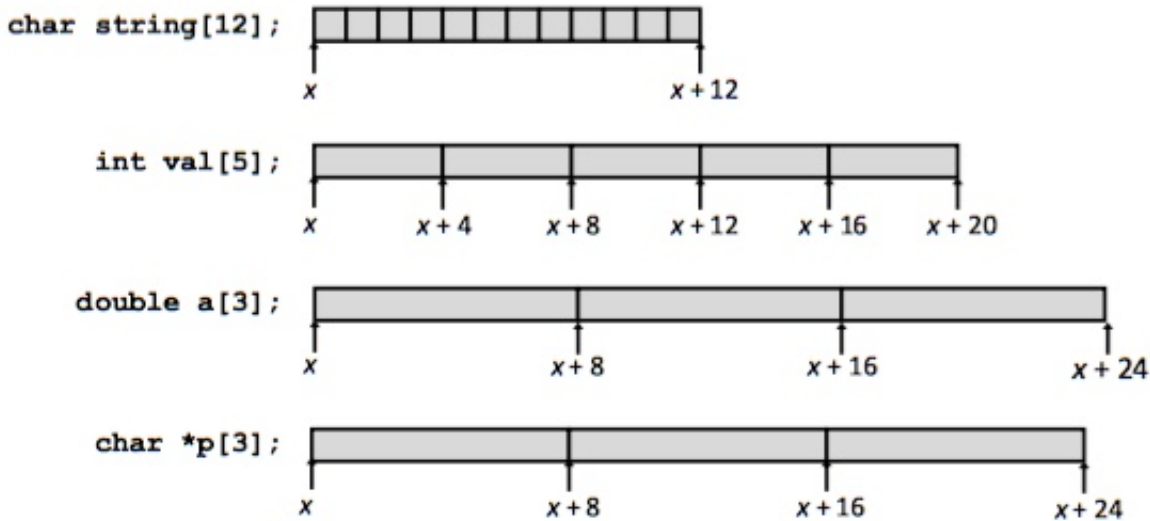


## Array Allocation

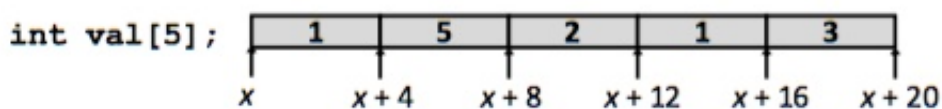
- `T A[L];`
  - array of data type `T` and length `L`
  - contiguously allocated region of `L * sizeof(T)` bytes in memory



- nested arrays: `A[n][n]`
  - static
    - compiler can optimize this very well
  - dynamic
    - hard for compiler to optimize
- multi-level
  - i.e. array of pointers to arrays (like hash tables)

## Array Access

- identifier `A` can be used as a pointer to array element 0: Type `T*`



Reference	Type	Value
<code>val[4]</code>	<code>int</code>	3
<code>val</code>	<code>int *</code>	$x$
<code>val+1</code>	<code>int *</code>	$x+4$
<code>&amp;val[2]</code>	<code>int *</code>	$x+8$
<code>val[5]</code>	<code>int</code>	??
<code>*(val+1)</code>	<code>int</code>	5
<code>val + i</code>	<code>int *</code>	$x+4i$

## Array Loop Example

```

1 void zincr(zip_dig z) {
2     size_t i;
3     for (i = 0; i < ZLEN; i++)
4         z[i]++;
5 }

```

```

1 ; %rdi = z
2
3 movl    $0, %eax
4 jmp     .L3
5
6 .L4:
7     addl    $1, (%rdi,%rax,4)
8     addq    $1, %rax
9 .L3:
10    cmpq    $4, %rax
11    jbe     .L4
12    ret

