Web search engines

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Two main difficulties

The Web:

Extracting “significant data” is difficult !!

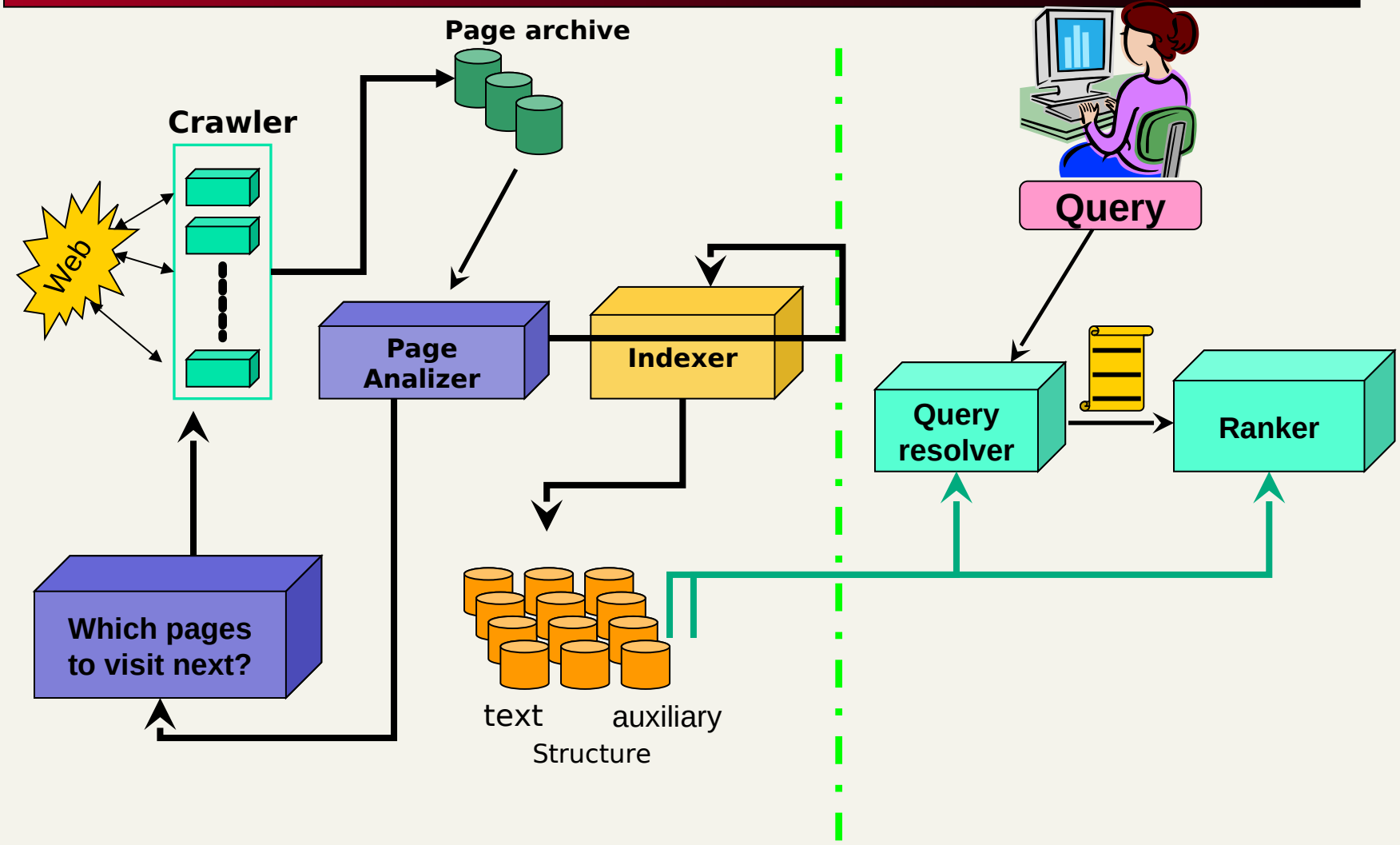
- Size: more than 1 trillion pages
- Language and encodings: hundreds…
- Distributed authorship: SPAM, format-less,…
- Dynamic: in one year 35% survive, 20% untouched

The User:

Matching “user needs” is difficult !!

- Query composition: short (2.5 terms avg) and imprecise
- Query results: 85% users look at just one result-page
- Several needs: Informational, Navigational, Transactional

