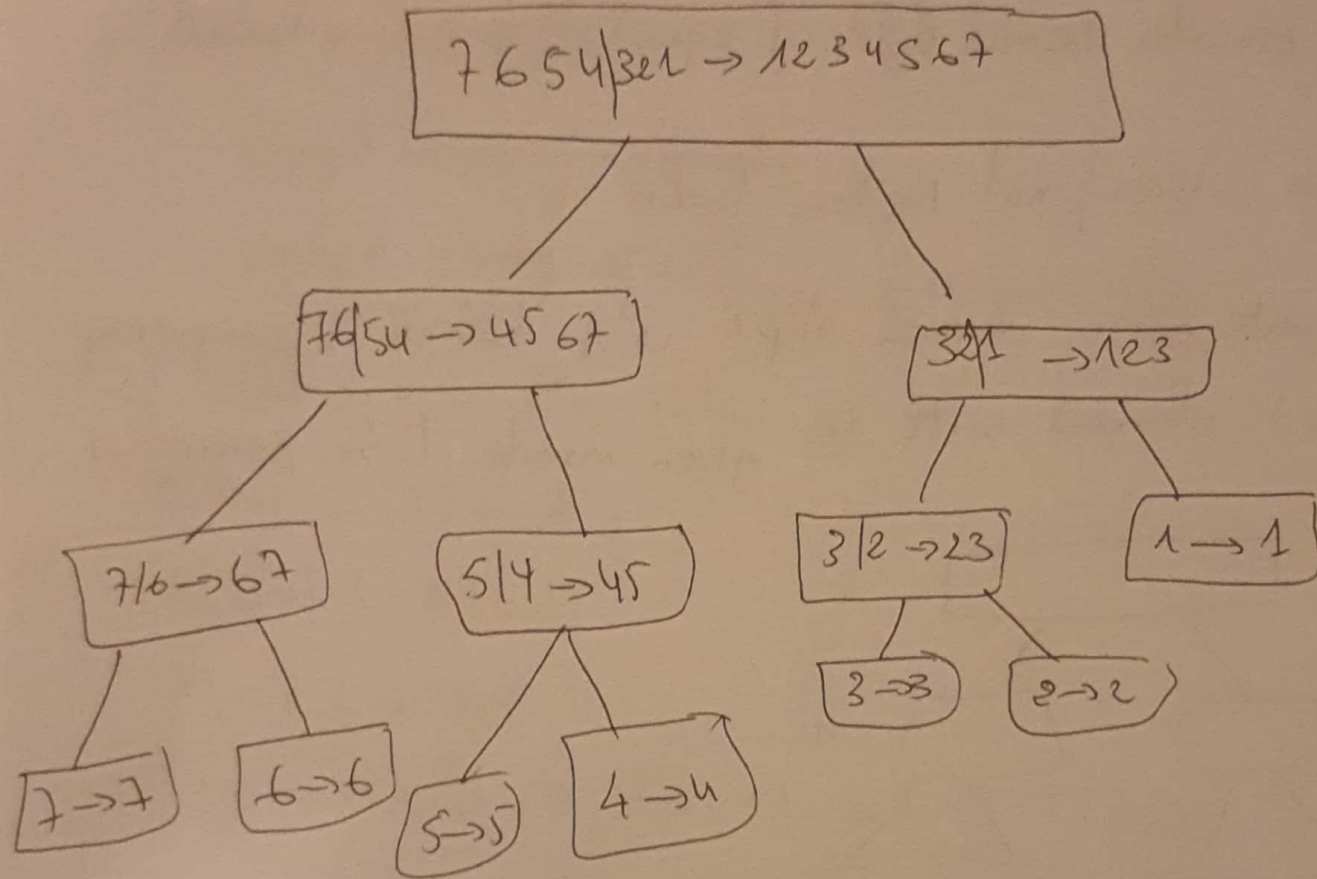


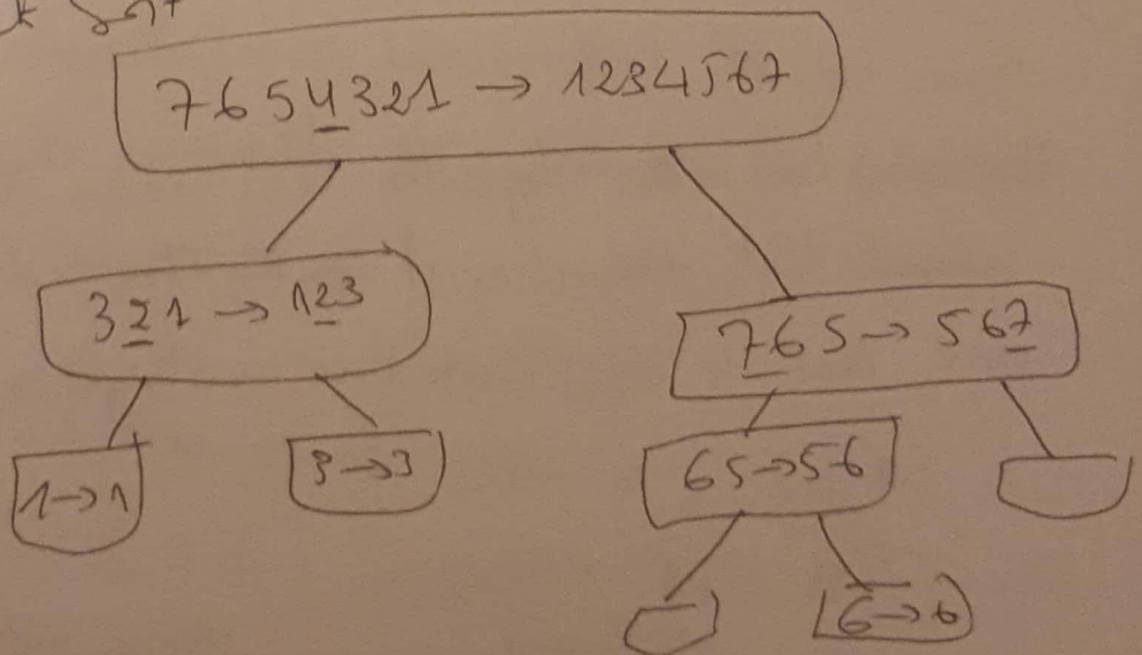
Lab 3B

④

1) using Merge Sort



2) using Quick Sort

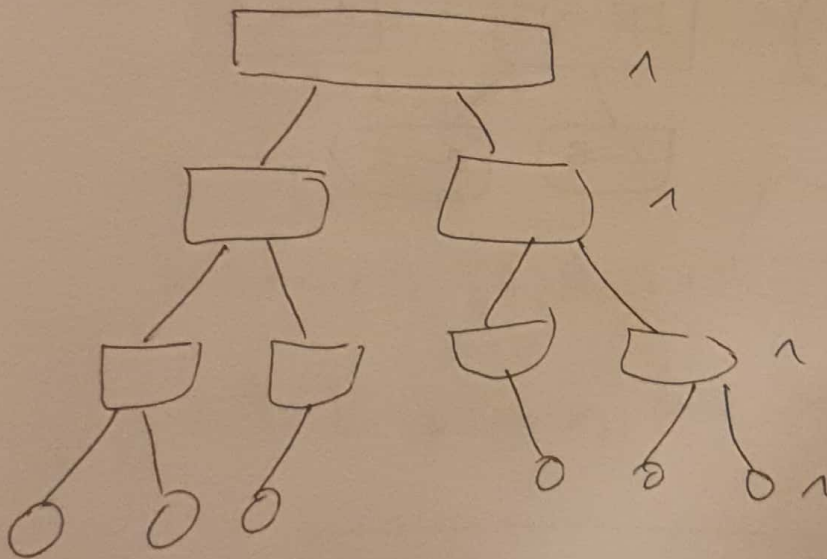


pb2: binary search algorithm only search at one side of the list

element is compared with the middle element if larger search right if smaller search left if equal found element

return true and if not return false.

note: work done at each step is 1 which is comparing the mid element with the given number to be searched.



the worst case in this algorithm is if we are looking for element that does not exist on the tree, at each level one there is $O(1)$ work done and height of the tree is

$(\log n)$

$$T(n) = O(1 \times \log n) \\ = O(\log n).$$

③

Problem 3: recursive algorithm to reverse order

Algorithm reverse order

Input: array of n elements

Output: array of n elements in reversed order

