SQLintersection

Thursday, 11.30 – 12.45

DBA Mythbusters (Level 2-300)

Paul S. Randal paul@SQLskills.com





This is me: Paul S. Randal

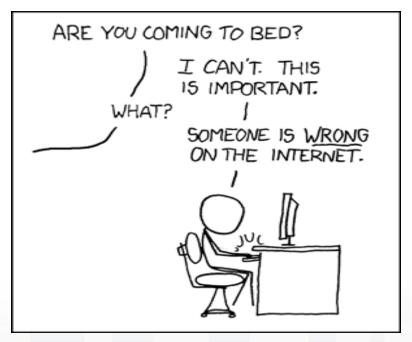
- Consultant/Trainer/Speaker/Author
- CEO, SQLskills.com
 - Email: Paul@SQLskills.com
 - Blog: http://www.SQLskills.com/blogs/Paul
 - Twitter: @PaulRandal
 - 5 years at DEC responsible for the VMS file-system and chkdsk
 - Almost 9 years as developer/manager in the SQL Storage Engine team through August 2007, ultimately responsible for Core Storage Engine
- Instructor-led training (US, UK, Australia), consulting (anything you need)
- Online training: pluralsight http://pluralsight.com/
- Become a SQLskills Insider: http://www.sqlskills.com/Insider







And This is My Life...



Source: http://xkcd.com/386/



Why Is This Important?

- Lots of myths and misconceptions have grown and persisted over the years about how SQL Server behaves
- Adherence to these misconceptions can lead to:
 - Bad practices
 - Wasted time and resources
 - Confusion
 - □ Arguments ☺
- Five years since I last did Mythbusters sessions and posts...
 - See http://www.sqlskills.com/blogs/paul/CommonSQLServerMyths.pdf
- Let's debunk some myths!



- Oracle is much better than SQL Server
- This one's obviously untrue! ◎
- On to the real ones...



- Page life expectancy should be around 300
- Page life expectancy measures (in seconds) how long a new page is expected to stay in the buffer pool
 - Also can be thought of as measure of memory pressure on the buffer pool
- Do you think that flushing your 100GB buffer pool every 5 minutes is a sign of a healthy SQL Server?
- That guidance is from 10 years ago!
- Correct PLE is whatever is normal for *your* system
 - If it drops and stays dropped, there's a problem





- Buffer Manager: Page life expectancy is the best counter to use
- If on a NUMA system, Buffer Manager PLE is the harmonic mean of PLEs from all Buffer Nodes
 - Buffer pool is split into partitions on NUMA system.
- Monitor all Buffer Node: Page life expectancy counters
- Example harmonic mean calculation
 - 4 buffer nodes with PLE = 4000, 4000, 4000, 2200
 - Buffer Manager PLE = 4 / (1/(1000x4000) + 1/(1000x4000) + 1/(1000x4000) + 1/(1000x3300)) / 1000 = 3321





- You can offload consistency checks to an availability group secondary
- The AG secondary is on a different I/O subsystem from the primary
- Running consistency checks on the secondary says nothing about the state of the primary
- You need to run consistency checks on the primary and ALL secondaries!
- Same argument holds for trying to offload consistency checks to a mirror, log shipping secondary, SAN mirror, etc. – you can't!



- AG readable secondaries don't cause performance problems
- All queries on a readable secondary are converted to snapshot isolation
- This means 14-byte versioning tags must be present on the secondary
- This means the tags have to be added on the primary
 - As the primary and secondaries are exact physical copies of each other
 - But the tags don't have to be filled in on the primary, just the space needs to be accounted for
- All changes to the primary database will incur a versioning tag which will start to cause page splits and index fragmentation!





