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	• [Ste97] W. Stevens "TCP congestion control"			
	• [Ste97] W. Stevens, "TCP congestion control", RFC 2001, Jan 1997.			
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	Overview			
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	> Layering			
	End-to-End Arguments			
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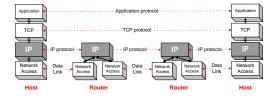
What is Layering?

 A technique to organize a network system into a succession of logically distinct entities, such that the service provided by one entity is solely based on the service provided by the previous (lower level) entity

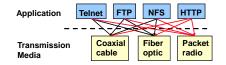
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Network Protocols

- Specify any function that requires cooperation between two or more network entities
 - specify the format of the information that is sent/received among routers and end-systems
 - specify timings and the actions that a node has to take when it receives special messages or special events occur



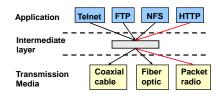
Why Layering?



 No layering: each new application has to be reimplemented for every network technology!

Why Layering?

Solution: introduce an intermediate layer that provides a unique abstraction for various network technologies



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Layering

- Advantages
 - Modularity protocols easier to manage and maintain
 - Abstract functionality –lower layers can be changed without affecting the upper layers
 - Reuse upper layers can reuse the functionality provided by lower layers
- Disadvantages
 - Information hiding inefficient implementations

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ISO OSI Reference Model

- ISO International Standard Organization
- OSI Open System Interconnection
- Started in 1978; first standard 1979
 - ARPANET started in 1969; TCP/IP protocols ready by 1974
- Goal: a general open standard
 - allow vendors to enter the market by using their own implementation

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ISO OSI Reference Model Seven layers - Lower three layers are peer-to-peer - Next four layers are end-to-end Application Application Presentation Presentation Session Session Transport Transport Network Network Network Datalink Datalink Physical Physical Physical Physical medium

Data Transmission • A layer can use only the service provided by the layer immediate below it • Each layer may change and add a header to data packet data data

OSI Model Concepts Service – says what a layer does

- Interface says how to access the service
- Protocol says how is the service implemented
- a set of rules and formats that govern the communication between two peers

Physical Layer (1)

- Service: move the information between two systems connected by a physical link
- Interface: specifies how to send a bit
- Protocol: coding scheme used to represent a bit, voltage levels, duration of a bit
- Examples: coaxial cable, optical fiber links; transmitters, receivers

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Encoding Schemes NRZ (non-return to zero) NRZI (non-return to zero intermediate) 4B/5B encode each 4 bits in a 5-bit code such that there is at most one leading 0 and at most two trailing zero and use NRZI encoding

Datalink Layer (2)

- Service:
 - framing, i.e., attach frames separator
 - send data frames between peers
 - others:
 - arbitrate the access to common physical media
 - ensure reliable transmission
 - provide flow control
- Interface: send a data unit (packet) to a machine connected to same physical media
- Protocol: layer addresses, implement Medium Access Control (MAC) (e.g., CSMA/CD)...

Marking the Beginning/Ending of a Frame

- Use a special bit sequence
- Problem: what happens if this sequence appears in the data payload?
- Use bit stuffing technique; e.g., assume both encoding, decoding bit sequences are 01111110
 - sender: inserts a 0 after five consecutive 1s
 - receiver: when it sees five 1s makes decision on next two bits
 - if next bit 0 (this is a stuffed bit), remove it
 - if next bit 1, look at the next bit
 - o if 0 this is end-of-frame (receiver has seen 01111110) o if 1 this is an error, discard the frame (receiver has seen 01111111)

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Network Layer (3)

- Service:
 - deliver a packet to specified destination
 - perform segmentation/reassembly
 - others:
 - packet scheduling
 - buffer management
- Interface: send a packet to a specified destination
- Protocol: define global unique addresses; construct routing tables

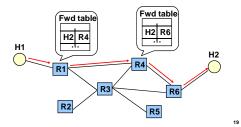
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Data and Control Planes

- Data plane: concerned with
 - packet forwarding
 - buffer management
 - packet scheduling
- Control Plane: concerned with installing and maintaining state for data plane

Example: Routing

- Data plane: use Forwarding Table to forward packets
- Control plane: construct and maintain Forwarding Tables (e.g., Distance Vector, Link State protocols)



