

1. Does Gaussian elimination preserve the column space of a matrix? Justify. (3 marks)
2. In the simplex method discussed in class, we made an assumption that the feasible region is bounded. Give a modified simplex algorithm that handles unbounded feasible regions. Explain in a few sentences why your modifications work. (5 marks)
3. Write the dual of the following LP: (4 marks)

$$\begin{array}{ll}
 \text{Maximize} & 8x_1 + 3x_2 - 2x_3 \\
 \text{Subject to} & x_1 - 6x_2 + x_3 \geq 2 \\
 & 5x_1 + 7x_2 - 2x_3 = -4 \\
 & x_1 \leq 0 \\
 & x_2 \geq 0
 \end{array}$$

4. Exhibit a primal-dual pair such that both are infeasible. (5 marks)
5. Consider the following problem. (8 marks)

$$\begin{array}{ll}
 \text{Maximize} & 10x_1 + 24x_2 + 20x_3 + 20x_4 + 25x_5 \\
 \text{Subject to} & x_1 + x_2 + 2x_3 + 3x_4 + 5x_5 \leq 19 \quad (C1) \\
 & 2x_1 + 4x_2 + 3x_3 + 2x_4 + x_5 \leq 57 \quad (C2) \\
 & x_1, x_2, x_3, x_4, x_5 \geq 0
 \end{array}$$

- (a) Write its dual with two variables  $w_1, w_2$  (corresponding to the constraints (C1) and (C2)) and verify that  $(w_1, w_2) = (4, 5)$  is a feasible solution.
- (b) Use complementary slackness to show that  $(w_1, w_2) = (4, 5)$  gives the optimal solution to the dual.
6. (a) Write the ILP for finding min-cost perfect matching in a bipartite graph with edge weights.
- (b) Show that in the relaxed LP, every non-integral feasible point can be expressed as a convex combination of two distinct feasible points.
- (c) Use (b) to show that every  $k$ -regular bipartite graph has a perfect matching<sup>1</sup>. (10 marks)
7. Here is the set cover problem:

**Input.** A universe  $D$  consisting of finite number of elements, and a family  $S_1, S_2, \dots, S_m$  of sets, with each  $S_i \subseteq D$ . Assume that  $\bigcup_{i \in \{1, \dots, m\}} S_i = D$ .

**Goal.** Find minimum size subset  $W \subseteq \{1, \dots, m\}$  such that  $\bigcup_{i \in W} S_i = D$

- (a) Write an ILP for the set cover problem. (3 marks)
- (b) Give the relaxed LP and its dual. (4 marks)
- (c) Design a primal-dual algorithm which gives an  $f$ -approximation where  $f$  is the maximum frequency of an element, i.e.  $f = \max_{e \in D} \{\text{Number of sets containing element } e\}$ . In other words, show that the answer given by the algorithm is at most  $f$  times the optimal solution. Describe the algorithm and its analysis. (8 marks)

<sup>1</sup>Graph theoretic answers not using (b) will not be accepted