

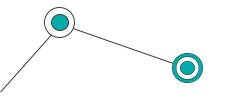




# Enhancing Non-Repudiation Protocols with Knowledge Graphs

STSM Research Funded by DKG Cost Action (CA19134)

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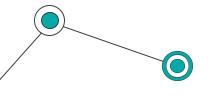




Let's imagine a dispute between Organization A and Organization B. Org. A states that Org. B accessed sensitive content at a given time, while Org. B refuses this statement.







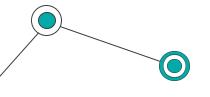


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Who's right? How this can be demonstrated?









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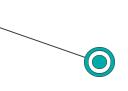
Who's right? How this can be demonstrated?





Who is the judge?



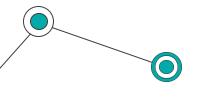




**Non-repudiation** is a security principle that ensures a party in a communication cannot deny the authenticity of their actions, such as sending a message or completing a transaction.

**Non-repudiation** in systems has always been a requirement for systems requesting a high level of security and auditability. These properties are particularly important for sensitive scenarios such as use cases 2.3 (legal data), 4.5 (healthcare), etc.

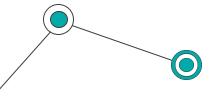
However, current non-repudiation protocols face critical challenges: there is no widely accepted standard for representing non-repudiation, particularly in **decentralized** environments, where two organizations communicate without having a **trust authority**.



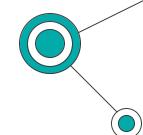
# DKGs for Non-Repudiation Benefits



- A standardized way for saving logs ensures consistent, **interoperable**, and tamper-proof records that can be easily verified and audited across different systems, independently from representation used (RDF, XML, JSON-LD, etc.)
- Signature applied over the graphs ensure **data integrity**, verifiability, and tamper-proof. Also in this case the signature is applied over the semantics rather than the single representation.
  - DKGs enables machine-to-machine (M2M) communication as well as user friendly visualization of relationships, by enabling seamless interaction and understanding of complex data connections, making the users more aware of actions performed.



# DKGs for Non-Repudiation Challenges



Application of DKGs for guaranteeing non-repudiation introduces huge advantages, anyway the implementation poses three main challenges.



Data Canonization and Signature

Private Key management

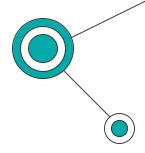




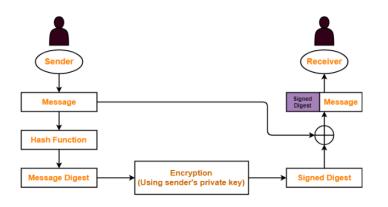
Decentralized Communication



#### **Data Canonization and Signature**



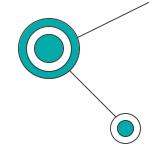
Typical message signature involves the hashing of the message and signature applied over the resulting digest. When translating this to KG we need to take into account canonization of data.



In general the operation uses the RDF
Dataset Normalization Algorithm transform
an input document into its canonical form.
The cannonical representation is then
hashed and signed with a detached
signature algorithm.



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#### Canonized KG



#### **Private Key Management**



Non-Repudiation algorithms rely Public Key Infrastructure (PKI) to guarantee signatures and checks over them. A secure storage is needed to preserve malicious usage of these keys.

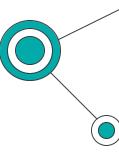
Aries Askar is a secure storage and a key management service suitable for use with <u>Hyperledger Aries</u> agents and possibly other digital trust agents



It leverages biometric authentication to protect private keys and offers all the crypto-suite needed to sign and encrypt data.



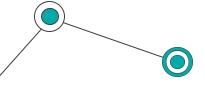
#### **Decentralized Communication**



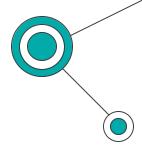
Decentralized communication is much more critical without a centralized authority responsible for manage certificates and create trust among partecipants.

DIDComm offers a strong security againist attackers, and it is also able to create a narrowed channel between communication parties.

