# Concurrent Skip Lists

Companion slides for
The Art of Multiprocessor
Programming
by Maurice Herlihy & Nir Shavit

# Set Object Interface

- Collection of elements
- No duplicates
- Methods
  - add() a new element
  - remove() an element
  - contains() if element is present



# Many are Cold but Few are Frozen

- Typically high % of contains() calls
- Many fewer add() calls
- And even fewer remove() calls
  - 90% contains()
  - 9% add()
  - 1% remove()
- Folklore?
  - Yes but probably mostly true



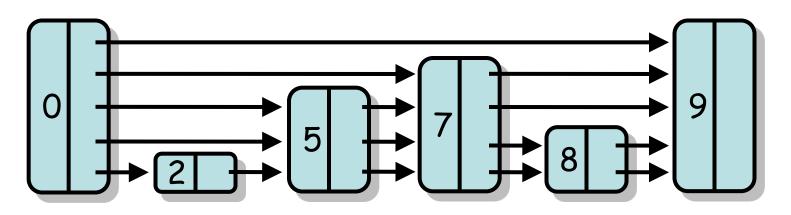
#### Concurrent Sets

- Balanced Trees?
  - Red-Black trees, AVL trees, ...
- Problem: no one does this well ...
- ... because rebalancing after add() or remove() is a global operation

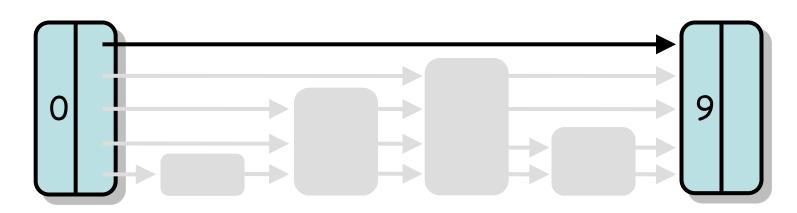


# Skip Lists

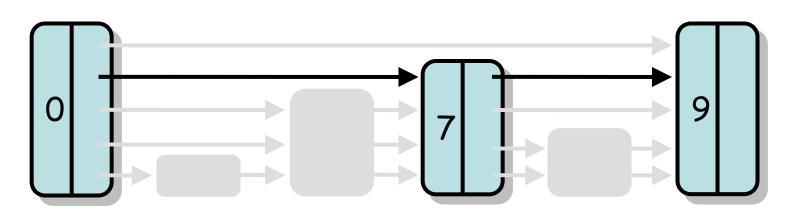
- Probabilistic Data Structure
- No global rebalancing
- · Logarithmic-time search



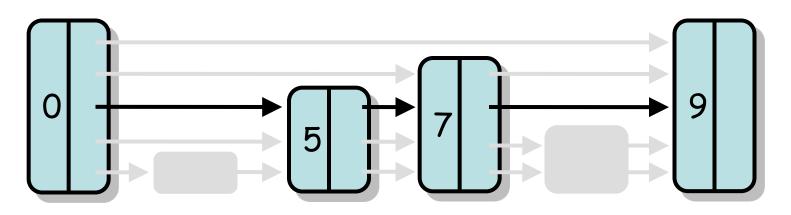




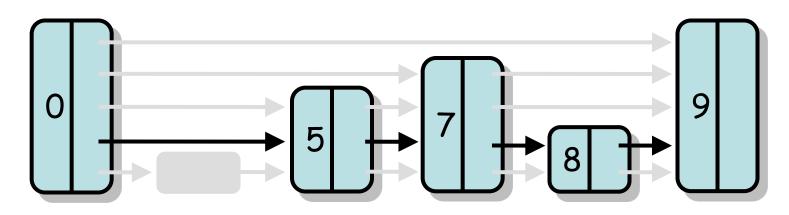






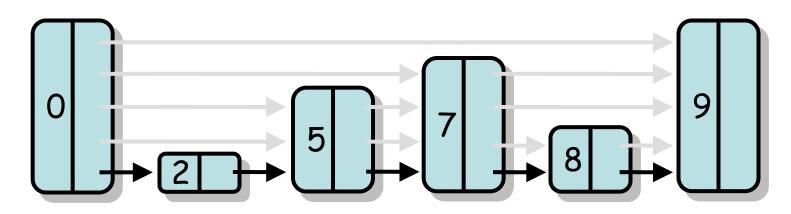






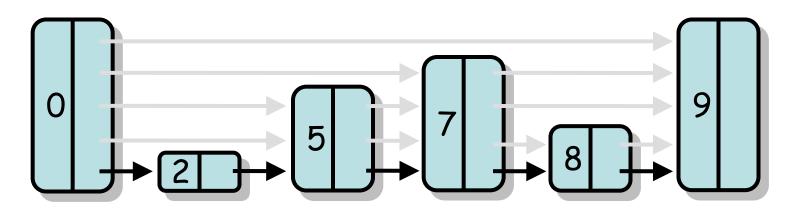


- Each layer is sublist of lower-levels
- · Lowest level is entire list

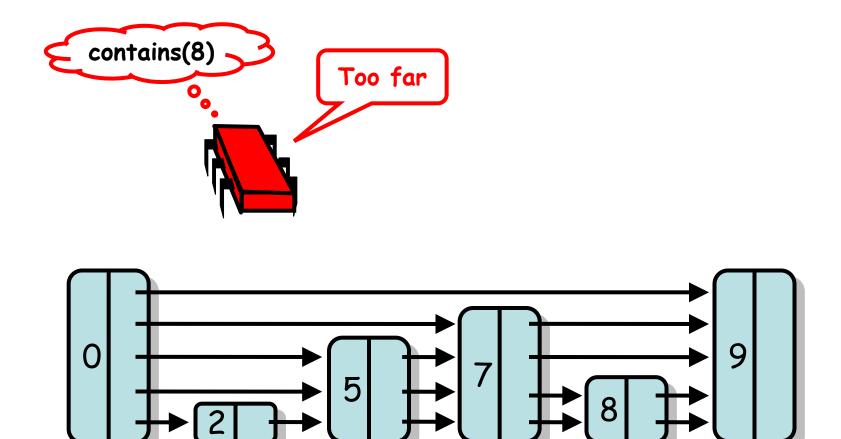




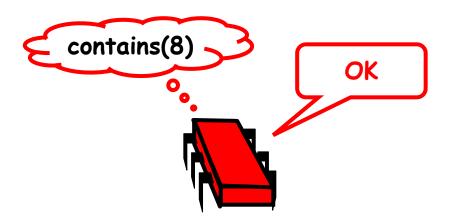
- Each layer is sublist of lower-levels
- Not easy to preserve in concurrent implementations ...

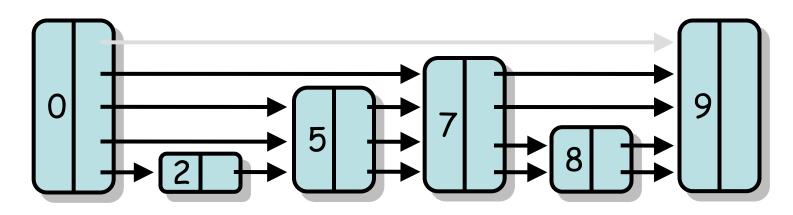




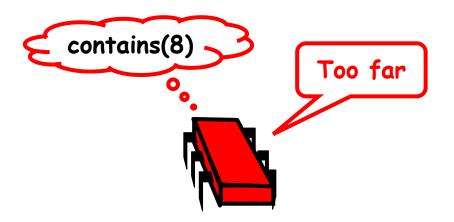


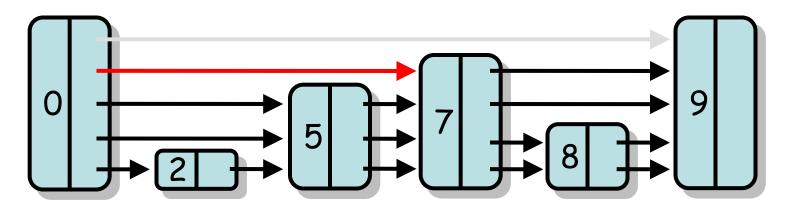




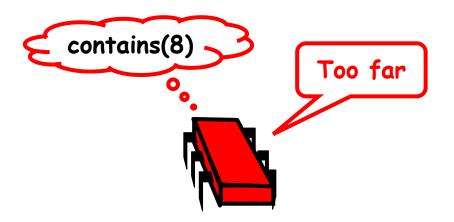


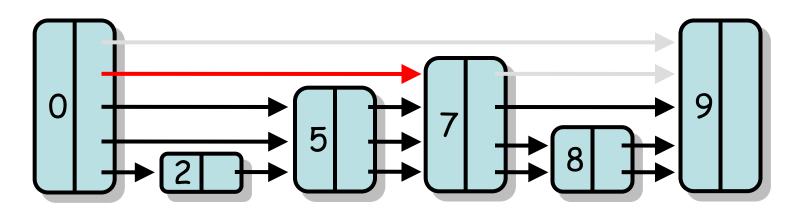




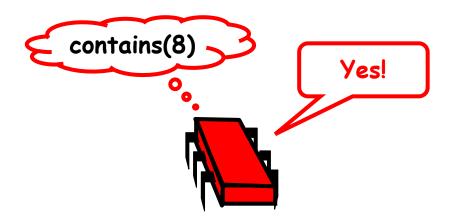


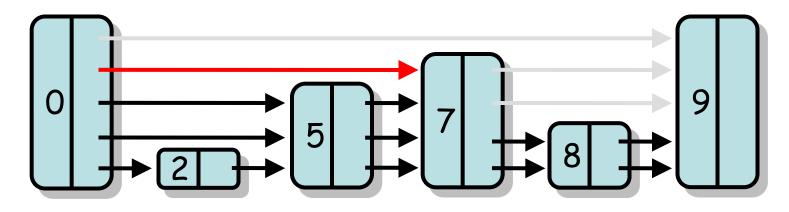




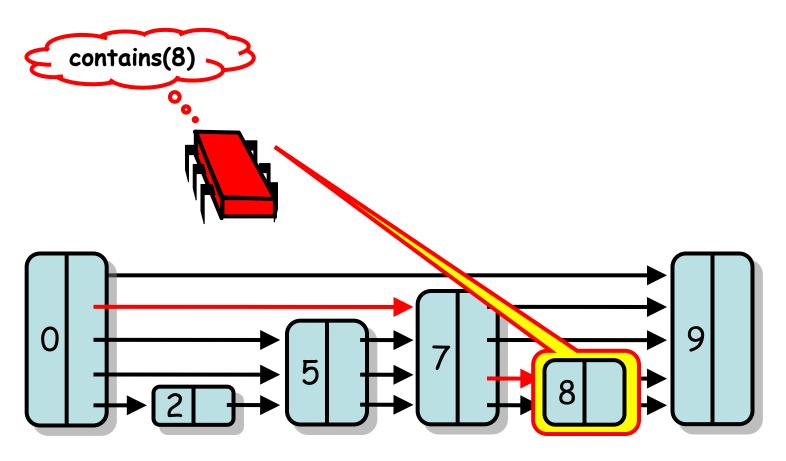






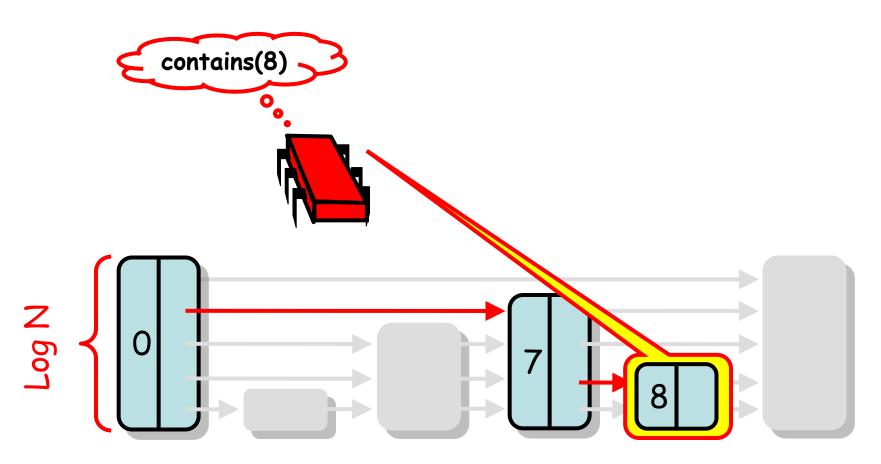








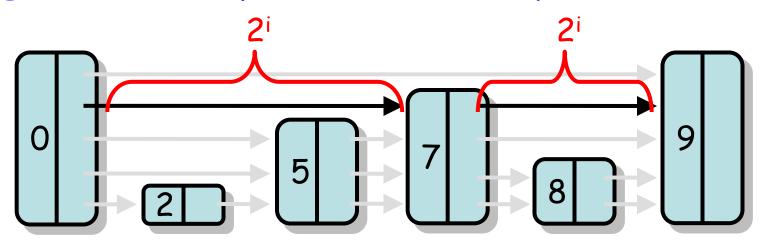
# Logarithmic





# Why Logarthimic

- Property: Each pointer at layer i jumps over roughly 2<sup>i</sup> nodes
- Pick node heights randomly so property guaranteed probabilistically





```
int find(T x, Node<T>[] preds, Node<T>[] succs) {
    ...
}
```

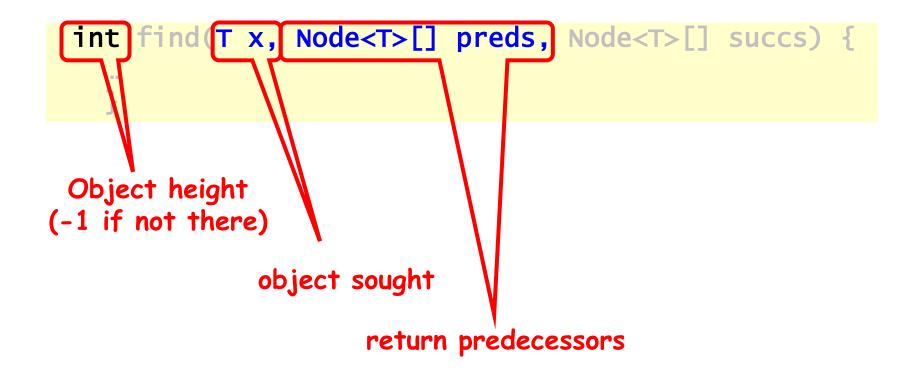


```
int find(T x, Node<T>[] preds, Node<T>[] succs) {
  object height
(-1 if not there)
```

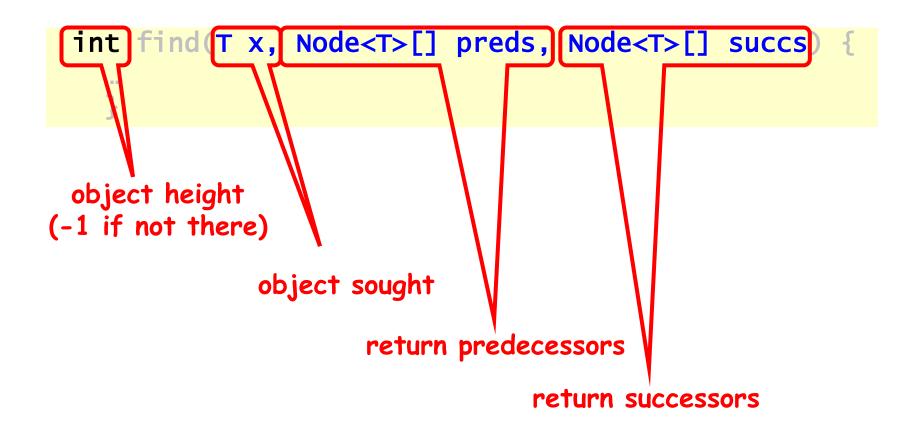


```
int find(T x, Node<T>[] preds, Node<T>[] succs) {
  object height
(-1 if not there)
  Object sought
```



