

5) D: prove bubble sort is correct

example 3, 2, 5, 4, 7, 1

base case $i = 0$

$i = 0$ 2, 3, 5, 4, 7, 1

$i = 1$ 2, 3, 5, 4, 7, 1

$i = 2$ 2, 3, 4, 5, 7, 1

$i = 3$ 2, 3, 4, 5, 7, 1

$i = 4$ 2, 3, 4, 5, 1, 7

$i = 5$ 2, 3, 4, 5, 1, 7

we are pushing the maximum elements to the last position of the array

So, after i^{th} pass elements of $a[n-1] \dots a[n-1-i]$ are in their final sorted order

$I(i)$ $a[n-1] \dots a[n-1-i]$ are sorted for $0 \leq i \leq n-1$

hypothesis

assume $I(i)$ holds for $0 \leq i \leq n-1$

induction step

If $i+1$ show that this holds

(2)

the algorithm will place the large elements

among $an[n-i-1] \dots an[0]$ in position $n-i-2$

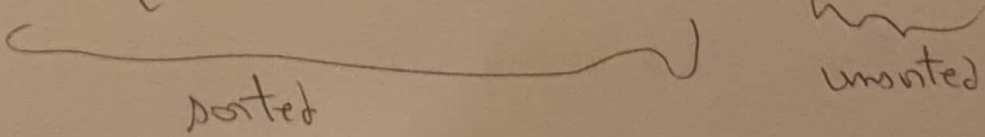
$I(i+1) : an[\dots n-i-2] \dots an[n-1]$ will be in its final sorted order

this shows $I(i)$ is true $0 \leq i \leq n-1$.

E) whenever again $an[i]$ and $an[i+1]$ in $0 \leq i \leq n-1$ in i th iteration $0 \leq i \leq n-1$ occur in inverted order, so that $an[i] > an[i+1]$

then bubble sort performs comparison and swap

~~after~~ after the i th step in the outer loop we will

have $an[n-1] \dots an[n-i-1] \dots an[0]$

sorted unsorted

As bubble sort algorithm attempt to place an element in $an[n-i-2]$ position it compares elements from $an[n-i-1]$ to $an[0]$

that will be placed in $an[n-i-2]$

There are as many inversions as there are comparisons in this algorithm

bubble sort is inversion bound.

1) ~~phase~~ $T(n) = O(n)$