

( Lectures 4 ) Divide-and-Conquer Algorithms: Polynomial Multiplication & FFT (Additional Slides)

**Given a Polynomial of Degree Bound 8**  
**Find 8 Distinct Points to Efficiently Evaluate it at**

$$A(x) = a_0 + a_1x + a_2x^2 + a_3x^3 + a_4x^4 + a_5x^5 + a_6x^6 + a_7x^7$$

$$\begin{aligned} A(x_0) &= a_0 + a_1x_0 + a_2(x_0)^2 + a_3(x_0)^3 + a_4(x_0)^4 + a_5(x_0)^5 + a_6(x_0)^6 + a_7(x_0)^7 \\ A(x_1) &= a_0 + a_1x_1 + a_2(x_1)^2 + a_3(x_1)^3 + a_4(x_1)^4 + a_5(x_1)^5 + a_6(x_1)^6 + a_7(x_1)^7 \\ A(x_2) &= a_0 + a_1x_2 + a_2(x_2)^2 + a_3(x_2)^3 + a_4(x_2)^4 + a_5(x_2)^5 + a_6(x_2)^6 + a_7(x_2)^7 \\ A(x_3) &= a_0 + a_1x_3 + a_2(x_3)^2 + a_3(x_3)^3 + a_4(x_3)^4 + a_5(x_3)^5 + a_6(x_3)^6 + a_7(x_3)^7 \end{aligned}$$

$x_4 = -x_0$	$A(x_4) = a_0 + a_1x_4 + a_2(x_4)^2 + a_3(x_4)^3 + a_4(x_4)^4 + a_5(x_4)^5 + a_6(x_4)^6 + a_7(x_4)^7$
$x_5 = -x_1$	$A(x_5) = a_0 + a_1x_5 + a_2(x_5)^2 + a_3(x_5)^3 + a_4(x_5)^4 + a_5(x_5)^5 + a_6(x_5)^6 + a_7(x_5)^7$
$x_6 = -x_2$	$A(x_6) = a_0 + a_1x_6 + a_2(x_6)^2 + a_3(x_6)^3 + a_4(x_6)^4 + a_5(x_6)^5 + a_6(x_6)^6 + a_7(x_6)^7$
$x_7 = -x_3$	$A(x_7) = a_0 + a_1x_7 + a_2(x_7)^2 + a_3(x_7)^3 + a_4(x_7)^4 + a_5(x_7)^5 + a_6(x_7)^6 + a_7(x_7)^7$

STRATEGY: Set  $x_{4+i} = -x_i$  for  $0 \leq i < 4$

**Given a Polynomial of Degree Bound 8**  
**Find 8 Distinct Points to Efficiently Evaluate it at**

$$A(x) = a_0 + a_1x + a_2x^2 + a_3x^3 + a_4x^4 + a_5x^5 + a_6x^6 + a_7x^7$$

$$\begin{aligned} A(x_0) &= a_0 + a_1x_0 + a_2(x_0)^2 + a_3(x_0)^3 + a_4(x_0)^4 + a_5(x_0)^5 + a_6(x_0)^6 + a_7(x_0)^7 \\ A(x_1) &= a_0 + a_1x_1 + a_2(x_1)^2 + a_3(x_1)^3 + a_4(x_1)^4 + a_5(x_1)^5 + a_6(x_1)^6 + a_7(x_1)^7 \\ A(x_2) &= a_0 + a_1x_2 + a_2(x_2)^2 + a_3(x_2)^3 + a_4(x_2)^4 + a_5(x_2)^5 + a_6(x_2)^6 + a_7(x_2)^7 \\ A(x_3) &= a_0 + a_1x_3 + a_2(x_3)^2 + a_3(x_3)^3 + a_4(x_3)^4 + a_5(x_3)^5 + a_6(x_3)^6 + a_7(x_3)^7 \end{aligned}$$

$x_4 = -x_0$	$A(-x_0) = a_0 - a_1x_0 + a_2(x_0)^2 - a_3(x_0)^3 + a_4(x_0)^4 - a_5(x_0)^5 + a_6(x_0)^6 - a_7(x_0)^7$
$x_5 = -x_1$	$A(-x_1) = a_0 - a_1x_1 + a_2(x_1)^2 - a_3(x_1)^3 + a_4(x_1)^4 - a_5(x_1)^5 + a_6(x_1)^6 - a_7(x_1)^7$
$x_6 = -x_2$	$A(-x_2) = a_0 - a_1x_2 + a_2(x_2)^2 - a_3(x_2)^3 + a_4(x_2)^4 - a_5(x_2)^5 + a_6(x_2)^6 - a_7(x_2)^7$
$x_7 = -x_3$	$A(-x_3) = a_0 - a_1x_3 + a_2(x_3)^2 - a_3(x_3)^3 + a_4(x_3)^4 - a_5(x_3)^5 + a_6(x_3)^6 - a_7(x_3)^7$

