

# Optimal flow problems, the Edmonds-Karp algorithm, and birdsong

Dan Stowell

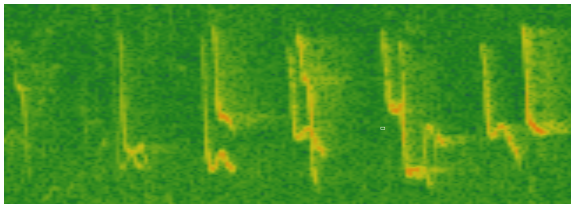
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Queen Mary, University of London

Oct 2014, Big-O

# Intro

My research:

- ▶ “Machine listening”  
= machine learning & signal processing  
for understanding sound
- ▶ Automatic analysis of bird sounds  
(which species, how many, how often...)



## Algorithm: Edmonds-Karp

- ▶ Graph theory, “flow networks”
- ▶ We'll consider the *maximum flow* problem
- ▶ Related to shortest-path problems  
(SatNav routing, travelling salesman)  
... but harder
- ▶ Imagine you have to send as many trains-per-hour as possible from London to Edinburgh, and that different segments of the route have different maximum capacities.

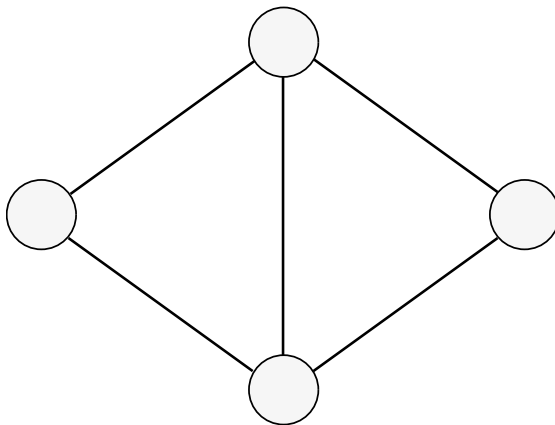
# Intro

What does that have to do with birdsong?



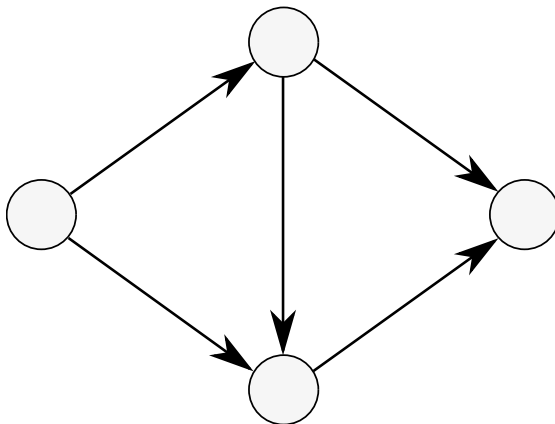
# Definitions

Undirected graph



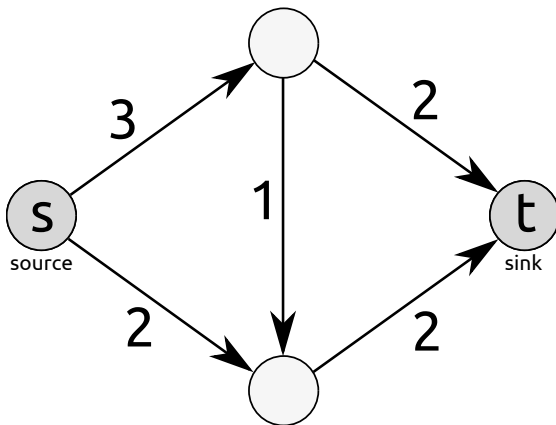
# Definitions

Directed graph (or “digraph”)



# Definitions

Network

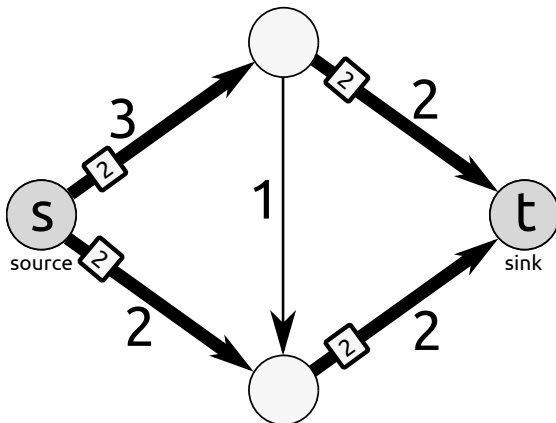


Capacities on each arc



# Definitions

## Flow



Flow values must balance at each node (except s, t)