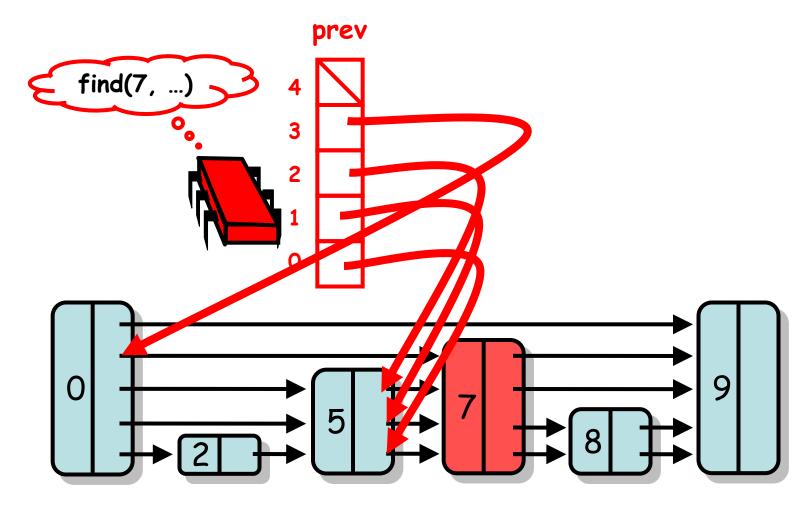






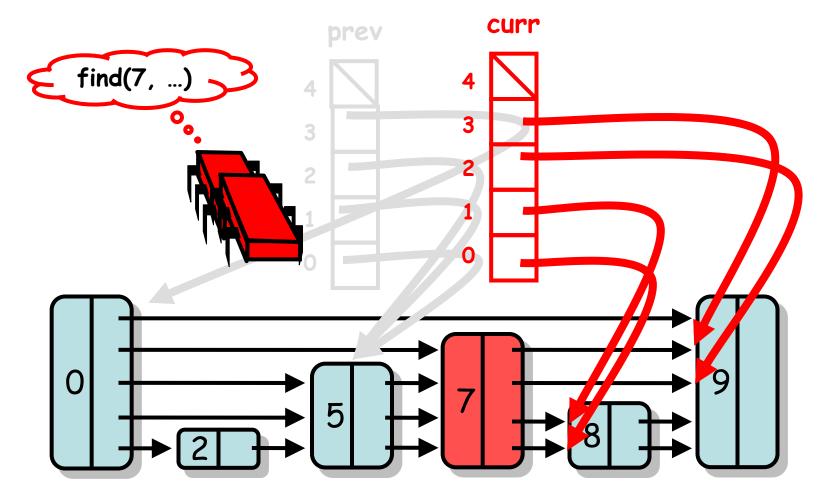


#### Successful Search



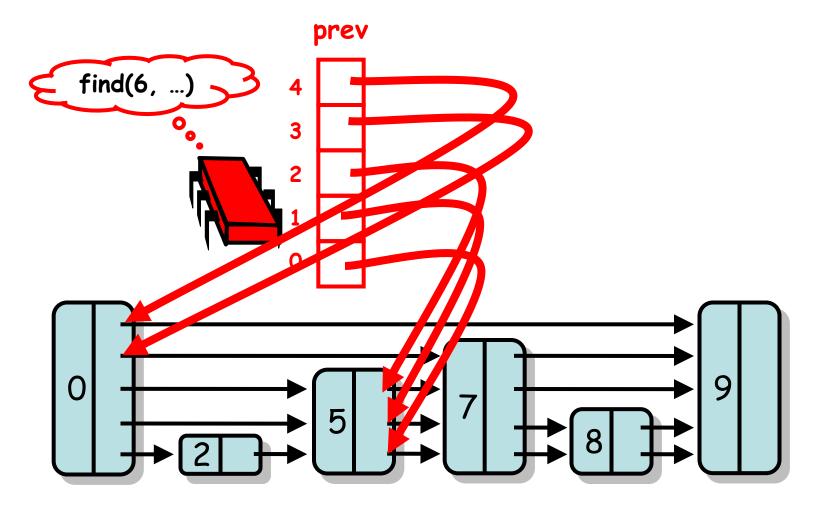


### Successful Search



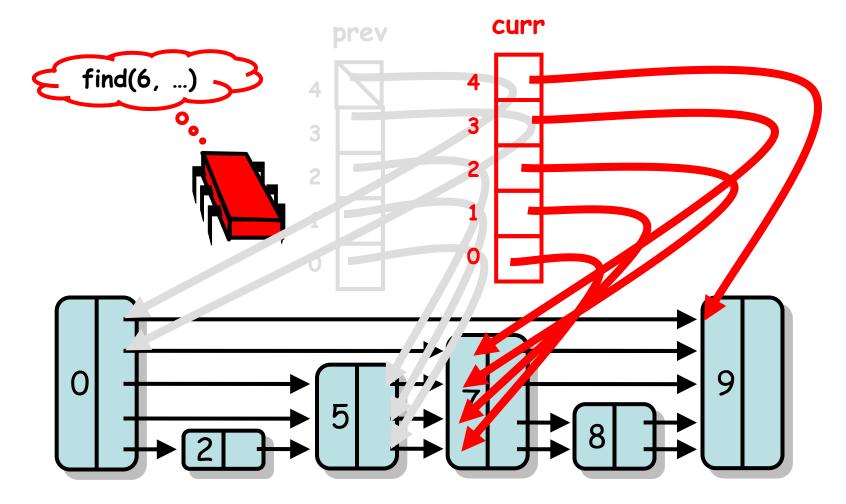


#### Unsuccessful Search





#### Unsuccessful Search

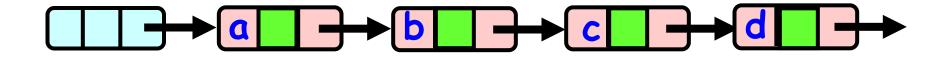




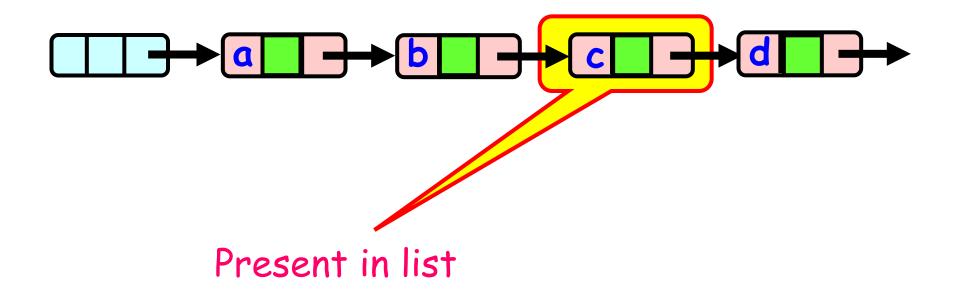
### Lazy Skip List

- Mix blocking and non-blocking techniques:
  - Use optimistic-lazy locking for add() and remove()
  - Wait-free contains()
- Remember: typically lots of contains() calls but few add() and remove()

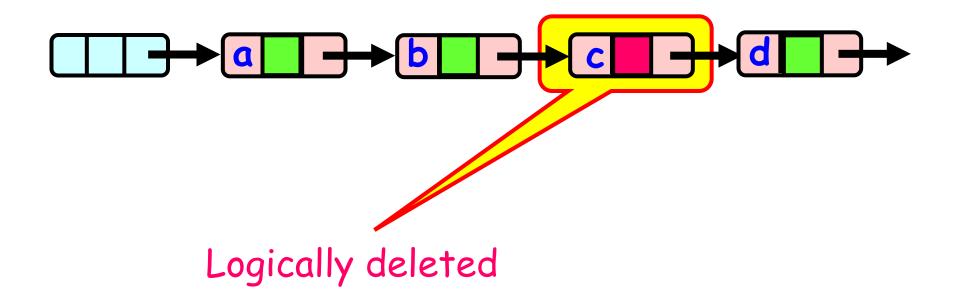




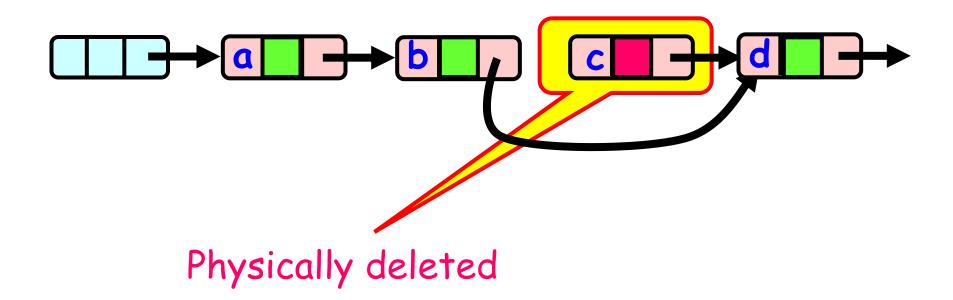








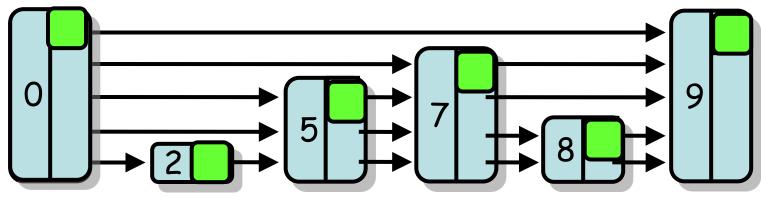






# Lazy Skip Lists

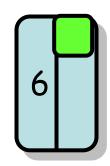
Use a mark bit for logical deletion

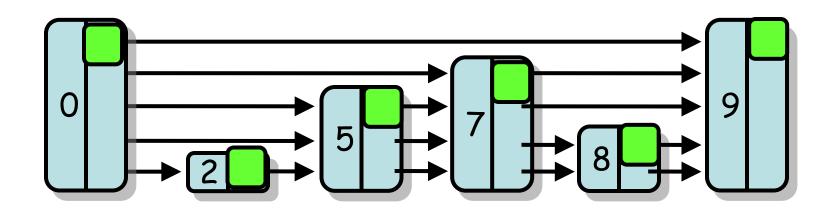




### add(6)

· Create node of (random) height 4

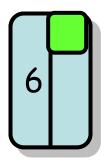


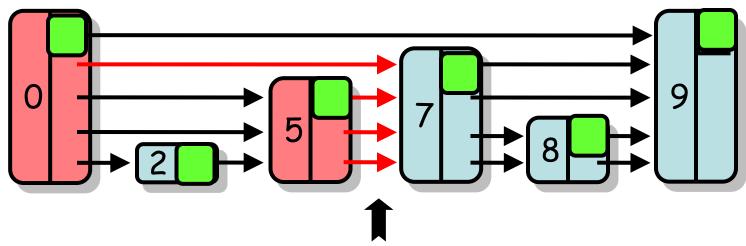




### add(6)

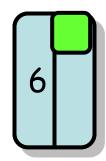
find() predecessors

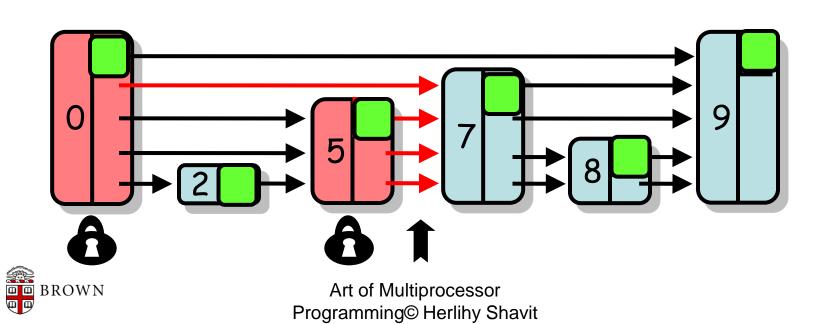






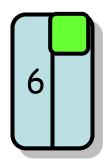
- find() predecessors
- · Lock them

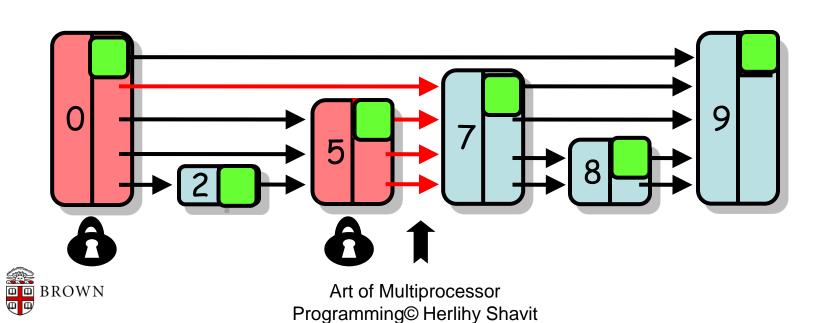




2007

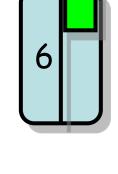
- find() predecessors
- · Lock them } Optimistic approach
- Validate

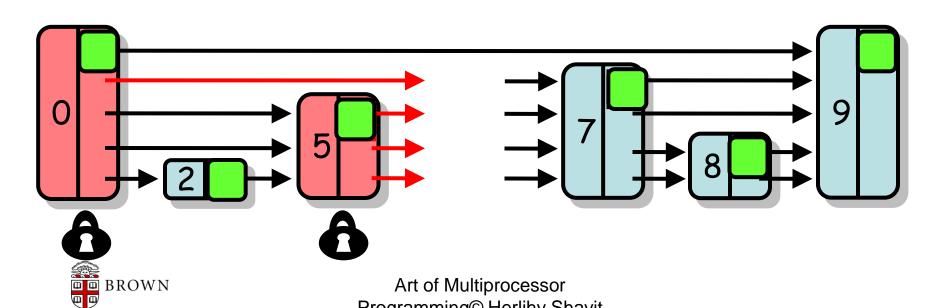




2007

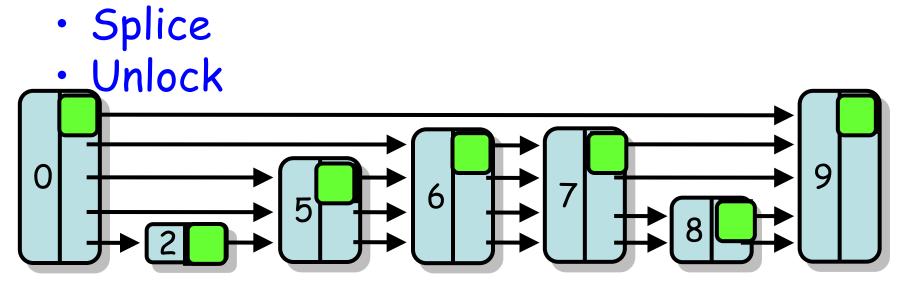
- find() predecessors
- · Lock them
- Validate
- Splice



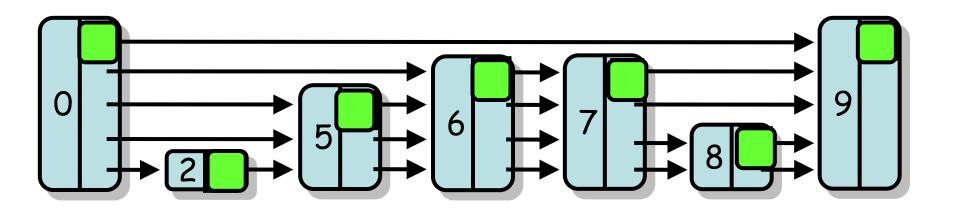


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- find() predecessors
- · Lock them
- Validate

