarrow .b$



State	Action			Goto	
	a	b	\$	A	S
I ₀	S ₃	S ₄		2	1
I ₁			Accept		
I ₂	S ₃	S ₄		5	
I ₃	S ₃	S ₄		6	
I ₄	R ₁	R ₁	R ₃		
I ₅	R ₁	R ₁	R ₁		
I ₆	R ₂	R ₂	R ₂		

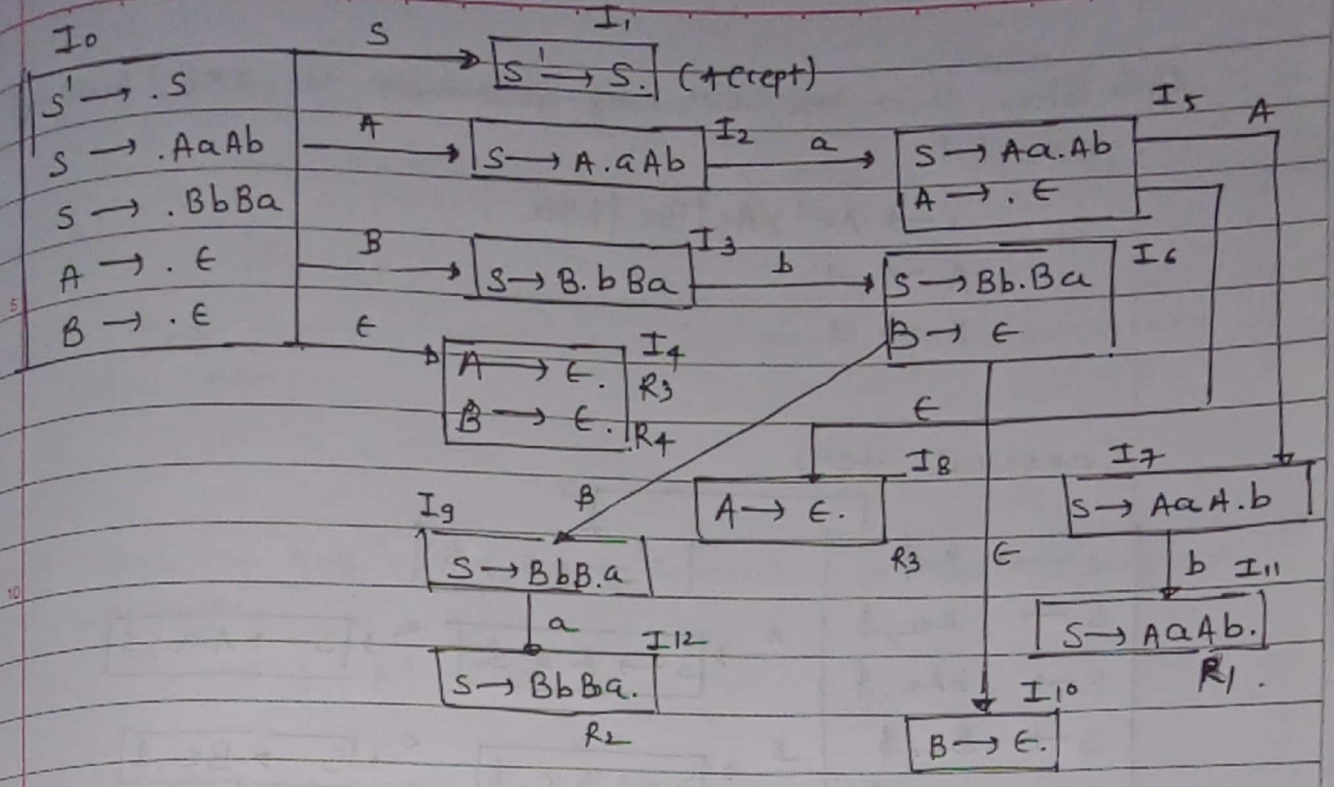
Q.4 Show that the following grammar is LL(1) but not SLR(1)

	first	follow
$S \rightarrow AaAb \mid BbBa$		
$A \rightarrow \epsilon$	$S \rightarrow \{a, b\}$	$S \rightarrow \{\$, \}$
$B \rightarrow \epsilon$	$A \rightarrow \{\epsilon\}$	$A \rightarrow \{a, b\}$
	$B \rightarrow \{\epsilon\}$	$B \rightarrow \{b, a\}$