STRATEGY: Set  $x_{4+i} = -x_i$  for  $0 \leq i < 4$

**Given a Polynomial of Degree Bound 8**  
**Find 8 Distinct Points to Efficiently Evaluate it at**

$$A(x) = a_0 + a_1x + a_2x^2 + a_3x^3 + a_4x^4 + a_5x^5 + a_6x^6 + a_7x^7$$

$$\begin{aligned} A(x_0) &= a_0 + a_2(x_0)^2 + a_4(x_0)^4 + a_6(x_0)^6 + a_1x_0 + a_3(x_0)^3 + a_5(x_0)^5 + a_7(x_0)^7 \\ A(x_1) &= a_0 + a_2(x_1)^2 + a_4(x_1)^4 + a_6(x_1)^6 + a_1x_1 + a_3(x_1)^3 + a_5(x_1)^5 + a_7(x_1)^7 \\ A(x_2) &= a_0 + a_2(x_2)^2 + a_4(x_2)^4 + a_6(x_2)^6 + a_1x_2 + a_3(x_2)^3 + a_5(x_2)^5 + a_7(x_2)^7 \\ A(x_3) &= a_0 + a_2(x_3)^2 + a_4(x_3)^4 + a_6(x_3)^6 + a_1x_3 + a_3(x_3)^3 + a_5(x_3)^5 + a_7(x_3)^7 \end{aligned}$$

$x_4 = -x_0$	$A(-x_0) = a_0 + a_2(x_0)^2 + a_4(x_0)^4 + a_6(x_0)^6 - a_1x_0 - a_3(x_0)^3 - a_5(x_0)^5 - a_7(x_0)^7$
$x_5 = -x_1$	$A(-x_1) = a_0 + a_2(x_1)^2 + a_4(x_1)^4 + a_6(x_1)^6 - a_1x_1 - a_3(x_1)^3 - a_5(x_1)^5 - a_7(x_1)^7$
$x_6 = -x_2$	$A(-x_2) = a_0 + a_2(x_2)^2 + a_4(x_2)^4 + a_6(x_2)^6 - a_1x_2 - a_3(x_2)^3 - a_5(x_2)^5 - a_7(x_2)^7$
$x_7 = -x_3$	$A(-x_3) = a_0 + a_2(x_3)^2 + a_4(x_3)^4 + a_6(x_3)^6 - a_1x_3 - a_3(x_3)^3 - a_5(x_3)^5 - a_7(x_3)^7$

STRATEGY: Set  $x_{4+i} = -x_i$  for  $0 \leq i < 4$

**Given a Polynomial of Degree Bound 8**  
**Find 8 Distinct Points to Efficiently Evaluate it at**

$$A(x) = a_0 + a_1x + a_2x^2 + a_3x^3 + a_4x^4 + a_5x^5 + a_6x^6 + a_7x^7$$

$$\begin{aligned} A(x_0) &= (a_0 + a_2(x_0)^2 + a_4(x_0)^4 + a_6(x_0)^6) + x_0(a_1 + a_3(x_0)^2 + a_5(x_0)^4 + a_7(x_0)^6) \\ A(x_1) &= (a_0 + a_2(x_1)^2 + a_4(x_1)^4 + a_6(x_1)^6) + x_1(a_1 + a_3(x_1)^2 + a_5(x_1)^4 + a_7(x_1)^6) \\ A(x_2) &= (a_0 + a_2(x_2)^2 + a_4(x_2)^4 + a_6(x_2)^6) + x_2(a_1 + a_3(x_2)^2 + a_5(x_2)^4 + a_7(x_2)^6) \\ A(x_3) &= (a_0 + a_2(x_3)^2 + a_4(x_3)^4 + a_6(x_3)^6) + x_3(a_1 + a_3(x_3)^2 + a_5(x_3)^4 + a_7(x_3)^6) \end{aligned}$$

$x_4 = -x_0$	$A(-x_0) = (a_0 + a_2(x_0)^2 + a_4(x_0)^4 + a_6(x_0)^6) - x_0(a_1 + a_3(x_0)^2 + a_5(x_0)^4 + a_7(x_0)^6)$
$x_5 = -x_1$	$A(-x_1) = (a_0 + a_2(x_1)^2 + a_4(x_1)^4 + a_6(x_1)^6) - x_1(a_1 + a_3(x_1)^2 + a_5(x_1)^4 + a_7(x_1)^6)$
$x_6 = -x_2$	$A(-x_2) = (a_0 + a_2(x_2)^2 + a_4(x_2)^4 + a_6(x_2)^6) - x_2(a_1 + a_3(x_2)^2 + a_5(x_2)^4 + a_7(x_2)^6)$
$x_7 = -x_3$	$A(-x_3) = (a_0 + a_2(x_3)^2 + a_4(x_3)^4 + a_6(x_3)^6) - x_3(a_1 + a_3(x_3)^2 + a_5(x_3)^4 + a_7(x_3)^6)$