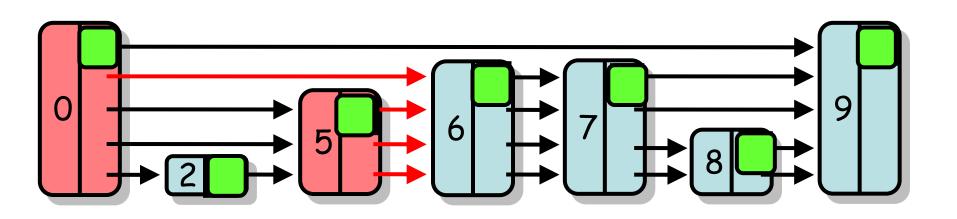






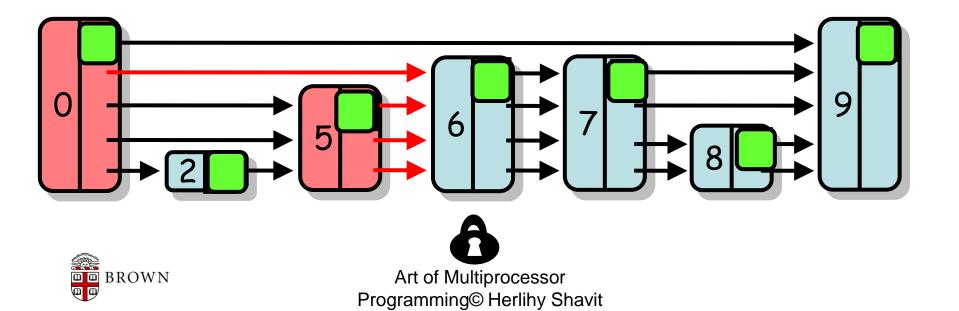


find() predecessors





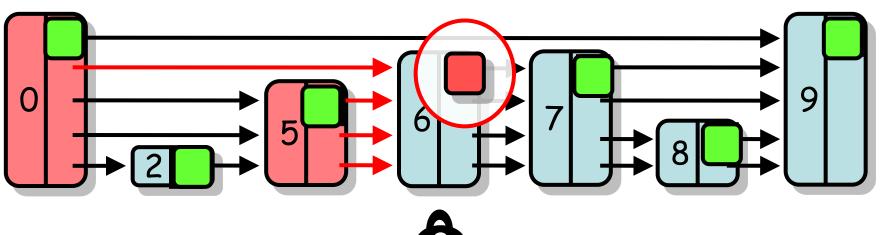
- find() predecessors
- Lock victim



2007

- find() predecessors
- Lock victim
- Set mark (if not already set)

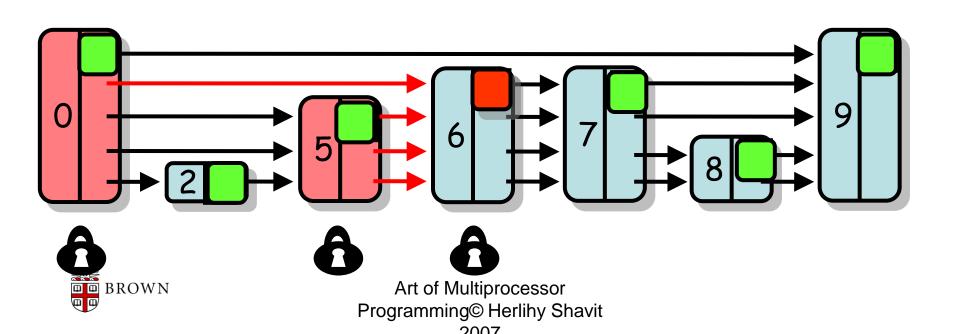
#### Logical remove...



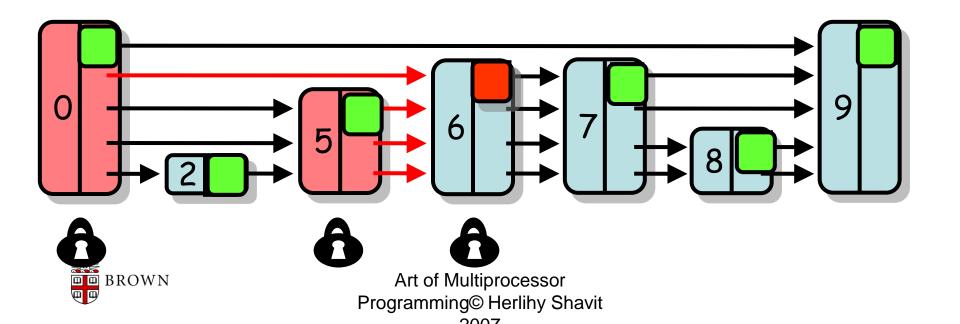


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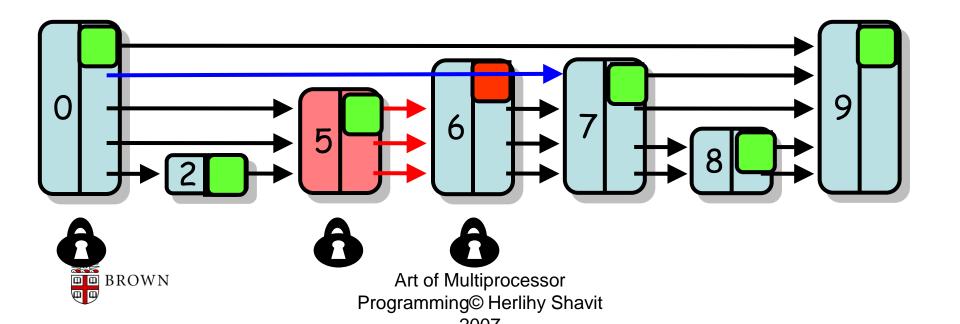
- find() predecessors
- Lock victim
- Set mark (if not already set)
- · Lock predecessors (ascending order) & validate



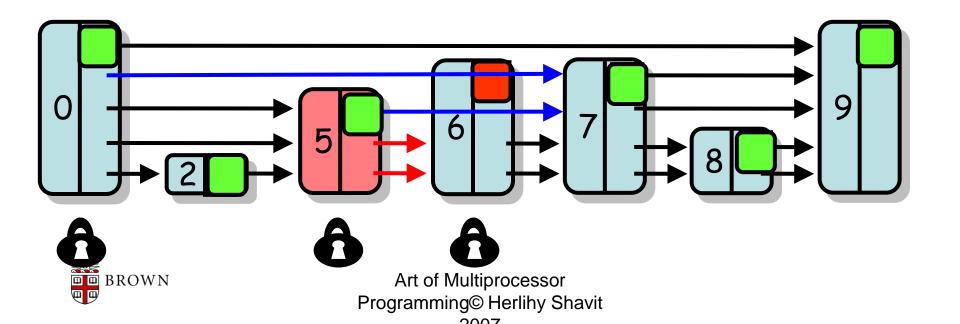
- find() predecessors
- Lock victim
- Set mark (if not already set)
- · Lock predecessors (ascending order) & validate
- · Physically remove



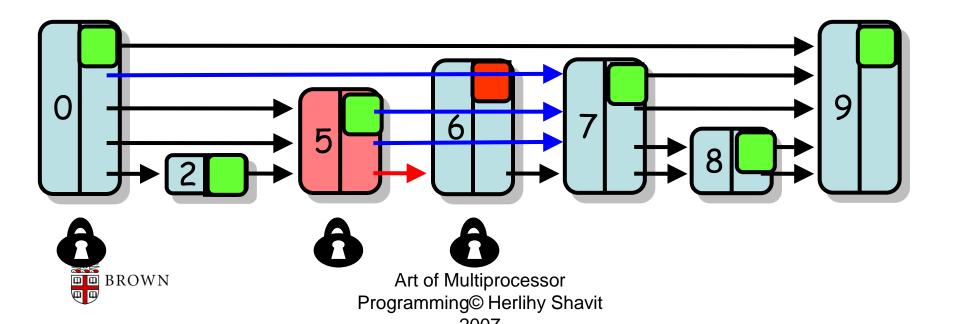
- find() predecessors
- Lock victim
- Set mark (if not already set)
- · Lock predecessors (ascending order) & validate
- · Physically remove



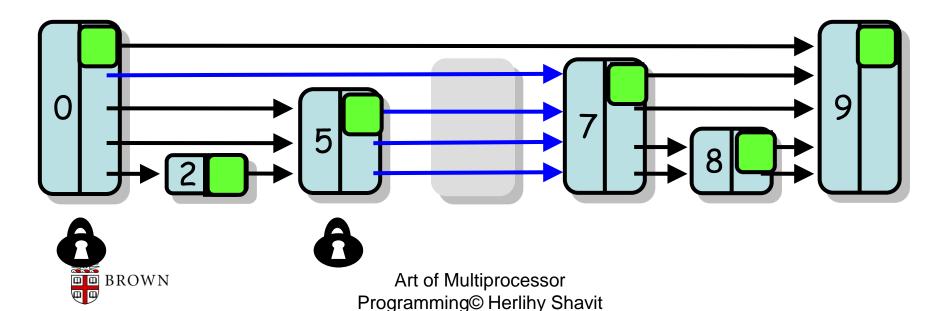
- find() predecessors
- Lock victim
- Set mark (if not already set)
- · Lock predecessors (ascending order) & validate
- · Physically remove



- find() predecessors
- Lock victim
- Set mark (if not already set)
- · Lock predecessors (ascending order) & validate
- · Physically remove

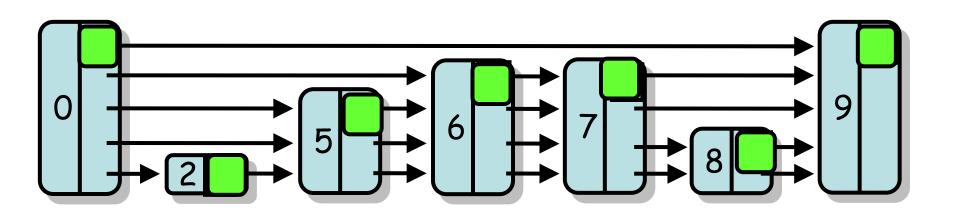


- find() predecessors
- Lock victim
- Set mark (if not already set)
- · Lock predecessors (ascending order) & validate
- · Physically remove



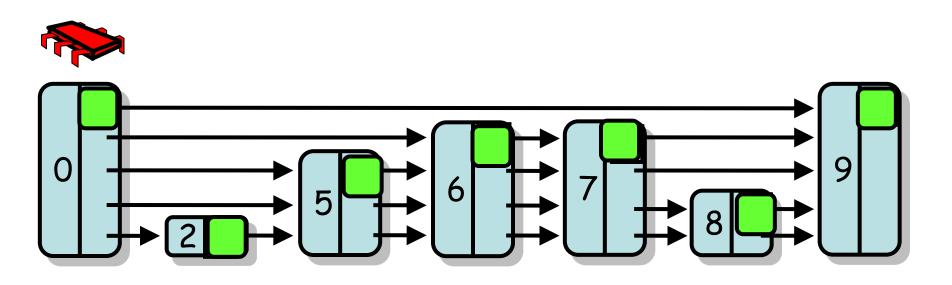
2007

Find() & not marked

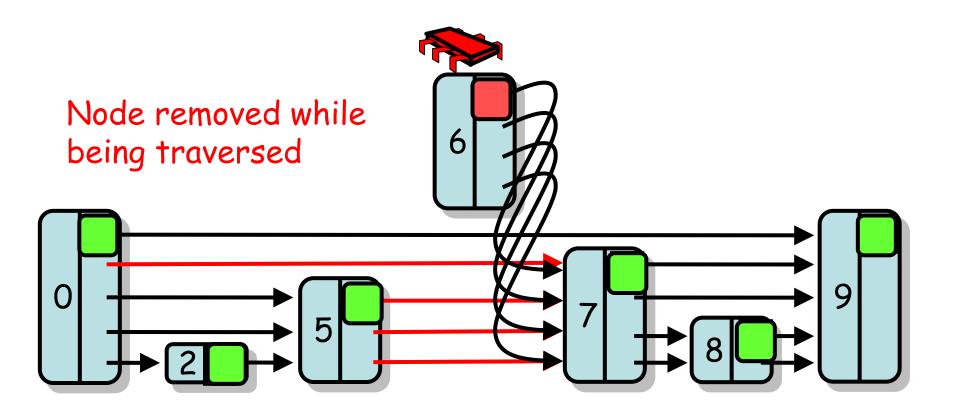




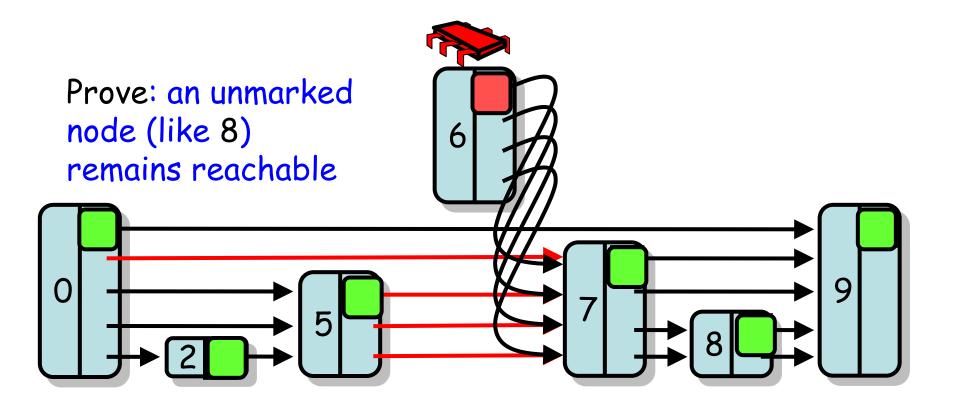
#### Node 6 removed while traversed









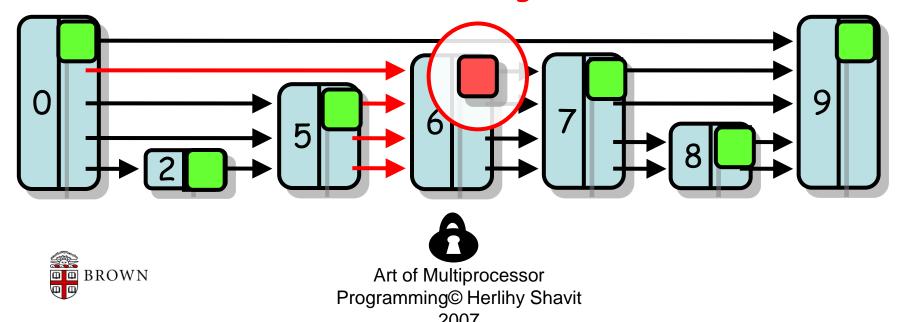




## remove(6): Linearization

· Successful remove happens when bit is set

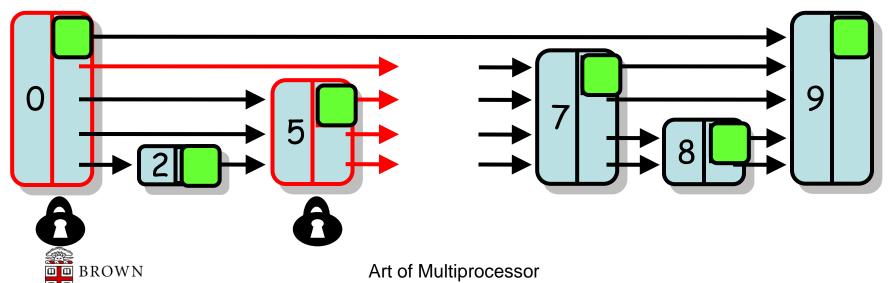
#### Logical remove...



