## OSFanatics - CS 270

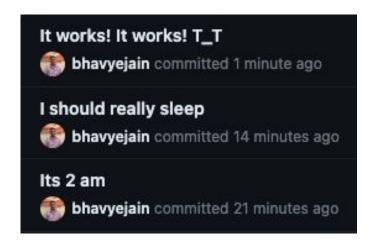
AltFS

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#### File-tastical follies

#### What we lost and learnt through our odyssey

- Code only as good as tests
- Logs for life!
- #SleepDeprivedButOurFileSystemThrived (hopefully:))

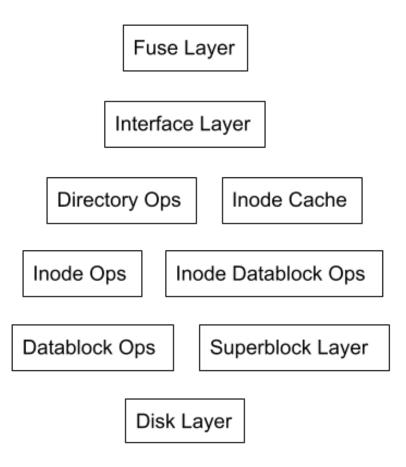


## Because 1 filesystem just wasn't enough...

We built & mounted our FS inside our FS and it works!!

```
[root@ip-172-31-19-208 mnt]# pwd
/home/cloud-user/AltFileSystem/disk/AltFileSystem/mnt
[root@ip-172-31-19-208 mnt]# mkdir dir1
[root@ip-172-31-19-208 mnt]# touch dir1/file.txt
[root@ip-172-31-19-208 mnt]# echo "Hello world!" > dir1/file.txt
[root@ip-172-31-19-208 mnt]# cat dir1/file.txt
Hello world!
[root@ip-172-31-19-208 mnt]#
```

### Code structure for AltFS



# Filesystem config parameters

Block size : 4096 B

Inode structure

```
Follows a structure similar to ext4.
(https://www.kernel.org/doc/html/latest/filesystems/ext4/inodes.html?highlight=inode)
struct inode
   mode_t i_mode; // permission mode
   id_t i_uid; // lower 16-bits of owner id
   id_t i_gid; // lower 16-bits of group id
   time_t i_status_change_time; // status change time
   time_t i_atime; // last access time
   time_t i_ctime; // last inode change time
   time_t i_mtime; // last data modification time
   nlink t i links count; // hard link count
   ssize_t i_file_size; // file size
   ssize_t i_blocks_num; // num of blocks the file has
   bool i_allocated; // flag to indicate if inode is allocated
   // TODO: Kept number of direct blocks as 12 in sync with ext4
   ssize_t i_direct_blocks[NUM_OF_DIRECT_BLOCKS];
   ssize_t i_single_indirect; // stores block num for single indirect block
   ssize_t i_double_indirect; // stores block num for double indirect block
   ssize t i triple indirect; // stores block num for triple indirect block
   nlink_t i_child_num; // stores the current number of entries for a directory (minimum 2) or 0 for others
    char padding;
```

Inode blocks

Inode cache capacity

: 1.5% of total blocks

: 100000 records

#### **Makefs**

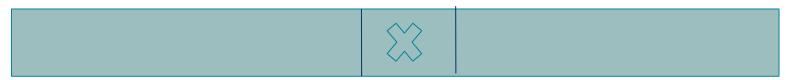
- Iteration 1: Erasing disk contents and formatting as part of initialization.
  - Erase required to prevent file corruption.
- Iteration 2: A separate mkaltfs utility and ability to load from disk on startup.
  - Persist data on unmount and remount.
- Iteration 3: Remove requirement for erasing disk while formatting!
  - Mkaltfs gets options!

#### **Directory records**

- No holes in a data block
- Record length = 0 => no directory entries from that point onwards

Record Length (2 B)	Inode number (8 B)	File name (variable)
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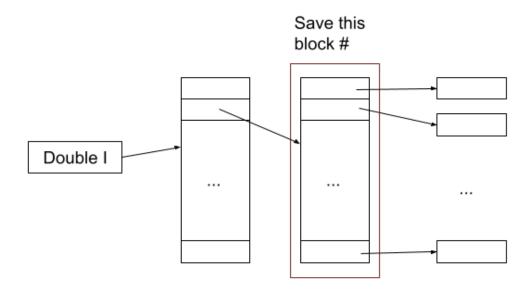
- Add new entry search for first record where record length = 0
- Remove entry Copy rest of the block starting from the next entry into buffer, delete the current entry and move the rest of the blocks left. This leaves no holes



 Get file position in directory - Check for each block until record len = 0 and if file name matches given name

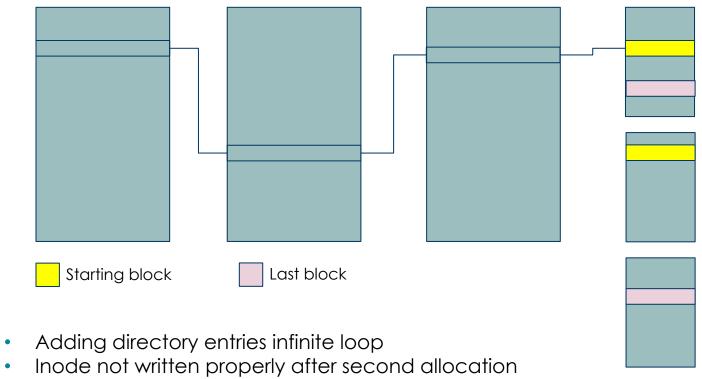
#### Inode data block traversal

- Fetching double and triple indirect blocks is expensive!
- Make it stateful!



# How our tests helped better our code

Remove data blocks from inode starting from a given logical number



- Pointer freeing in cache flush
- operating vim and gcc writing more bytes than requested

## Limitations

- Data locality too many seeks!
- Finding next free inode is expensive.
- Permissions are partially supported.
- Symlinks not supported.

# UC SANTA BARBARA