STRATEGY: Set  $x_{4+i} = -x_i$  for  $0 \leq i < 4$

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**Given a Polynomial of Degree Bound 8**  
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$$A(x) = a_0 + a_1x + a_2x^2 + a_3x^3 + a_4x^4 + a_5x^5 + a_6x^6 + a_7x^7$$

$$A_{\text{even}}(x) = a_0 + a_2x + a_4x^2 + a_6x^3$$

$$A_{\text{odd}}(x) = a_1 + a_3x + a_5x^2 + a_7x^3$$

$$\begin{aligned} A(x_0) &= (a_0 + a_2(x_0^2) + a_4(x_0^2)^2 + a_6(x_0^2)^3) + x_0(a_1 + a_3(x_0^2) + a_5(x_0^2)^2 + a_7(x_0^2)^3) \\ A(x_1) &= (a_0 + a_2(x_1^2) + a_4(x_1^2)^2 + a_6(x_1^2)^3) + x_1(a_1 + a_3(x_1^2) + a_5(x_1^2)^2 + a_7(x_1^2)^3) \\ A(x_2) &= (a_0 + a_2(x_2^2) + a_4(x_2^2)^2 + a_6(x_2^2)^3) + x_2(a_1 + a_3(x_2^2) + a_5(x_2^2)^2 + a_7(x_2^2)^3) \\ A(x_3) &= (a_0 + a_2(x_3^2) + a_4(x_3^2)^2 + a_6(x_3^2)^3) + x_3(a_1 + a_3(x_3^2) + a_5(x_3^2)^2 + a_7(x_3^2)^3) \end{aligned}$$

$$\begin{aligned} x_4 = -x_0 \quad A(-x_0) &= (a_0 + a_2(x_0^2) + a_4(x_0^2)^2 + a_6(x_0^2)^3) - x_0(a_1 + a_3(x_0^2) + a_5(x_0^2)^2 + a_7(x_0^2)^3) \\ x_5 = -x_1 \quad A(-x_1) &= (a_0 + a_2(x_1^2) + a_4(x_1^2)^2 + a_6(x_1^2)^3) - x_1(a_1 + a_3(x_1^2) + a_5(x_1^2)^2 + a_7(x_1^2)^3) \\ x_6 = -x_2 \quad A(-x_2) &= (a_0 + a_2(x_2^2) + a_4(x_2^2)^2 + a_6(x_2^2)^3) - x_2(a_1 + a_3(x_2^2) + a_5(x_2^2)^2 + a_7(x_2^2)^3) \\ x_7 = -x_3 \quad A(-x_3) &= (a_0 + a_2(x_3^2) + a_4(x_3^2)^2 + a_6(x_3^2)^3) - x_3(a_1 + a_3(x_3^2) + a_5(x_3^2)^2 + a_7(x_3^2)^3) \end{aligned}$$

STRATEGY: Set  $x_{4+i} = -x_i$  for  $0 \leq i < 4$

**Given a Polynomial of Degree Bound 8**  
**Find 8 Distinct Points to Efficiently Evaluate it at**

$$A(x) = a_0 + a_1x + a_2x^2 + a_3x^3 + a_4x^4 + a_5x^5 + a_6x^6 + a_7x^7$$

$$A_{\text{even}}(x) = a_0 + a_2x + a_4x^2 + a_6x^3$$

$$A_{\text{odd}}(x) = a_1 + a_3x + a_5x^2 + a_7x^3$$

$$\begin{aligned} A(x_0) &= A_{\text{even}}(x_0^2) + x_0 A_{\text{odd}}(x_0^2) \\ A(x_1) &= A_{\text{even}}(x_1^2) + x_1 A_{\text{odd}}(x_1^2) \\ A(x_2) &= A_{\text{even}}(x_2^2) + x_2 A_{\text{odd}}(x_2^2) \\ A(x_3) &= A_{\text{even}}(x_3^2) + x_3 A_{\text{odd}}(x_3^2) \end{aligned}$$

$$\begin{aligned} x_4 = -x_0 \quad A(-x_0) &= A_{\text{even}}(x_0^2) - x_0 A_{\text{odd}}(x_0^2) \\ x_5 = -x_1 \quad A(-x_1) &= A_{\text{even}}(x_1^2) - x_1 A_{\text{odd}}(x_1^2) \\ x_6 = -x_2 \quad A(-x_2) &= A_{\text{even}}(x_2^2) - x_2 A_{\text{odd}}(x_2^2) \\ x_7 = -x_3 \quad A(-x_3) &= A_{\text{even}}(x_3^2) - x_3 A_{\text{odd}}(x_3^2) \end{aligned}$$