### Add: Linearization

6

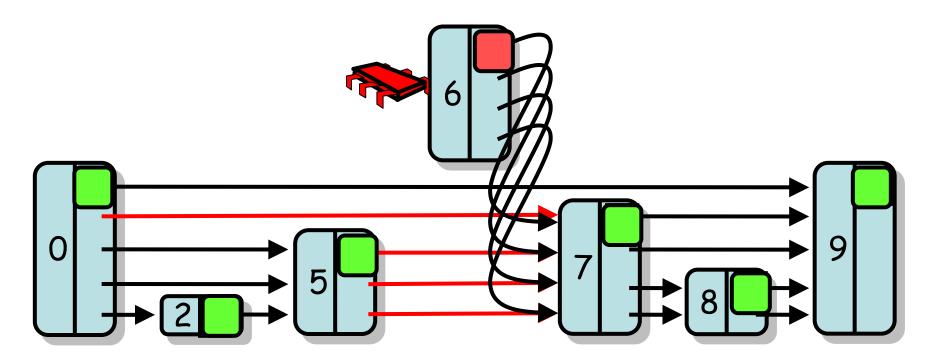
- Successful add() at point when fully linked
- · Add fullyLinked bit to indicate this
- Bit tested by contains()



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# Unsuccessful add(6) Linearization

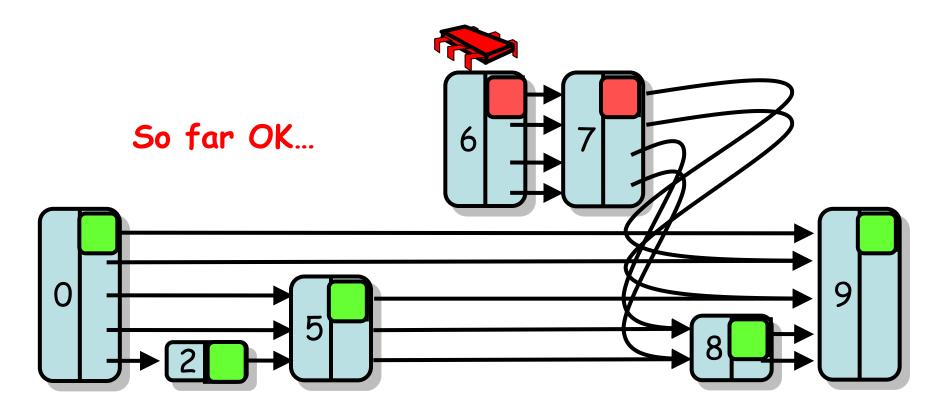
- · When not fully-linked unmarked node found
- · Pause while fullyLinked bit sets





# contains(7): Linearization

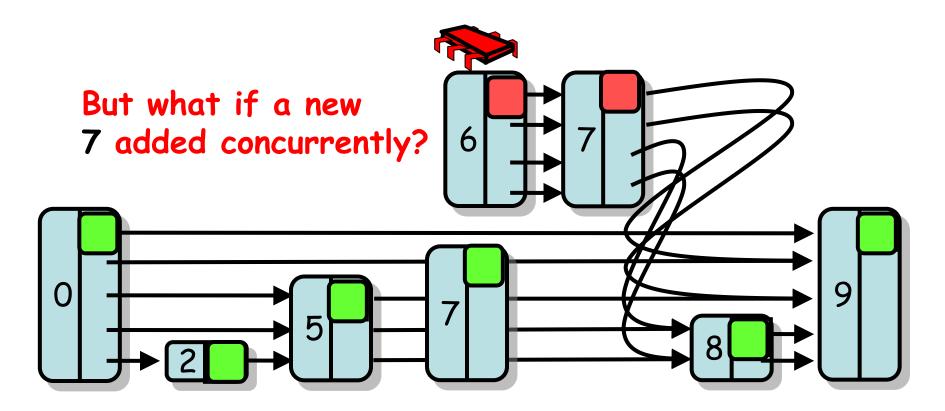
· When do we linearize unsuccessful Search?





### contains(7): Linearization

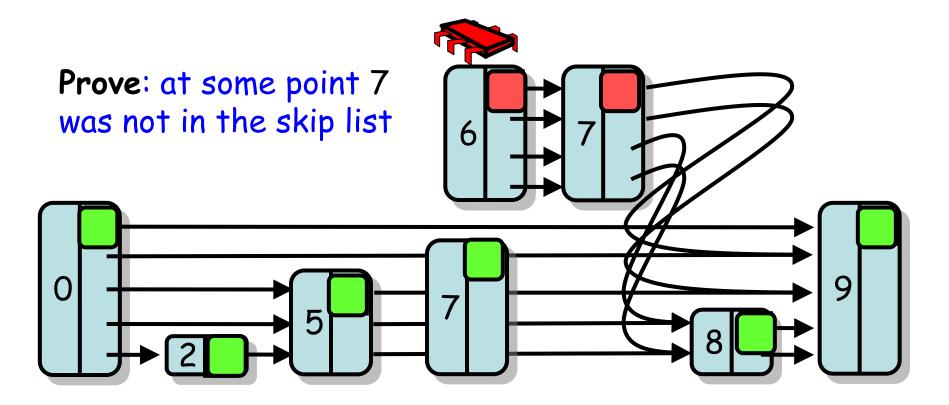
· When do we linearize unsuccessful Search?





### contains(7): Linearization

When do we linearize unsuccessful Search?





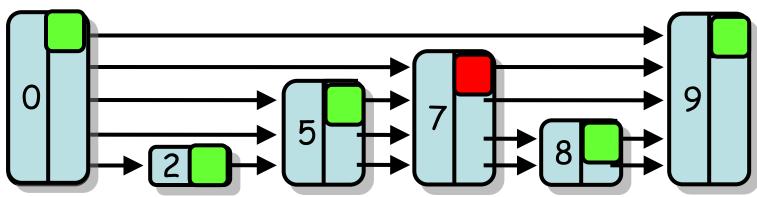
# Look-Free Skip List

- We can't manipulate multiple reference at the same time
  - Can't maintain the SkipList property
  - Each list could be not a sublist of the list at levels below
- The Abstract set is defined by the bottom-level
  - A key is in the set if there is a node with that key whose next reference is unmarked in the bottom level list Art of Multiprocessor

# LockFree Skip List

Use a mark bit for logical deletion

Did I mark it?



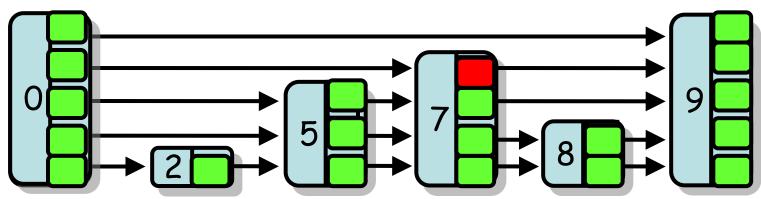


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# LockFree Skip List

Use a mark bit for logical deletion

Use a mark bit for each linked list





### How to Add/Remove?

- Each level of the list is treated as a LockFree list
- Use CAS() to insert
- Mark the next reference to remove
- find() method cleans up marked nodes
  - It's possible for a node to be physically removed while it's being linked at lower level



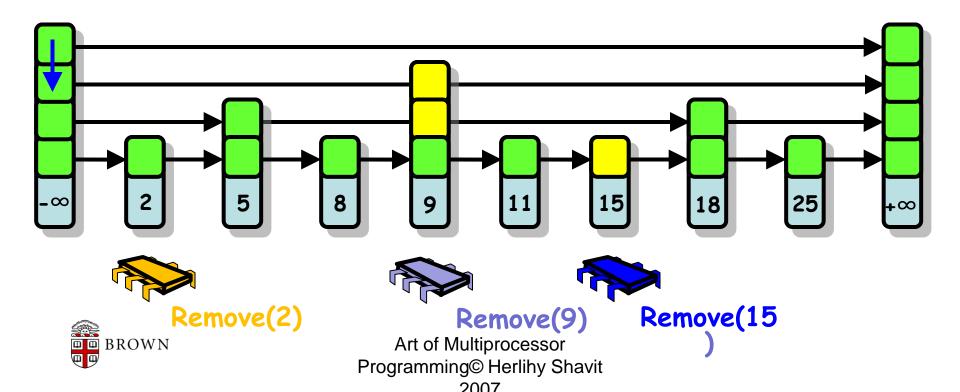
# Add()

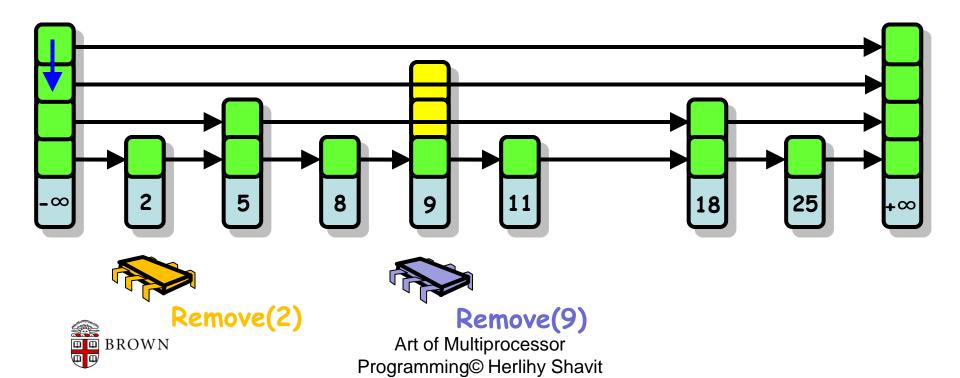
- Call find()
  - to determine whether a node is already in the list, find its set of predecessors and successors
- · A new node is prepared
  - with randomly chosen topLevel, and its next references are directed to the potential successors
- · Add to the bottom level LockFreeList
  - Successively links the node in higher levels

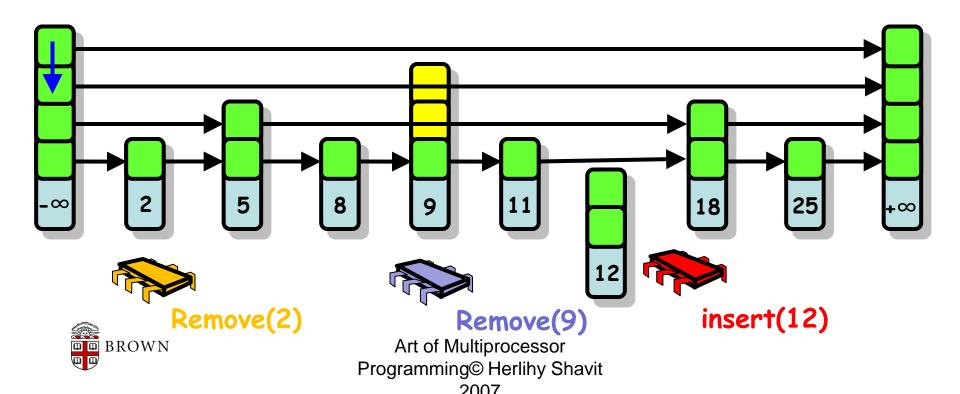


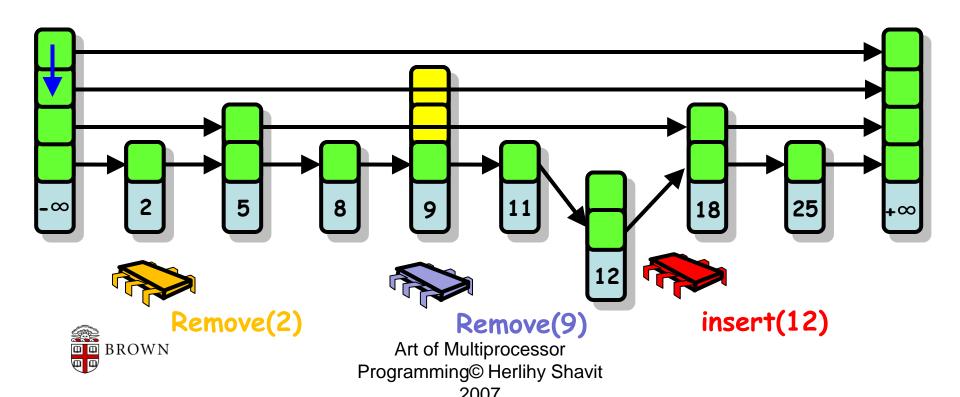


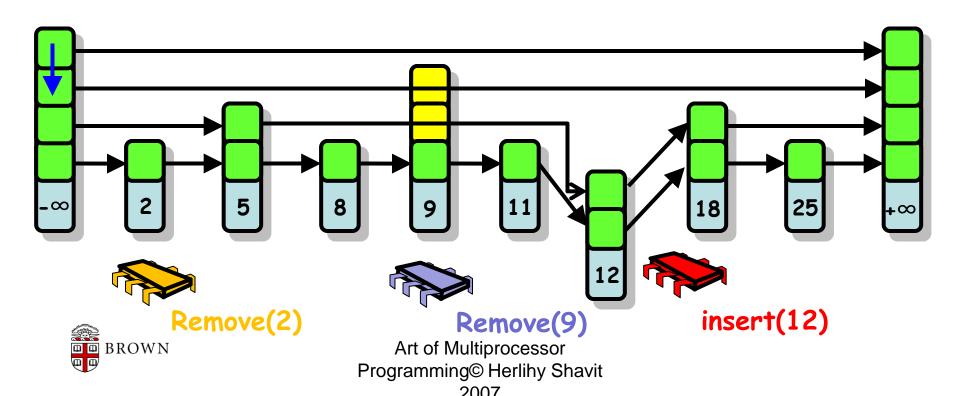
find(12)

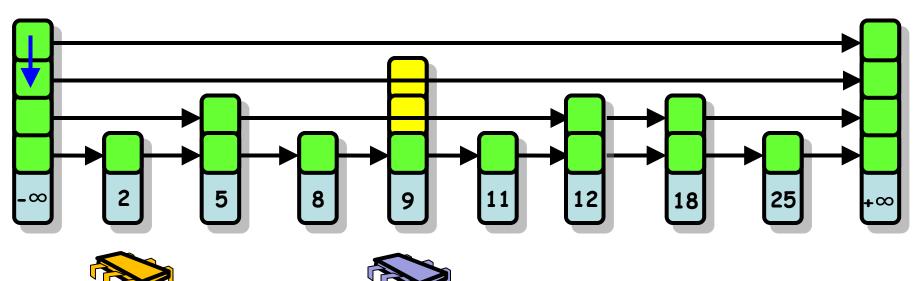










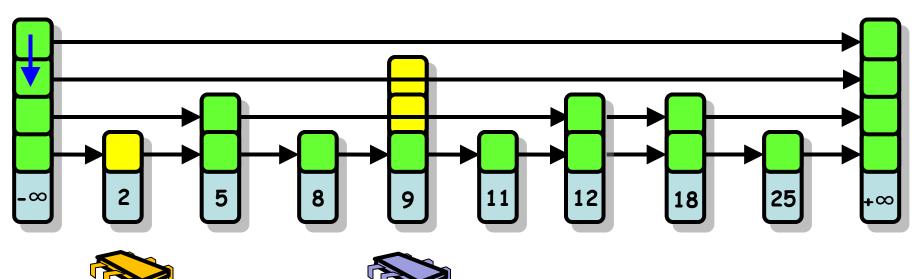






Remove(9)