if $\text{start} > \text{end}$ then return $+c$

~~temp~~ $\leftarrow A[\text{start}]$

$A[\text{start}] = A[\text{end}]$

$A[\text{end}] \leftarrow \text{temp}$

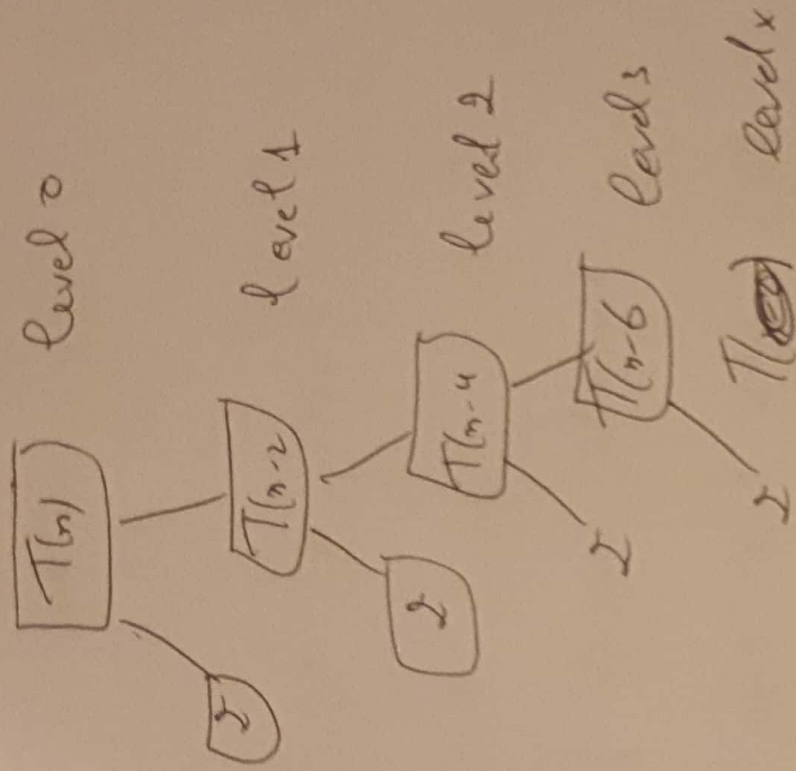
reverseOrder($A, \text{start}+1, \text{end}-1$) $-T(n-1)$

compute the running time using count self calls

in the given algorithm we can observe that the next method recursive call is smaller by 2, then the first one because $\text{start}+1$ and $\text{end}-1$

$$T(n) = \begin{cases} c & n=0 \\ T(n-2) + c & \text{otherwise} \end{cases}$$

because we don't have formula to resolve this we will use the tree method



problem 4. algorithm fib it (n)

Input : non negative integer (n)

output the fibonacci value at n^{th} sequence

fib ← 1

fib ← 1

for $i \in 2 \text{ to } n-1$ do

temp ← fib

pfib ← fib + pfib

pfib ← temp

return fib

$T(n) = 4(n) = O(n)$