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#### Online Gradient Descent

#### Solution

a) The pythagorean inequality becomes apparent when we expand by  $\hat{y}$ :

$$||y - u||^2 = ||y - \hat{y} + \hat{y} - u||^2$$
$$= ||(y - \hat{y}) + (\hat{y} - u)||^2$$

This way we can apply the parallelogram law  $||a+b||^2 = ||a||^2 + ||b||^2 + 2\langle a,b\rangle$ :

$$= ||y - \hat{y}||^2 + ||\hat{y} - u||^2 + 2\langle y - \hat{y}, \hat{y} - u\rangle$$

Since the  $\langle a, b \rangle \geq 0$ :

$$\geq ||y - \hat{y}||^2 + ||\hat{y} - u||^2$$

b) Consider the distance between the  $w_{t+1}$  and u:

$$||w_{t+1} - u||^2$$

We can expand  $w_{t+1}$  with our step rule before applying  $P_U$ 

$$= ||w_t - \eta \nabla f_t(w_t) - u||^2 - \epsilon$$

Where  $\epsilon$  describes the nonnegative length lost in the projection. Without  $\epsilon$  clearly:

$$||w_{t+1} - u||^{2} \le ||w_{t} - \eta \nabla f_{t}(w_{t}) - u||^{2}$$

$$= ||(w_{t} - u) + (-\eta \nabla f_{t}(w_{t}))||^{2}$$

$$= ||(w_{t} - u)||^{2} + ||(-\eta \nabla f_{t}(w_{t}))||^{2} + 2\langle w_{t} - u, -\eta \nabla f_{t}(w_{t})\rangle$$

$$= ||(w_{t} - u)||^{2} + \eta^{2}||\nabla f_{t}(w_{t})||^{2} - 2\eta\langle w_{t} - u, \nabla f_{t}(w_{t})\rangle$$

Remember the term we started with:

$$||w_{t+1} - u||^2 \le ||(w_t - u)||^2 + \eta^2 ||\nabla f_t(w_t)||^2 - 2\eta \langle w_t - u, \nabla f_t(w_t) \rangle$$
$$2\eta \langle w_t - u, \nabla f_t(w_t) \rangle \le ||(w_t - u)||^2 + \eta^2 ||\nabla f_t(w_t)||^2 - ||w_{t+1} - u||^2$$

Divide by  $2\eta$ :

$$\langle w_t - u, \nabla f_t(w_t) \rangle \le \frac{||(w_t - u)||^2 - ||w_{t+1} - u||^2}{2\eta} + \frac{\eta}{2} ||\nabla f_t(w_t)||^2$$

c) From convexity we know that a (differentiable) function stays above any linear local approximation, that is to say for x, y in the domain of f:

$$f(x) \ge f(y) + \langle \nabla f(y), x - y \rangle$$
  

$$\Leftrightarrow -\langle \nabla f(y), x - y \rangle \ge f(y) - f(x)$$
  

$$\Leftrightarrow \langle \nabla f(y), y - x \rangle \ge f(y) - f(x)$$

Translating to our problem, we obtain:

$$f_t(w_t) - f_t(u) \le \langle \nabla f_t(w_t), w_t - u \rangle$$

We can replace the right side with our inequality from b):

$$f_t(w_t) - f_t(u) \le \frac{||(w_t - u)||^2 - ||w_{t+1} - u||^2}{2\eta} + \frac{\eta}{2}||\nabla f_t(w_t)||^2$$

The sum  $\sum_{t=1}^{T} f_t(w_t) - f_t(u)$  can then be decomposed to examine both sums on the right hand side individually:

$$\sum_{t=1}^{T} \frac{||(w_t - u)||^2 - ||w_{t+1} - u||^2}{2\eta} = \frac{1}{2\eta} \sum_{t=1}^{T} ||(w_t - u)||^2 - ||w_{t+1} - u||^2$$

It can be seen that this sum is telescopic, that is, it expands to:

$$\frac{1}{2\eta}(||w_1-u||^2-||w_2-2||^2+||w_2-u||^2-||w_3-u||^2+||w_3-u||^2+\dots)$$

leaving only the first and last term.