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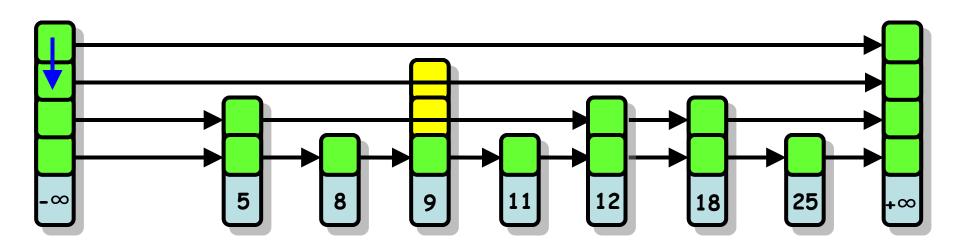




Remove(9)

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# add(12)





Remove(9)

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### Remove()

- Call find()
  - to determine whether an unmarked node with the target key is in the bottom level list
- Logical remove
  - mark from the top level to the bottom level
- physical remove
  - is done either by itself or by the other find()
     traversing the link

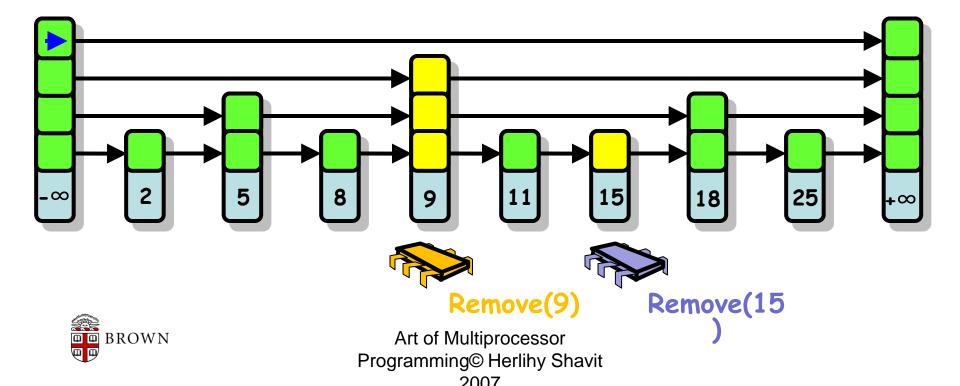


- Can't use find()
  - trying to remove marked nodes might generate too much contention
- Can't use LockFreeList's contains()
  - might skip nodes reachable from the bottom level
- Use find()
  - without actually doing physical remove



