#### **USCC**

# Programming Assignment 1 – Recursive Descent Parsing

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#### Introduction

In this assignment, you will write large portions of a recursive descent parser for the USC programming language. Before you begin working on this assignment, you should read the USC language reference, which is in the same directory as this document. You will need to repeatedly refer to the grammar definitions while working on this assignment, so you may want to print out those pages from the language reference.

The scanner for USCC has already been provided to you. It is implemented via flex, and the input file is scan/usc.1. The list of tokens are in scan/Tokens.def, and they are defined in an enum. Provided your code is using namespace uscc::scan, the tokens can be accessed via a Token:: prefix. For example, Token::Key\_while would access the while keyword token. The Parser class, which uses the flex generated scanner, is declared in parse/Parse.h, and implemented in parse/Parse.cpp, parse/ParseExpr.cpp, and parse/ParseStmt.cpp.

You will notice that the Parser class has a large number of functions prefixed with parse, such as parseProgram, parseStmt, parseCompoundStmt, and so on. These functions are the recursive descent parsing functions, most of which you will implement in this assignment. Each of these functions returns a std::shared\_ptr to a specific type of abstract syntax tree (or AST) node. The AST is used as the initial intermediate representation of the source program. Each of these nodes is declared in parse/ASTNodes.h. For this assignment, you do not need to worry about the implementation of any of the nodes. You simply need to ensure that you are constructing the correct types of nodes during the parse.

The parse functions are speculative when invoked. For example, if parseWhileStmt is called, it does not mean that there definitely is a while statement upcoming. Rather, parseWhileStmt will first determine whether or not it thinks the upcoming token matches a while statement — specifically, if the current token is the while keyword. If the current token is not a while token, then parseWhileStmt will simply return an empty (or null) shared\_ptr to denote that a while statement is not a match for the current token stream.

However, suppose that parseWhileStmt encounters a while keyword. That means that the token stream can only be valid if the rest of the while statement is correct. However, if something in the while statement is malformed, parseWhileStmt should throw an exception to denote there was an error in the parse. The most common exception you will manually throw is a ParseExceptMsg, though there are other types of exceptions as defined in parse/ParseExcept.h. These exceptions should then be caught somewhere, ideally as deep in the call stack as possible. This is because ideally, the parser should catch as many errors as possible in one pass. It would be very annoying for end users if it simply stopped on the first error in a source file that had ten errors.

In this particular example, if you look at the partial implementation of parseStmt (in ParseStmt.cpp), you will see that there is a try/catch that catches any exceptions of type ParseExcept and then uses reportError to actually report the error. Errors that are reported are then written to stderr upon conclusion of the parse.

There are also several helper functions that are provided to aid with the parsing. The peekToken function simply returns the current token in the token stream. If there are multiple possibilities you want to peek for, you can use peekIsOneOf, which takes in an initializer\_list of tokens to choose from. Similarly, getTokenTxt returns the actual C-style string for the current token, which is important for tokens that could be a variety of strings, such as identifiers. The consumeToken function "eats" the current token and then moves on to the next token that's not a new line or comment. Since a very common operation is to peek for a specific token and then consume it if it is that token, this combined functionality is in the peekAndConsume function.

For cases where you absolutely require a specific token, or sequence of tokens, you can use matchToken or matchTokenSeq. These functions will match and consume either one or a sequence of tokens. The difference between these functions and peekAndConsume is that these functions will throw a TokenMismatch exception in the event of a mismatch. Returning to the while statement example, if you peekAndConsume a while token, you know for a fact that the next token must be a open parenthesis in a valid program. So you can use matchToken to process the parenthesis.

The last set of helper functions that have some use in error correction is the two consumeUntil functions. These functions mostly will be used by the catch blocks. If you look at the code for parseStmt again, you will see consumeUntil is used to skip all tokens until the next semi-colon. This way, if there's a major error in a statement, the parser can attempt to continue on the subsequent statement.

In this assignment, you should only have to edit ParseStmt.cpp and ParseExpr.cpp. Furthermore, a handful of the functions in these files have already been implemented for you, but you will have to implement the rest. Specifically, parseDecl, parseAssignStmt, parseExpr, parseExprPrime, and parseIdentFactor should not need to be modified.

In order to validate that your parser works correctly, there is a parsing test suite. To run the test suite run the following command from inside the tests directory:

#### \$ python testParse.py

When you start this assignment, only one of the 23 tests (Err\_Parse03) will pass. The tests are broken down into two types. There are some tests which are intended to cause parse errors, and these tests confirm that your parser identifies the same errors. The other set of tests should parse and generate an AST – the expected AST is compared against the AST your program generated. Note that several of these parse tests are actually semantically invalid USC programs, but in PA1 we aren't checking for semantic validity.

Upon conclusion of this assignment, all of the tests in the test suite should pass.

You also can run USCC directly from the command line to look at a specific test case. In order to output the AST of a program, you can run the command with the -a switch.

For example, if you wanted to output the AST of test002.usc, you could run the following command (again, from the tests directory):

```
$ ../bin/uscc -a test002.usc
```

Once you've implemented the first three items in the implementation section, this command will output:

Alternatively, if there are errors you will see a list of errors. For example, if you try to compile parse03e.usc, with the following command:

# Implementation

It is recommended that you implement the parsing functions in the order outlined in these instructions. This will allow you to verify the functions work as you go along.