The structure of a search Engine



The web graph: properties

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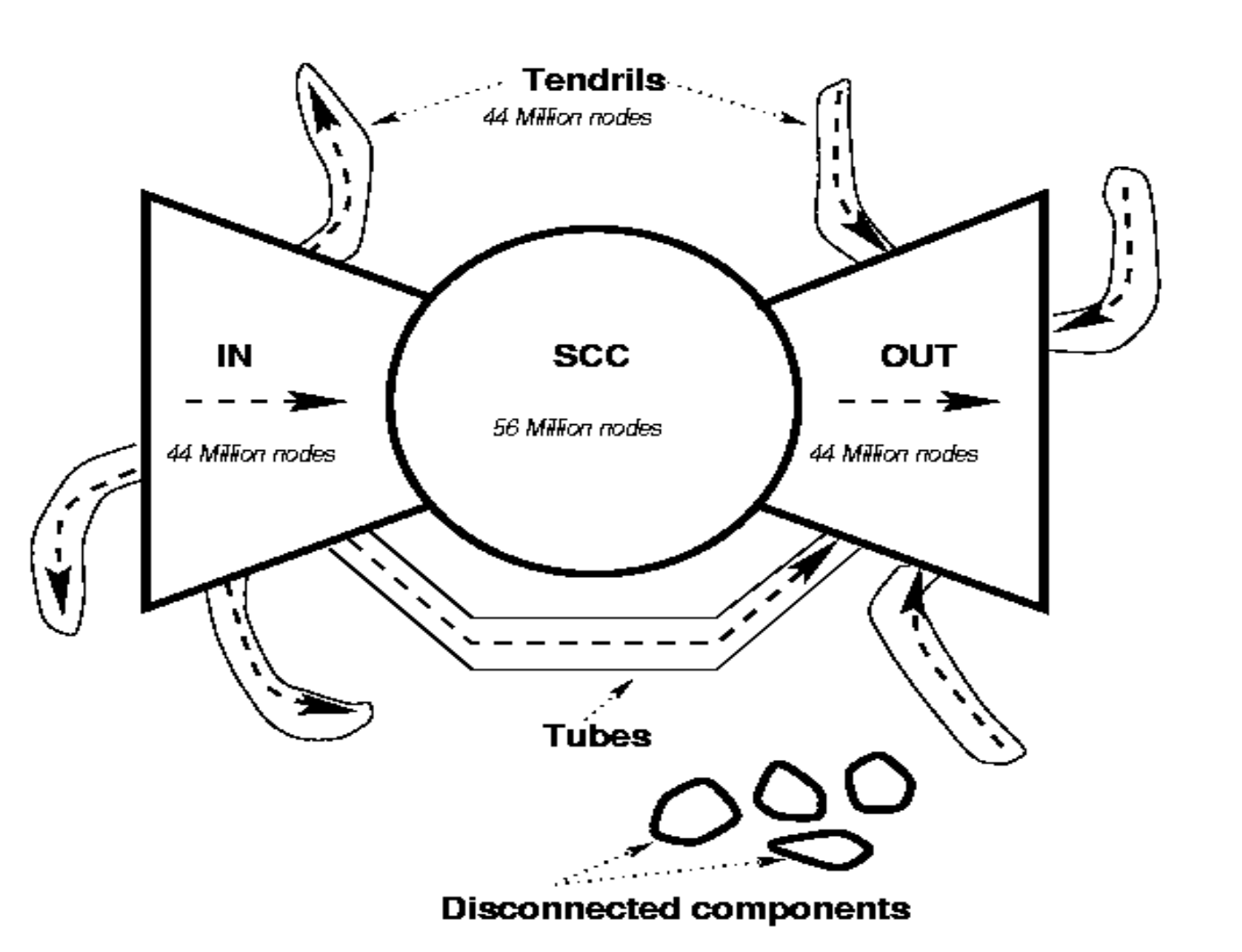
The Web’s Characteristics

- It’s a graph whose size is
  - 1 trillion of pages is available
    - 50 billion pages crawled (09/15)
  - 5-40K per page => terabytes & terabytes
  - Size grows every day!!
  - Actually the web is infinite… Calendars…

- It’s a dynamic graph with
  - 8% new pages, 25% new links change weekly
  - Life time of about 10 days

The Bow Tie


Tendrils
44 Million nodes
IN
44 Million nodes
SCC
56 Million nodes
OUT
44 Million nodes
Tubes
Disconnected components

Crawling

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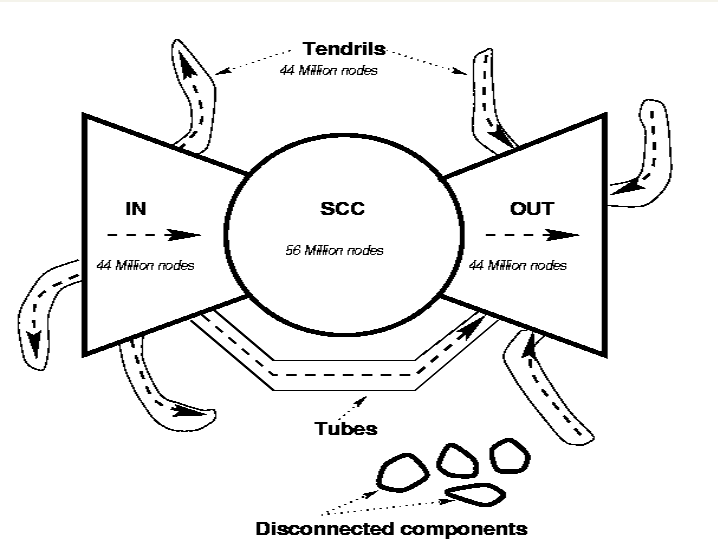
Spidering