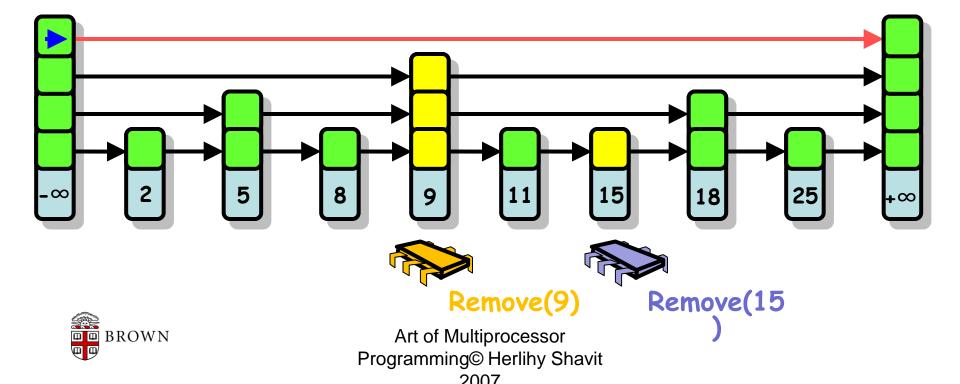






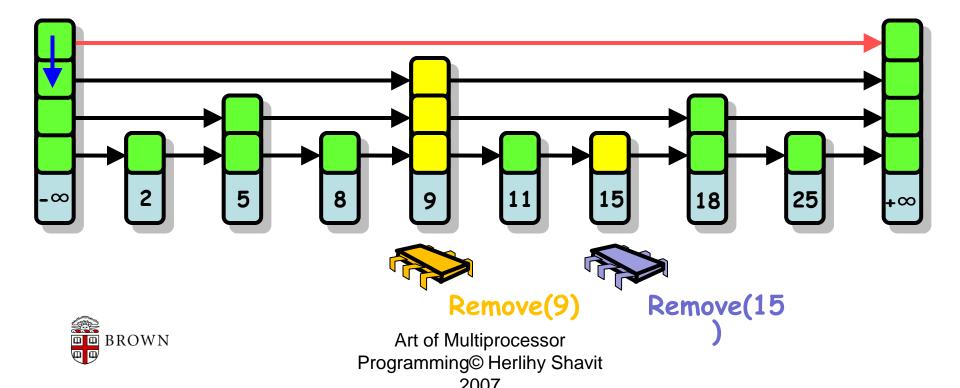




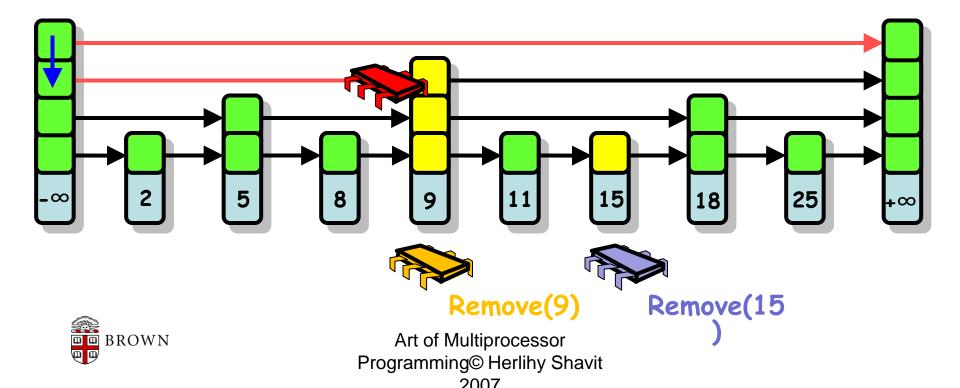






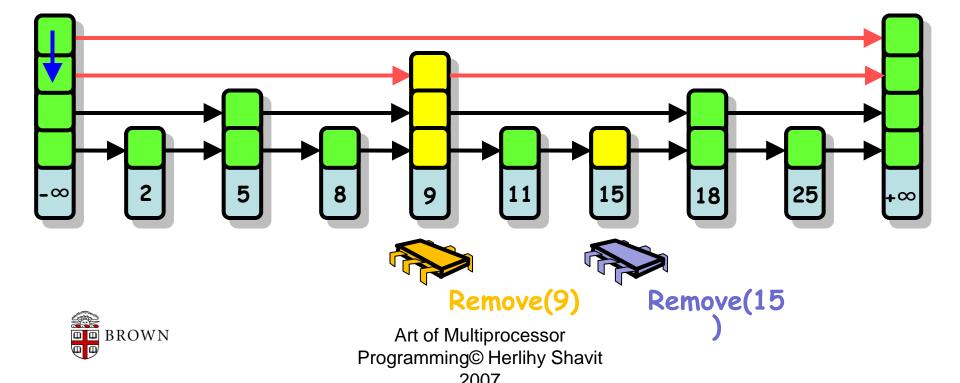






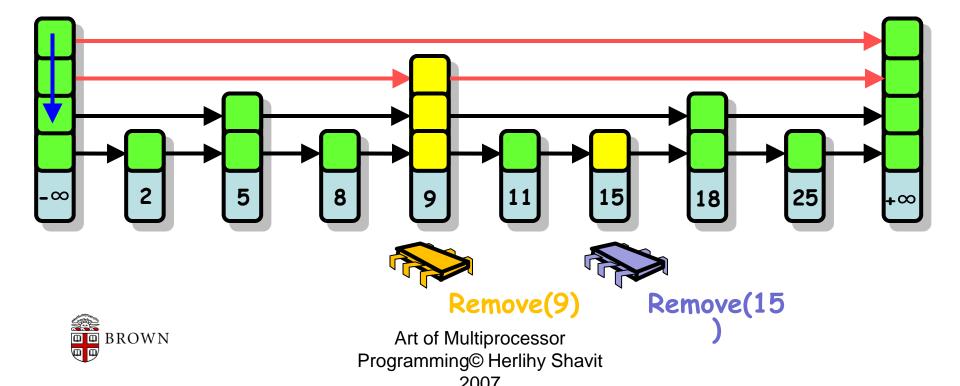




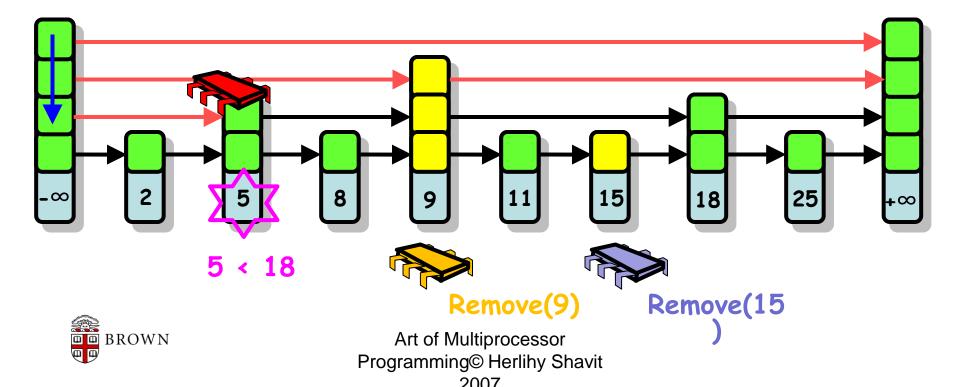




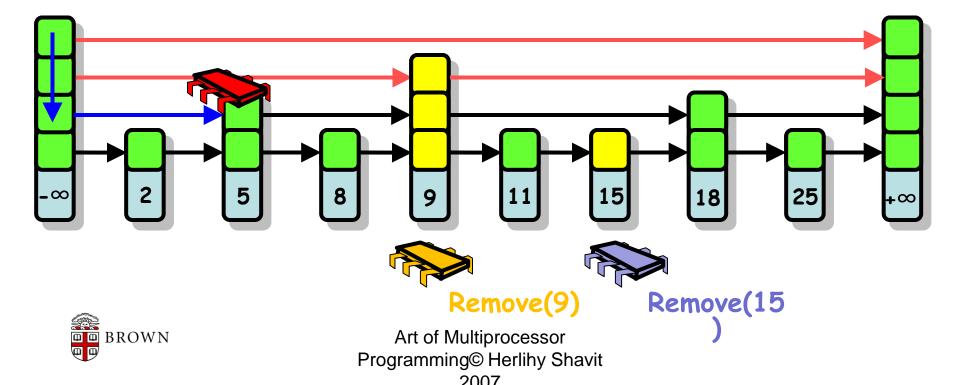




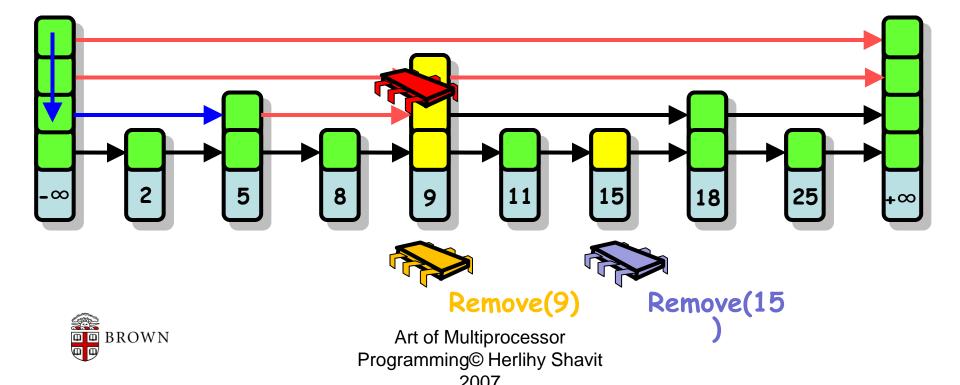




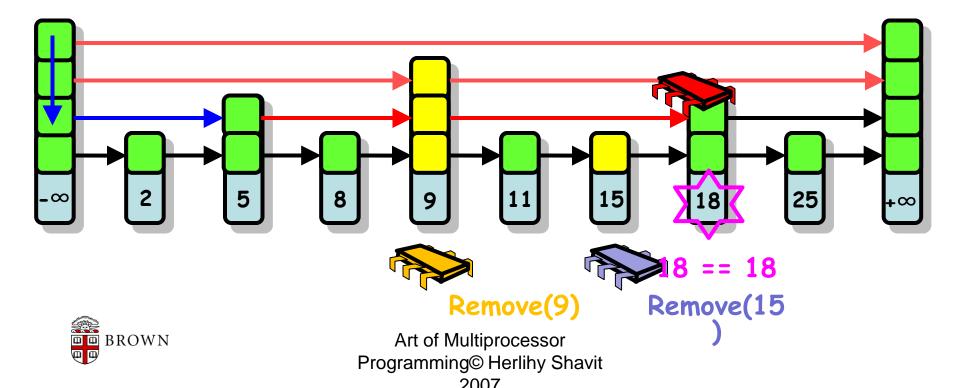






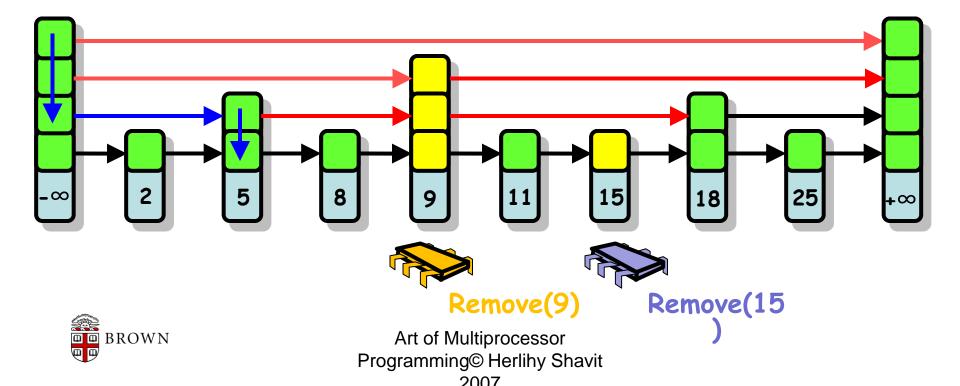




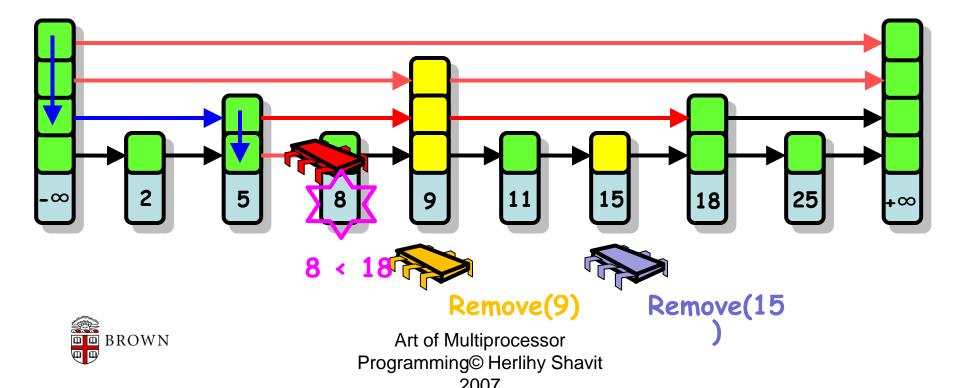




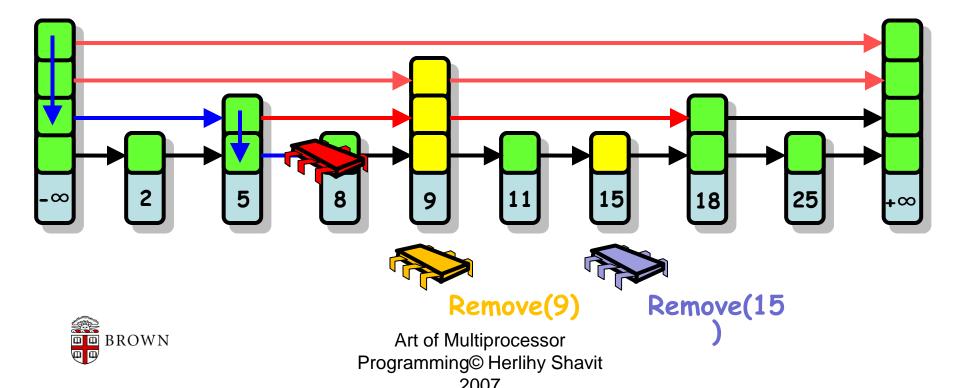




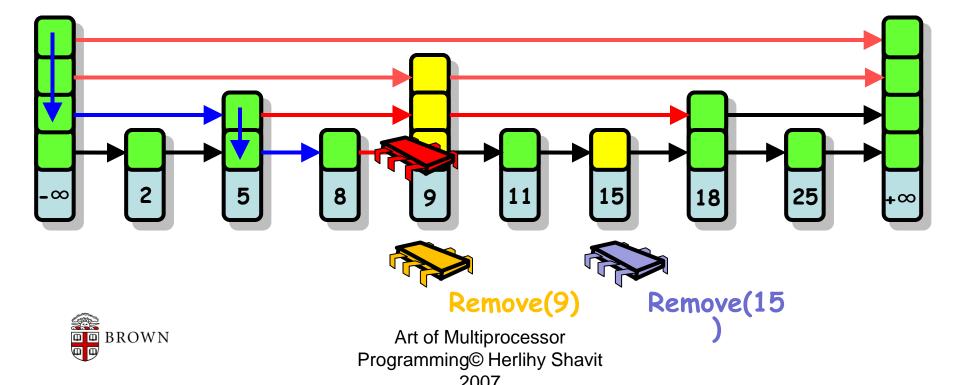




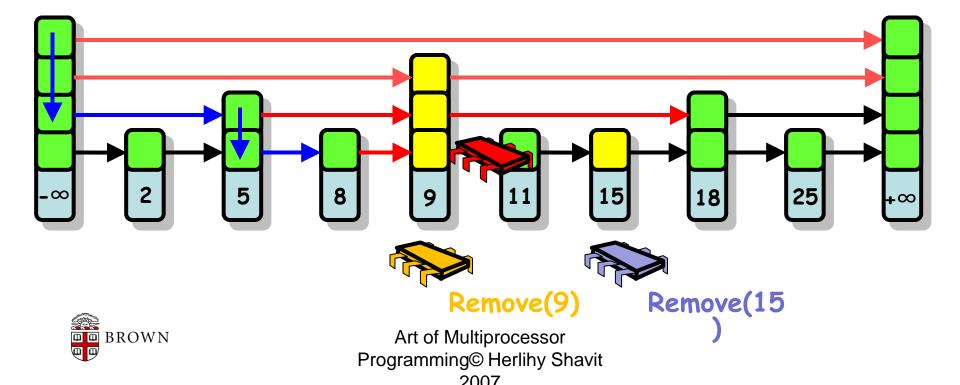




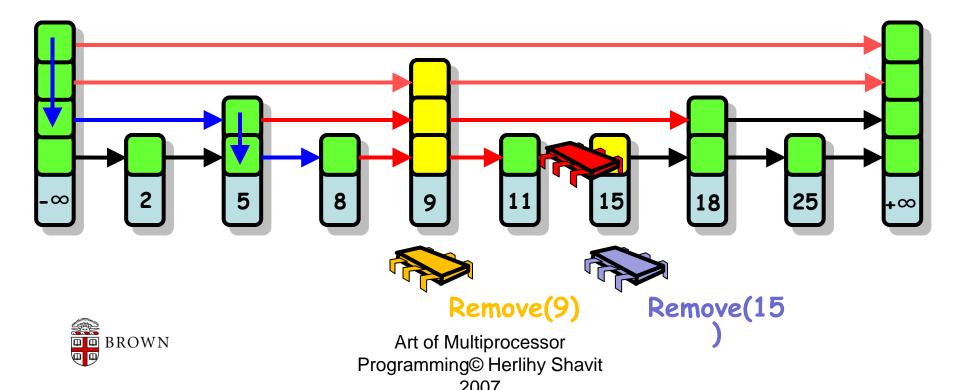






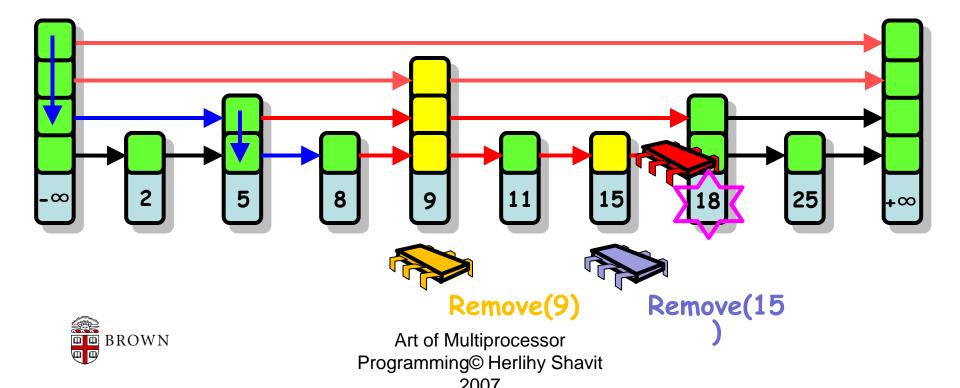






Contain(18) returns TRUE





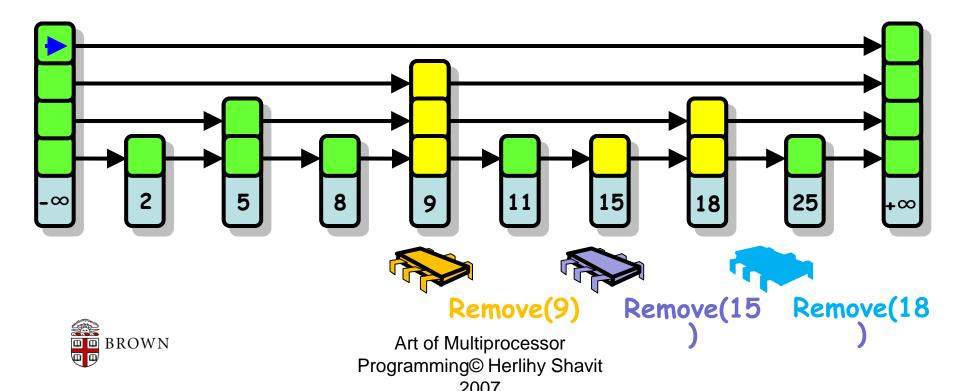
# When is contains() linearized?

 It might be possible contains() returns false even though the item has been inserted!



