

به نام خدا



# نظریه بازی‌ها

تمرین سوم نظری

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## سوال ۱

**Shoutout to Ali Safarafard for helping me out.**

action  $a^*$  evolutionary stable iff  $(a^*, a^*)$  is NE and for every action,  $b$ , that is a best response to  $a^*$ ,  $u(a^*, b) > u(b, b)$ .

Now lets see, the only strategy which could be evolutionary stable in our case would be the third strategy because  $(a^*, a^*)$  need to be a NE and the only place where both strategies by player I and player II are the same and NE is (1,1).

Lets check the second condition,  $u(3rd, b) > u(b, b)$ , now we could see that the best response to the 3rd action is only the 3rd action so the second condition is also satisfied and we could say that **the 3rd strategy is evolutionary stable**.

Now for the mixed strategy, a mixed strategy,  $\alpha$  is evolutionary stable if  $(\alpha, \alpha)$  is a NE and for every strategy,  $\beta$  that is a best response to  $\alpha$ ,  $u(\alpha, \beta) > u(\beta, \beta)$ .

Lets find a mixed strategy in our game.

$$U_{II}(1st, \{p_1, p_2, p_3\}) = 3p_2$$

$$U_{II}(2nd, \{p_1, p_2, p_3\}) = p_1$$

$$U_{II}(3rd, \{p_1, p_2, p_3\}) = p_3$$

$$3p_2 = p_1 = p_3, p_1 + p_2 + p_3 = 1 \rightarrow p_1 = p_3 = \frac{3}{7}, p_2 = \frac{1}{7}$$

Game is symmetric so  $U_{II} = U_I$ .

Now we see that our mixed strategy is a NE so lets see the second condition.

$$U(\beta, \beta) = U(\{p_1, p_2, p_3\}, \{p_1, p_2, p_3\}) = p_2p_1 + 3p_1p_2 + p_3^2 = p_3^2 + 4p_1p_2$$

$$U(\alpha, \beta) = U(\{\frac{3}{7}, \frac{1}{7}, \frac{3}{7}\}, \{p_1, p_2, p_3\}) = \frac{p_1}{7} + \frac{9p_2}{7} + \frac{p_3}{7} = \frac{1+8p_2}{7}$$

Now :  $\frac{1+8p_2}{7} > p_3^2 + 4p_1p_2$ , if for example  $p_1 \approx 0.339811, 0 \leq p_2 \leq 0.4438, p_3 \approx 0.0714286(9.24264 - 14p_2)$ , then the inequality would not hold, hence the mixed strategy is **not an ESS**.

This means that there exists an element on row  $i$  column  $i$  which is **bigger than all other elements on row  $i$**  and since the game is **symmetric**, this element is also **greater than all elements on column  $i$**  so we could say that this is the best response for both player 1 and player 2 and we know that in a symmetric game, if  $(a_i, a_i)$  is a **strict NE**, then it is also evolutionary stable strategy.

So to sum things up, we showed that since this game is symmetric and by the questions hypothesis,  $a_{ii}$  is a strict NE and we know that strict NE in symmetric games where both strategies are the same is an ESS.

	$x$	$y$	$z$
$x$	$-\infty, -\infty$	$\textcircled{3}, 1$	$0, \textcircled{2}$
$y$	$1, \textcircled{3}$	$0, 0$	$\textcircled{1}, -\infty$
$z$	$\textcircled{2}, 0$	$-\infty, \textcircled{1}$	$0, 0$

As you could see, there is no pure NE.

Now to find mixed NEs:

$$U_1(x, P) = -\infty * p_1 + 3 * p_2$$

$$U_1(y, P) = p_1 + p_3$$

$$U_1(z, P) = 2p_1 - \infty * p_2$$

We see that  $\infty$  is making a problem in our equations and we cant solve them, simply because infinity is considered unknown and that is why in an infinity utility game, there is a possibility which we will not find any equilibriums.

Now for the second part think as if we remove the profiles with  $-\infty$  and randomize between the other strategies,

each player now has expected of  $6/6 = 1$  payoff.