STRATEGY: Set  $x_{4+i} = -x_i$  for  $0 \leq i < 4$

**Given a Polynomial of Degree Bound 8**  
**Find 8 Distinct Points to Efficiently Evaluate it at**

$$A(x) = a_0 + a_1x + a_2x^2 + a_3x^3 + a_4x^4 + a_5x^5 + a_6x^6 + a_7x^7$$

$$A_{\text{even}}(x) = a_0 + a_2x + a_4x^2 + a_6x^3$$

$$A_{\text{odd}}(x) = a_1 + a_3x + a_5x^2 + a_7x^3$$

$$\begin{aligned} A(x_0) &= A_{\text{even}}(x_0^2) + x_0 A_{\text{odd}}(x_0^2) \\ A(x_1) &= A_{\text{even}}(x_1^2) + x_1 A_{\text{odd}}(x_1^2) \\ A(x_2) &= A_{\text{even}}(x_2^2) + x_2 A_{\text{odd}}(x_2^2) \\ A(x_3) &= A_{\text{even}}(x_3^2) + x_3 A_{\text{odd}}(x_3^2) \\ A(-x_0) &= A_{\text{even}}(x_0^2) - x_0 A_{\text{odd}}(x_0^2) \\ A(-x_1) &= A_{\text{even}}(x_1^2) - x_1 A_{\text{odd}}(x_1^2) \\ A(-x_2) &= A_{\text{even}}(x_2^2) - x_2 A_{\text{odd}}(x_2^2) \\ A(-x_3) &= A_{\text{even}}(x_3^2) - x_3 A_{\text{odd}}(x_3^2) \end{aligned}$$

STRATEGY: Set  $x_{4+i} = -x_i$  for  $0 \leq i < 4$

We save roughly half the work.

**Given a Polynomial of Degree Bound 2**  
**Find 2 Distinct Points to Efficiently Evaluate it at**

$$A(x) = a_0 + a_1x$$

$$\begin{aligned} A(x_0) &= a_0 + a_1x_0 \\ A(x_1) &= a_0 + a_1x_1 \end{aligned}$$

STRATEGY: Set  $x_{1+i} = -x_i$  for  $0 \leq i < 1$

( Lectures 4 ) Divide-and-Conquer Algorithms: Polynomial Multiplication & FFT (Additional Slides)

**Given a Polynomial of Degree Bound 2**  
**Find 2 Distinct Points to Efficiently Evaluate it at**

$$A(x) = a_0 + a_1x$$

$$\begin{array}{lcl} A(x_0) & = & a_0 + a_1x_0 \\ x_1 = -x_0 & A(-x_0) & = a_0 - a_1x_0 \end{array}$$

STRATEGY: Set  $x_{1+i} = -x_i$  for  $0 \leq i < 1$



**Given a Polynomial of Degree Bound 2**  
**Find 2 Distinct Points to Efficiently Evaluate it at**

$$A(x) = a_0 + a_1x$$

$$\begin{array}{lcl} x_0 = 1 & A(x_0) & = a_0 + a_1 \\ x_1 = -1 & A(x_1) & = a_0 - a_1 \end{array}$$

STRATEGY: We will evaluate any polynomial of degree bound 2 at  
 $x_0 = 1$   
 $x_1 = -1$



**Given a Polynomial of Degree Bound 4**  
**Find 4 Distinct Points to Efficiently Evaluate it at**

$$A(x) = a_0 + a_1x + a_2x^2 + a_3x^3$$

$$\begin{array}{lcl} A(x_0) & = & a_0 + a_1x_0 + a_2(x_0)^2 + a_3(x_0)^3 \\ A(x_1) & = & a_0 + a_1x_1 + a_2(x_1)^2 + a_3(x_1)^3 \\ x_2 = -x_0 & A(x_2) & = a_0 + a_1x_2 + a_2(x_2)^2 + a_3(x_2)^3 \\ x_3 = -x_1 & A(x_3) & = a_0 + a_1x_3 + a_2(x_3)^2 + a_3(x_3)^3 \end{array}$$

STRATEGY: Set  $x_{2+i} = -x_i$  for  $0 \leq i < 2$



**Given a Polynomial of Degree Bound 4**  
**Find 4 Distinct Points to Efficiently Evaluate it at**

$$A(x) = a_0 + a_1x + a_2x^2 + a_3x^3$$

$$\begin{array}{lcl} A(x_0) & = & a_0 + a_1x_0 + a_2(x_0)^2 + a_3(x_0)^3 \\ A(x_1) & = & a_0 + a_1x_1 + a_2(x_1)^2 + a_3(x_1)^3 \\ x_2 = -x_0 & A(-x_0) & = a_0 - a_1x_0 + a_2(x_0)^2 - a_3(x_0)^3 \\ x_3 = -x_1 & A(-x_1) & = a_0 - a_1x_1 + a_2(x_1)^2 - a_3(x_1)^3 \end{array}$$