```

## Multidimensional (Nested) Arrays

### Declaration

`T A[R][C]`

- 2D array of data type T
- R rows, C columns
- type T element requires K bytes

### Array Size

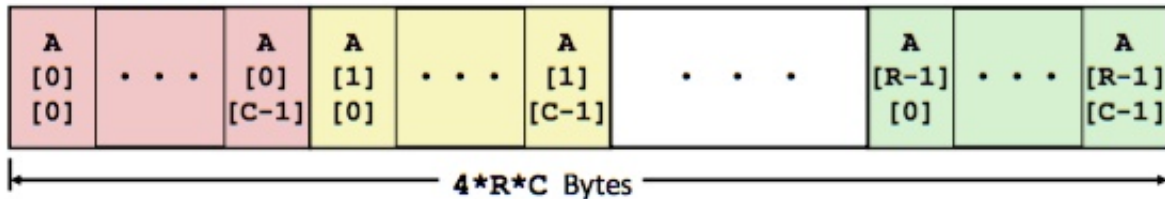
- $R * C * K$  bytes

### Arrangement

- row-major ordering

$$\begin{bmatrix} A[0][0] & \dots & A[0][C-1] \\ \vdots & & \vdots \\ A[R-1][0] & \dots & A[R-1][C-1] \end{bmatrix}$$

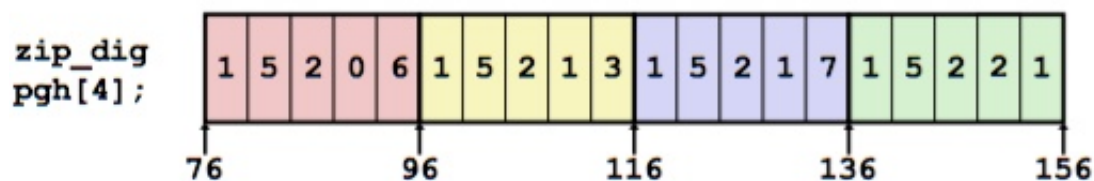
```
int A[R][C];
```



- each color is a separate row

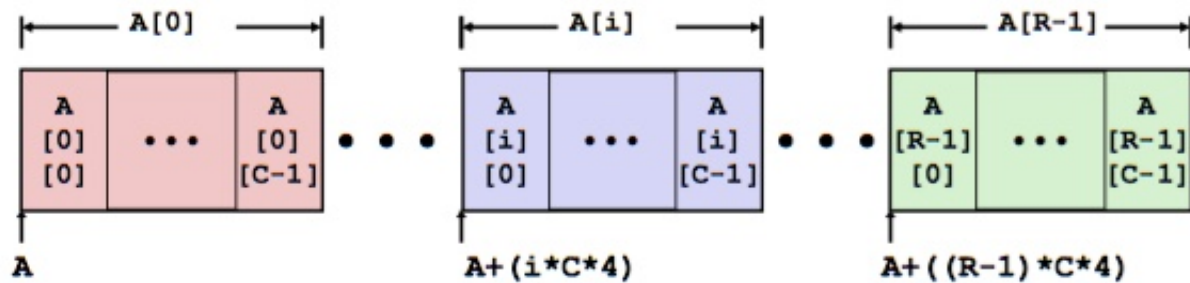
## Nested Array Example

```
1 #define PCOUNT 4
2 zip_dig pgh[PCOUNT] =
3     {{1, 5, 2, 0, 6},
4      {1, 5, 2, 1, 3},
5      {1, 5, 2, 1, 7},
6      {1, 5, 2, 2, 1}};
```



## Nested Array Row Access

- `A[i]` is array of `C` elements
- each element of type `T` requires `K` bytes
- $A + i * (C * K)$



