

# Event Ticketing Platform

## Solution Architecture Assessment

February 2026

# Problem & Scenarios

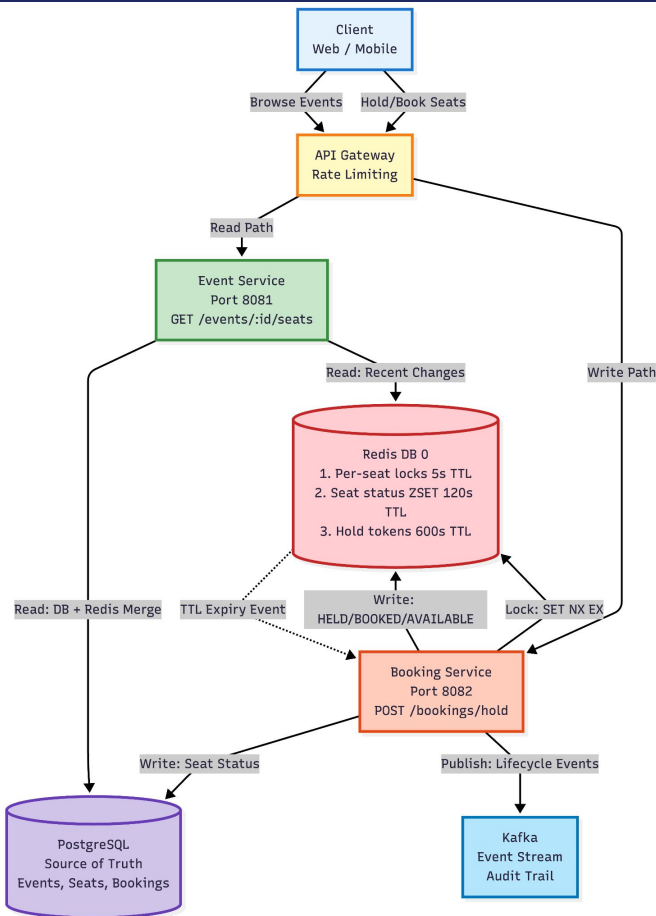
Category	Scenario	Strategic Rationale
Write Path	Seat Hold (e.g 10-Minute TTL)	Addresses high-concurrency locking, distributed coordination, expiry handling, and event-driven workflows
Read Path	Event Discovery with Caching	Optimizes search latency, availability reads, and cache consistency strategies

Challenge	Business Risk / Impact
Double Booking	Revenue loss and reputational damage due to concurrent seat allocation conflicts
Abandoned Holds	Inventory lock leading to reduced conversion and revenue leakage
Lack of Audit Trail	Operational risk in dispute resolution and regulatory traceability

## Success Criteria (North Star Outcomes)

Zero double bookings | Automated seat release via deterministic TTL enforcement | End-to-end booking lifecycle auditability

# Architectural Overview



## Architecture Style

Microservices-based, event-driven system separating Event and Booking domains for scalability and fault isolation.

## Concurrency Control

Redis distributed locks + ACID DB transactions ensure seat integrity under high contention.

## Event-Driven Workflow

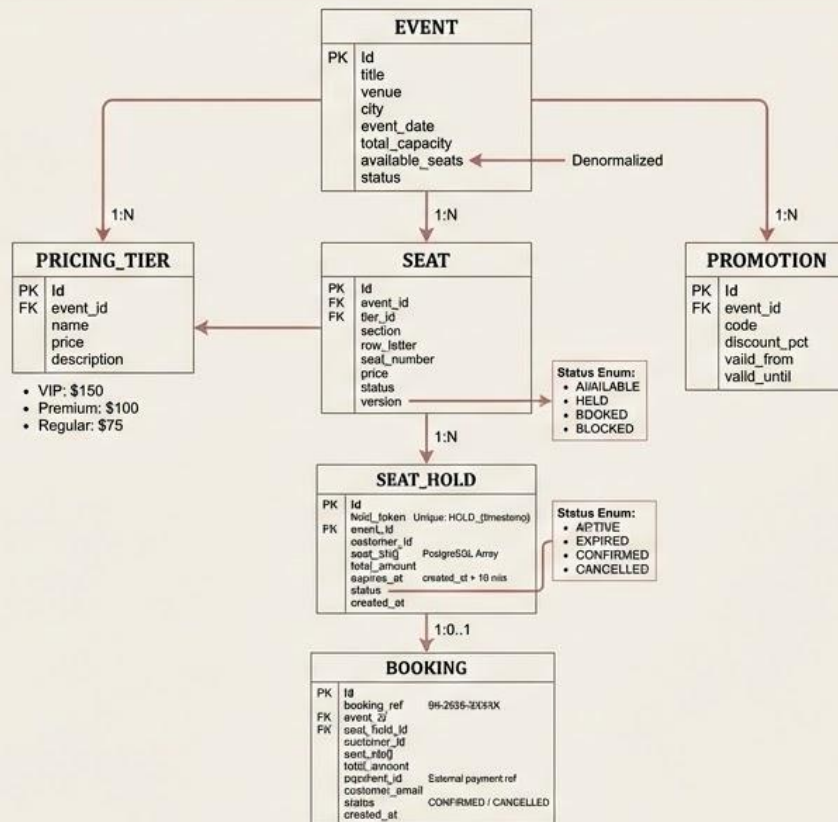
Kafka publishes lifecycle events:

SEAT-HOLD-CREATED, SEAT-HOLD-CONFIRMED, SEAT-HOLD-CANCELLED,  
SEAT-HOLD-EXPIRED, BOOKING-CONFIRMED

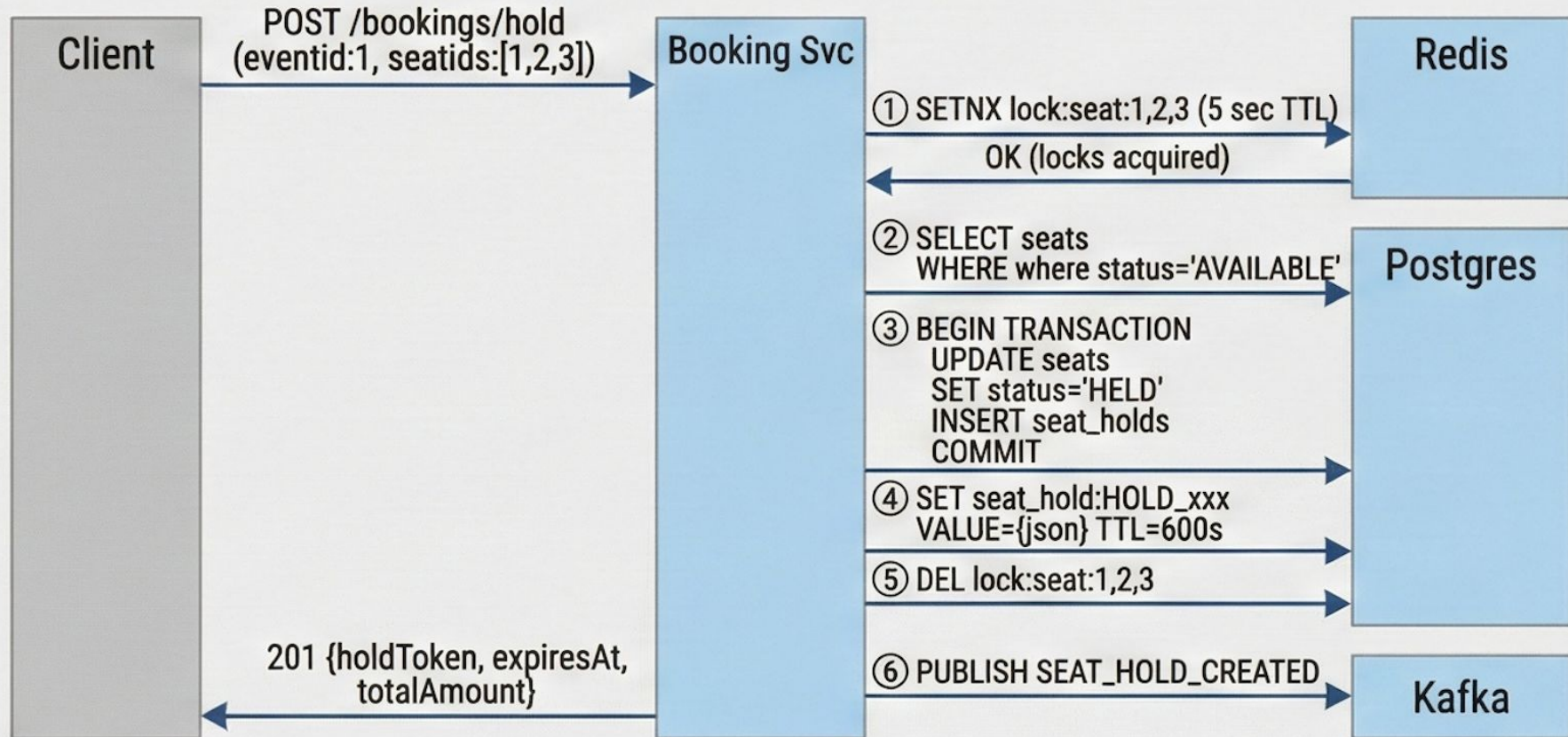
## Resilience & Recovery

TTL-based auto release, idempotent APIs, and service isolation prevent deadlocks and duplicate bookings.