STRATEGY: Set  $x_{2+i} = -x_i$  for  $0 \leq i < 2$

( Lectures 4 ) Divide-and-Conquer Algorithms: Polynomial Multiplication & FFT (Additional Slides)

**Given a Polynomial of Degree Bound 4**  
**Find 4 Distinct Points to Efficiently Evaluate it at**

$$A(x) = a_0 + a_1x + a_2x^2 + a_3x^3$$

	$A(x_0)$	$=$	$a_0$	$+$	$a_2(x_0)^2$	$+$	$a_1x_0$	$+$	$a_3(x_0)^3$
	$A(x_1)$	$=$	$a_0$	$+$	$a_2(x_1)^2$	$+$	$a_1x_1$	$+$	$a_3(x_1)^3$
$x_2 = -x_0$	$A(-x_0)$	$=$	$a_0$	$+$	$a_2(x_0)^2$	$-$	$a_1x_0$	$-$	$a_3(x_0)^3$
$x_3 = -x_1$	$A(-x_1)$	$=$	$a_0$	$+$	$a_2(x_1)^2$	$-$	$a_1x_1$	$-$	$a_3(x_1)^3$

STRATEGY: Set  $x_{2+i} = -x_i$  for  $0 \leq i < 2$



**Given a Polynomial of Degree Bound 4**  
**Find 4 Distinct Points to Efficiently Evaluate it at**

$$A(x) = a_0 + a_1x + a_2x^2 + a_3x^3$$

	$A(x_0)$	$=$	$(a_0 + a_2(x_0)^2)$	$+$	$x_0(a_1 + a_3(x_0)^2)$
	$A(x_1)$	$=$	$(a_0 + a_2(x_1)^2)$	$+$	$x_1(a_1 + a_3(x_1)^2)$
$x_2 = -x_0$	$A(-x_0)$	$=$	$(a_0 + a_2(x_0)^2)$	$-$	$x_0(a_1 + a_3(x_0)^2)$
$x_3 = -x_1$	$A(-x_1)$	$=$	$(a_0 + a_2(x_1)^2)$	$-$	$x_1(a_1 + a_3(x_1)^2)$

STRATEGY: Set  $x_{2+i} = -x_i$  for  $0 \leq i < 2$



**Given a Polynomial of Degree Bound 4**  
**Find 4 Distinct Points to Efficiently Evaluate it at**

$$A(x) = a_0 + a_1x + a_2x^2 + a_3x^3$$

$$A_{\text{even}}(x) = a_0 + a_2x^2$$

$$A_{\text{odd}}(x) = a_1 + a_3x$$

	$A(x_0)$	$=$	$(a_0 + a_2(x_0)^2)$	$+$	$x_0(a_1 + a_3(x_0)^2)$
	$A(x_1)$	$=$	$(a_0 + a_2(x_1)^2)$	$+$	$x_1(a_1 + a_3(x_1)^2)$
$x_2 = -x_0$	$A(-x_0)$	$=$	$(a_0 + a_2(x_0)^2)$	$-$	$x_0(a_1 + a_3(x_0)^2)$
$x_3 = -x_1$	$A(-x_1)$	$=$	$(a_0 + a_2(x_1)^2)$	$-$	$x_1(a_1 + a_3(x_1)^2)$

STRATEGY: Set  $x_{2+i} = -x_i$  for  $0 \leq i < 2$



**Given a Polynomial of Degree Bound 4**  
**Find 4 Distinct Points to Efficiently Evaluate it at**

$$A(x) = a_0 + a_1x + a_2x^2 + a_3x^3$$

$$A_{\text{even}}(x) = a_0 + a_2x^2$$

$$A_{\text{odd}}(x) = a_1 + a_3x$$

	$A(x_0)$	$=$	$A_{\text{even}}(x_0^2)$	$+$	$x_0 A_{\text{odd}}(x_0^2)$
	$A(x_1)$	$=$	$A_{\text{even}}(x_1^2)$	$+$	$x_1 A_{\text{odd}}(x_1^2)$
$x_2 = -x_0$	$A(-x_0)$	$=$	$A_{\text{even}}(x_0^2)$	$-$	$x_0 A_{\text{odd}}(x_0^2)$
$x_3 = -x_1$	$A(-x_1)$	$=$	$A_{\text{even}}(x_1^2)$	$-$	$x_1 A_{\text{odd}}(x_1^2)$