# Maximum flow

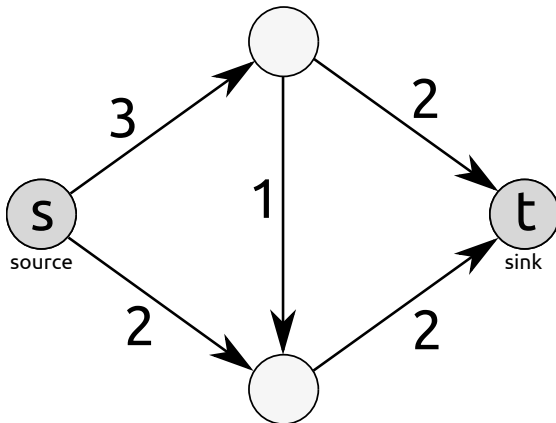
The **maximum flow** problem:

*For a given network,  
find the flow of maximum value  
that is compatible with the network capacities  
(a flow that is “feasible”)*

Flow might be integer or real-valued; I'll focus on integer.

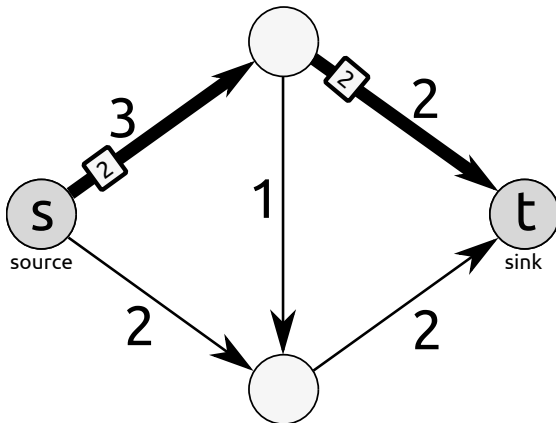
# Maximum flow

Greedy approach: use a single-path algorithm, and iterate



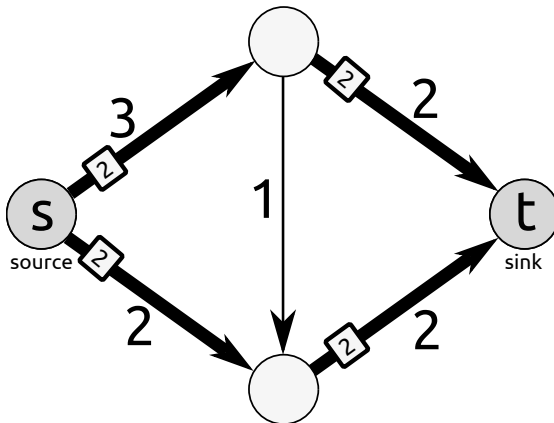
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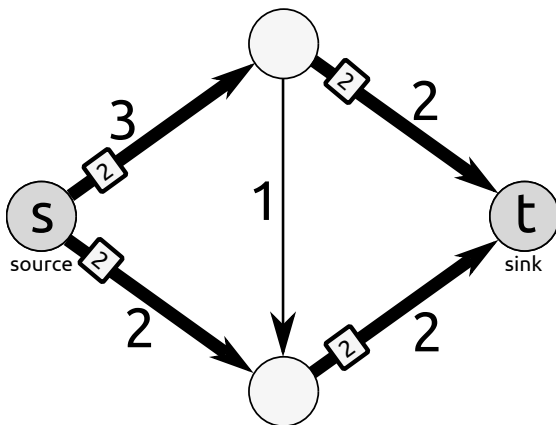
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# Maximum flow

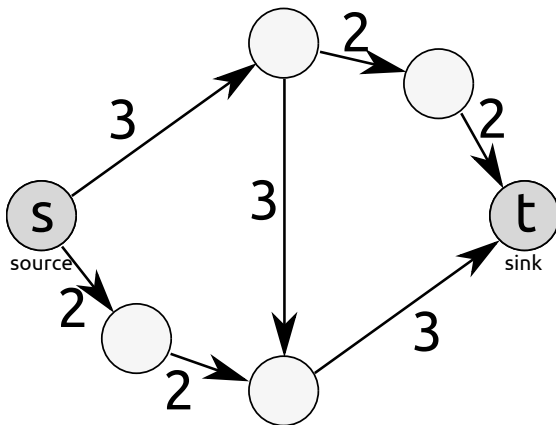
Greedy approach: use a single-path algorithm, and iterate