We could also remove the profiles in which both players get 0,0 so that each players payoff would be  $6/4 = 1.5$  and no player would deviate.

Consider this scenario that there is a third party which could draw (y,x) (z,x) (x,y) (x,z) and then recommend a strategy to each player, we know that the players won't deviate from the recommended strategy.

First we will see for what range of  $x$ , the game would have pure ESS.

Action  $a$  would be ESS iff  $(a,a)$  is NE and for every  $b$  that is best responses to  $a$ ,  $u(a,b) > u(b,b)$ .

	$i$	$ii$
$i$	$x, x$	$0.5x, 0.5$
$ii$	$0.5, 0.5x$	$1, 1$

1.  $i$  is ESS

$(x,x)$  is NE  $\rightarrow 0.5 < x \rightarrow ESS$

$x = 0.5 \rightarrow u(i,i) = u(ii,i) \rightarrow u(i,ii) < u(ii,ii) \rightarrow$  not ESS.

2.  $ii$  is ESS

$(1,1)$  is NE  $\rightarrow 0.5x < 1 \rightarrow x < 2 \rightarrow ESS$

$x = 2 \rightarrow u(ii,ii) = u(i,ii) \rightarrow u(ii,i) < u(i,i) \rightarrow$  not ESS.

Now if either  $x > 0.5$  or  $2 > x$  then the game would have ESS.

سوال ۵

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We know that  $P$  is the matrix of probabilities which hold the probabilities for mixed nash equilibrium, so  $RP = cI \xrightarrow{RR^{-1}=I} P = cR^{-1}$

We know by the property of matrix inversion, a matrix could only have at best, one inverse so we know that  $P$  is unique.

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This equilibrium exists if  $R$  is inversable.

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The strategy is best response for both players due to  $A$  being symmetric, so this strategy would always be an ESS.

If  $V - c > \alpha V \rightarrow (protect, protect) = strictNE \rightarrow protect = ESS$

We also know that in this case, the  $(leave, leave)$  option would be strictly dominated so it could not be an NE, hence  $leave$  is not ESS in this case.

If  $\alpha V - c < 0 \rightarrow (leave, leave) = strictNE \rightarrow ESS$

We also know that in this case, the  $(protect, protect)$  option would be strictly dominated so it could not be an NE, hence  $protect$  is not ESS in this case.

**We proved that in the first and second case, there would only exist one ESS.**

First lets see protect, we could see that  $(protect, protect)$  is a strict NE because :

$$V - c > (1 - \alpha)V - c > \alpha V - c \xrightarrow{c > \alpha V} V - c > \alpha V \rightarrow protect = ESS$$

Also leave is a strict NE because :

$$\alpha V < c \rightarrow \alpha V - c < 0 \rightarrow leave = ESS$$

So we saw that protect and leave are both ESS because they are **both strict NEs**.

In this state we could see that neither  $(protect, protect)$  or  $(leave, leave)$  is not a strict NE so they could not be ESS, hence we need to find the mixed NE and then find if its a mixed ESS.

$$U_1(protect, P) = p_1 * (V - c) + (1 - p_1)(\alpha V - c)$$

$$U_1(leave, P) = p_1 * \alpha V$$

$$p_1 \alpha V = p_1 * (V - c) + (1 - p_1)(\alpha V - c) \rightarrow p_1 = \frac{\alpha V - c}{(2\alpha - 1)V}$$

Now lets see if its also ESS, to do this we should check that for every mixed strategy,  $\beta$   
 $u(NE, \beta) > u(\beta, \beta)$  :

$$\begin{aligned} & \frac{\alpha V - c}{(2\alpha - 1)V} * p * (V - c) + (1 - \frac{\alpha V - c}{(2\alpha - 1)V}) * p * \alpha V + \frac{\alpha V - c}{(2\alpha - 1)V} * (1 - p) * (\alpha V - c) > \\ & p^2 * (V - c) + p * (1 - p) * (\alpha V - c) + p * (1 - p) * \alpha V \\ & \rightarrow True \rightarrow ESS \end{aligned}$$