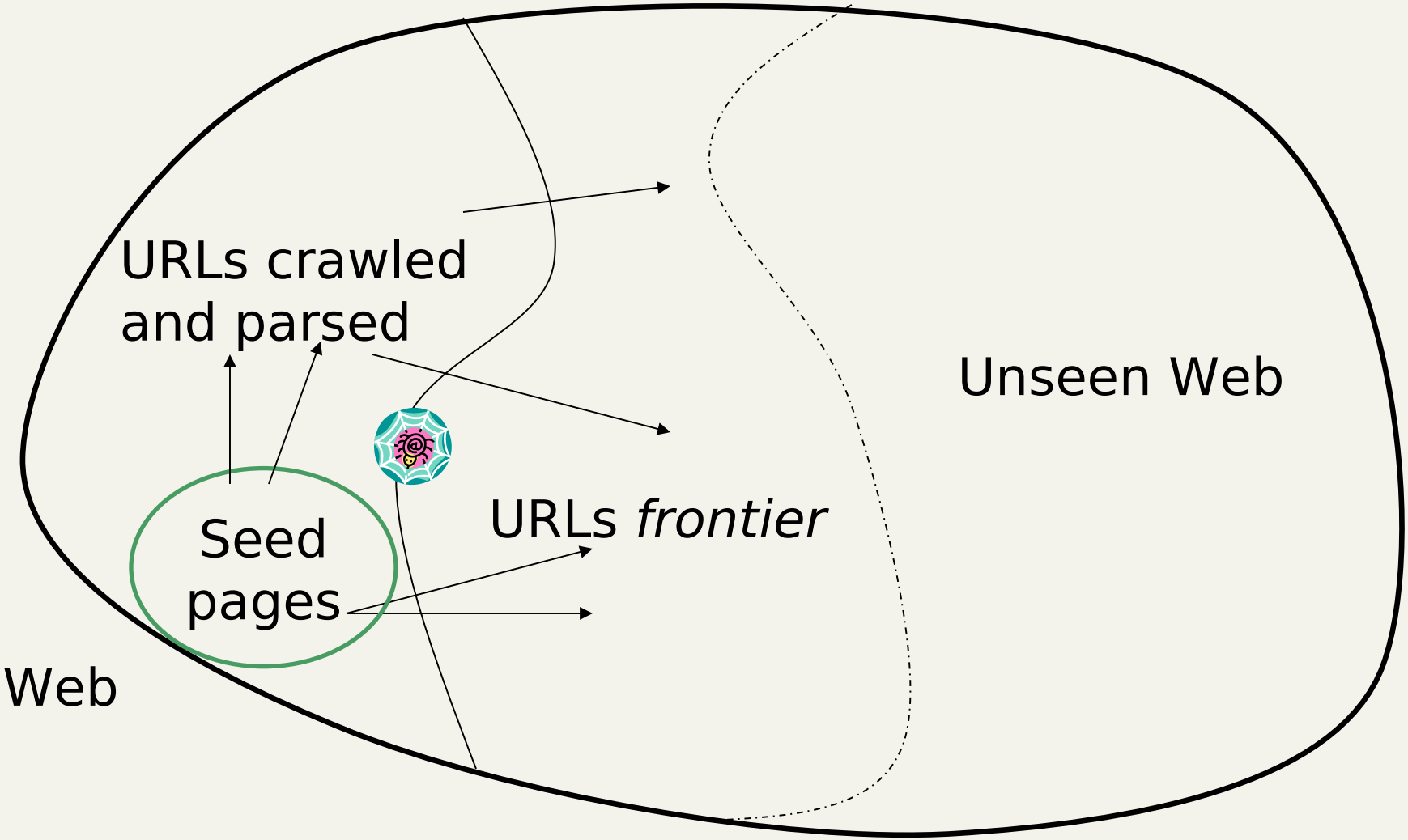
- 24h, 7days “walking” over the web graph
- Recall that:
  - Direct graph G = (N, E)
  - N changes (insert, delete) >> trillion nodes
  - E changes (insert, delete) > 10 links per node
    - Trillion entries in posting lists
    - Many more if we consider also the word positions in every document where it occurs.