It is based on Elliptic Curve Diffie Hellman (ECDH) exchange of XChaCha20Poly1305 key used for the encryption.

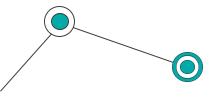


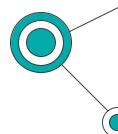
#### **Decentralized Communication**



It is based on Elliptic Curve Diffie Hellman (ECDH) exchange of XChaCha20Poly1305 key used for the encryption.

The public key, which are managed by users individually, are published in a DIDDocument under the control of the user itself.





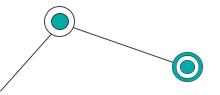
A concrete application of the proposed protocol has been investigated during STSM period. Such a protocol has been applied to decentralized authentication, fully leveraging decentralized identifiers and KGs.

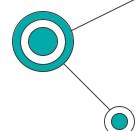
Parties are identified by their respective DIDs, each of them mapped into a DIDDocument.

#### Resolution

| $DID_U = did:web:user \mapsto (g^u, pk(sk_U))$     |
|----------------------------------------------------|
| $DID_A = did:web:app \mapsto (g^a, pk(sk_A))$      |
| $DID_I = did:web:issuer \mapsto (g^i, pk(sk_I))$   |
| $DID_V = did:web:verifier \mapsto (g^v, pk(sk_V))$ |

Table 1: Mapping of DIDs using Web method resolution.





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#### Resolution

```
DID_U = did:web:user \mapsto (g^u, pk(sk_U))

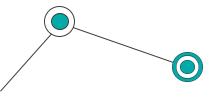
DID_A = did:web:app \mapsto (g^a, pk(sk_A))

DID_I = did:web:issuer \mapsto (g^i, pk(sk_I))

DID_V = did:web:verifier \mapsto (g^v, pk(sk_V))
```

Table 1: Mapping of DIDs using Web method resolution.

```
**Special Control of the Control of
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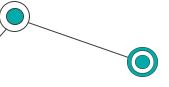


The first step of the proposed protocol is the authentication request from the Org. 1 (taget) to the user (agent), which is intermediated by the Org. 2 (client).

