

Summation:

$$\llbracket \sum_i \pi_i.R_i \rrbracket^d \stackrel{\text{def}}{=} d?(n).(d!\langle n+1 \rangle \mid \text{new}(l).(l!\langle true \rangle \mid \prod_i \llbracket \pi_i.R_i \rrbracket_{n,l}^d))$$

Sending:

$$\llbracket c!\langle \bar{V} \rangle.P \rrbracket_{n,l}^d \stackrel{\text{def}}{=} \text{new}(ack).(c!\langle n, l, ack, \bar{V} \rangle \mid ack?(x).\text{if } x = true \text{ then } \llbracket P \rrbracket^d \\ \text{else (if } x = retry \text{ then } c!\langle n, l, ack, \bar{V} \rangle \text{ else } stop))$$

Receiving:

$$\llbracket c?(\bar{X}).P \rrbracket_{n,l}^d \stackrel{\text{def}}{=} \text{rec } q.(c?(m, r, ack, \bar{X}).(\\ \text{if } n = m \text{ then } (ack!\langle retry \rangle \mid q) \text{ else } (\\ \text{if } n < m \text{ then } (l, r)?d\text{-lock}.\llbracket P \rrbracket^d \text{ else } (r, l)?rd\text{-lock}.\llbracket P \rrbracket^d)))$$

(l, r)?d-lock.P means:

$l?(x).(\text{if } x = true$
 then $r?(y).(\text{if } y = true$
 then $l!\langle false \rangle \mid r!\langle false \rangle \mid ack!\langle true \rangle \mid P$
 else $l!\langle true \rangle \mid r!\langle false \rangle \mid ack!\langle false \rangle \mid q)$
 else $l!\langle false \rangle \mid ack!\langle retry \rangle)$

(r, l)?rd-lock.P means:

$r?(x).(\text{if } x = true$
 then $l?(y).(\text{if } y = true$
 then $l!\langle false \rangle \mid r!\langle false \rangle \mid ack!\langle true \rangle \mid P$
 else $l!\langle false \rangle \mid r!\langle true \rangle \mid ack!\langle retry \rangle)$
 else $r!\langle false \rangle \mid ack!\langle false \rangle \mid q)$