- Automatic page repair can instantly fix broken pages
- Applies to database mirroring and availability groups
- The secondary/mirror cannot send the requested page back to the primary unless the page is known to be at the right point in time
 - What is the LSN in the primary when it requested the page image?
 - Secondary redo queue must be replayed to that LSN
 - Otherwise the secondary may be sending back the wrong version of the page





- Using temp tables for intermediate query results is always a good idea
- Creating a temp table to hold intermediate results forces SQL Server to interrupt the data pipeline through a query to persist the results to disk
- Sometimes just doing one query rather than pre-aggregating or presorting can be way more efficient and lead to far lower run time and tempdb usage
- Always compare the methods before production
- And if using temp tables, use minimum amount of data and correct indexes



- "Fully logged" means you'll always see one log record for each part of an operation
- Consider a rebuild of a 100,000 row index
 - You would expect to see 100 thousand LOP_INSERT_ROW log records, right?
 - Wrong it will log about LOP_FORMAT_PAGE log records instead with full page images with the net effect of all the inserts on
- "Fully logged" simply means the transaction log contains enough information to reconstitute the transaction after a crash or restore
- What about TRUNCATE TABLE?





- NOLOCK / READ UNCOMMITTED means no locks
- First off, they're the same thing
- And they do have to acquire some locks:
 - Schema-stability locks (Sch-S) to prevent the structure of the table/index changing
 - BULK_OPERATION locks on heaps to prevent reading of unformatted pages
- And they still have to take latches to access the physical page images in memory, so there's still some potential for blocking at the latch level



- You should always plan a backup strategy
- Always plan a *restore* strategy
- Then plan what backups you need to take
- The other way can result in disaster
- Let me tell you a story...





- The best thing to put on SSDs are always tempdb and transaction logs
- Don't fall into the trap of listening to other people
- Investigate where your biggest I/O subsystem bottleneck is
 - Try to solve it within SQL Server
 - □ If not, put that on your SSD
- Or design a new I/O subsystem layout to take advantage of the SSD
- What about the RAID level to use?



- Using SSDs means you don't have to care about index fragmentation
- Index fragmentation has two forms:
 - Logical fragmentation that stops efficient readahead
 - Low page density that wastes space
- SSDs make reads faster, but still a trip down/up I/O stack for each one
- SSDs don't stop page splits from happening
 - Lots of extra transaction log
- SSDs don't stop low page density from happening
 - Wasted disk space, wasted buffer pool memory



- Adding more memory is always a good idea
- Consider some of the potential problems
 - Shutting down the instance will take longer
 - P.O.S.T. of the server will take longer
 - Allocating buffer pool memory may take longer (see KB article 2819662)
 - Warming up the buffer pool will take longer
 - Could lead to complacency



- Shrinking tempdb can cause corruption
- KB article 307487 was updated in 2014
- No problems with shrinking tempdb since SQL Server 2000
- However, just because you can, doesn't mean you should...
- And be aware of what happens to tempdb size on a server restart



- DBCC CHECKDB runs when SQL Server starts up
- The messages in the error log are confusing:
 - 2015-05-15 13:16:20.07 spid7s CHECKDB for database 'master' finished without errors on 2015-05-01 09:59:42.447 (local time). This is an informational message only; no user action is required.
- This is just reporting the time that DBCC CHECKDB last completed without finding any errors
 - Stored in the boot page of the database (file 1, page 9)
 - Check with DBCC TRACEON (3604); DBCC DBINFO;





- Rebuilding indexes solves performance problems even when there's no index fragmentation
- It's the query plan recompilation that 'fixes' the performance problem
- Rebuilding an index causes plan recompilation for plans on that table
- If a poor query plan had resulted (e.g. from parameter sniffing), the next plan to be compiled might be better
- Also, updating statistics in 2012+ doesn't invalidate plans if no table rows changed (this is a good thing!)