STRATEGY: Set  $x_{2+i} = -x_i$  for  $0 \leq i < 2$

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$$A(x) = a_0 + a_1x + a_2x^2 + a_3x^3$$

$$A_{\text{even}}(x) = a_0 + a_2x$$
$$A_{\text{odd}}(x) = a_1 + a_3x$$

$A(x_0)$	=	$A_{\text{even}}(x_0^2)$	+	$x_0 A_{\text{odd}}(x_0^2)$	
$A(x_1)$	=	$A_{\text{even}}(x_1^2)$	+	$x_1 A_{\text{odd}}(x_1^2)$	
$x_2 = -x_0$	$A(-x_0)$	=	$A_{\text{even}}(x_0^2)$	-	$x_0 A_{\text{odd}}(x_0^2)$
$x_3 = -x_1$	$A(-x_1)$	=	$A_{\text{even}}(x_1^2)$	-	$x_1 A_{\text{odd}}(x_1^2)$

STRATEGY: Set  $x_{2+i} = -x_i$  for  $0 \leq i < 2$

**Given a Polynomial of Degree Bound 4**  
**Find 4 Distinct Points to Efficiently Evaluate it at**

$$A(x) = a_0 + a_1x + a_2x^2 + a_3x^3$$

$$A_{\text{even}}(x) = a_0 + a_2x$$
$$A_{\text{odd}}(x) = a_1 + a_3x$$

$A(x_0)$	=	$A_{\text{even}}(x_0^2)$	+	$x_0 A_{\text{odd}}(x_0^2)$	
$A(x_1)$	=	$A_{\text{even}}(x_1^2)$	+	$x_1 A_{\text{odd}}(x_1^2)$	
$x_2 = -x_0$	$A(-x_0)$	=	$A_{\text{even}}(x_0^2)$	-	$x_0 A_{\text{odd}}(x_0^2)$
$x_3 = -x_1$	$A(-x_1)$	=	$A_{\text{even}}(x_1^2)$	-	$x_1 A_{\text{odd}}(x_1^2)$

Observe that we evaluate both  $A_{\text{even}}(x)$  and  $A_{\text{odd}}(x)$  at  $x = x_0^2$  and  $x = x_1^2$ .  
But we decided to always evaluate polynomials of degree bound 2 at  $x = 1$  and  $x = -1$ .  
So,  $x_0^2 = 1 \Rightarrow x_0 = 1$  and  $x_1^2 = -1 \Rightarrow x_1 = \sqrt{-1} = i$ .

**Given a Polynomial of Degree Bound 4**  
**Find 4 Distinct Points to Efficiently Evaluate it at**

$$A(x) = a_0 + a_1x + a_2x^2 + a_3x^3$$

$$A_{\text{even}}(x) = a_0 + a_2x$$
$$A_{\text{odd}}(x) = a_1 + a_3x$$

$A(x_0)$	=	$A_{\text{even}}(x_0^2)$	+	$x_0 A_{\text{odd}}(x_0^2)$	
$A(x_1)$	=	$A_{\text{even}}(x_1^2)$	+	$x_1 A_{\text{odd}}(x_1^2)$	
$x_2 = -x_0$	$A(-x_0)$	=	$A_{\text{even}}(x_0^2)$	-	$x_0 A_{\text{odd}}(x_0^2)$
$x_3 = -x_1$	$A(-x_1)$	=	$A_{\text{even}}(x_1^2)$	-	$x_1 A_{\text{odd}}(x_1^2)$

So, we evaluate any polynomial of degree bound 4 at  
 $x_0 = 1, x_1 = i$   
and  
 $x_2 = -x_0 = -1, x_3 = -x_1 = -i$

**Given a Polynomial of Degree Bound 8**  
**Find 8 Distinct Points to Efficiently Evaluate it at**

$$A(x) = a_0 + a_1x + a_2x^2 + a_3x^3 + a_4x^4 + a_5x^5 + a_6x^6 + a_7x^7$$

$$A_{\text{even}}(x) = a_0 + a_2x + a_4x^2 + a_6x^3$$
$$A_{\text{odd}}(x) = a_1 + a_3x + a_5x^2 + a_7x^3$$

