

# Error Correction

Error Correction codes are used to detect and correct the errors when data is transmitted from the sender to the receiver.

Error Correction can be handled in two ways:

- **Backward error correction:** Once the error is discovered, the receiver requests the sender to retransmit the entire data unit.
- **Forward error correction:** In this case, the receiver uses the error-correcting code which automatically corrects the errors.

A single additional bit can detect the error, but cannot correct it.

For correcting the errors, one has to know the exact position of the error. For example, If we want to calculate a single-bit error, the error correction code will determine which one of seven bits is in error. To achieve this, we have to add some additional redundant bits.

## Hamming Code

**Parity bits:** The bit which is appended to the original data of binary bits so that the total number of 1s is even or odd.

**Even parity:** To check for even parity, if the total number of 1s is even, then the value of the parity bit is 0. If the total number of 1s occurrences is odd, then the value of the parity bit is 1.

**Odd Parity:** To check for odd parity, if the total number of 1s is even, then the value of parity bit is 1. If the total number of 1s is odd, then the value of parity bit is 0.

## Relationship b/w Error position & binary number.

| Error Position | Binary Number |
|----------------|---------------|
| 0              | 000           |
| 1              | 001           |
| 2              | 010           |
| 3              | 011           |
| 4              | 100           |
| 5              | 101           |
| 6              | 110           |
| 7              | 111           |

## Determining the position of the redundant bits

The number of redundant bits is 3. The three bits are represented by  $r_1$ ,  $r_2$ ,  $r_4$ . The position of the redundant bits is calculated with corresponds to the raised power of 2. Therefore, their corresponding positions are  $1$ ,  $2^1$ ,  $2^2$ .

1. The position of  $r_1 = 1$
2. The position of  $r_2 = 2$
3. The position of  $r_4 = 4$

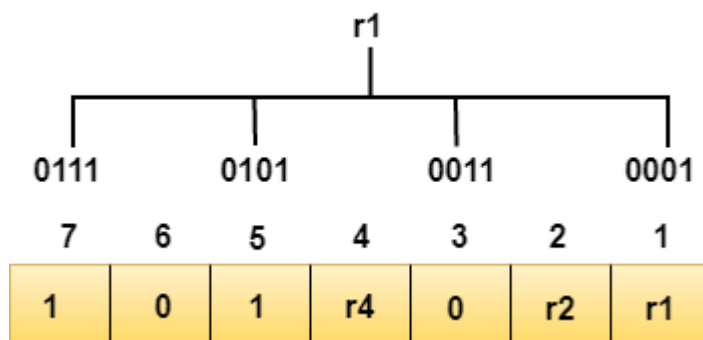
## Representation of Data on the addition of parity bits:

| 7 | 6 | 5 | 4     | 3 | 2     | 1     |
|---|---|---|-------|---|-------|-------|
| 1 | 0 | 1 | $r_4$ | 0 | $r_2$ | $r_1$ |

## Determining the Parity bits

### Determining the $r_1$ bit

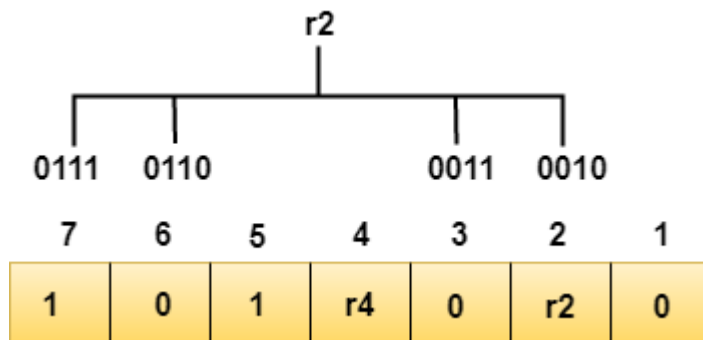
The  $r_1$  bit is calculated by performing a parity check on the bit positions whose binary representation includes 1 in the first position.



We observe from the above figure that the bit positions that includes 1 in the first position are 1, 3, 5, 7. Now, we perform the even-parity check at these bit positions. The total number of 1 at these bit positions corresponding to  $r_1$  is **even, therefore, the value of the  $r_1$  bit is 0.**

### Determining $r_2$ bit

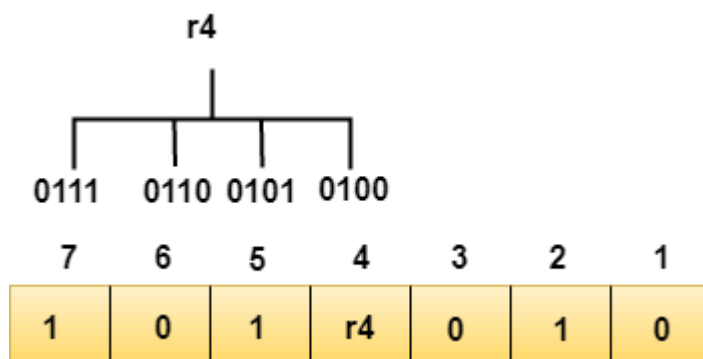
The r2 bit is calculated by performing a parity check on the bit positions whose binary representation includes 1 in the second position.