- Adding an extra file to tempdb will help solve contention issues
- Adding an extra file means SQL Server can alternate between the files
- But allocation also takes into account proportional fill
 - It will allocate proportionally more from files with more free space
- If the existing file is quite full, the new file becomes allocation hot spot
 - No alleviation of contention issues!
- Make sure to take that into account when working with tempdb





- Lots of OLEDB waits always means linked-server problems
- OLEDB waits mean that the OLE-DB protocol is being used
 - OLE-DB is not just used by linked servers
- How long are the waits?
 - \Box 0 1-2ms = not linked servers
 - □ 10s or 100s of ms are likely to be linked servers





- ASYNC_NETWORK_IO waits always means network problems
- Rare for it to be the network
- The word NETWORK is horribly misleading in the name
- More likely the application doing RBAR processing





- CXPACKET waits mean disable parallelism
- Check you expect parallelism for that query
- Check you don't have skewed parallelism
- Consider altering MAXDOP for query
- Consider setting server MAXDOP (8 or number of cores in a NUMA node)
- Consider using Resource Governor
- Better: increase 'cost threshold for parallelism'





- Checkpoints only write committed changes to disk
- Checkpoint writes all pages marked dirty regardless of whether the change was made by a committed or uncommitted transaction
 - Crash recovery takes care of fixing things up if there's a crash
- Use sys.dm_os_buffer_descriptors to examine the relative proportion of dirty vs. clean pages in the buffer pool





- DBCC DROPCLEANBUFFERS flushes the buffer pool
- Look at the name carefully drop clean buffers
- It doesn't flush dirty pages
- You have to do a checkpoint for that, and then DROPCLEANBUFFERS





- Tempdb data files should be 1:1 with processor cores
- SQL Server 2000: rule was #files = #logical processor cores, and TF 1118
 - □ E.g. my laptop CPU has 4 physical cores plus hyperthreading = 8 logical cores
- SQL Server 2005 onwards: Microsoft guidance was same until 2011
 - □ Everyone else said to start with ¼ to ½ the number of logical processor cores
- Universal guidance now in KB article 2154845
 - < 8 cores, start with #files = #cores</p>
 - □ > 8 cores, start with #files = 8
 - Increase in blocks of 4 if still seeing contention





- Multiple log files will help performance
- SQL Server will always use log files sequentially
- You may see them all having I/Os, but that's just updating the file header pages
- The only time another log file is needed is if the first one fills up and cannot grow, you cannot take a log backup, and you do not want to break the log backup chain
- Remove additional log files once you don't need them





- The log should always be as small as possible
- The log needs to be as big as it needs to be
- Do not regularly shrink the transaction log
 - It'll just have to grow again, and can't use instant initialization
- How big should the log be?
 - Single largest transaction (ETL, large index rebuild, large update)
 - Asynchronous database mirroring/AG SEND queue
 - How long is the longest data backup?
 - Transactional replication (beware of CDC too)



- It depends
- The answer to all questions about SQL Server that do not have obvious yes/no answers always starts with 'it depends'

YES, IT REALLY DOES!!

- The trick is then to explain *why* it to pe ds *v are it explain to explain *when* it de to ds
 - n no control of the property o
- One exception: should auto-shrink be enabled?



Plenty More Myths Around...

Repair

- It can fix everything
- Safe to repair system databases
- SQL Server causes corruptions
- Corruptions can disappear

Performance

- You can't override MAXDOP
- Always use data compression
- Nested transactions exist

The transaction log

- Log records can move
- The log is zeroed when cleared
- BULK LOGGED lowers backup size

High availability

- Just use a cluster
- Replication isn't an HA solution





Summary and Resources

- Make sure you corroborate what you read online
- If something sounds fishy, try it yourself!
- Blog:
 - http://www.sqlskills.com/blogs/paul/category/misconceptions/
- Pluralsight
 - SQL Server: Myths and Misconceptions



Questions?

Don't forget to complete an online evaluation on EventBoard!

DBA Mythbusters

Your evaluation helps organizers build better conferences and helps speakers improve their sessions.



Thank you!