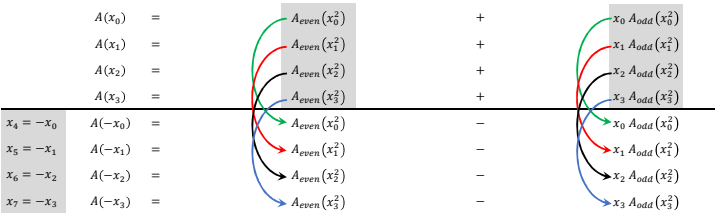
$A(x_0)$	=	$A_{\text{even}}(x_0^2)$	+	$x_0 A_{\text{odd}}(x_0^2)$	
$A(x_1)$	=	$A_{\text{even}}(x_1^2)$	+	$x_1 A_{\text{odd}}(x_1^2)$	
$A(x_2)$	=	$A_{\text{even}}(x_2^2)$	+	$x_2 A_{\text{odd}}(x_2^2)$	
$A(x_3)$	=	$A_{\text{even}}(x_3^2)$	+	$x_3 A_{\text{odd}}(x_3^2)$	
$x_4 = -x_0$	$A(-x_0)$	=	$A_{\text{even}}(x_0^2)$	-	$x_0 A_{\text{odd}}(x_0^2)$
$x_5 = -x_1$	$A(-x_1)$	=	$A_{\text{even}}(x_1^2)$	-	$x_1 A_{\text{odd}}(x_1^2)$
$x_6 = -x_2$	$A(-x_2)$	=	$A_{\text{even}}(x_2^2)$	-	$x_2 A_{\text{odd}}(x_2^2)$
$x_7 = -x_3$	$A(-x_3)$	=	$A_{\text{even}}(x_3^2)$	-	$x_3 A_{\text{odd}}(x_3^2)$

Observe that we evaluate both  $A_{\text{even}}(x)$  and  $A_{\text{odd}}(x)$  at  $x = x_0^2, x = x_1^2, x = x_2^2$  and  $x = x_3^2$ .  
But we decided to always evaluate polynomials of degree bound 4 at  $x = 1, x = i, x = -1$  and  $x = -i$ .  
So,  $x_0^2 = 1 \Rightarrow x_0 = 1, x_1^2 = i \Rightarrow x_1 = \frac{1+i}{\sqrt{2}}, x_2^2 = -1 \Rightarrow x_2 = i$ , and  $x_3^2 = -i \Rightarrow x_3 = \frac{-1+i}{\sqrt{2}}$ .

( Lectures 4 ) Divide-and-Conquer Algorithms: Polynomial Multiplication & FFT (Additional Slides)

Given a Polynomial of Degree Bound 8  
Find 8 Distinct Points to Efficiently Evaluate it at

$A(x) = a_0 + a_1x + a_2x^2 + a_3x^3 + a_4x^4 + a_5x^5 + a_6x^6 + a_7x^7$   
 $A_{\text{even}}(x) = a_0 + a_2x + a_4x^2 + a_6x^3$   
 $A_{\text{odd}}(x) = a_1 + a_3x + a_5x^2 + a_7x^3$



So, we evaluate any polynomial of degree bound 8 at  
 $x_0 = 1, x_1 = \frac{1+i}{\sqrt{2}}, x_2 = i, x_3 = \frac{-1+i}{\sqrt{2}}$   
and  
 $x_4 = -x_0 = -1, x_5 = -x_1 = -\frac{1+i}{\sqrt{2}}, x_6 = -x_2 = -i, x_7 = -x_3 = -\frac{-1+i}{\sqrt{2}}$

Given a Polynomial of Degree Bound  $n = 2^k$   
Find  $n = 2^k$  Distinct Points to Efficiently Evaluate it at

degree bound	how did we find the points to evaluate the polynomial at?	the points	point property
$2^1$	...	1, -1	all $2^{\text{nd}}$ roots of unity
$2^2$	take positive and negative square roots of points used for degree bound $2^1$ which are already the $2^{\text{nd}}$ roots of unity	1, i, -1, -i	all $4^{\text{th}}$ roots of unity
$2^3$	take positive and negative square roots of points used for degree bound $2^2$ which are already the $4^{\text{th}}$ roots of unity	1, $\frac{1+i}{\sqrt{2}}$ , i, $\frac{-1+i}{\sqrt{2}}$ , -1, $\frac{1+i}{\sqrt{2}}$ , -i, $\frac{-1+i}{\sqrt{2}}$	all $8^{\text{th}}$ roots of unity
$2^4$	take positive and negative square roots of points used for degree bound $2^3$ which are already the $8^{\text{th}}$ roots of unity	1, $\frac{\sqrt{2}+\sqrt{2}}{2} + i\frac{\sqrt{2}-\sqrt{2}}{2}$ , ..., -1, $-\frac{\sqrt{2}+\sqrt{2}}{2} + i\frac{\sqrt{2}-\sqrt{2}}{2}$	all $16^{\text{th}}$ roots of unity
...	...	...	...
$2^{k-1}$	take positive and negative square roots of points used for degree bound $2^{k-2}$ which are already the $2^{k-2}$ th roots of unity	...	all $2^{k-1}$ th roots of unity
$n = 2^k$	take positive and negative square roots of points used for degree bound $2^{k-1}$ which are already the $2^{k-1}$ th roots of unity	...	all $2^k$ th roots of unity (i.e., $n^{\text{th}}$ roots of unity)

How to Find all  $n^{\text{th}}$  Roots of Unity

The  $n^{\text{th}}$  roots of unity are:  $1, \omega_n, (\omega_n)^2, (\omega_n)^3, \dots, (\omega_n)^{n-1}$ ,  
where  $\omega_n = \cos \frac{2\pi}{n} + i \sin \frac{2\pi}{n} = e^{\frac{2\pi i}{n}}$  is known as the primitive  $n^{\text{th}}$  roots of unity.  
The result above can be derived using Euler's Formula.

