

## Parameterized Algorithms

### Quiz 1

The allotted time is 180 minutes. Solve all problems. All problems carry equal marks.

### Problems

1. Recall that a language  $L$  is *decidable* if there is an algorithm which (i) halts on every input, and (ii) correctly reports whether an arbitrary input string belongs to  $L$ .
  - (i) Define the following:
    - A parameterized language
    - A fixed-parameter tractable (FPT) algorithm
    - A kernelization algorithm, and a kernel
  - (ii) Let  $L$  be a decidable parameterized language. Prove that  $L$  is fixed-parameter tractable if and only if  $L$  has a kernel.
2. Recall that a subset  $S \subseteq V(G)$  of vertices of a graph  $G$  is a *vertex cover* of  $G$  if every edge in  $G$  has at least one of its two end-vertices in  $S$ . The standard parameterization of the VERTEX COVER problem is as follows:

VERTEX COVER

Parameter:  $k$

Input: A graph  $G$  and an integer  $k$ .

Question: Does  $G$  have a vertex cover of size at most  $k$ ?

- (i) Write pseudocode for the  $O^*(2^k)$  FPT algorithm for VERTEX COVER which we discussed in class.

Your pseudocode should consist of a function `vertexCover()`, invoked as `vertexCover(G,k)`, which returns true if graph  $G$  has a vertex cover of size at most  $k$ , and false otherwise. You should write *what* the algorithm should do, and omit the fine details of *how* it should be done. In particular, you should assume that you have access to a Graph API using which you can manipulate and query graphs; you should *not* implement a graph object of your own from scratch. Your pseudocode for `vertexCover(G,k)` should not be much more than ten or fifteen lines.
  - (ii) Prove that your algorithm correctly solves VERTEX COVER, and that it runs in  $O^*(2^k)$  time.
3. In this problem we improve on the previous algorithm.



- (i) Write the pseudocode for an algorithm which solves VERTEX COVER in  $O^*(1.5^k)$  time.

Once again, your pseudocode should consist of a function `vertexCover()`, invoked as `vertexCover(G,k)`, which returns true if graph  $G$  has a vertex cover of size at most  $k$ , and false otherwise. It should also have the other properties of pseudocode as stated in Problem (2), except that your function may have more than fifteen lines this time.

- (ii) Prove that your algorithm correctly solves VERTEX COVER, and that it runs in  $O^*(1.5^k)$  time.

4. Recall that a subset  $S \subseteq V(G)$  of vertices of a graph  $G$  is a *feedback vertex set* (FVS) of  $G$  if every *cycle* in  $G$  has at least one of its vertices in  $S$ . The standard parameterization of the FEEDBACK VERTEX SET problem is as follows:

FEEDBACK VERTEX SET

Parameter:  $k$

Input: A graph  $G$  and an integer  $k$ .

Question: Does  $G$  have a feedback vertex set of size at most  $k$ ?

- (i) Write pseudocode for the  $O^*(5^k)$  FPT algorithm for FEEDBACK VERTEX SET which we discussed in class.

As for previous problems, you should assume that you have access to a Graph API using which you can manipulate and query graphs.

- (ii) Prove that your algorithm correctly solves FEEDBACK VERTEX SET, and that it runs in  $O^*(5^k)$  time.