

IR Lowering

Expressions Statements

CONST(n)	MOVE(e_{dst}, e)
TEMP(t)	EXP(e)
OP(e_1, e_2)	SEQ(\vec{s})
MEM(e)	JUMP(e, l)
CALL(e, \vec{t})	CJUMP(e, l_1, l_2)
NAME(ℓ)	LABEL(ℓ)
ESEQ(s, e)	RETURN(e)
	CALL(e_f, \vec{e})

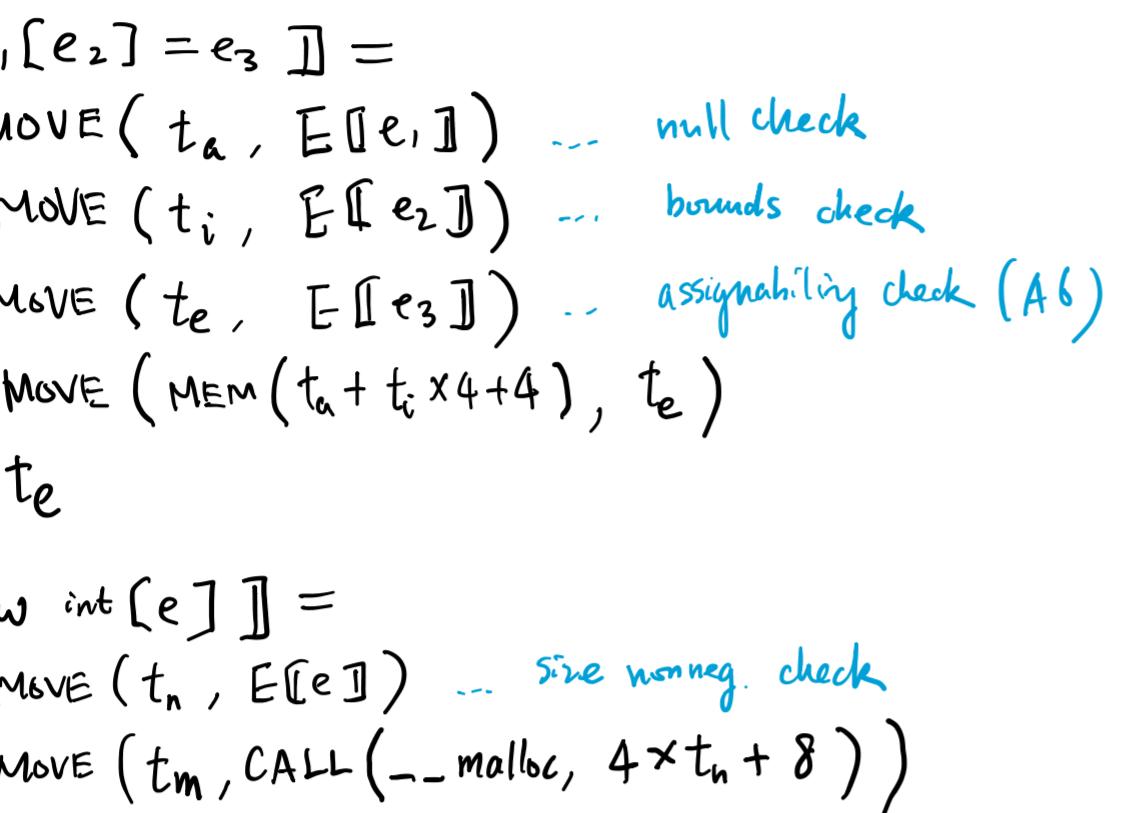
Recap

- TIR
- Expressions and statements have mutually recursive syntax
- Translating Joos expressions and statements to IR:

$E[e] = e'$ and $S[s] = s'$: mutually recursive functions

- Flatten ESEQ and SEQ to a single top-level SEQ
- expressions don't have side effects
- ≤ 1 side effect in each statement
- no two-way jumps

$E[e_1[e_2]] = ESEQ(MOVE(t_a, E[e_1]); CJUMP(EQ(t_a, 0), \ell_{err}, \ell_{notnull}, \ell_{err}); // null check \ell_{err}: CALL(NAME(__exception)); \ell_{notnull}: MOVE(t_i, E[e_2]); CJUMP(LTU(t_i, MEM(t_a - 4)), \ell_{inbound}, \ell_{err}); // bounds check \ell_{inbound}: MEM(t_a + t_i \times 4 + 4))$



$E[e_1[e_2] = e_3] =$

- MOVE($t_a, E[e_1]$) ... null check
- MOVE($t_i, E[e_2]$) ... bounds check
- MOVE($t_e, E[e_3]$) ... assignability check (A6)
- MOVE($MEM(t_a + t_i \times 4 + 4), t_e$)

t_e

$E[\text{new int}[e]] =$

- MOVE($t_n, E[e]$) ... size memory check
- MOVE($t_m, CALL(\text{--malloc}, 4 \times t_n + 8)$)
- MOVE($MEM(t_m), t_n$)
- MOVE($MEM(t_m + 4), \dots$)
- ... loop to zero-initialize array

$t_m + 4$

Lower TIR to canonical-form IR.

