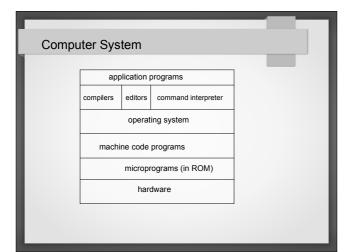


# software for using the hardware computer resources controls shares program development environment kernel = operating system



Operating System

direct access to hardware not allowed

user mode x kernel mode

hides difficulties of using hardware

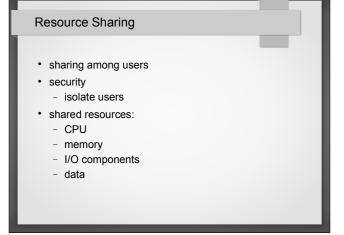
interface between user and hardware

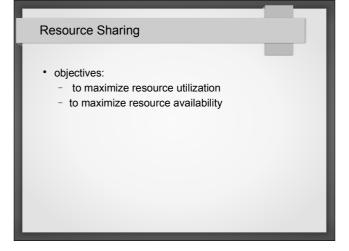
system calls

# for user programs to interact with operating system get operating system to perform a task for them a library routine for every system call user program uses library routine

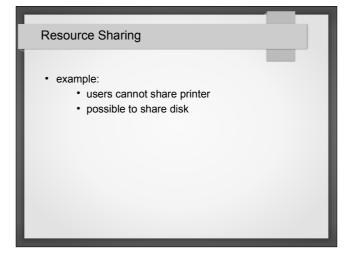
Operating System Responsibilities

• resource sharing
• virtual machine

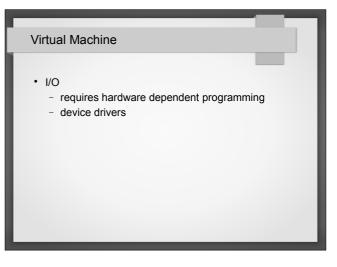


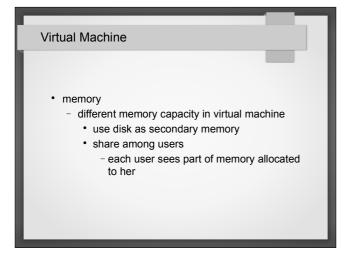


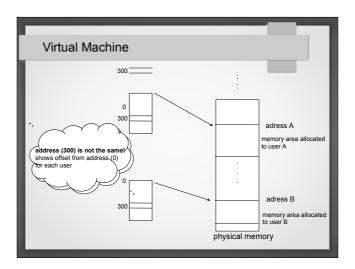
# Provided services: - define user interface - system calls - sharing and usage control of resources in multi-use systems - prevent race for resources - mutual exclusion - allow users to share data (shared memory) - resource scheduling - I/O scheduling - error handling



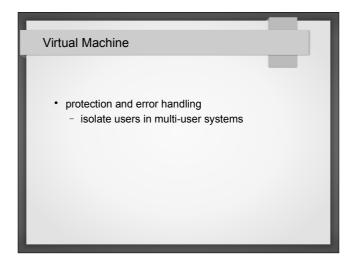
# Virtual Machine as if single user resource sharing transparent to user virtual machine may be different from actual physical machine: I/O memory file system protection and error handling program interaction program control



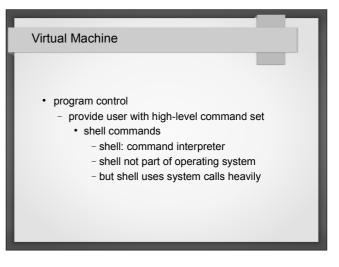




# • file system - for longterm storage of program and data - on disk - use symbols to acces info instead of physical addresses • naming - all accessed as files in UNIX



# Program interaction in runtime for example one program may use output of another program as input



# Types of Operating Systems

- · mainframe operating systems
- · server operating systems
- multi-processor operating systems
- PC operating systems
- real-time operating systems
- embedded operating systems
- · smartcard operating systems