Complexity:  $O(|V| + |A|)$  per iteration

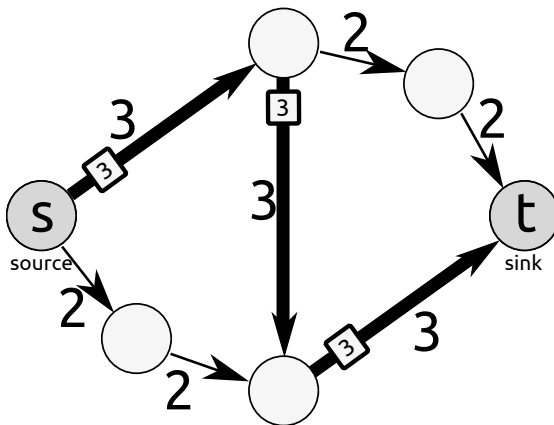
## Greedy approach may fail

Many networks lead to local optima



## Greedy approach may fail

Many networks lead to local optima





# Edmonds-Karp algorithm

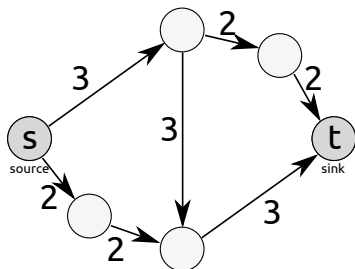
The **Edmonds-Karp algorithm** is guaranteed to find the exact solution.

Based on the Ford-Fulkerson algorithm (plus shortest-path).

Key concept: the *residual network*.