```
a, sk_A, g^u, pk_U, g^i, pk_I, g^v, pk_V
                                                                    v, sk_V, g^u, pk_U, g^a, pk_A, g^i, pk_I
                     App
                                                                                     CSS
new sa. sav: SecretKev
                                                               new sc: SecretKev
new k0, k2, k3, k10: SymmetricKey
                                                               new iv_1, iv_4, iv_{11}, n_c: nonce
                                                               new k1, k4, k11: SymmetricKey
new iv0, iv2, iv3, iv10: nonce
                                                               RULE := (DID_{II}, DID_{I}, URI, DID_{A})
  m_1 := sig(DIDDoc(g^{sac}), sk_A)
                              \{k_0\}_{(g^v)^{sav}}, \{g^{sav}\}_{pk(sk_v)}, \{m_1\}_{F(iv_0,k_0)}, iv_0
                                                                    if check(m_1, s_1, pk(sk_A)) then
                                                                   m_2 := sig(DIDDoc(g^{sc}), sk_V)
                             \{k_1\}_{(g^{sar})^{\infty}}, \{g^{sc}\}_{pk(sk_{+})}, \{m_2\}_{F(iv_1,k_1)}, iv_1
if check(m_2, s_2, pk(sk_V)) then
m_3 := ack
m_4 := (DID_U, DID_I, URI, DID_A)
                             \{k_2\}_{(g^{sc})^{tav}}, \{g^{sav}\}_{pk(sk_V)}, \{m_3\}_{F(iv_2,k_2)}, iv_2
                             \{k_3\}_{(g^{sc})^{sav}}, \{g^{sav}\}_{pk(sk_V)}, \{m_4\}_{F(iv_1,k_1)}, iv_3
                                                            if (DID'_{U}, DID'_{I}, URI', DID'_{A}) == RULE then
                            \{k_4\}_{(g^{tatv})^{sc}}, \{g^{sc}\}_{pk(sk_A)}
      if RULE' == m_4 then
      VPR^2 = (VPR, scope)
       m_5 := sig(VPR^2, sk_A)
```

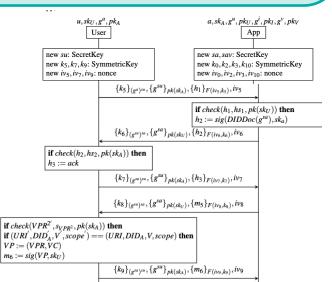
```
"presentation_definition": {
  "id": "32f54163-7166-48f1-93d8-ff217bdb0653",
  "input_descriptors": [
                        "$.credentialSubject.degree",
"$.credentialSubject.claims.degree"
   "challenge": "zFWfzGWq8UtAuZOiwKwV8g==",
```

Note: The VPR, which is the request of credentials, is a Knowledge Grarph.



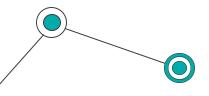


Before to exchange the VPR a connection is established. The VPR is then signed by the client and forwarded to the agent.



