

Intermediate Representation

Statements $S ::=$

MOVE(e_{dst}, e_{src})
 \hookrightarrow either MEM(e) or TEMP(t)

EXP(e)
EXP(CALL(e_f, e))

SEQ(s_1, s_2, \dots, s_n) $s_1; s_2; \dots; s_n$

JUMP(e)

CSJUMP(e, l_1, l_2)

LABEL(l) $l:$

RETURN(e)

OP ::= ADD | SUB | MUL | DIV | MOD |
AND | OR | XOR | NOT |
RSHIFT | LSHIFT | ARSHIFT |
EQ | NEQ | LT | LEQ | ...

Translating Joss to TIR.

$$E[e] = e' \quad S[s] = s' \quad F[f] = s$$
$$\begin{aligned} E[n] &= n & E[\text{true}] &= 1 \\ E[\text{false}] &= 0 \end{aligned}$$
$$E[e_1 + e_2] = \text{Add}(E[e_1], E[e_2])$$
$$[[x]] = \text{TEMP}(x)$$
$$E[f(e_1, \dots, e_n)] = \text{CALL}(\text{NAME}(f), E[e_1], \dots, E[e_n])$$
$$S[\text{int } x = e;] = \text{MOVE}(x, E[e])$$
$$S[\langle s_1; \dots; s_n \rangle] = \text{SEQ}(S[\langle s_1 \rangle], \dots, S[\langle s_n \rangle])$$

$S \llbracket \text{if}(e) s \rrbracket =$ $\text{cJump}(E \llbracket e \rrbracket, l_t, l_f);$
 $l_t:$ $S \llbracket s \rrbracket;$
 $l_f:$

$$S[\text{return } e] = \text{RETURN}(E[e])$$
$$F[\llbracket \tau \ f(\tau_1 x_1, \dots, \tau_n x_n) \{s\} \rrbracket =$$

$$f: \text{MOVE}(x_1, \text{arg}_0); \dots; \text{MOVE}(x_n, \text{arg}_{n-1}); S[s]$$
$$E[e_1 \& e_2] = \text{AND}(E[e_1], E[e_2])$$

```

MOVE(t, o);
CJUMP(Ell, l_t, l_f);
l_t: MOVE(t, Ell);
l_f:
t

```

$$S \llbracket \text{if } (e_1 \&\& e_2) \text{ s} \rrbracket = \text{cJump} \left(E \llbracket e_1 \&\& e_2 \rrbracket, l_t, l_f \right)$$

$l_t: S \llbracket s \rrbracket$
 $l_f:$

Better:

$$\begin{aligned} S \llbracket (e_1 \&\& e_2) \rrbracket s &= \\ c_{\text{Jump}}(\llbracket e_1 \rrbracket, l, l_f); \\ l : c_{\text{Jump}}(\llbracket e_2 \rrbracket, l_t, l_f); \\ l_t : S \llbracket s \rrbracket \\ l_f : \end{aligned}$$

$C[e, lb, lf] = \text{IR statement that}$
 - jumps to lb , if e is true
 - lf , otherwise

$$S[\llbracket f(e) \rrbracket s] = C[\llbracket e, lt, lf \rrbracket;$$

$$lt: S[s];$$

$$lf;$$
$$C \llbracket \text{true}, l_t, l_f \rrbracket = \text{JUMP}(l_t)$$

$C \models \text{false}, l_t, l_f \models \text{Jump}(l_f)$

$$C \llbracket e, l_t, l_f \rrbracket = C \llbracket e, l_f, l_t \rrbracket$$
$$C[e_1 \& e_2, l_t, l_f] = \text{Jump}(E[e_1], l, l_f);$$

$$l: \text{Jump}(E[e_2], l_t, l_f);$$

Better:

$$C[e_1 \& e_2, l_t, l_f] = C[e_1, l, l_f];$$

$$l: C[e_2, l_t, l_f]$$
$$C[e_1 \| e_2, l_t, l_f] = C[e_1, l_t, l];$$

$$l: C[e_2, l_t, l_f]$$

fall-back case

$$C \llbracket e, l_t, l_f \rrbracket = \text{Jump}(E \llbracket e \rrbracket, l_t, l_f)$$
$$\begin{aligned} S[\llbracket \text{while}(e) \ s \rrbracket] = & \quad l_w: C[\llbracket e, l_t, l_f \rrbracket] \\ & \quad l_t: S[\llbracket s \rrbracket]; \text{Jump}(l_w); \\ & \quad l_f: \end{aligned}$$

Arrays

```
int [] a = new int [10];
```

$$E[e_1 [e_2]] =$$

```
MOVE(tn, E[er]);
```

null check

$$\text{CJump}(\text{EQ}(t_0, 0), l_{\text{notnull}}, l_{\text{err}})$$

terr: CALL (__exception)

$\ell_{\text{not null}}$:

MOVE(t_i , $[l_2]$); bounds check

$$C_{\text{Jump}}(\text{LTU}(t_i, \text{MEM}(t_{a-4})), l_{\text{inbound}}, l_{\text{err}})$$

linkword:

$$\text{MEM}(t_n + t_1 \times 4 + 4)$$