Predictive Parsing table

	a	b	\$
S	$S \rightarrow AaAb$	$S \rightarrow BbBa$	
A	$A \rightarrow \epsilon$	$A \rightarrow \epsilon$	
B	$B \rightarrow \epsilon$	$B \rightarrow \epsilon$	

In predictive parsing table, no production rules are overlapping. So it is an LL(1) grammar



States	Actions				Goto		
	a	b	ϵ	\$	S	A	B
I ₀			S ₄		1	2	3
I ₁				Accept			
I ₂	S ₅	S ₆					
I ₃		S ₆					
I ₄	R ₃ /R ₄	R ₃ /R ₄					
I ₅			S ₈			7	
I ₆			S ₁₀				9
I ₇		S ₁₁					
I ₈	R ₃	R ₃					
I ₉	S ₁₂						
I ₁₀	R ₄	R ₄					
I ₁₁				R ₁			
I ₁₂				R ₂			

There is Reduce-Reduce conflict, so this is not SLR(1) grammar

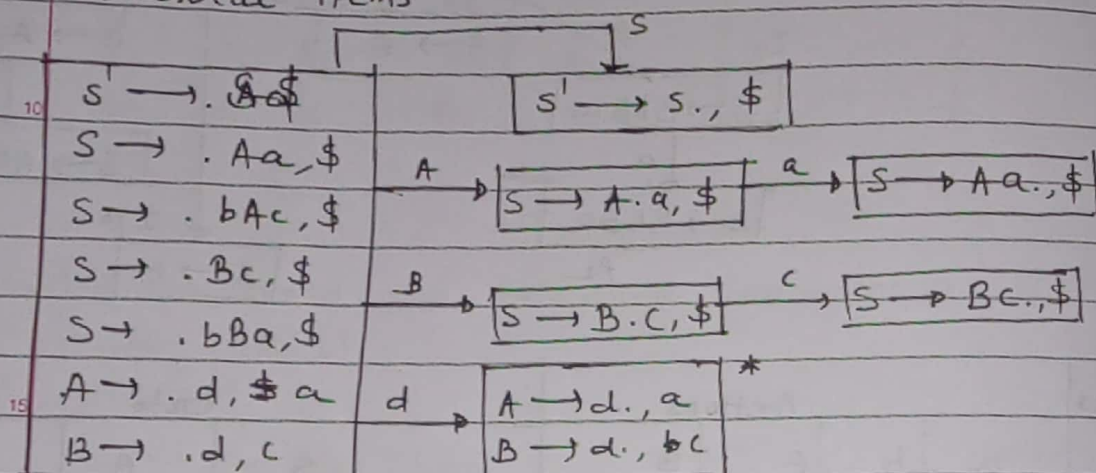
Q5 Show that the following grammar is LR(1) but not LALR(1)

$$S \rightarrow Aa | bAc | Bc | bBa$$

$$A \rightarrow d$$

$$B \rightarrow d$$

Canonical items



for LR(1)	a	b	c	d	\$
I_0	R5		R5		
I_3	R6		R6		

Here, there is no RR conflict, hence given grammar is LR(1)

for LALR(1), combining state 2 & state 3

state	a	b	c	d	\$
2, 3	R5/R6		R5/R6		

Here Reduce-Reduce occurs, hence given grammar is not LALR(1)

Q.6 Parse the give string using operator precedence parser
id + id * id

$$E \rightarrow E + E \mid E * E \mid id$$

Operator Precedence Table

	id	+	*	\$
id		>	>	>
+	<	>	<	>
*	<	>	>	>
\$	<	<	<	Accept

Input: id + id * id \$	Action: shift	stack: \$
+ id * id \$	Reduce { $E \rightarrow id$ }	\$ id
id * id \$	shift	\$ E
id * id \$	shift	\$ E +
* id \$	shift	\$ E + id
* id \$	Reduce { $E \rightarrow id$ }	\$ E + E
id \$	shift	\$ E + E *
\$	shift	\$ E + E * id
\$	Reduce { $E \rightarrow id$ }	\$ E + E * E
\$	Reduce { $E \rightarrow E * E$ }	\$ E + E
\$	Reduce { $E \rightarrow E + E$ }	\$ E
\$	accept	\$ E