## Mainframe Operating Systems

- for heavily I/O bound tasks
- · three main services:
  - batch mode
    - · non-interactive, routine tasks
  - e.g. preparing employee paychecks
  - transaction processing
    - · e.g. airline reservation systems
  - time-sharing
    - · multiple remote users running tasks
    - e.g. database
  - e.g.: OS/390

# Server Operating Systems

- · on servers
  - PCs with high resource capacities
  - workstations
  - mainframe systems
- · services for multi-users over a network
  - hardware and software sharing
  - e.g: printer services, file sharing, web access
- · e.g.: UNIX, Windows 2000

# Multi-Processor Operating Systems

- · for multi-processor systems
- · to increase computing power
- · based on interconnection between processors:
  - parallel systems
  - networked computers
  - multi-processor computers
- · special operating system features required
  - design objectives similar to server operating systems
  - extra features for interconnection and communication between processors

# PC Operating Systems

- · efficient and easy to use interface
- · office applications
- e.g.:
  - Windows 98, 2000, XP
  - Macintosh
  - Linux

# Real-Time Operating Systems

- · time constraints important
- · industrial control systems
  - feedback
- two types:

  - hard real-timetime constraints compulsory
    - e.g. robots in car production line
  - soft-real-time
    - possible not to obey some constraints
      e.g. multimedia systems
- · örnek: VxWorks ve QNX

# **Embedded Operating Systems**

- · palm computers and embedded systems
- · limited operation
- · special purpose
- e.g.: TV, microwave oven, cell phones, ...
- in some systems, size, memory and power consumption constraints
- e.g.: PalmOS, Windows CE

# **Smart-Card Operating Systems**

- smallest operating system
- · on credit card sized cards with processor
- · strict memory and CPU constraints
- · some are dedicated e.g. elektronic payments
- · some may have several functionalities
- usually special purpose operating systems developed by card companies
- some Java based
  - possible to load and execute small JAVA programs (applet)
  - some may execute more than one applet
    - multi-programming, scheduling, resource sharing and protection

# Main Kernel Architectures

- · monolithic
- modular
- layered
- · virtual machine
- exo-kernel
- · server-client model

# Monolithic

- · no general structure
- all services and functionalities included in operating system
- all functional procedures
  - at the same level
  - may interact with each other
- large

### Modular

- · minimal kernel
- services added to kernel at runtime as they are needed
  - e.g. device drivers
- small kernel size
- slower

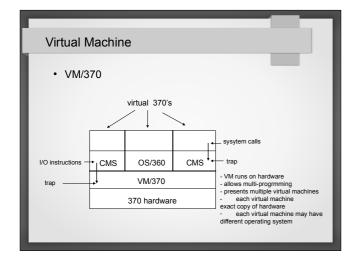
# Layered

- · layered structure
  - hierarchical
- e.g.: THE operating system

5	operator	
4	user programs	• layer 0:
3	I/O control	•
2	operator – process interaction	Each layer layers bel
1	memory and disk control	e.g.: for la
0	CPU sharing and multi-programming	

· layer 0: processor · layer 1: memory management

Each layer independent of operations of layers below. e.g.: for layer 2 operations, data may be on memory or disk



### **Exo-Kernel**

- · developed at MIT
- · similar to the virtual machine concept
  - · copy of system
  - difference: each virtual machine gets subset of system resources
- · external kernel
  - controls that virtual machines do not exceed thir allocated resources
- each virtual machine may have different operating systems

### Server-Client Model

- · minimal kernel micro-kernel
- · most of operating system in user mode
- · server and client processes
  - e.g. file read operation
    - · client process asks from server process
    - server carries out operation
    - · gives reply to client
- kernel coordinates communication and interaction between servers and clients

# Server-Client Model

- · servers in user mode
  - file server
  - process server
  - terminal server
  - memory server
- · operating system consists of many smaller sub-units:
  - easy to manage
  - error in one does not affect others (units do not access hardware directly)
  - implementation problems: not possible to implement especially some I/O device drivers at user mode
- · suitable for distributed systems