## Access Code

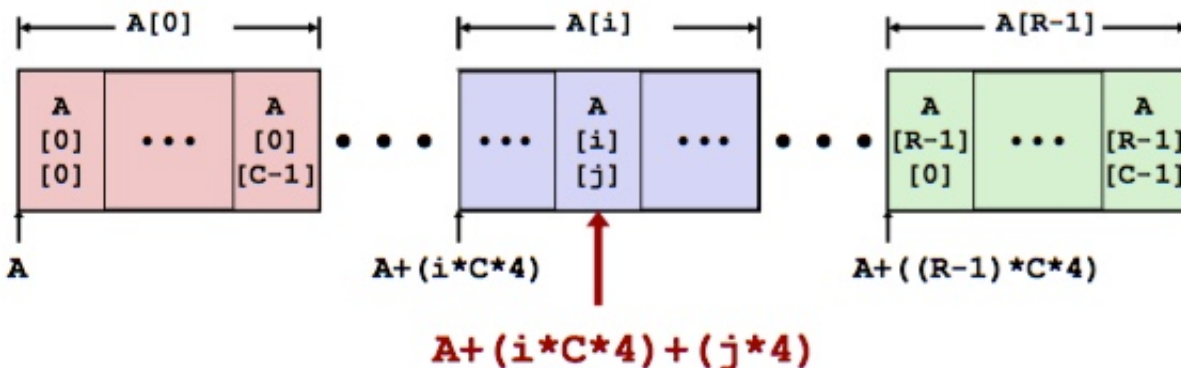
```
1 int *get_pgh_zip(int index) {
2     return pgh[index];
3 }
```

```
1 ; %rdi = index
2 ; leaq doesn't dereference the contents at the address
3 ; just puts memory location into %rax
4 leaq    (%rdi,%rdi,4), %rax    ; 5 * index into %rax
5 leaq    pgh(,%rax,4), %rax     ; pgh + (20 * index) into %rax
```

## Nested Array Element Access

- $A[i][j]$  is element of type  $T$ , which requires  $K$  bytes
- **address:**  $A + (i * (C * K)) + (j * K) = A + (i * C + j) * K$

```
int A[R][C];
```

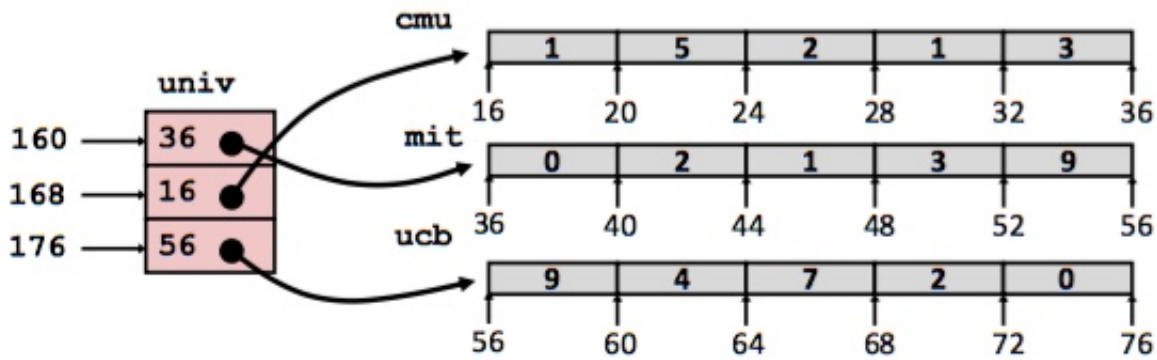


## Access Code

```
1 int get_pgh_digit(int index, int dig) {
2     return pgh[index][dig];
3 }
```

```
1 leaq    (%rdi,%rdi,4), %rax    ; 5*index into %rax
2 addl    %rax, %rsi             ; (5*index) + dig
3 movl    pgh(,%rsi,4), %eax     ; M[pgh + 4*(5*index)+dig]
```