Transport Layer (4)

- Service:
 - provide an error-free and flow-controlled end-to-end connection
 - multiplex multiple transport connections to one network connection
 - split one transport connection in multiple network connections
- Interface: send a packet to specify destination
- Protocol: implement reliability and flow control
- Examples: TCP and UDP

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Session Layer (5)

- Service:
 - full-duplex
 - access management, e.g., token control
 - synchronization, e.g., provide check points for long transfers
- Interface: depends on service
- Protocols: token management; insert checkpoints, implement roll-back functions

Presentation Layer (6)

- Service: convert data between various representations
- Interface: depends on service
- Protocol: define data formats, and rules to convert from one format to another

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Application Layer (7)

- Service: any service provided to the end user
- Interface: depends on the application
- Protocol: depends on the application
- Examples: FTP, Telnet, WWW browser

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OSI vs. TCP/IP

- OSI: conceptually define: service, interface, protocol
- Internet: provide a successful implementation

Application
Presentation
Session
Transport
Network
Datalink
Physical

Application	Telnet FTP DNS
Transport	TCP UDP
Internet	IP
Host-to- network	LAN Packet radio

Key Design Decision

How do you divide functionality across the layers?

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Overview

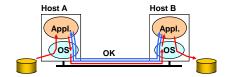
- Layering
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End-to-End Argument

- Think twice before implementing a functionality that you believe that is useful to an application at a lower layer
- If the application can implement a functionality correctly, implement it at a lower layer only as a performance enhancement

Example: Reliable File Transfer



- Solution 1: make each step reliable, and then concatenate them
- Solution 2: end-to-end check and retry

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Discussion

- Solution 1 not complete
 - What happens if the sender or/and receiver misbehave?
- The receiver has to do the check anyway!
- Thus, full functionality can be entirely implemented at application layer; no need for reliability from lower layers
- Is there any need to implement reliability at lower layers?

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Discussion

- Yes, but only to improve performance
- Example:
 - assume a high error rate on communication network
 - then, a reliable communication service at data link layer might help

Trade-offs

- Application has more information about the data and the semantic of the service it requires (e.g., can check only at the end of each data unit)
- A lower layer has more information about constraints in data transmission (e.g., packet size, error rate)
- Note: these trade-offs are a direct result of layering!

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Rule of Thumb

 Implementing a functionality at a lower level should have minimum performance impact on the application that do not use the functionality

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Other Examples

- Secure transmission of data
- Duplicate message suppression
- Transaction management
- RISC vs. CISC

Overview

- Layering
- End-to-End Arguments
- > A Case Study: the Internet

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Internet & End-to-End Argument

- Network layer provides one simple service: best effort datagram (packet) delivery
- Only one higher level service implemented at transport layer: reliable data delivery (TCP)
 - performance enhancement; used by a large variety of applications (Telnet, FTP, HTTP)
 - does not impact other applications (can use UDP)
- Everything else implemented at application level

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Internet Architecture

- Packet-switched datagram network
- IP is the glue
- · Hourglass architecture
 - all hosts and routers run IP
- Stateless architecture
 - no per flow state inside network



Key Advantages

- The service can be implemented by a large variety of network technologies
- Does not require routers to maintain any fine grained state about traffic. Thus, network architecture is
 - Robust
 - Scalable
 - Stateless

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What About Other Services?

- Multicast?
 - Multicast group management at routers?
 - End host based multicast?
- Quality of Service (QoS)?
 - Per flow state: resource reservation, classification, buffer management, scheduling
 - Stateless quality of service?

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Summary: Layering

- Key technique to implement communication protocols; provides
 - Modularity
 - Abstraction
 - Reuse
- Key design decision: what functionality to put in each layer?

Summary: End-to-End Arguments

- If the application can do it, don't do it at a lower laver
 - the application knows the best what it needs anyway
 - add functionality in lower layers iff it is
 - (1) used and improves performance of a large number of applications, and
 - (2) does not hurt other applications
- Success story: Internet

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Summary

 Challenge of building a good (network) system: find the right balance between:

Reuse, implementation effort (apply layering concepts)



No universal answer: the answer depends on the goals and assumptions!