$$= \frac{||w_1 - u||^2 - ||w_{T+1} - u||^2}{2\eta}$$

Since  $w_1 = 0$ ,  $||u||^2 \le D$  we obtain:

$$\leq \frac{D^2 - ||w_{T+1} - u||^2}{2\eta} \\ \leq \frac{D^2}{2\eta}$$

Returning to the original expression and examining the other summand:

$$\frac{\eta}{2} \sum_{t=1}^{T} ||\nabla f_t(w_t)||^2$$

Since that  $\nabla |f_t| \leq G$ :

$$\leq \frac{\eta}{2} \sum_{t+1}^{T} G^2$$
$$= \frac{\eta}{2} T G^2$$

This concludes the proof that:

$$\max_{u \in U} \sum_{t=1}^{T} f_t(w_t) - f_t(u) \le \frac{D^2}{2\eta} + \frac{\eta}{2} TG^2$$

d) To find the optimal  $\eta$  we investigate the upper bound for critical points.

$$\frac{d}{d\eta} \left( \frac{D^2}{2\eta} + \frac{\eta}{2} T G^2 \right) \stackrel{!}{=} 0$$

$$\frac{-D^2}{2\eta^2} + \frac{T G^2}{2} = 0$$

$$\frac{T G^2}{2} = \frac{D^2}{2\eta^2}$$

$$\eta^2 T G^2 = D^2$$

$$\eta^2 = \frac{D^2}{T G^2}$$

$$\eta = \frac{D}{\sqrt{T} G}$$

This is a minimum, since the second derivative is  $\frac{D^2}{6\eta^3}$ , which is positive. Inserting this into the original equation yields a worst case regret of:

$$\frac{D^2}{2\left(\frac{D}{\sqrt{T}G}\right)} + \frac{\left(\frac{D}{\sqrt{T}G}\right)}{2}TG^2$$

$$= \frac{D^2\sqrt{T}G}{2D} + \frac{D}{2\sqrt{T}G}TG^2$$

$$= \frac{D\sqrt{T}G}{2} + \frac{D\sqrt{T}G}{2}$$

$$= \frac{D\sqrt{T}G}{2} + \frac{D\sqrt{T}G}{2}$$

## Formal proof of the existence of Lagrange Multipliers

Let  $f, g\mathbb{R}^d \to \mathbb{R}$ .

#### Solution

We want to show that the problem  $\min_x f(x)$  under the constraint g(x) = 0 is solved by  $\nabla f(x) + \nu g(x) = 0$ . For this we first introduce the parametrisation function h(t), such that the image of h is  $\{x|g(x)=0\}$ . For this function we know that  $g(h(t)) = 0 \forall t$ . In particular, this means that it is a constant function and thererfore by definition  $\frac{dg(h(t))}{dt} = 0$ . Writing this out with the chain rule we obtain:

$$\nabla g(h(t))Dh(t) = 0.$$

On the other hand f(h(t)) also fulfills the condition that it only is defined over all x for which g(x) = 0. Taking the derivative and setting to zero to find the minimum we find:

$$\nabla f(h(t*))Dh(t*) = 0,$$

with h(t\*) = x\*, the minimal point of the optimization problem. This means that  $\nabla f(h(t))$  and as  $\nabla g(h(t))$  are perpendicular to range(Dh(t)). But we can assume that Dh(t) has full rank the space perpendicular to Dh(p) has only dimension d - (d-1) = 1 and therefore  $\nabla g(h(t)) = \nu \nabla f(h(t*))$  and thus the proposition:

$$\nabla f(x*) + \nu g(x*) = 0$$

holds true.

# Differentiable approximation of $\mathcal{L}^1$ -approximation

### Solution

a) Given:

$$\min_{x,y} ||y||_1$$

Subject to 
$$y - Ax + b = 0$$

We can rearrange into the Lagrangian:

$$\mathcal{L}(y, x, y) = \min_{x, y} ||y||_1 + \lambda(y - Ax + b)$$

The dual is then:

$$g(\lambda) = \inf_{y,x} \mathcal{L}(y,x,\lambda)$$

To obtain our constrained version we can differentiate this wrt x and y. Starting with y:

$$\inf_{y} ||y||_1 + \lambda^t y$$

It can be seen that this is lower bounded by 0 if the all entries of  $\lambda$  are in [-1,1], as any negative terms are canceled out by the norm. Without this restriction, however, it is unbounded. Therefore:

$$= \begin{cases} 0 & \text{if} |\lambda_i| \le 1 \ \forall i \\ -\infty & otherwise \end{cases}$$

Differentiating wrt x yields:

$$\inf_{x} - \lambda^t Ax$$

For dual feasibilty we require  $\lambda A = 0$ , yielding our second constraint. Substituting this for the infimum into the dual problem leaves:

$$g(\lambda) = \max_{\lambda} \lambda^t b$$

subject to 
$$|\lambda_i| \leq 1 \ \forall i, \ \lambda^t A = 0$$

- b)
- c)