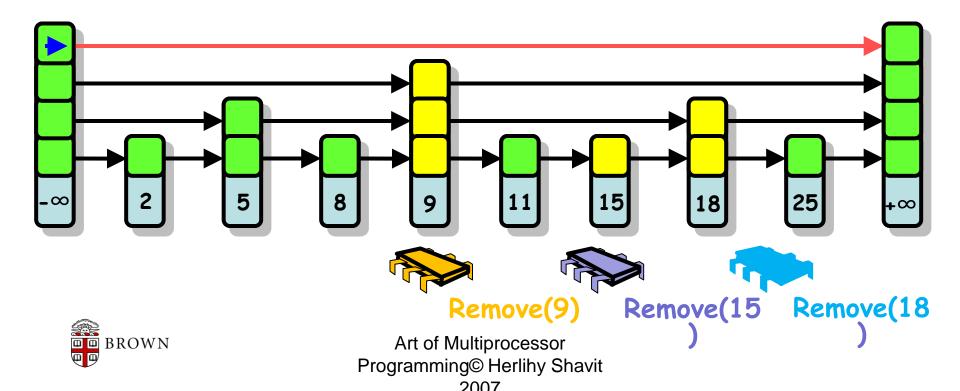






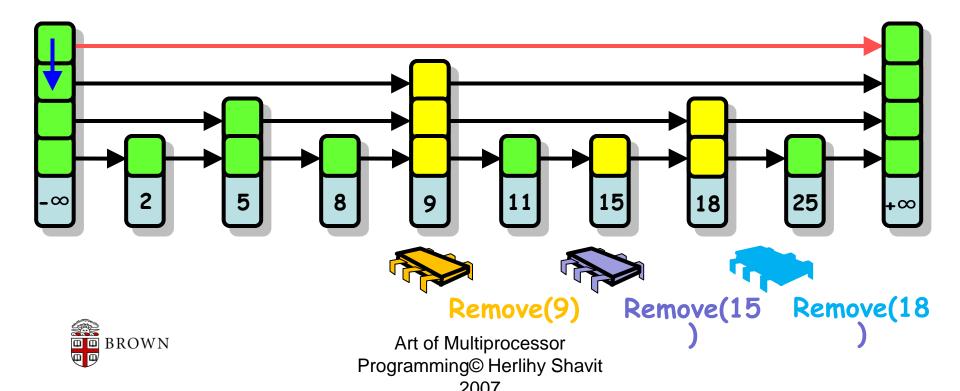




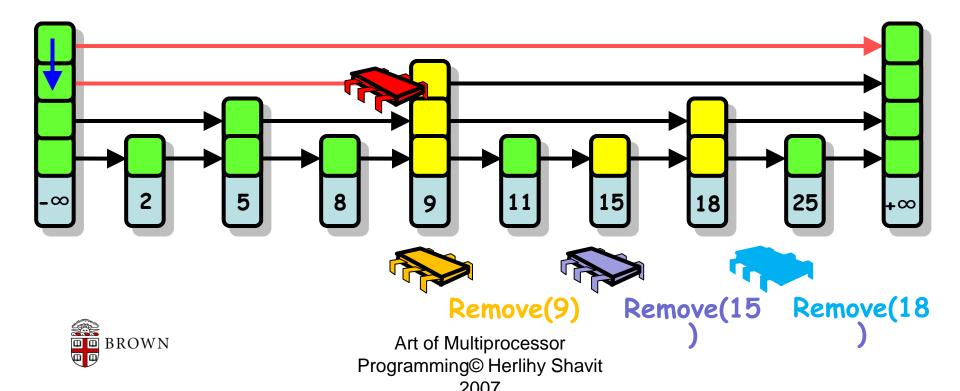






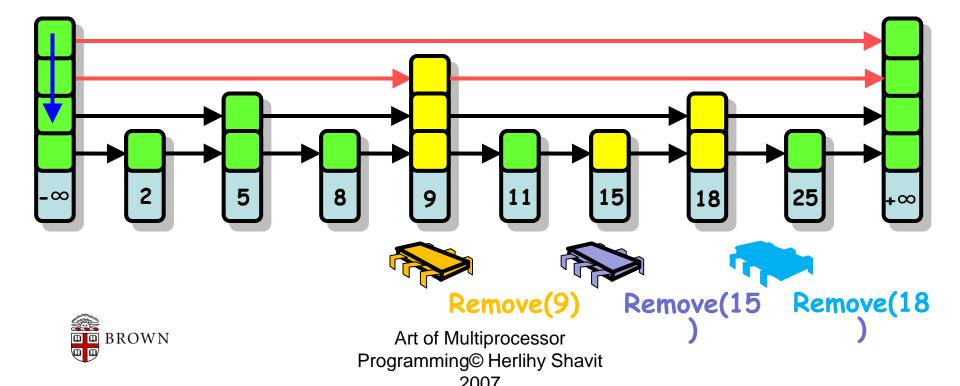


current
prev

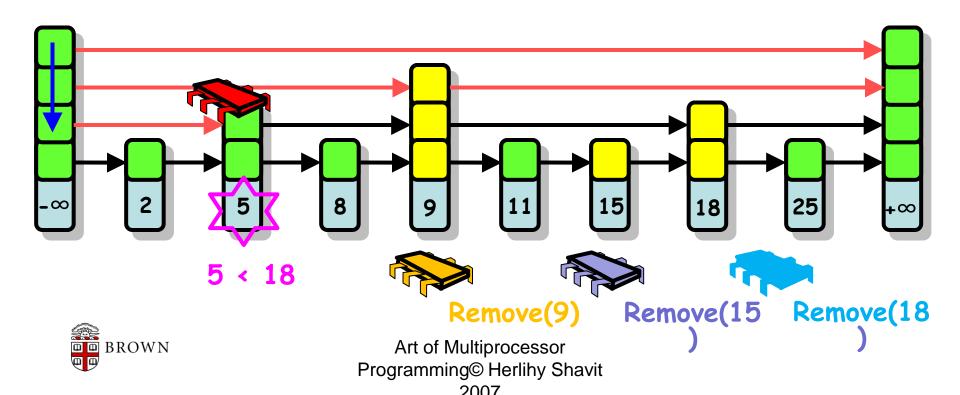




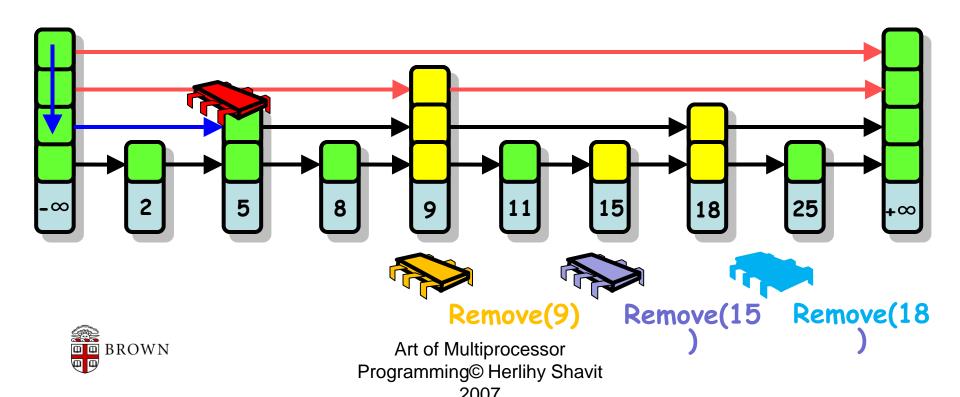




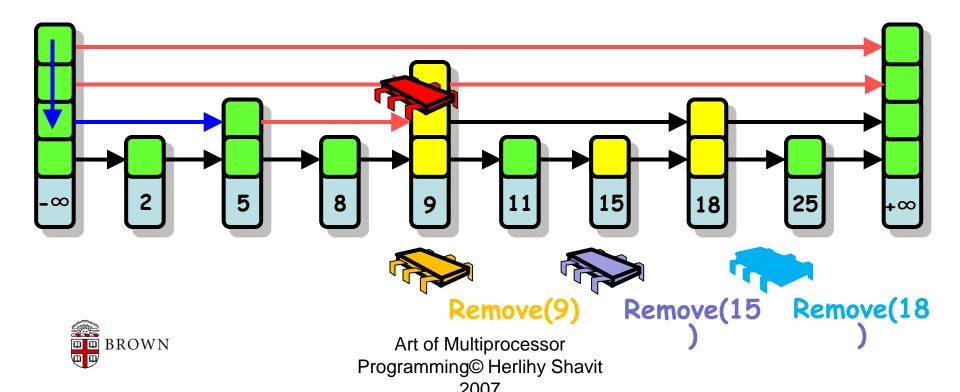




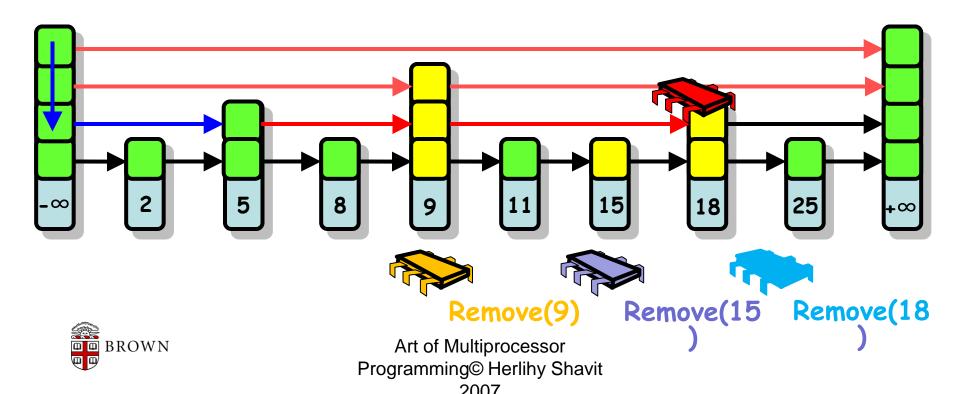




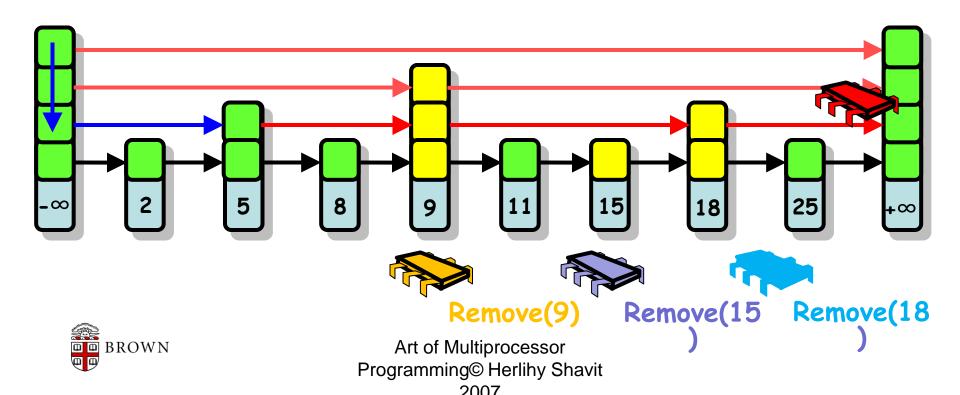




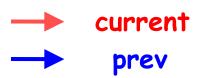


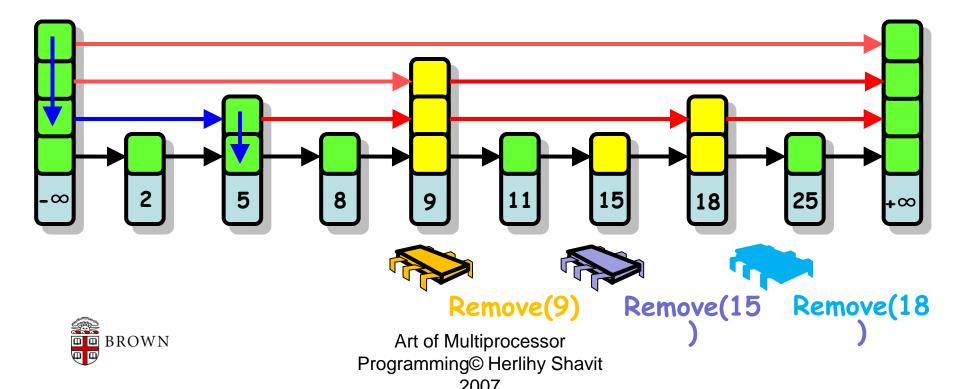




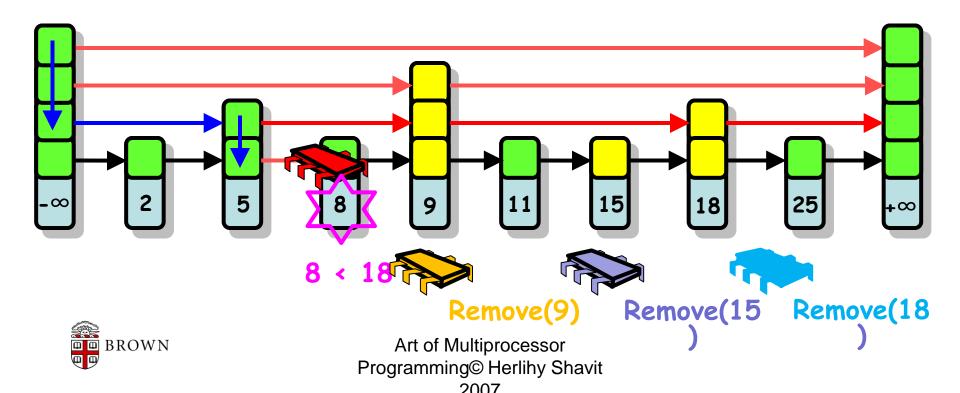




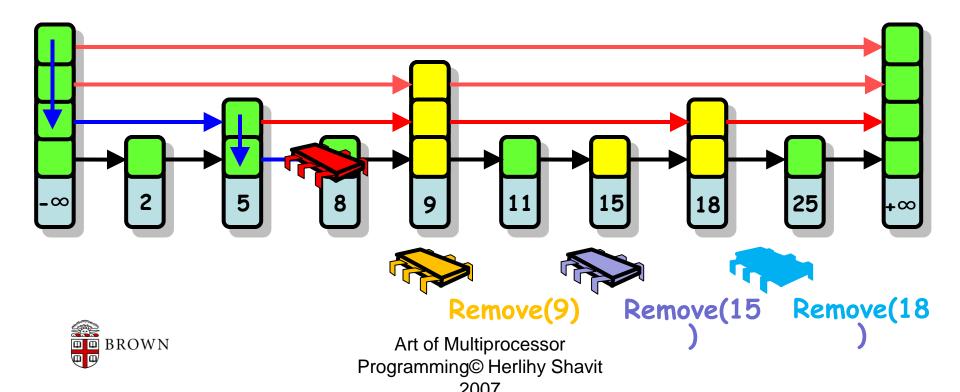




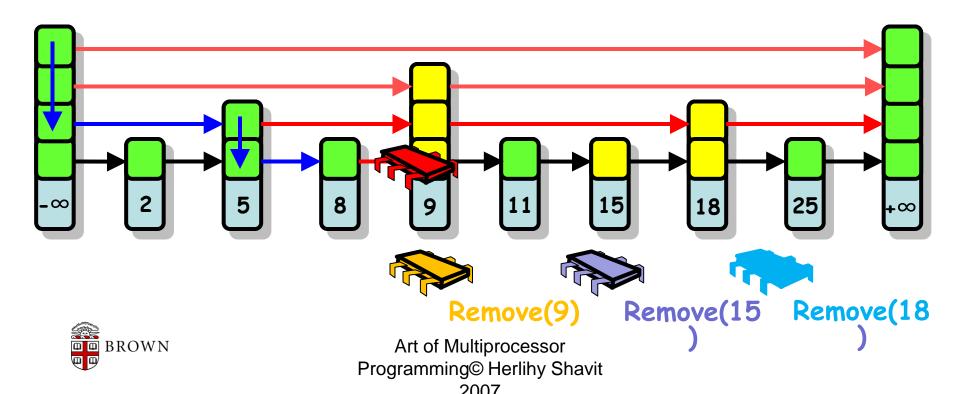




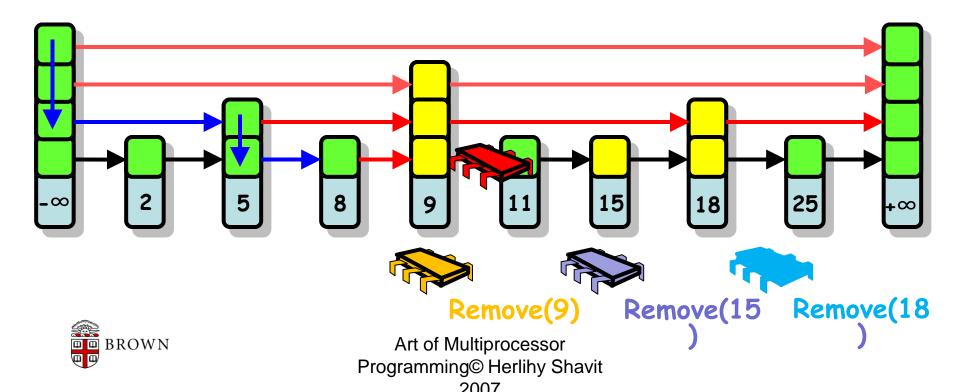






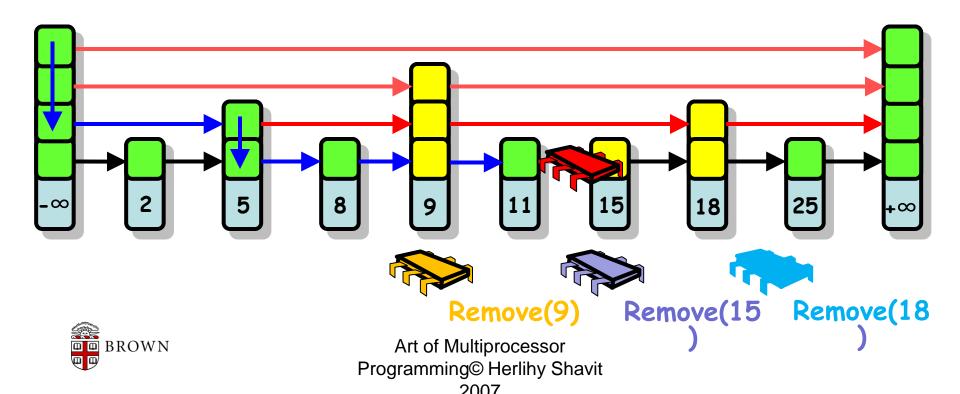






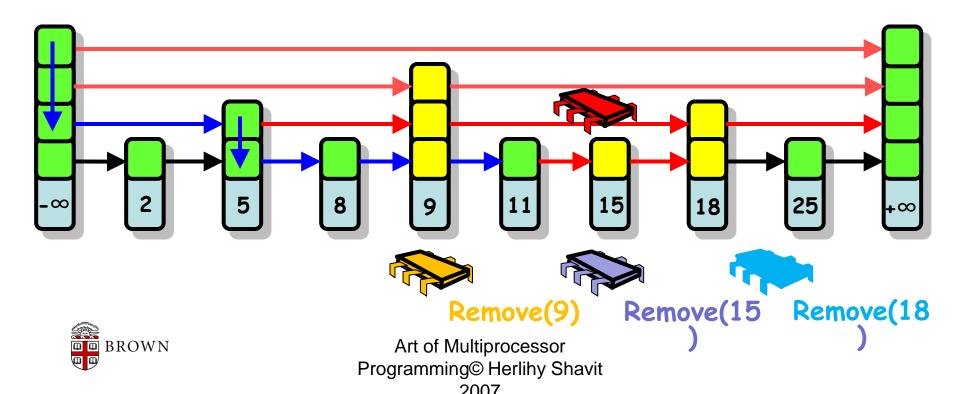
Contain(18





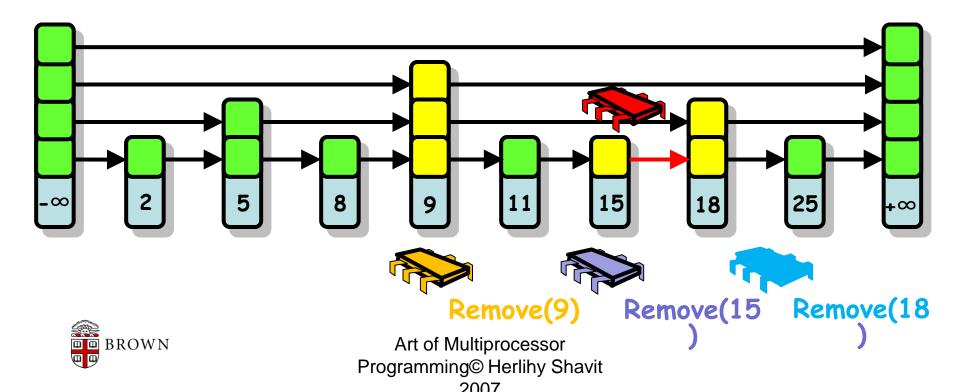
Contain(18





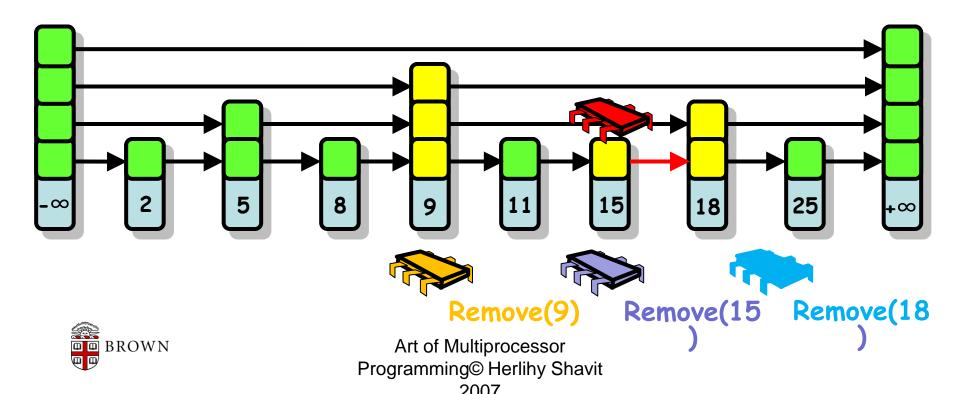
Contain(18
)
Add(18





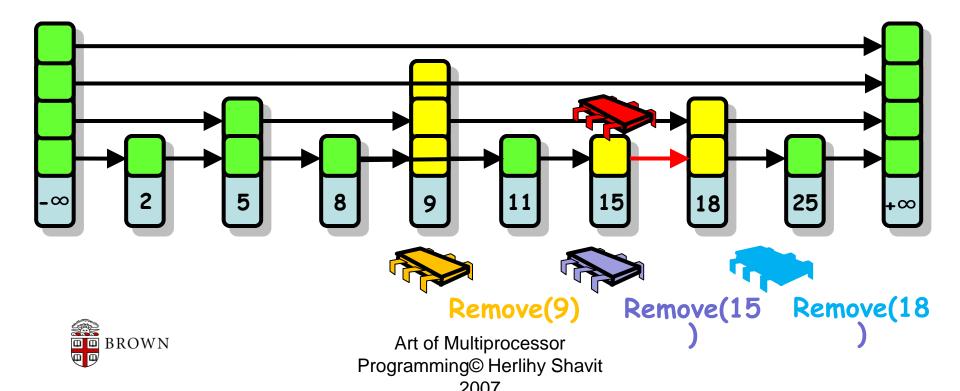
```
Contain(18
)
Add(18) -> find(18)
```





Contain(18 ) Add(18) -> find(18)



