



← Notes

▲ Introduction To Dynamic Programming

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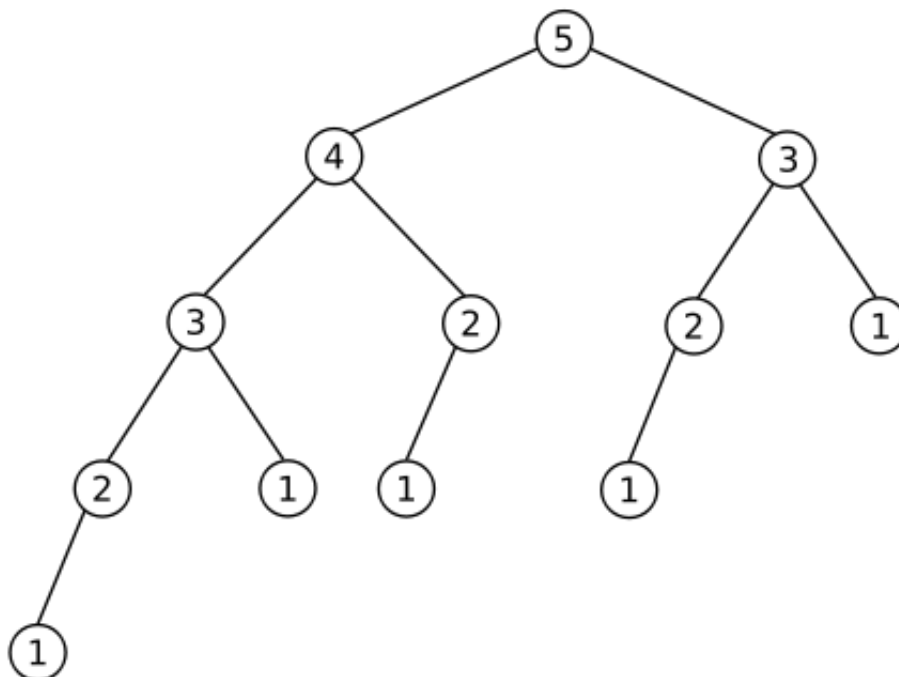
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Introduction To Dynamic Programming

DP is a very useful and effective technique for finding optimal solutions for problems having exponential time complexities($O(n!)$ or $O(2^n)$) as it may bring down the complexity to $O(n^2)$ or $O(n^3)$.

DP applies to the problems having overlapping sub-problems. DP computes sub-problems before the problem itself and combines value of sub-problem to generate solution to the problem. Since problems are overlapping we may end up computing same sub-problems over and over again. So, DP speeds up the process by storing those results in a table and fetch result instead of computing it again.

For ex : Take an example of Fibonacci series, $\text{fib}(4)$ and $\text{fib}(3)$ are sub-problems to $\text{fib}(5)$ and as we can see from image we'll compute $\text{fib}(3)$ twice one while computing $\text{fib}(5)$ and one while computing $\text{fib}(4)$, that's a lot of wastage, in DP technique we'll store $\text{fib}(3)$ and use this result when required.



Let's take a simple example: Problem: Given an array of integers of size n . We've to answer Q queries. Each query consists of 2 integers L and R and we've to print sum of elements of array from L to R (inclusive).

input: 5 1 2 3 4 5 2 1 5 2 5 output: 15 14

Naive: For each query iterate through L to R and calculate the sum

```
for q<-1 upto Q:
    sum = 0
    for i<-L upto R:
        sum = sum + Arr[i]
    print sum
```

∴ for single Query takes $O(n)$ time and for answering Q queries it takes $O(Q*n)$ time.

DP : Find and store the cumulative frequency in a table(array)

```
table[1...n] = {0}
table[0] = Arr[0]
for i<-1 upto n:
    table[i] = table[i-1] + Arr[i]
```

this pre-computation requires $O(n)$ time. Now we can answer each query in $O(1)$ time as follows:

```
for q<-1 upto Q:
    print (table[R] - table[L-1])
```

∴ for Q queries complexity becomes $O(1*Q) = O(Q)$

So, overall complexity of problem using DP is: $O(Q+n)$ which is linear in time. Though DP requires Extra space of order $O(n)$, but the reduction in the time complexity covers for it.

Exercise: Write an algorithm to compute Fibonacci series using DP.

Note : C++ codes related to this article can be found on <https://github.com/techmon/Posts-Related-Codes>

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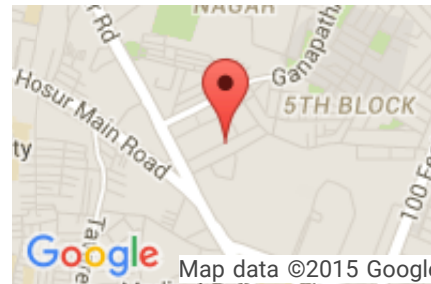
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