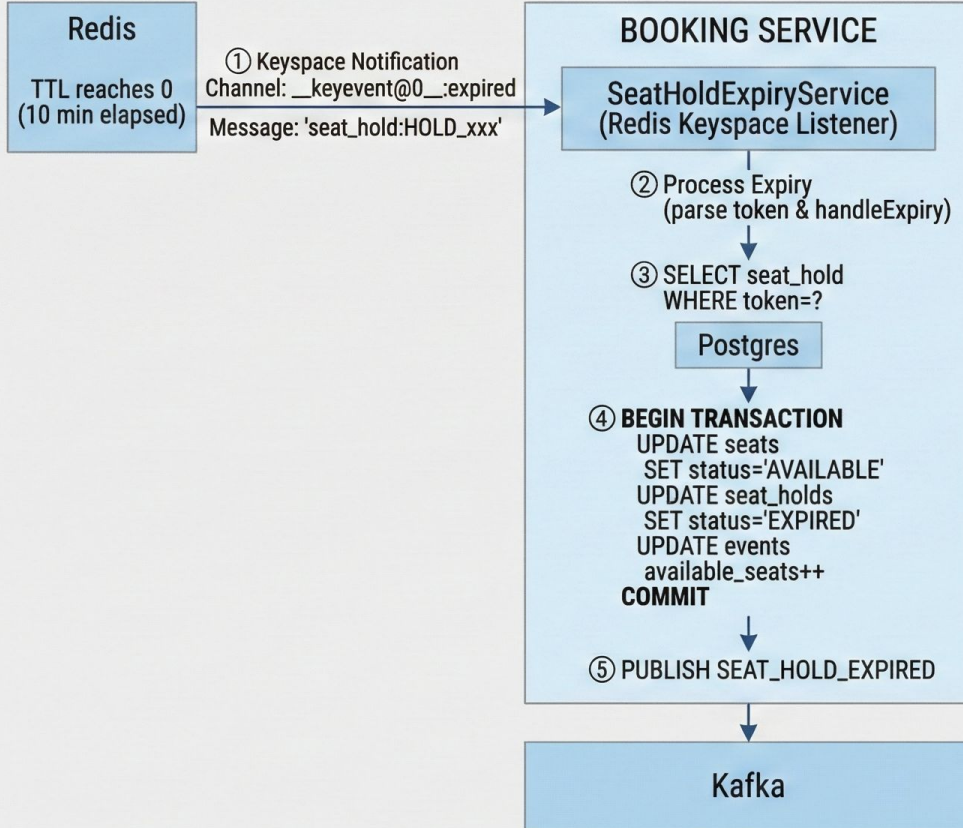
# Data Entity Design



## Scenario Flow - Seat Hold (Write)



## Scenario Flow - Auto-Expiry (Timeout)



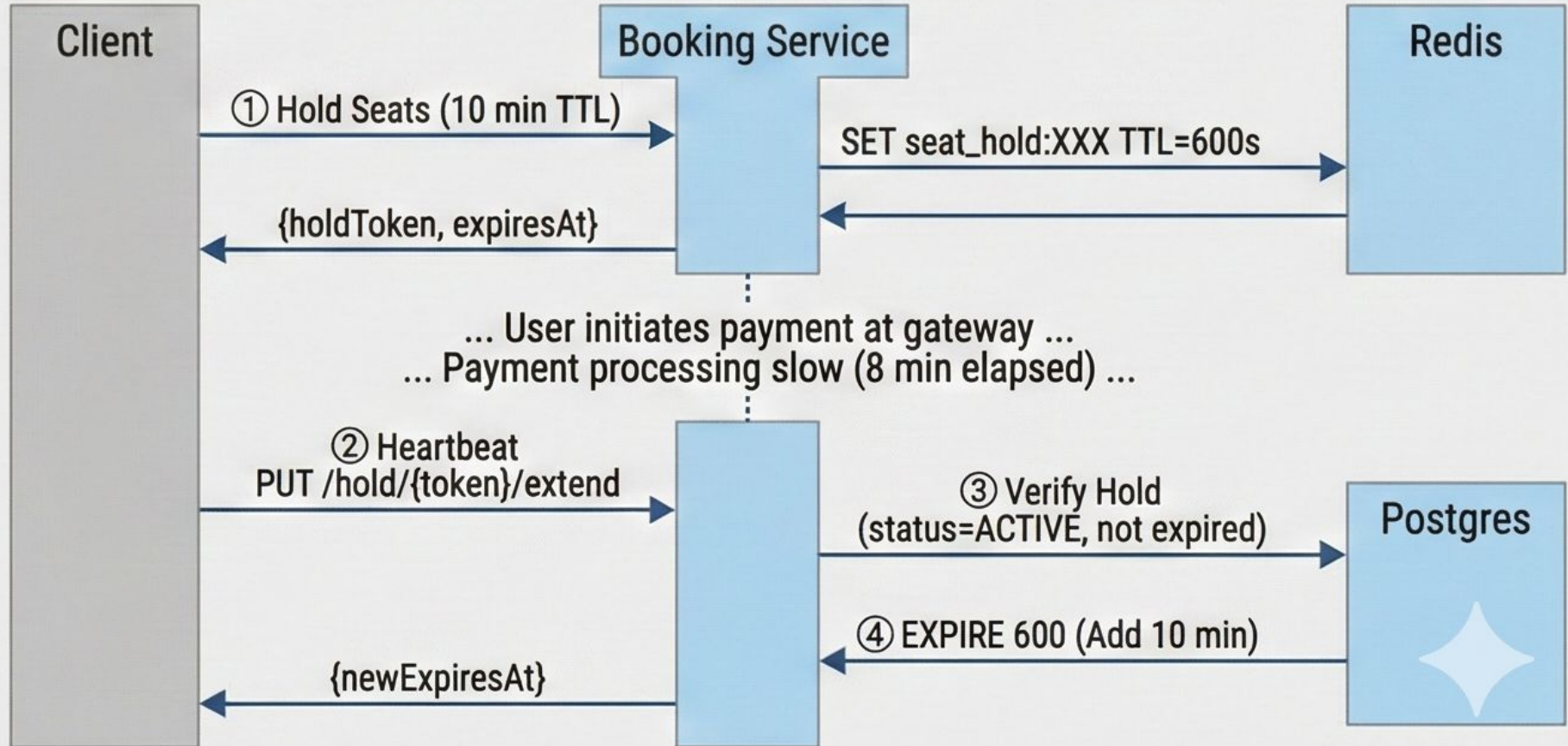
Seats now AVAILABLE for other customers

## Assumption :

Booking not completed within Hold-Window

Redis key pattern: `seat:{eventId}:{seatId}:HELD`

## Sequence 2: Heartbeat Extension for Slow Payment Gateways



# Core Architectural Decision

## Database as Source of Truth

All booking confirmations validated via **conditional DB update** Prevents double booking under race conditions Redis used only for coordination, not durability Strong consistency enforced at booking boundary

## Deterministic Expiry & Self-Healing

Seat holds enforced with TTL  
Delayed expiry validation via background consumer  
Periodic reconciliation job (Redis ↔ DB drift correction)  
System remains consistent even after crashes

## Distributed Locks

**SET NX EX** with unique lock token (UUID)  
Lua-based safe unlock (ownership verification)  
Short lock TTL (e.g., 30s) to avoid deadlocks  
Designed to handle high-contention flash-sale scenarios

## Idempotent by Design State Machine

The holdToken acts as a natural, single-use idempotency key for the confirm flow. A retry against an already-confirmed hold is rejected by the isActive() pre-check and the WHERE status='HELD' guard. domain state itself prevents duplicate effects