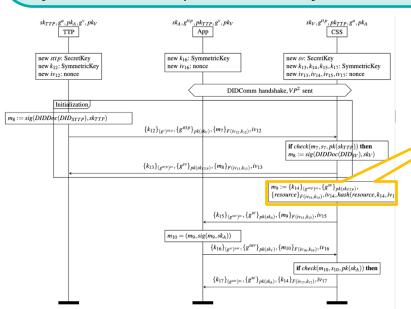
I received the following VPR (Connection Record: 8987fdfa-7203-4a56-928b-3d8362a1cff5 with DID did:web:secureapp.solidcommunity.net:public and label SecureS olidpppppppppp):
ACP Policy:

For the target: http://localhost:3000/my-pod/test-folder/test-resource.txt with the client (app): did:web:secureapp.solidcommunity.net:public



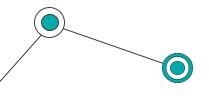


The resource is delivered in an encrypted and signed way. This operation is needed to guarantee non-repudiation of origin (NRO).



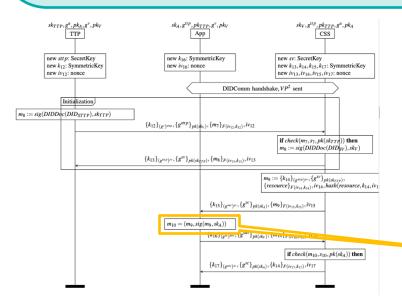
Non-repudiation of origin (NRO)

{"protected":"eyJlbmMi0iJ4Y2hhY2hhMjBwb2x5MTMwNV9pZXRmIiwidHlwIjoiSldMLzEuMCIsImFsZyI6IkFub25jc
JDMHp50WIweW1UUzV3aFdoMEVVT0tBbE9yQkdZYi1XdM5qd1NpNVZhSkp6WmlfRVcyNHJDLXIyNUInSlJHS0tqTnNtTEh5
Z2N1pMNnZ0R0d1d2h0cHNnaGJBckdQenR0c25DM2JDb3Y3QINIcTRpMXc1SyJ9fV19", "ciphertext":"j9TWK01JyQaSK
76Qm9XekJt6QlCNlDfZtSrNUFu43nHCD851LwzgeQx1Ei-i5qx2bB3BZ6tEPdSjpQ1rfcKzaqT93J19MbRzdBxmgeYcQaCP
GTF0kdpz0Zm\_bjqEc4CYmeerkZT2sq6-UJ01QWzB4ImT5qXuohS7oQQcml6alf9e5Kbz3xGF2GnocHq32qj56iynPmZsmbC
0sfjHL\_gvbfAdy39SWg\_hBClIIpw0Sbo2Q0KR7d4oQ300-pQpFkJ50YfdPKetskudqnDkX1iIsV0HrdjTipbsCLoN2P03NR
P6koy93Bt19GFFHqIRUGxT8\_tpw72fcDd4WBqGWJrAX9pWBs0ZAjtylyjlQut6vDrolSPokHfdgj5-UN9ZrLSTdZJSSjPmF
76UakR0QVP\_sRHGG09Jj7j2y01HxvSlzRwE4PHDhqmYVbxa9EPMeR4JrbzYQ","iv":"fySs3Zk2exgrK1YY","tag":"UZ
765c5e172fb13077bf0edfbc724b03449"}

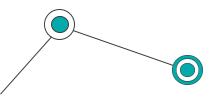




To obtain access to decrypted resource, the client must sign a message which represents the non-repudiation of reciept (NRR) that demonstrates the app really accessed the resource. At this point the server release the key for the decryption.

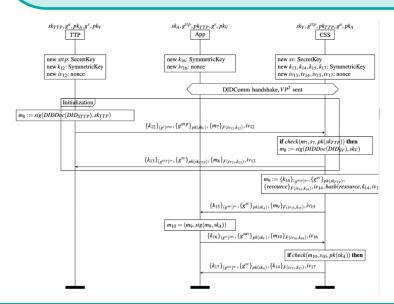


Non-repudiation of reciept (NRR)



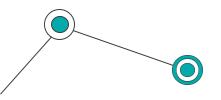


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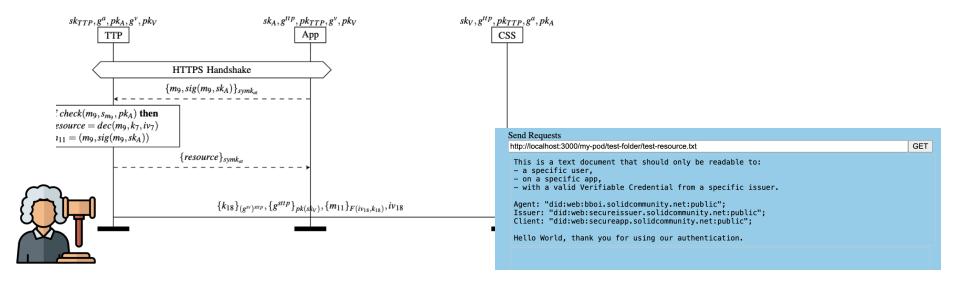
And if the server is malicious and sends the wrong key? Then we delivered NRR without really accessing the resource!!!!

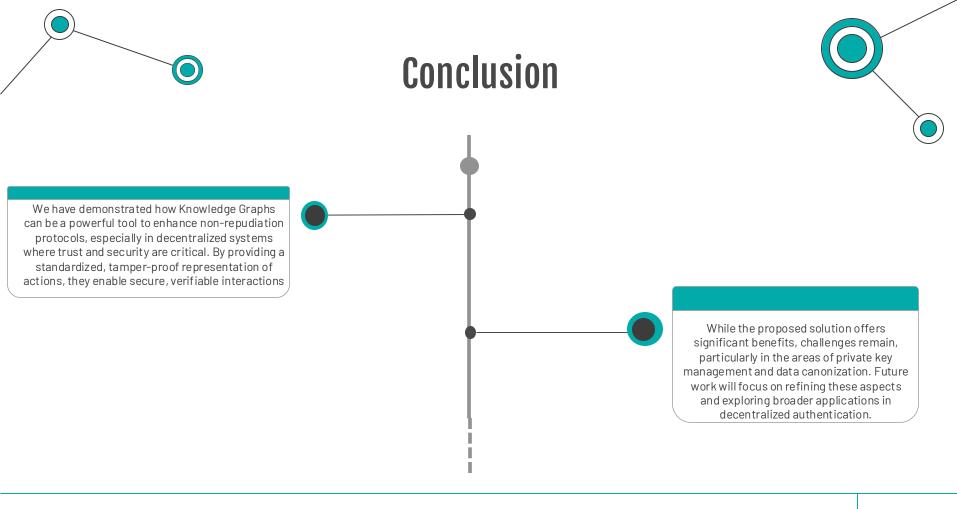






To overcome this problem, the key used for ciphering the resource is shared among CSS and TTP, meaning that the App can ask to the TTP to decrypt the message in case the CSS acts malicously.





# Thank you! Questions?