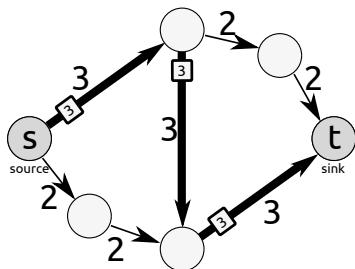
# Edmonds-Karp algorithm

Network



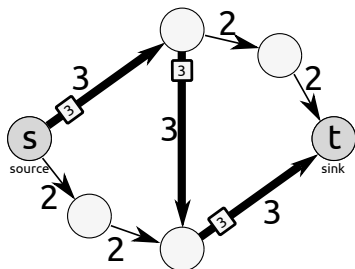
# Edmonds-Karp algorithm

Network

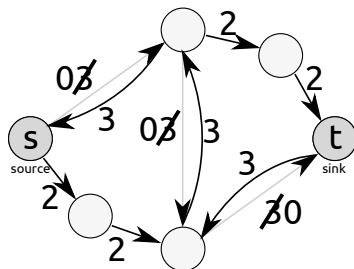


# Edmonds-Karp algorithm

Network

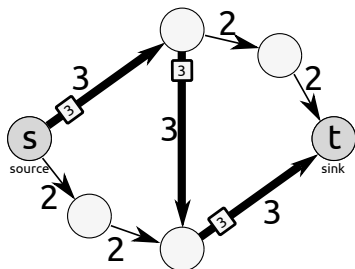


Residual network

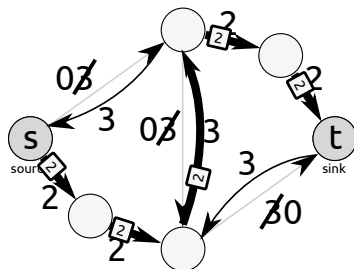


# Edmonds-Karp algorithm

Network

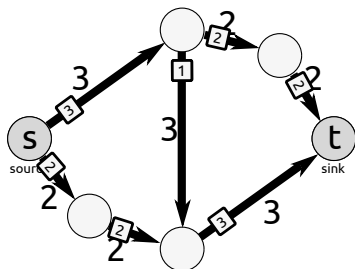


Residual network



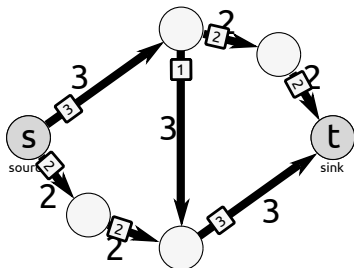
# Edmonds-Karp algorithm

Network

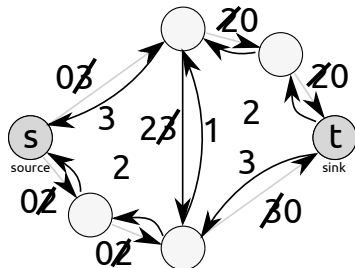


# Edmonds-Karp algorithm

Network



Residual network



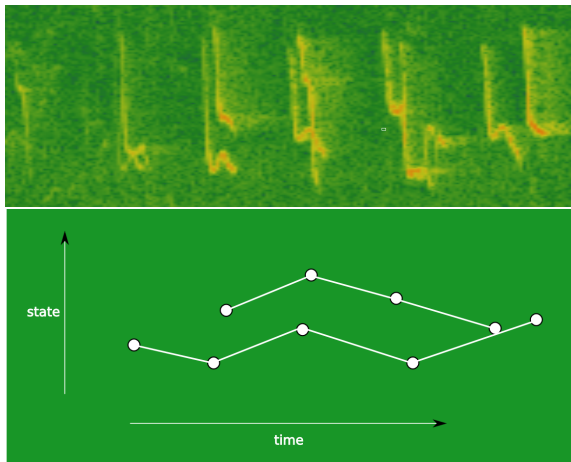
# Edmonds-Karp algorithm

- ▶ Guaranteed to have no local optima
- ▶ Guaranteed to terminate in bounded number of iterations (as long as all capacities are rational)
- ▶ Complexity:  $O(|V||A|^2)$

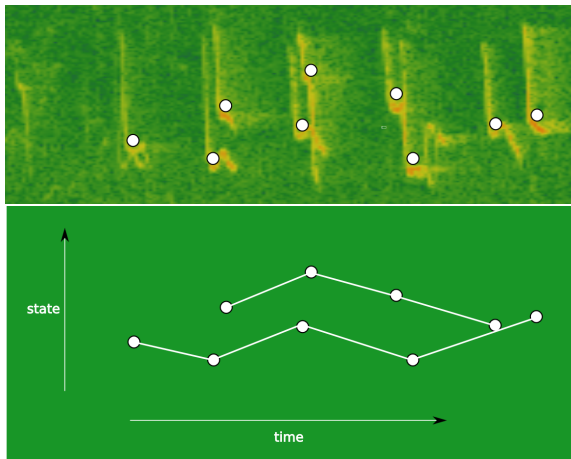


# Pause

# Disentagling birds



# Disentagling birds

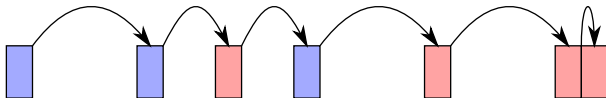


# Modelling an intermittent source

Markov renewal process (“MRP”):

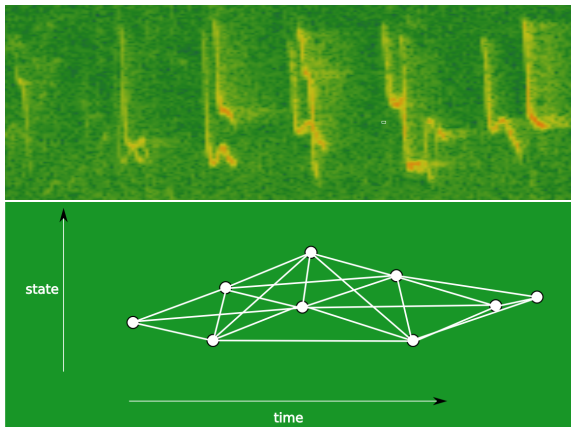
$$\begin{aligned} P(\tau_{n+1} \leq t, X_{n+1} = j \mid (X_1, T_1), \dots, (X_n = i, T_n)) \\ = P(\tau_{n+1} \leq t, X_{n+1} = j \mid X_n = i) \end{aligned}$$

where  $\tau_{n+1}$  is the time difference  $T_{n+1} - T_n$ .



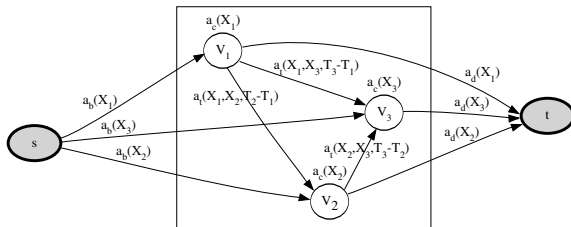
# Multiple MRPs

Problem sketch: assume multiple MRPs, plus potential “clutter”.