#### parseCompoundStmt

The interior of a compound statement can begin with zero or more declarations and is followed by zero or more statements. So you'll want to use parseDecl and parseStmt to find the declarations and statements. The ASTCompoundStmt node has addDecl and addStmt to append the declaration/statement nodes as children. Remember that because nodes handled as shared pointers, you have to use std::make\_shared to dynamically create instances of a node.

#### parseReturnStmt

Next, implement parseReturnStmt. In USCC, return statements can either be void or return a value. When you implement this, you need to update parseStmt so that it calls parseReturnStmt, as well.

#### parseConstantFactor and parseStringFactor

If you look at the current implementation of parseAndTerm in ParseExpr.cpp, you will see that it directly calls parseFactor. This is actually not correct as per the grammar, but for the moment we just want to skip all the other grammar rules and jump directly to factor.

You should implement parseConstantFactor (which is for constant numeric expressions) and parseStringFactor (which is for string expressions). Then, you should update parseFactor so that after checking for an ident factor, it checks for constants and strings, as well.

You should now have three additional tests pass: AST\_002, Err\_Parse01, and Err\_Parse02. Note that the error tests will fail if your error messages are not identical to the expected ones, so you may have to edit your error messages so that they match the test cases.

#### parseParenFactor

Next, implement parseParenFactor and hook it up to the parseFactor function. Now the AST\_003 test should also pass.

#### parseIncFactor and parseDecFactor

Now implement parseIncFactor and parseDecFactor, and hook these up to parseFactor, as well. You should now also pass the AST\_004 test case.

#### parseValue

You should now implement parseValue. Once you've implemented parseValue, you should change parseAndTerm so it now calls parseValue instead of parseFactor. Now the AST\_005 test should also pass.

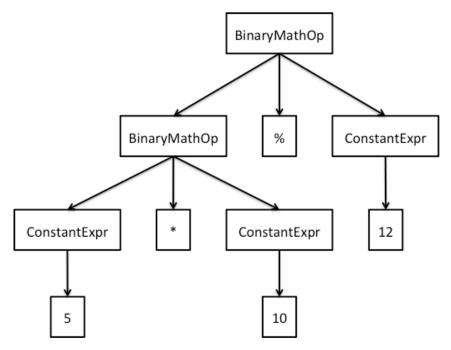
#### parseTerm and parseTermPrime

Now implement parseTerm and parseTermPrime. The reason there is a function for term prime is because the term grammar rule is left-recursive. So you will need to transform the grammar so it is right-recursive instead, and then implement these two rules.

There's one aspect of parseTermPrime that may be a little confusing. Suppose parseTermPrime sees a \* token. In this case, it means that a valid program must have an RHS value. So you can parse the value for the RHS operand. However, it's possible that after you get the RHS operand, it is followed by another term prime, since it's possible the expression might be along the lines of:

#### 5 \* 10 % 12

In the above case, this branch of the AST should actually have the modulus node at the top:



You need this *rightmost derivation* to ensure that binary operators with the same precedence are evaluated left to right in a post-order traversal of the AST. However, by default a recursive parser will produce a leftmost derivation. To solve this issue, after grabbing the RHS value you need to call parseTermPrime again, and see if there's another term prime. If there is another term prime, you should actually return that ASTBinaryMathOp and not the one you originally found. This guarantees a rightmost derivation as above.

Next, you should update parseAndTerm so that it calls parseTerm instead of parseValue. You should now pass AST\_006.

## parseNumExpr/Prime

You should now implement the rules for parseNumExpr/Prime. Since these are also binary operators, you need to handle the leftmost derivation as with parseTermPrime. Then update parseAndTerm so that it calls parseNumExpr instead. You should now also pass AST\_007 and AST\_013.

# parseRelExpr/Prime

These are very similar to parseNumExpr/Prime, except they are relational operators. Once you update parseAndTerm so that it calls parseRelExpr, you should additionally pass AST\_008.

#### parseAndTerm/Prime

One of the last expression rules to implement is parseAndTerm/Prime. Same rules as all other binary operators apply. Once you've implemented this, you will also pass AST\_009.

### parseWhileStmt

Now back in ParseStmt.cpp, implement parseWhileStmt and update parseStmt so that it also checks for while loops. You should now also pass AST 010.

## parseExprStmt and parseNullStmt

Now implement these two types of statements, and hook them up to parseStmt. You should then also pass AST\_011, AST\_015, and AST\_016.

## parselfStmt

The last statement type to parse is if statements. Remember that ifs may or may not have an else associated with them. Once hooked up to parseStmt, you should now pass every test other than Err\_06.

## parse Addr Of Array Factor

The last thing to implement is back in ParseExpr.cpp. Implement parseAddrOfArrayFactor and hook up parseFactor to check for this as well. You should now pass all the tests.

## Conclusion

If you pass all the tests, you now have a working parser for the USC language. However, if you look at the test programs, you will see that the majority of them are actually not valid USC code. This is because the parser only checks whether or not the file conforms to the context-free grammar for the language. In the next assignment, you will add semantic checks which will ensure that the program conforms to all of the rules of the language.