Joos AST \rightarrow TIR \rightarrow -- IR

$L[s] = s_1; s_2; \dots; s_n$

$L[e] = \underbrace{s_1; s_2; \dots; s_n}_{\text{side effects}} \mid \underbrace{e'}_{\text{pure}}$

$L[l] = l$:

$L[s_1; s_2; \dots; s_n] = L[s_1]; L[s_2]; \dots; L[s_n]$

$L[e] = \vec{s} \mid e'$

$L[\text{Jump}(e)] = \vec{s}; \text{JUMP}(e')$

$L[\text{Exp}(e)] = \vec{s}$

$L[e] = \vec{s} \mid e'$

$L[\text{RETURN}(e)] = \vec{s}; \text{RETURN}(e')$

$L[n] = \bullet \mid n$

$L[t] = \bullet \mid t$

$L[\ell] = \bullet \mid \ell$

$L[e] = \vec{s} \mid e'$

$L[MEM(e)] = \vec{s} \mid MEM(e')$

$L[ESEQ(s, e)] = \vec{s}; \vec{s'} \mid e'$

$L[e_1] = \vec{s}_1 \mid e'_1 \quad L[e_2] = \vec{s}_2 \mid e'_2$

$L[OP(e_1, e_2)] = \vec{s}_1; \vec{s}_2 \mid OP(e'_1, e'_2)$

When does s_2 have side effects that affect e'_1 ?

• s_2 writes into temp read by e'_1 easy to detect.
 $t_3 = t_4$?

• s_2 --- mem --- undecidable
 $MEM(e_3) = MEM(e_4)$?

$L[e] = \vec{s}_1 \mid e'_1 \quad L[e_2] = \vec{s}_2 \mid e'_2$

$L[MOVE(MEM(e_1), e_2)] = \vec{s}_1; MOVE(t, e'_1); \vec{s}_2 \mid OP(t, e'_2)$

$\forall i \quad L[e_i] = \vec{s}_i \mid e'_i$

$L[CALL(e, \vec{t})] = \vec{s}_0; MOVE(t_0, e'_0); \vec{s}_1; MOVE(t_1, e'_1); \dots; \vec{s}_n; MOVE(t_n, e'_n); \vec{s}_{n+1} \mid \vec{s}_{n+2} \mid \dots \mid \vec{s}_m \mid \vec{s}_{m+1} \mid \dots \mid \vec{s}_p$

TEMP(RET)

$L[e_1] = \vec{s}_1 \mid e'_1 \quad L[e_2] = \vec{s}_2 \mid e'_2$

$L[MOVE(MEM(e_1), e_2)] = \vec{s}_1; MOVE(t, e'_1); \vec{s}_2 \mid MOVE(MEM(t), e'_2)$

Eliminate two-way jumps.

$CJUMP(e, l_t, l_f) \rightsquigarrow CJUMP(e, l_t); \text{Jump}(l_f);$

Idea: reorder statements so that l_f is the next statement

ordinary starts $\rightarrow \circlearrowleft \rightarrow \circlearrowright$

jump $\rightarrow \circlearrowleft \rightarrow \circlearrowright$

return $\rightarrow \circlearrowleft \rightarrow \circlearrowright$

labels $\rightarrow \circlearrowleft \rightarrow \circlearrowright$

Basic block: a sequence $s_1; \dots; s_n$ st.

• all but s_1 and s_n are ordinary

• s_1 can be a label statement

• s_n can be ordinary or jump/return.

$\ell: s_2;$

\vdots

s_{n-1}

$s_n / \text{jump/return}$

Maximal basic block.

Example:

$l_0: CJUMP(e, l_1, l_2)$

$l_1: MOVE(x, y)$

$l_2: MOVE(x, y + z)$

$\text{JMP}(l_1)$

$l_3: CALL(f)$

RETURN

CFG

Vertices = maximal BBs.

Edges = control flow.

$l_0: CJUMP(\dots)$

$l_2: \text{Move}$

$\text{Jump}(l_1)$

$l_3: \text{CALL}$

RET

$l_1: \text{Move}$

$\text{Jump}(l_1)$

$l_1: \text{Move}$

$\text{Jump}(l_1)$

• delete jumps to next BB

• invert jumps as needed, drop args

• append jumps to BB as needed

• delete unneeded labels

$\text{NOT}(e) l_3$

$l_0: CJUMP(e, l_1, l_2)$

$l_2: MOVE(x, y + z)$

$\text{Jump}(l_1)$

$l_0: MOVE(x, y)$

$\text{Jump}(l_2)$

$l_3: CALL(f)$

RETURN

Trace: a CFG path

of distinct BBs.

Reorder. (Greedy)

1. Init all BBs as "unmarked"

2. Repeat until all BBs are marked.

• Choose an unmarked block

• Mark a maximal unmarked trace

• Append blocks in that trace to output code

(while simplifying/fixing Jumps & CJumps)