Given transition probabilities, find the most likely set of paths. (Max 1 path per node)

# Convert max-likelihood problem to a flow problem



Convert likelihood expression to flow “costs”:

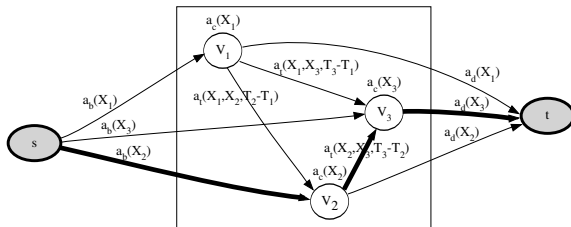
$$a_b(X) = -\log p_b(X)$$

$$a_d(X) = -\log p_d(X)$$

$$a_t(X, X', \tau) = -\log f_X(X', \tau)$$

$$a_c(X) = \log p_c(X)$$

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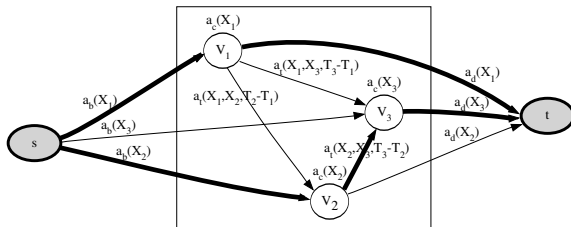
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# Minimum cost flow

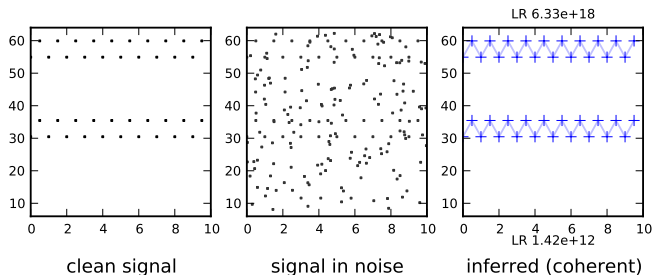
**Minimum cost flow** problem, with binary capacities

is solved the same way as a

**Maximum flow** problem

- ▶ Optimal minimum-cost flow: Edmonds-Karp algorithm, asymptotic time complexity  $O(|V||A|^2)$ .
- ▶ Or use inexact (greedy) algorithm:  $O(|V||A|)$  or lower.

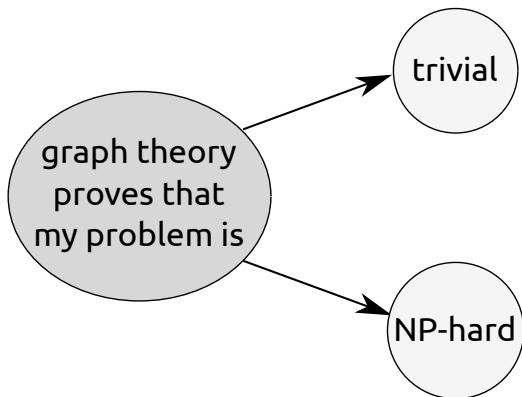
# Synthetic example



# Lesson

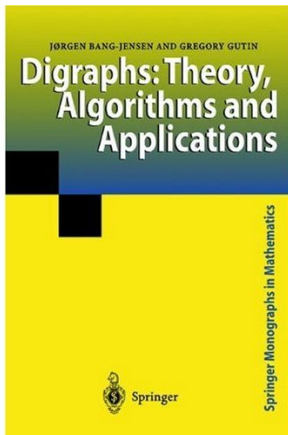
# Lesson

If your problem can be represented on a graph/network, you can probably transform it into a well-studied problem.



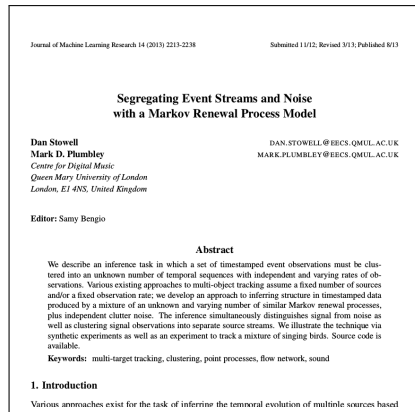
## Further reading

A good textbook:  
Bang-Jensen and Gutin



(free PDF on website)

Journal paper: JMLR,  
Stowell & Plumbley (2013)



(free, open-access) plus...



Funded **PhD position** available next year!

Working with me at QMUL  
on machine learning, signal processing and bird sounds

Drop me a line:  
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