We observe from the above figure that the bit positions that includes 1 in the second position are 2, 3, 6, 7. Now, we perform the even-parity check at these bit positions. The total number of 1 at these bit positions corresponding to r2 is **odd, therefore, the value of the r2 bit is 1.**

## Determining r4 bit

The r4 bit is calculated by performing a parity check on the bit positions whose binary representation includes 1 in the third position.



We observe from the above figure that the bit positions that includes 1 in the third position are 4, 5, 6, 7. Now, we perform the even-parity check at these bit positions. The total number of 1 at these bit positions corresponding to r4 is **even, therefore, the value of the r4 bit is 0.**

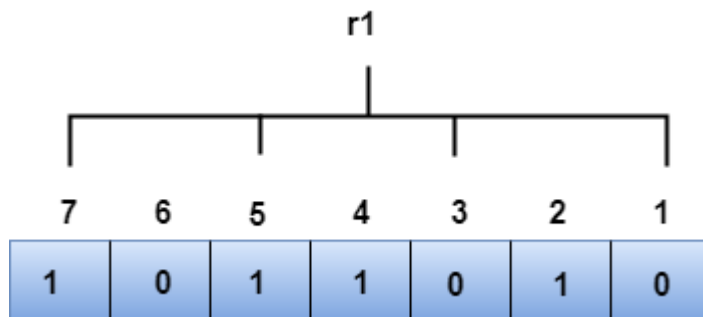
**Data transferred is given below:**

| 7 | 6 | 5 | 4 | 3 | 2 | 1 |
|---|---|---|---|---|---|---|
| 1 | 0 | 1 | 0 | 0 | 1 | 0 |

Suppose the 4<sup>th</sup> bit is changed from 0 to 1 at the receiving end, then parity bits are recalculated.

## R1 bit

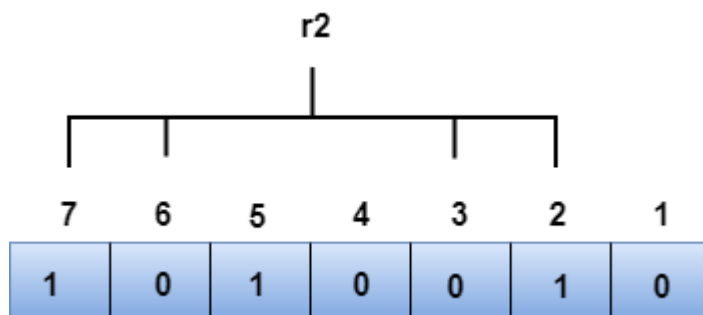
The bit positions of the r1 bit are 1,3,5,7



We observe from the above figure that the binary representation of r1 is 1100. Now, we perform the even-parity check, the total number of 1s appearing in the r1 bit is an even number. Therefore, the value of r1 is 0.

## R2 bit

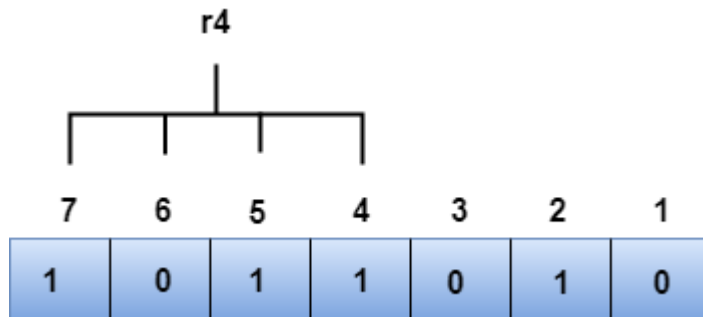
The bit positions of r2 bit are 2,3,6,7.



We observe from the above figure that the binary representation of r2 is 1001. Now, we perform the even-parity check, the total number of 1s appearing in the r2 bit is an even number. Therefore, the value of r2 is 0.

## R4 bit

The bit positions of r4 bit are 4,5,6,7.



We observe from the above figure that the binary representation of  $r4$  is 1011. Now, we perform the even-parity check, the total number of 1s appearing in the  $r4$  bit is an odd number. Therefore, the value of  $r4$  is 1.

- *The binary representation of redundant bits, i.e.,  $r4r2r1$  is 100, and its corresponding decimal value is 4. Therefore, the error occurs in a 4<sup>th</sup> bit position. The bit value must be changed from 1 to 0 to correct the error*