# Key Architectural Trade-offs

## PostgreSQL for Transactional Data

Strong ACID guarantees for bookings

Trade-off: Scaling requires partitioning

## Redis for Distributed State

Fast locks, TTL-based holds, Real-Time data

Trade-off: Additional infrastructure complexity

## Pessimistic Locking + Distributed Locks

Prevent double booking at multiple layers, supports Redis unavailability

Trade-off: Retry logic, potential contention

## Microservices Architecture

Independent scaling, fault isolation

Trade-off: Increased operational complexity

## Event-Driven Async Processing

Decouple notifications, analytics

Trade-off: Eventual consistency

# API Design - Complete User Journey

## GET /api/events/search

Browse events by city, date, category

200 OK with paginated results

## GET /api/events/{id}/seats

View seat layout and availability

200 OK with seat map

## POST /api/bookings/hold-seats

Hold selected seats (10 min TTL)

200 OK with hold ID

## POST /api/bookings/{hold-token}/confirm

Create booking from hold

201 Created with booking ref

## GET /api/bookings/{id}

Get booking details

200 OK with full details

## POST /api/bookings/hold/{hold-token}?customerid={castID}

Cancel hold & return to AVAILABLE pool

200 OK with clear hold

# Failure Scenario

## 1. Service Crash After Lock Acquisition

- Scenario: Lock acquired in Redis, but DB write fails/crashes.
- Risk: Seat locked in Redis, not in DB (state drift).

### Mitigation:

1. Persist DB immediately after lock (fail-fast).
2. On DB failure → Safe unlock using Lua script:  
if redis.call("GET", KEYS[1]) == ARGV[1] then  
3. return redis.call("DEL", KEYS[1])  
4. end
5. Reconciliation job (every 5 min):
  - Find Redis locks without DB records → Delete lock.
  - Find DB HELD records without Redis locks → Recreate/release.

## 2. Redis Crash / Restart

- Scenario: All locks lost after Redis restart.
- Risk: DB and Redis state drift (DB=HELD, no lock).

### Mitigation:

1. DB is source of truth; Redis is derived/cache.
2. Periodic reconciliation job:
  - DB=HELD & no Redis lock → Recreate lock or mark expired.
  - Redis lock & DB=AVAILABLE → Delete stale lock.
3. Use Redis persistence (AOF/RDB) for faster recovery.
4. Fallback: Use DB row locks if Redis is down (degraded mode).

# Failure Scenario continued..

## 3. Lock Expiry During Payment Processing

**Scenario:** The distributed lock (e.g., Redis lock) expires before the full payment transaction is finalized.

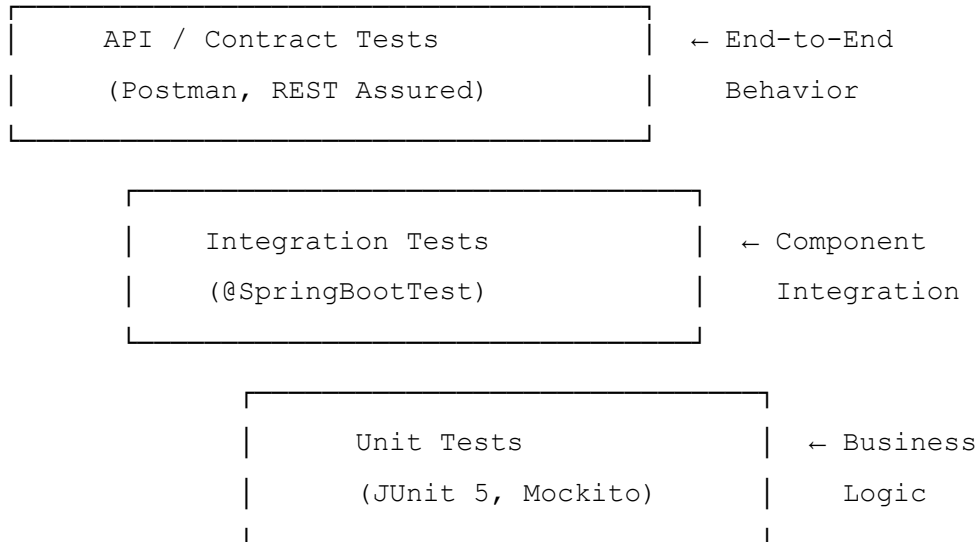
### **Mitigation:**

- Enforce mandatory database conditional confirmation:-

```
UPDATE seats SET status='BOOKED' WHERE id IN (1,2,3) AND status='HELD'
```

- Optionally implement a Time-To-Live (TTL) heartbeat mechanism to extend the lock duration for protracted payment processes.
- **The booking attempt must be rejected if the conditional database update fails**

# Testing Strategy and Quality Assurance



**Thank You**