## Multi-Level Array Example



```
1 zip_dig cmu = { 1, 5, 2, 1, 3 };
2 zip_dig mit = { 0, 2, 1, 3, 9 };
3 zip_dig ucb = { 9, 4, 7, 2, 0 };
```

```
1 #define UCOUNT 3
2 int *univ[UCOUNT] = {mit, cmu, ucb};
```

```
1 int get_univ_digit (size_t index, size_t digit) {
2     return univ[index][digit];
3 }
```

```
1 salq    $2, %rsi                ; 4*digit
2 addq    univ(,%rdi,8), %rsi      ; p = univ[index] + 4*digit
3 movl    (%rsi), %eax             ; return *p
4 ret
```

### Must do two memory reads

1. to get pointer to row array
2. then access element within array

Less efficient than nested memory access, which only requires ONE memory load

## N x N Matrix Access

TAKE ANOTHER LOOK AT THESE

```

/* Get element a[i][j] */
int var_ele(size_t n, int a[n][n], size_t i, size_t j)
{
    return a[i][j];
}

```

```

# n in %rdi, a in %rsi, i in %rdx, j in %rcx
imulq    %rdx, %rdi          # n*i
leaq     (%rsi,%rdi,4), %rax  # a + 4*n*i
movl     (%rax,%rcx,4), %eax  # a + 4*n*i + 4*j
ret

```

## Questions

1. Comparison between nested arrays, multi-level arrays
2. N x N matrix

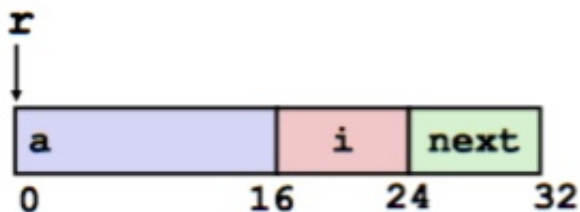
# Structures

## Structure Representation

```

1 struct rec {
2     int a[4];
3     size_t i;
4     struct rec *next;
5 };

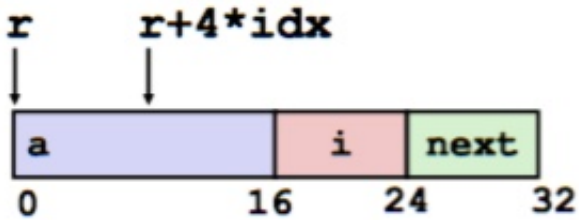
```



- structure represented as block of memory
  - big enough to hold all of the fields
- fields ordered according to declaration
  - even if another ordering could yield a more compact representation
- compiler determines overall size + position of fields
  - machine-level program has no understanding of the structures in the source code
- can access different elements of struct using pointer arithmetic

- data laid out linearly, the same way every time

## Generating Pointer to Structure Member



```
1 int *get_ap (struct rec *r, size_t idx) {
2     return r->a[idx];
3 }
```

```
1 ; r in %rdi, idx in %rsi
2 leaq (%rdi, %rsi, 4), %rax
3 ret
```

## Generating Pointer to Array Element

- offset of each structure member determined at compile time
- compute as `r + 4*idx`

## Following Linked List

```
1 void set_val (struct rec *r, int val) {
2     while (r) {
3         int i = r->i;
4         r->a[i] = val;
5         r = r->next;
6     }
7 }
```

```
1 ; r in %rdi, val in %rsi
2 .L11:                                ; loop:
3     movslq 16(%rdi), %rax            ; i = M[r+16]
4     movl   %esi, (%rdi,%rax,4)       ; M[r+4*i] = val
5     movq   24(%rdi), %rdi           ; r = M[r+24]
6     testq  %rdi, %rdi               ; Test r
7     jne    .L11                     ; if != 0 goto .L11
```