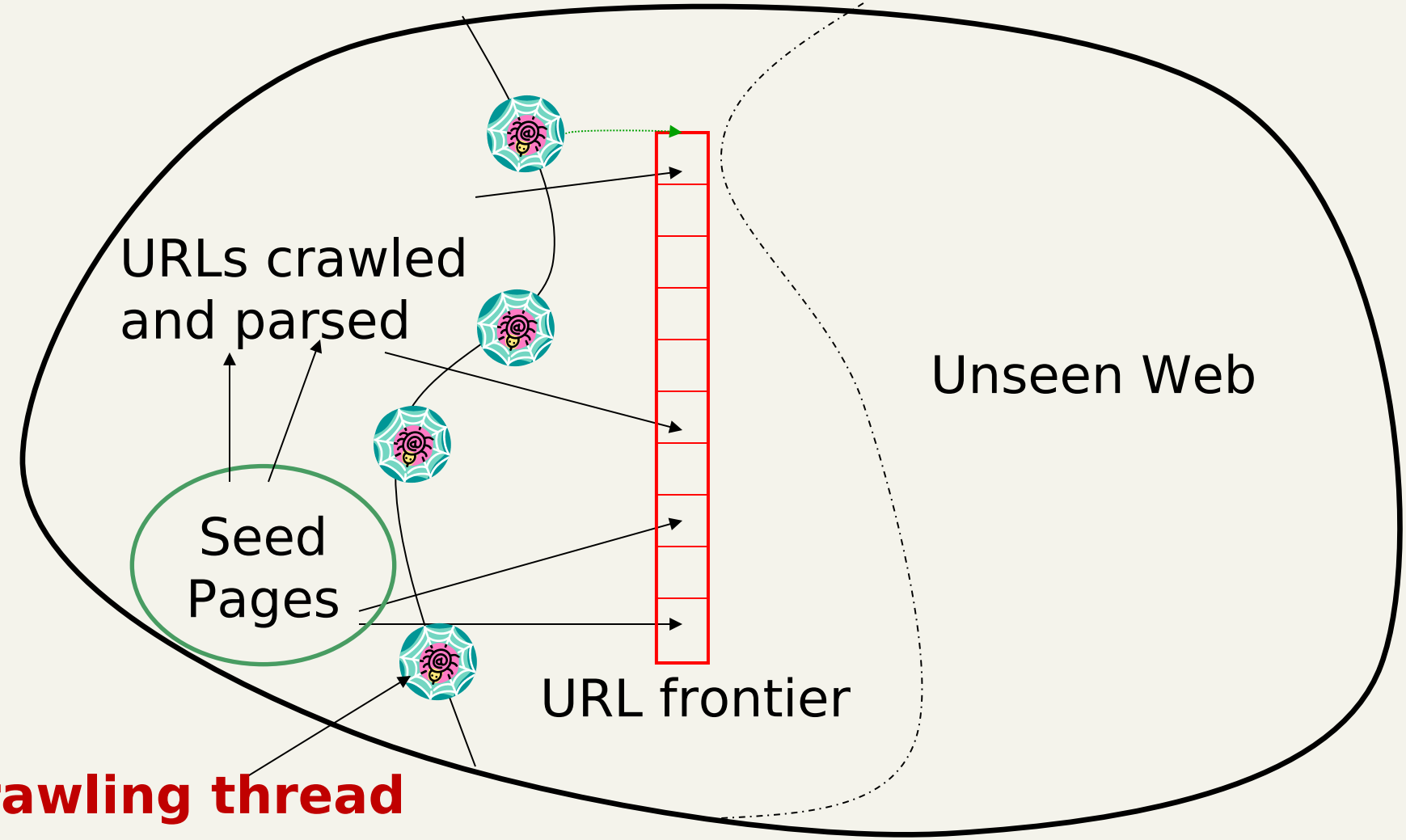
Crawling Issues

- How to crawl?
  - *Quality*: “Best” pages first
  - *Efficiency*: Avoid duplication (or near duplication)
  - *Etiquette*: Robots.txt, Server load concerns (Minimize load)
  - *Malicious pages*: Spam pages, Spider traps – incl dynamically generated
- How much to crawl, and thus index?
  - *Coverage*: How big is the Web? How much do we cover?
  - *Relative Coverage:* How much do competitors have?
- How often to crawl?
  - *Freshness*: How much has changed?
  - *Frequency*: Commonly insert time gap btw host requests