**Euler's Formula:** For any real number  $\alpha$ ,  $\cos \alpha + i \sin \alpha = e^{i\alpha}$

Euler's formula follows very easily from the following three power series each of which holds for  $-\infty < \alpha < +\infty$ :

$\cos \alpha = 1 - \frac{\alpha^2}{2!} + \frac{\alpha^4}{4!} - \frac{\alpha^6}{6!} + \frac{\alpha^8}{8!} - \dots$   
 $\sin \alpha = \alpha - \frac{\alpha^3}{3!} + \frac{\alpha^5}{5!} - \frac{\alpha^7}{7!} + \frac{\alpha^9}{9!} - \dots$   
 $e^\alpha = 1 + \alpha + \frac{\alpha^2}{2!} + \frac{\alpha^3}{3!} + \frac{\alpha^4}{4!} + \frac{\alpha^5}{5!} + \frac{\alpha^6}{6!} + \frac{\alpha^7}{7!} + \frac{\alpha^8}{8!} + \dots$

How to Find all  $n^{\text{th}}$  Roots of Unity

Observe that for (any) real numbers  $\alpha$  and  $p$ ,  
 $(\cos \alpha + i \sin \alpha)^p = (e^{i\alpha})^p = e^{i(p\alpha)} = \cos(p\alpha) + i \sin(p\alpha)$   
Also observe that for any integer  $k$ ,  $\cos(k \times 2\pi) + i \sin(k \times 2\pi) = 1 + i \times 0 = 1$   
Then the  $n^{\text{th}}$  root of 1 (unity) is  $1^{\frac{1}{n}} = (\cos(k \times 2\pi) + i \sin(k \times 2\pi))^{\frac{1}{n}} = \cos\left(k \times \frac{2\pi}{n}\right) + i \sin\left(k \times \frac{2\pi}{n}\right)$   
Observe that  $\cos\left(k \times \frac{2\pi}{n}\right) + i \sin\left(k \times \frac{2\pi}{n}\right)$  takes  $n$  distinct values for  $0 \leq k < n$ , and then simply repeats those values for  $k < 0$  and  $k \geq n$ .  
When  $k = 1$ , we have  $\cos\left(k \times \frac{2\pi}{n}\right) + i \sin\left(k \times \frac{2\pi}{n}\right) = \cos\left(\frac{2\pi}{n}\right) + i \sin\left(\frac{2\pi}{n}\right) = \omega_n$  = primitive  $n^{\text{th}}$  root of 1.  
Clearly, for any  $k$ ,  $\cos\left(k \times \frac{2\pi}{n}\right) + i \sin\left(k \times \frac{2\pi}{n}\right) = \left(\cos\left(\frac{2\pi}{n}\right) + i \sin\left(\frac{2\pi}{n}\right)\right)^k = (\omega_n)^k$   
Hence,  $1^{\frac{1}{n}} = \cos\left(k \times \frac{2\pi}{n}\right) + i \sin\left(k \times \frac{2\pi}{n}\right) = (\omega_n)^k$ , for  $k = 0, 1, 2, \dots, n-1$ .  
In other words, the  $n^{\text{th}}$  roots of 1 (unity) are:  $1, \omega_n, (\omega_n)^2, (\omega_n)^3, \dots, (\omega_n)^{n-1}$

Coefficient Form  $\Rightarrow$  Point-Value Form

```
Rec-FFT ( ( a0, a1, ..., an-1 ) )   { n = 2k for integer k ≥ 0 }
1.  if n = 1 then
2.    return ( a0 )
3.  ωn ← e2πi/n
4.  ω ← 1
5.  yeven ← Rec-FFT ( ( a0, a2, ..., an-2 ) )
6.  yodd  ← Rec-FFT ( ( a1, a3, ..., an-1 ) )
7.  for j ← 0 to n/2 - 1 do
8.    yj ← yeven + ω yodd
9.    yn/2+j ← yeven - ω yodd
10.   ω ← ω ωn
11. return y
```

Running time:

$$T(n) = \begin{cases} \Theta(1), & \text{if } n = 1, \\ 2T\left(\frac{n}{2}\right) + \Theta(n), & \text{otherwise.} \end{cases}$$
$$= \Theta(n \log n)$$