## Structures & Alignment

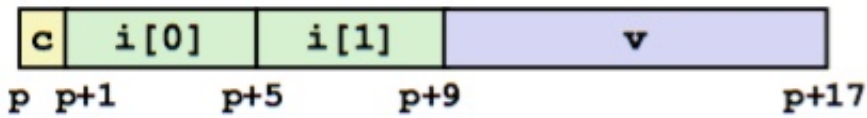
```
1 struct S1 {
2     char c;        // 1 1-byte quantity
```

```

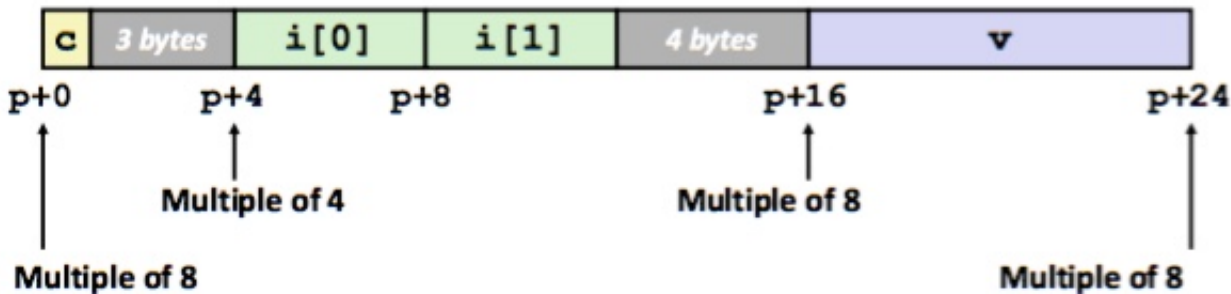
3     int i[2];    // 2 4-bytes quantities
4     double v;    // 8-byte quantities
5 }

```

## Unaligned



## Aligned



## Aligned Data

- primitive data type requires K bytes
- address must be multiple of K
  - char requires 1 byte, can be anywhere
  - int requires 4 bytes, must be at address multiple of 4
  - double requires 8 bytes, must be at address multiple of 8
- required on some machines; advised on x86-64

## Motivation for Aligning Data

- memory access by (aligned) chunks of 4 or 8 bytes (system dependent)
  - inefficient to load or store datum that spans quad word boundaries
  - virtual memory trickier when datum spans 2 pages

## Compiler

- inserts gaps in structure to ensure correct alignment of fields

# Satisfying Alignment with Structures

## Within Structure

- must satisfy each element's alignment requirement

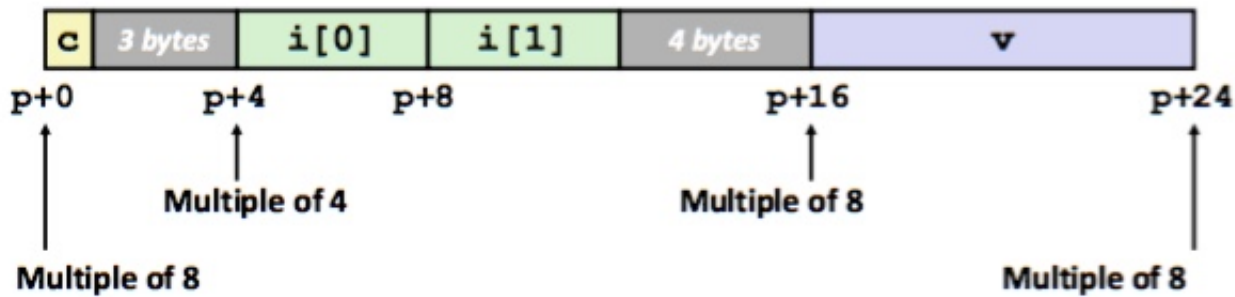


## Overall Structure Placement

- each structure has alignment requirement **K**
  - **K** = largest alignment of any element
- initial address & structure length must be multiples of **K**