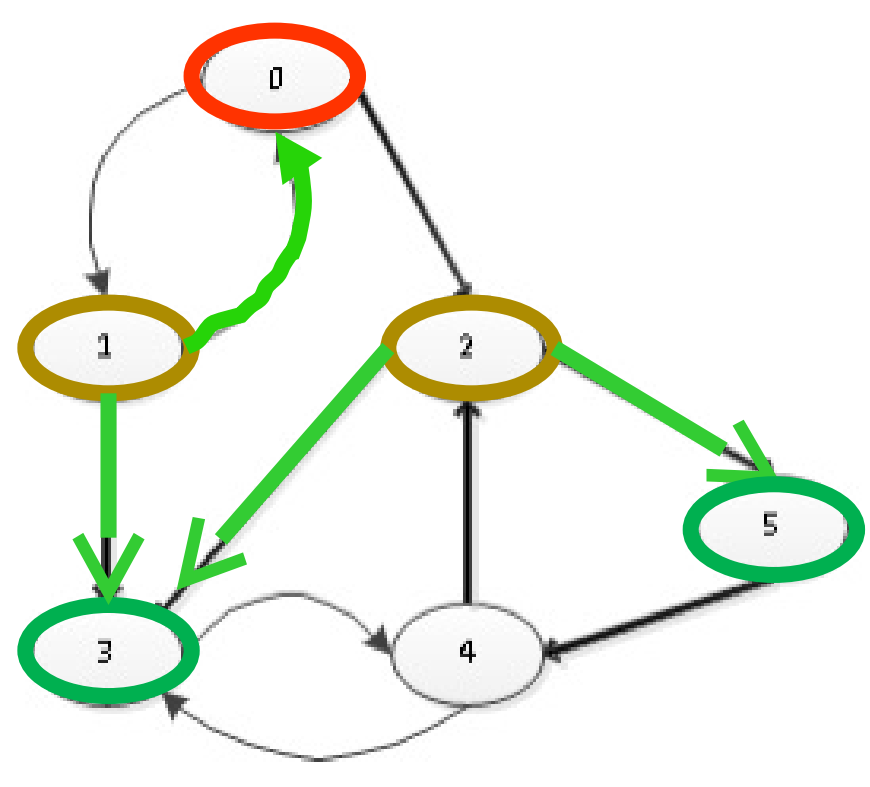
Crawling picture

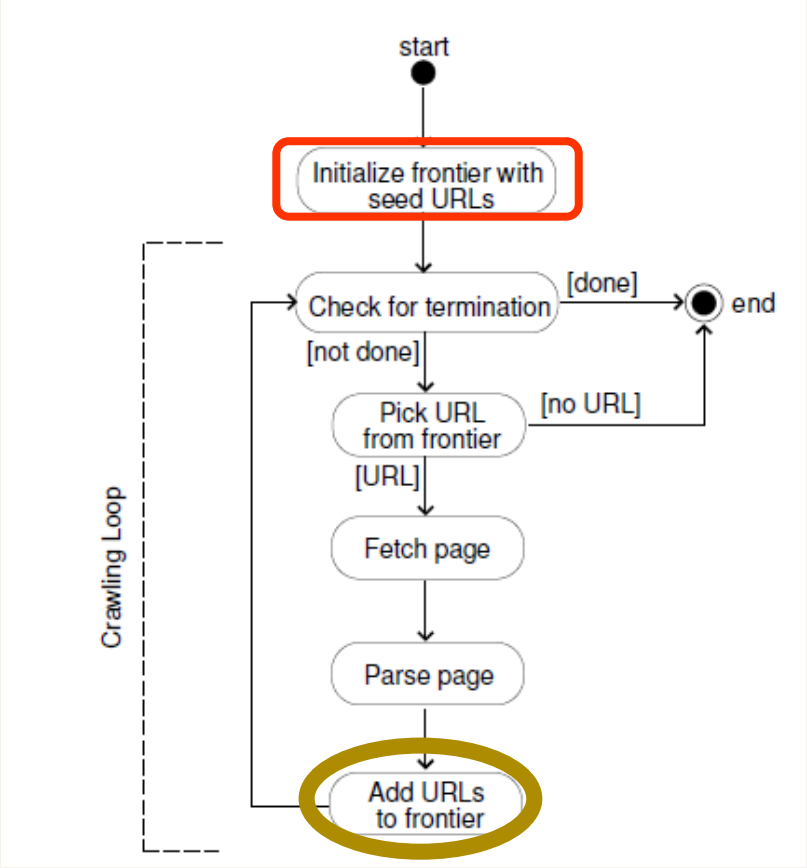


Updated crawling picture

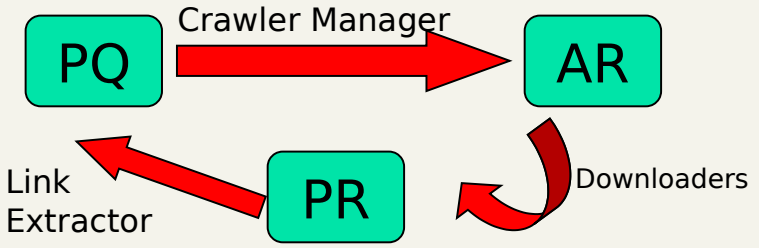


A small example





Crawler “life cycle”



One Link Extractor per page:

while(<**Page Repository** is not empty>){
  <take a page *p (check if it is new)*>
  <extract links contained in p within *href*>
  <extract links contained in *javascript*>
  <extract .....
  <insert these links into the Priority Queue>
}

One Downloader per page:

while(<**Assigned Repository** is not empty>){
  <extract url *u*>
  <download *page(u)*>
  <send *page(u)* to the **Page Repository**>
  <store *page(u)* in a proper archive,
    possibly compressed>
}

One single Crawler Manager:

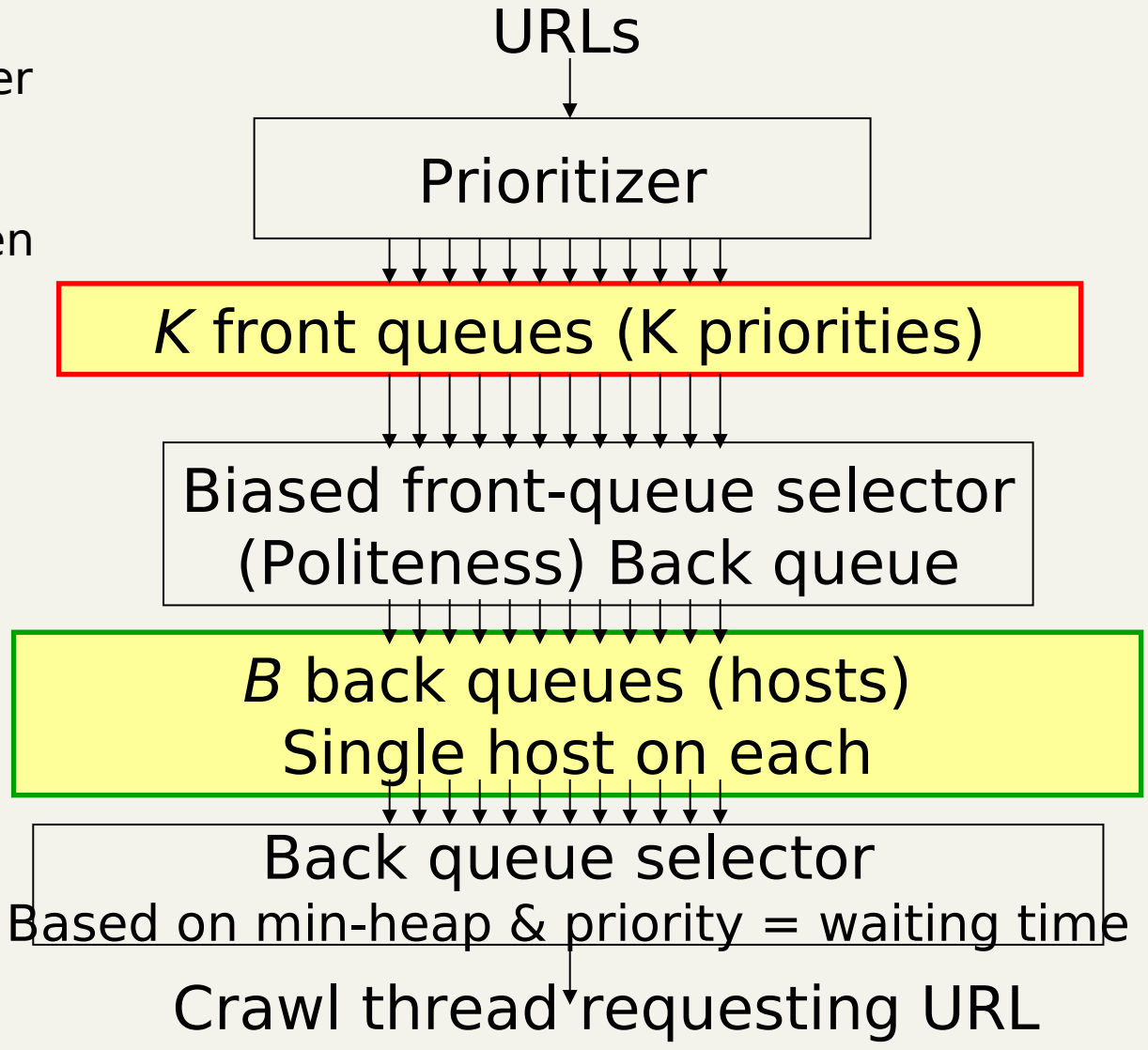
while(<**Priority Queue** is not empty>){
  <extract some URL *u* having the highest priority>
  foreach u extracted {
    if ((u ∉ “Already Seen Page”) ||
      (u ∈ “Already Seen Page” && <u’s version on the Web is more recent>)
) {
      <resolve u wrt DNS>
      <send u to the **Assigned Repository**>
    }
  }
}

URL frontier visiting

- Given a page P, define how “good” P is.
- Several *metrics* (via priority assignment):
  - BFS, DFS, Random
  - Popularity driven (*PageRank*, full vs partial)
  - Topic driven or focused crawling
  - Combined
- How to fast check whether the URL is new ?

