

Privacy Engineering (70018)

Zero-Knowledge Proofs

Problem 1. Sequential Composition. Consider an arbitrary interactive proving system $\pi = (p, v)$. Let $\Pi_N = (P_N, V_N)$ be an interactive proving system in which π is executed N independent times in sequence such that V_N accepts iff v accepts when invoked on the same common input in all N runs, and P_N invokes p on the same common input in all N runs.

- If the completeness and soundness errors of π are the constants c and s respectively, what are the completeness and soundness errors, C_N and S_N respectively, of Π_N ?
- If p achieves perfect zero-knowledge, what level, if any, of zero-knowledge does P_N achieve for all N ?
- If v has a knowledge error equal to the constant k , what is the knowledge error, K_N , of V_N ?
- If $s = \frac{1}{2}$, what is the minimum value of N required for S_N to be strictly less than 10^{-40} ?

Problem 2. Non-interactive Arguments. Assume that $\pi = (p, v)$ is the proving system for graph isomorphism knowledge defined in section 3.2 on page 11 of the lecture slides.

- What is the expected number of attempts for a cheating prover without knowledge of ϕ to construct a valid argument in π ?
- Let $\Pi_N = (P_N, V_N)$ be the sequentially composed version of N invocations to π as done in Problem 1. Specify Π_N in a pseudo-code style similar to that of π .
- Using Π_N as a basis, specify a non-interactive argument version of it $\bar{\Pi}_N = (\bar{P}_N, \bar{V}_N)$ in pseudo-code. Make sure $\bar{\Pi}_N$ achieves the same soundness, completeness and knowledge error values as Π_N for all values of $N \leq 256$.
- What is the expected number of attempts, in terms of N , for a cheating prover without knowledge of ϕ to construct a valid argument in $\bar{\Pi}_N$?
- Does \bar{P}_N reveal *any* knowledge?