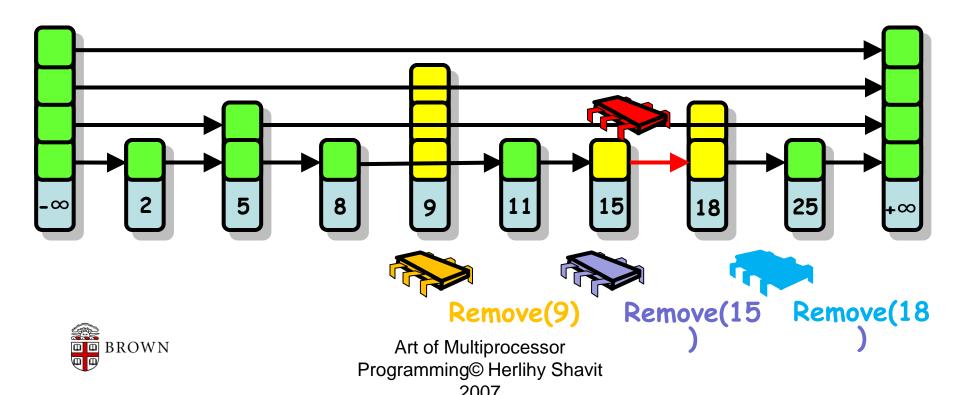


```
Contain(18

Current

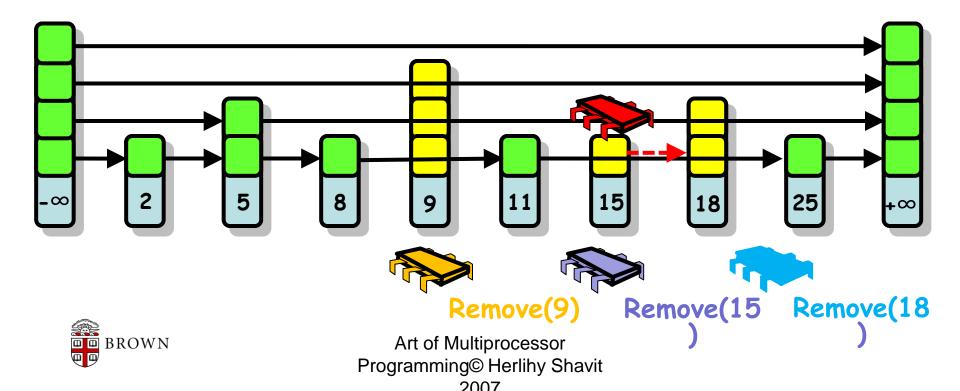
prev

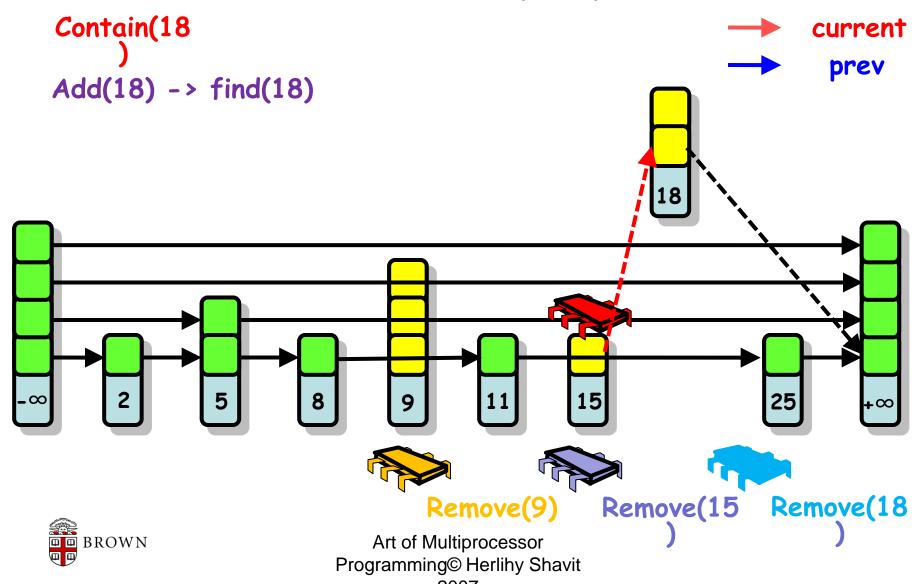
Add(18) -> find(18)
```

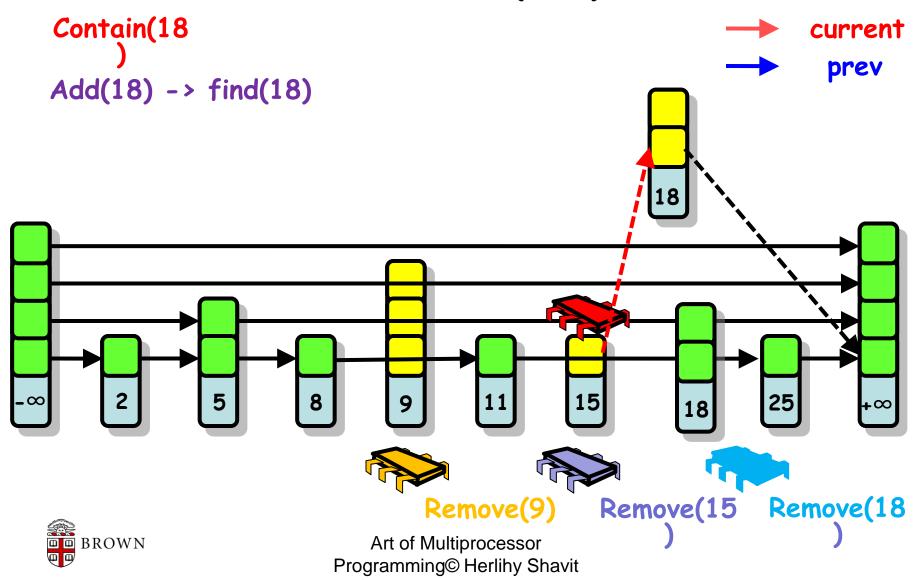


```
Contain(18
)
Add(18) -> find(18)
```

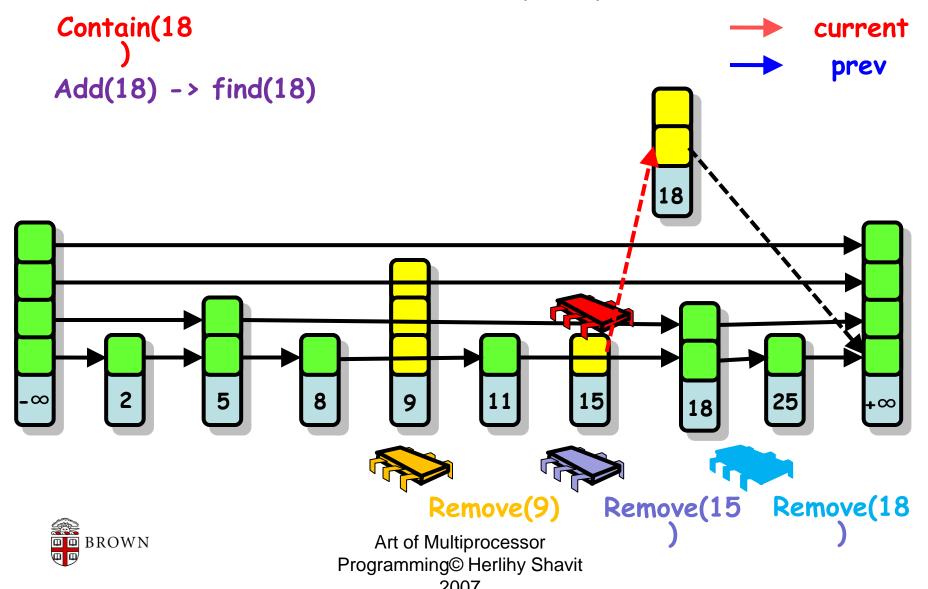


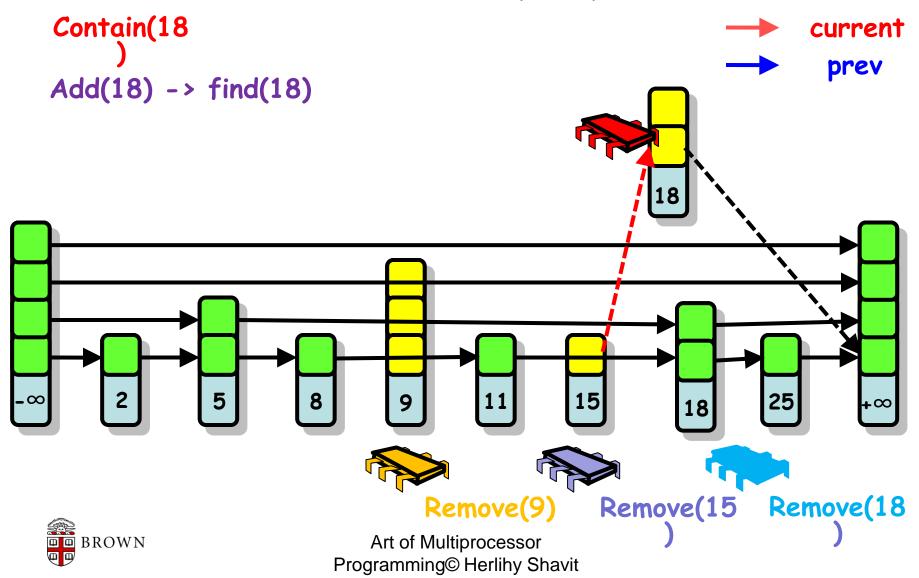


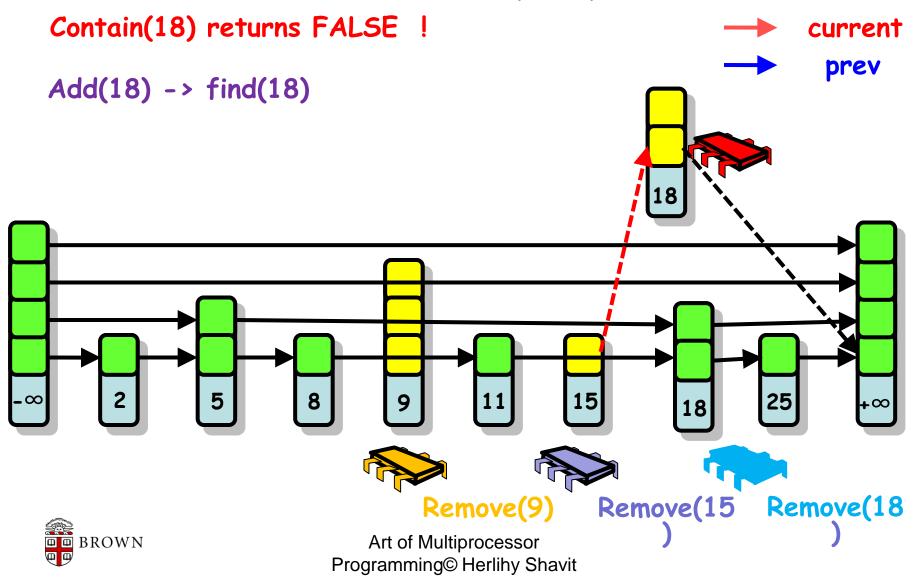




2007





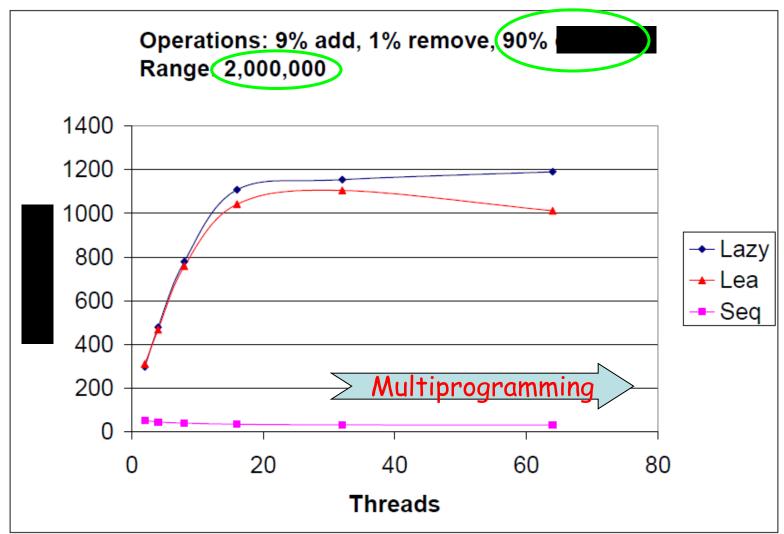


## A Simple Experiment

- · Each thread runs 1 million iterations, each either:
  - add()
  - remove()
  - contains()
- Item and method chosen in random from some distribution

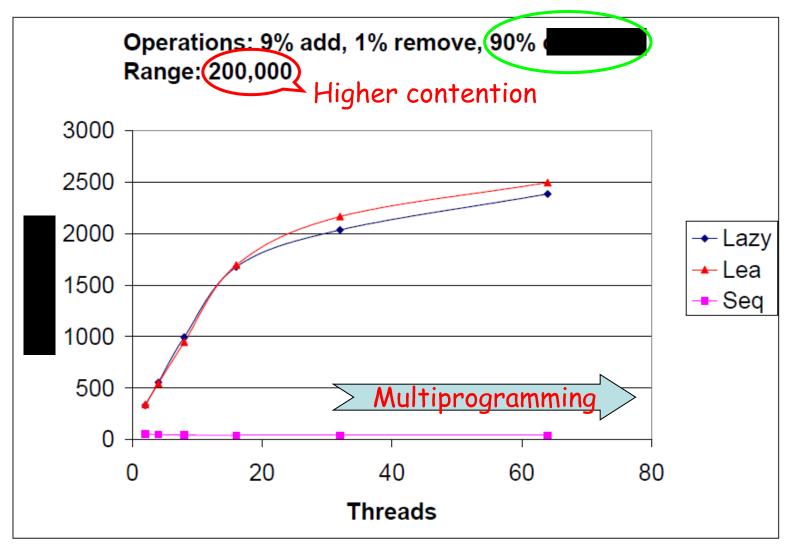


## Lazy Skip List: Performance



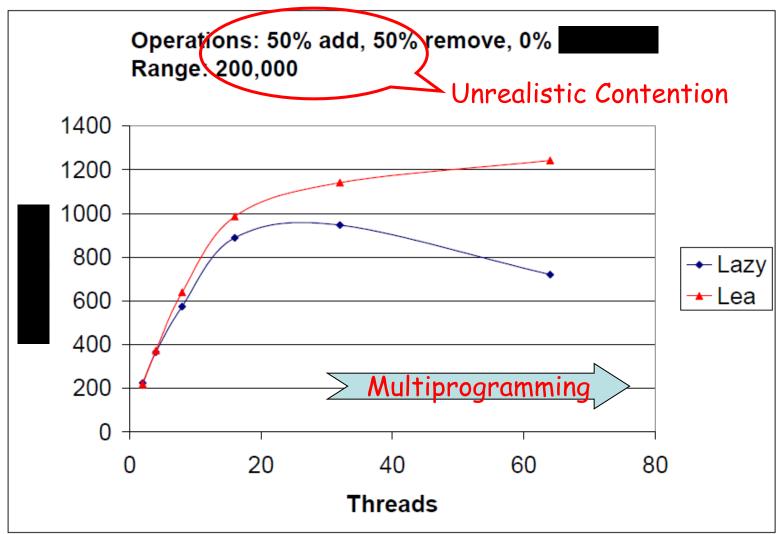


## Lazy Skip List: Performance





## Lazy Skip List: Performance





### Summary

- Lazy Skip List
  - Optimistic fine-grained Locking

 Performs as well as the lock-free solution in "common" cases

· Simple





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