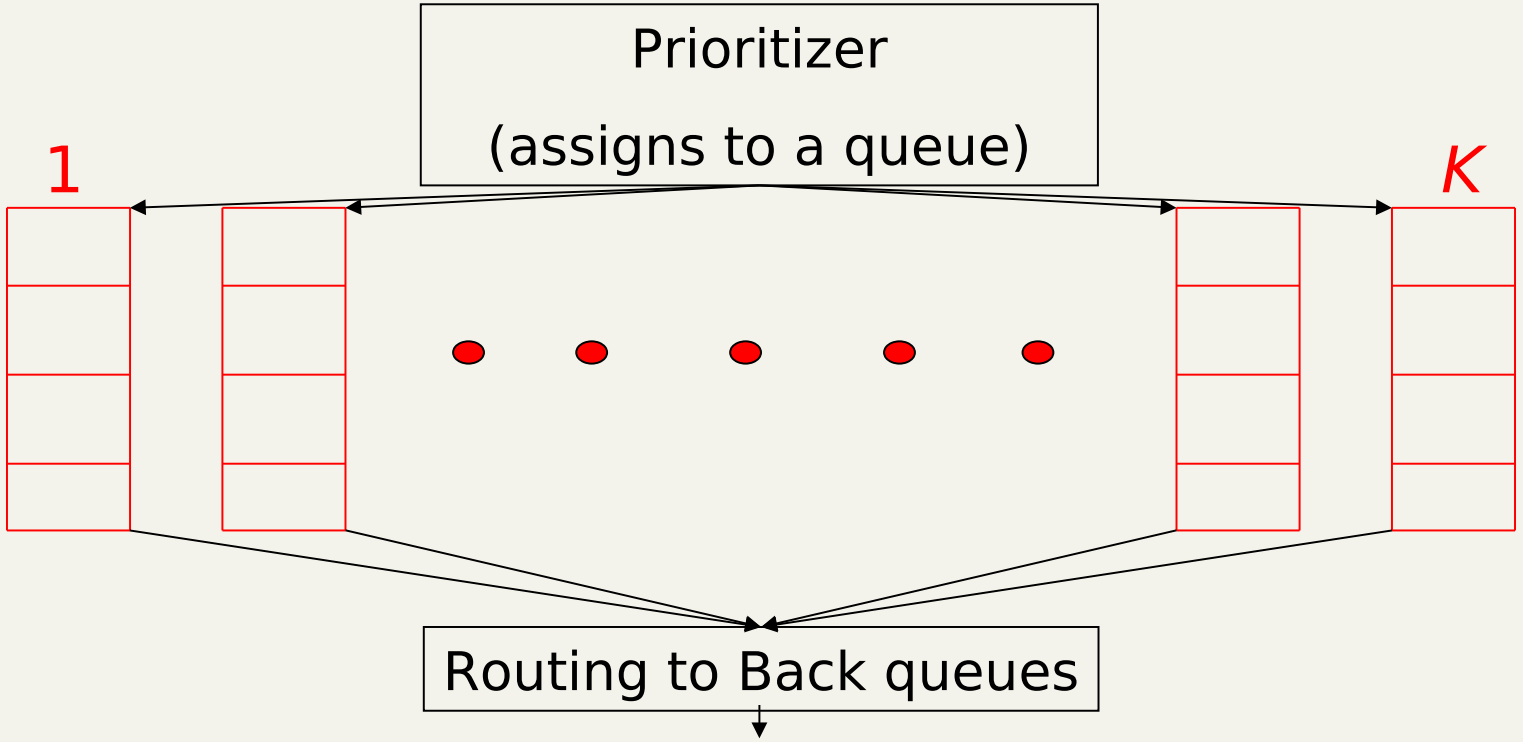
Mercator

1. Only one connection per host is open at a time;
2. a waiting time of a few seconds occurs between successive requests to the same host;
3. high-priority pages are crawled preferentially.

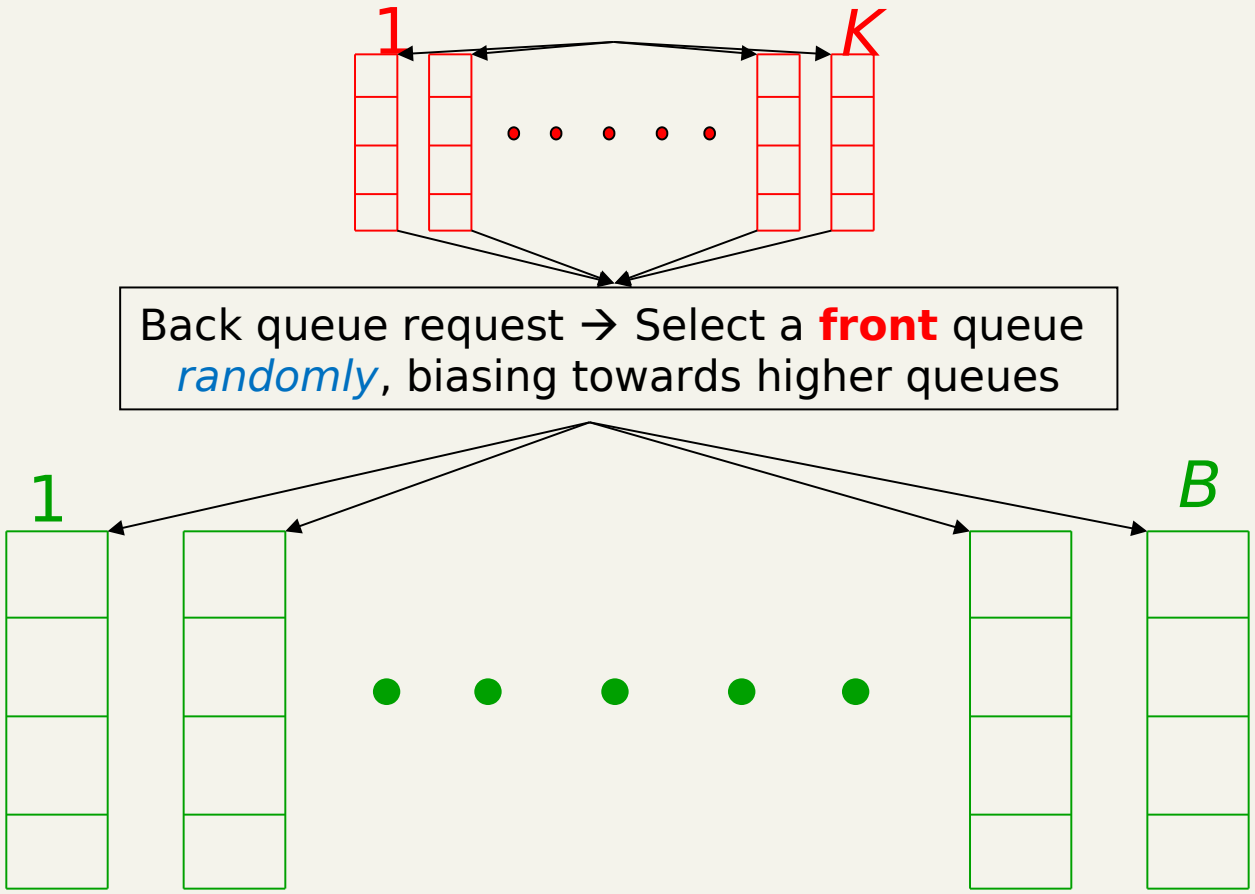


Front queues

- Front queues manage **prioritization**:
  - Prioritizer assigns to an URL an integer priority (refresh, quality, application specific) between 1 and *K*
  - Appends URL to corresponding queue, according to priority