## Example

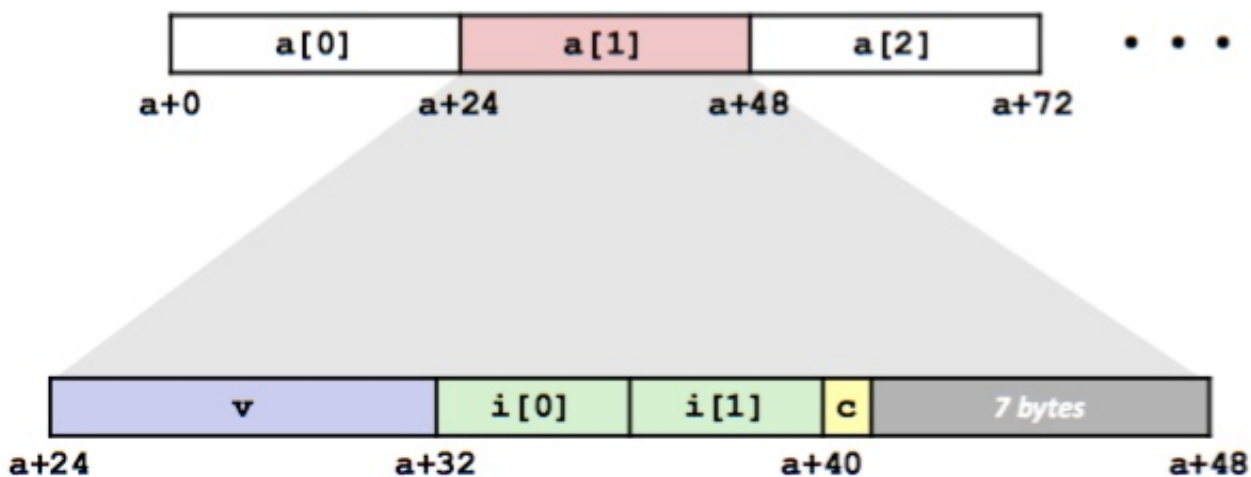
- $K = 8$ , due to `double` element



## Arrays of Structures

- overall structure length multiple of **K**
  - **array has 7 bytes of padding to meet  $K=8$  requirement**
- satisfy alignment requirement for every element

```
1 struct S2 {  
2     double v;  
3     int i[2];  
4     char c;  
5 } a[10];
```



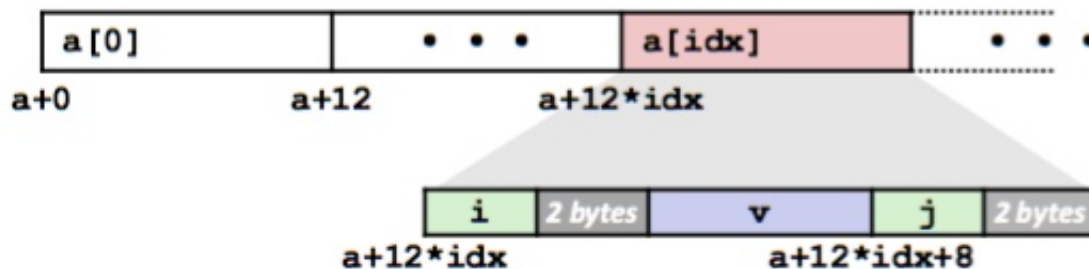
## Accessing Array Elements

```

1 struct S3 {
2     short i;
3     float v;
4     short j;
5 } a[10];

```

- compute array offset  $12 \cdot \text{idx}$ 
  - `sizeof(S3)`, including alignment spacers
- element `j` is at offset 8 within structure
- assembler gives offset `a+8`
  - resolved during linking



```

1 short get_j (int idx) {
2     return a[idx].j;
3 }

```

```

1 ; %rdi = idx
2 leaq    (%rdi,%rdi,2), %rax    ; 3*idx
3 movzwl  a+8(,%rax,4), %eax

```

## Saving Space

- put large data types first

```

struct S4 {
    char c;
    int i;
    char d;
} *p;

```



```

struct S5 {
    int i;
    char c;
    char d;
} *p;

```

