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## ▲ Lucas' Theorem. Wilson's Theorem

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CodeMonk

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### Lucas' Theorem

*Statement:*

$$C(N, K) \% MOD = (C(n_0, k_0) * C(n_1, k_1) * ... * C(n_{m-1}, k_{m-1})) \% MOD$$

$n_0, n_1, \dots, n_{m-1}$  and  $k_0, k_1, \dots, k_{m-1}$  are representations of the numbers  $N$  and  $K$  in the scale of notation with base  $MOD$ . In other words:

$$N = n_0 * MOD^0 + n_1 * MOD^1 + \dots + n_{m-1} * MOD^{m-1}$$

$$K = k_0 * MOD^0 + k_1 * MOD^1 + \dots + k_{m-1} * MOD^{m-1}$$

$C(N, K)$  is [Binomial coefficient](#) (number of ways to choose  $K$  elements from a set of  $N$  elements).

*Conditions:*  $MOD$  is a prime number (look at the end of the article to know what can we do with not prime  $MOD$ ), and you should be able to calculate  $C(n_i, k_i) \% MOD$ , where  $(0 \leq n_i, k_i < MOD)$ .

*Advices:* this theorem is very useful in case  $N \geq MOD$ , otherwise it's better to use formula  $C(N, K) = N! / ((N - K)! * K!)$  and tricks #2 or #3 from [there](#). If  $N \geq MOD$  then  $N! \% MOD = 0$ , when  $C(N, K) \% MOD$  is not necessary equals to 0.

*Realization:* let's see how can we get representation of some number  $N$  in the scale of notation with base  $MOD$ :

```
vector<int> getRepresentation(int N) {
    vector<int> res;
    while (N > 0) {
        res.push_back(N % MOD);
        N /= MOD;
    }
    return res;
}
```

Let  $n$  will be representation of  $N$  and  $k$  will be representation of  $K$ . They are not necessary have the same length. If  $K > N$  we can easily say that  $C(N, K) = 0$ . Otherwise  $k$  has less or equal length than  $n$ . To make them the same length we can add some extra zeroes to  $k$  and make them both of length of  $n$ , or we can take only some first elements of  $n$  and make them both of length of  $k$ . The second way has more sense because  $C(n_i, 0) = 1$ .

So the main part of code looks like:

```
vector<int> n = getRepresentation(N);
vector<int> k = getRepresentation(K);
long long res = 1;
for (int i = 0; i < k.size(); ++i) {
    res = (res * C(n[i], k[i])) % MOD;
}
```

Let's talk about function  $C(n[i], k[i])$  in more detail. It's easy to see that  $(0 \leq n[i], k[i] < \text{MOD})$ , so we can use formula  $C(N, K) = N! / ((N - K)! * K!)$  and trick #3 from [there](#):

```
int C(int N, int K) {
    if (K > N) {
        return 0;
    }
    return (((fact[N] * binpow(fact[N - K], MOD - 2)) % MOD) * bi
}
```

Let's precalc all possible factorials modulo  $\text{MOD}$  and store them in the array `fact`:

```
long long fact[MOD];
fact[0] = 1;
for (int i = 1; i < MOD; ++i) {
    fact[i] = (fact[i - 1] * i) % MOD;
}
```

Function `binpow` is just [Fast exponentiation](#), it can calculate  $A^N \% \text{MOD}$  in  $O(\log(N))$  time:

```
int binpow(int a, int n) {
    long long res = 1;
    while (n > 0) {
        if (n % 2 != 0) {
```

```

        res = (res * a) % MOD;
    }
    a = ((long long)a * a) % MOD;
    n /= 2;
}
return (int)res;
}

```

If  $n[i]$  and  $k[i]$  are small enough instead of using formulas and tricks we can just precalc [Pascal's triangle](#) and then get  $C(n[i], k[i])$  in  $O(1)$ :

```

int C[MOD][MOD];
for (int i = 0; i < MOD; ++i) {
    for (int j = 0; j <= i; ++j) {
        if (i == 0 || j == 0) {
            C[i][j] = 1;
        } else {
            C[i][j] = (C[i - 1][j - 1] + C[i - 1][j]) % MOD;
        }
    }
}
}

```

*Trick with not prime MOD:* let's factorize  $MOD = \text{mod}_1^{q_1} * \text{mod}_2^{q_2} * \dots * \text{mod}_m^{q_m}$  and calculate  $C(N, K) \% \text{mod}_1, C(N, K) \% \text{mod}_2, \dots C(N, K) \% \text{mod}_m$  using Lucas' Theorem. Now we can use [Chinese remainder theorem](#) to restore  $C(N, K) \% MOD$ .

## Wilson's Theorem

*Statement:*

**Natural number  $N$  is a prime number if and only if  $(N - 1)! + 1$  is divisible by  $N$ .**

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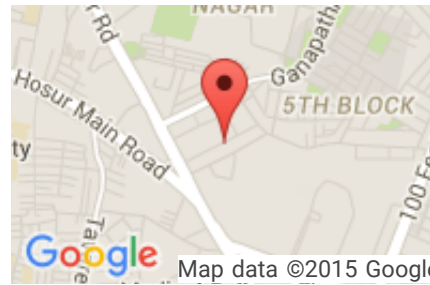
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