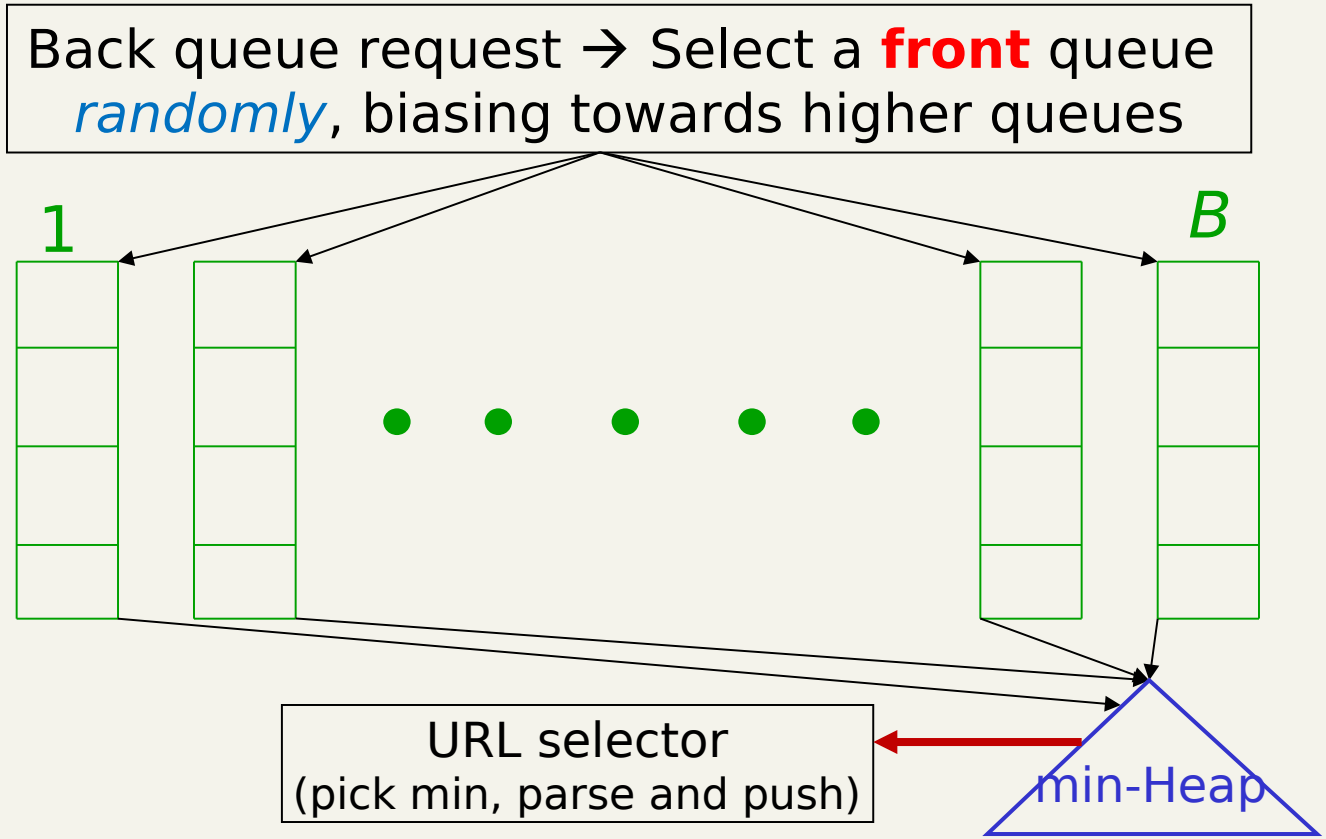


Back queues


1
K
Back queue request → Select a **front** queue
randomly, biasing towards higher queues
1
B

Back queues

- Back queues enforce **politeness:**
  - Each back queue is kept non-empty
  - Each back queue contains only URLs from a **single** host



The **min-heap**

- It contains one entry per back queue
- The entry is the earliest time t_e at which the host corresponding to the back queue can be "hit again"
- This earliest time is determined from
  - Last access to that host
  - Any time buffer heuristic we choose

The crawl thread

- A crawler seeks a URL to crawl:
  - Extracts the root of the heap: So it is an URL at the head of some back queue *q (and then removes it)*
  - Waits the indicated time t_{url}
  - Parses URL and adds its out-links to the Front queues
- If back queue *q* gets empty, pulls a **URL *v*** from some front queue *(more prob for higher queues)*
  - If there's already a back queue for *v*'s host, append *v* to it and repeat until q gets not empty;
  - Else, make *q* the back queue for *v's* host
- If back queue *q* is non-empty, pick URL and add it to the min-heap with priority = waiting time t_{url}

Keep crawl threads busy (B = 3 x threads)