

GRAPHS

MSCI 240: Algorithms & Data Structures

lecture summary

lots of definitions

more definitions

graph data structures

how to represent graphs in code

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2

Topic	Building Java Programs	Algorithms (Sedgewick)
classes, ADTs	chapter 8	1.2
arrays	chapter 7	
ArrayList<T>	chapter 10	1.3
Stack/Queue	chapter 14, (11)	1.3
LinkedList	chapter 16	1.3
Complexity		1.4
Searching	chapter 13	pp. 46-47
Sorting		chapter 2.1-2.3
Recursion	chapter 12	1.1 (p. 25)
Binary Trees	chapter 17	chapter 3.1-3.2
Dictionaries	chapter 18.1	chapter 3.4
Graphs	N/A (Wikipedia good)	chapter 4.1
Heaps/Priority Queues	chapter 18.2	chapter 2.4

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3

graph definitions

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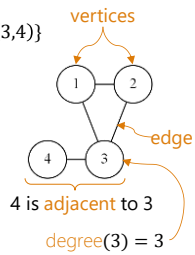
4

a **graph** is a set of vertices (a.k.a. nodes), V , and a set of edges, E , where each edge connects a pair of vertices

e.g., $V = \{1, 2, 3, 4\}$, $E = \{(1, 2), (1, 3), (2, 3), (3, 4)\}$

adjacent: u is adjacent to $v \Leftrightarrow (u, v) \in E$

degree (of a vertex $v \in V$): the number of edges connected to a vertex



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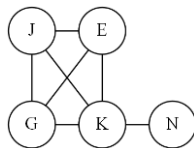
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e.g. (social),

$V = \{\text{Jerry, Elaine, George, Kramer, Newman}\}$

adjacent nodes?

degree of J, E, G, K, N?

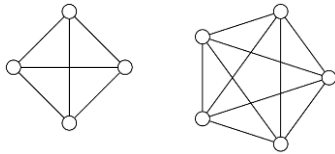


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complete graph: a graph where each vertex is connected to every other vertex

e.g., complete graphs of size 4 and 5



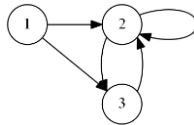
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directed graph: each edge in E is an **ordered pair** has a direction

in-degree (of $v \in V$): number of edges entering a vertex
in-degree of 1,2,3?

out-degree (of $v \in V$): number of edges leaving a vertex
out-degree of 1,2,3?



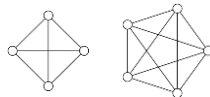
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how many edges are in a **complete** (undirected) graph?

let $n = |V|$ (number of vertices)

let $m = |E|$ (number of edges)



$$m = (n-1) + \dots + 2 + 1 = \sum_{i=1}^{n-1} i = \frac{n \cdot (n-1)}{2} \in O(n^2)$$

$$\text{alternatively, } m = \binom{n}{2} = \frac{n!}{2!(n-2)!} = \frac{n \cdot (n-1)}{2} \in O(n^2)$$

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more definitions

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dense graph

when m is in $\Theta(n^2)$ (e.g., almost all the edges)

sparse graph

when $m \ll n^2$ (e.g., only 1-2 edges per vertex, max, $\therefore m \in O(n)$)
(\ll means **way** less than)

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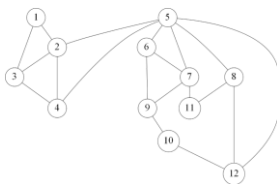
11

path: sequence of vertices connected by edges

e.g., $[1,2,5,7,11,8,12]$ is a path from 1 to 12

length of path: number of edges along the path

e.g., $[1,2,5,7,11,8,12]$ has length 6



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12

graph data structures

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13

there are two primary ways to represent graphs:

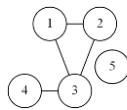
adjacency list
good for sparse graphs

adjacency matrix
good for dense graphs

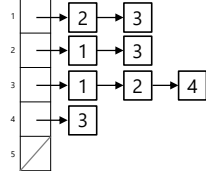
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$V = \{1, 2, 3, 4, 5\}$
 $E = \{(1, 2), (1, 3), (2, 3), (3, 4)\}$



adjacency list:



adjacency matrix:

	1	2	3	4	5
1	0	1	1	0	0
2	1	0	1	0	0
3	1	1	0	1	0
4	0	0	1	0	0
5	0	0	0	0	0

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IntGraphList class (undirected)

```
public class IntGraphList {
    private HashMap<Integer, LinkedList<Integer>> adjacencyList;

    public IntGraphList() {
        adjacencyList = new HashMap<>();
    }

    // precondition: value has not been added before
    public void addVertex(int value) {
        adjacencyList.put(value, new LinkedList<>());
    }

    // precondition: first and second must already be added as vertices
    public void addEdge(int first, int second) {
        adjacencyList.get(first).add(second);
        adjacencyList.get(second).add(first);
    }
}
```

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16

IntGraphMatrix class (undirected)

```
public class IntGraphMatrix {
    private boolean[][] adjMatrix;
    private int numVertices;

    public IntGraphMatrix(int numVertices) {
        adjMatrix = new boolean[numVertices][numVertices];
        this.numVertices = numVertices;
    }

    // no addVertex method (fixed number of vertices, numbered 1 to n)

    // precondition: 0 <= first, second < numVertices
    public void addEdge(int first, int second) {
        adjMatrix[first][second] = true;
        adjMatrix[second][first] = true;
    }
}
```

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how much space does an **adjacency list** use?

how much space does an **adjacency matrix** use?

sparse vs. dense graphs?

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18

graph summary

graphs have vertices and edges (pairs of vertices)

definitions: adjacent, degree (in/out), complete, sparse, dense, path

graphs can be **undirected** or **directed**

graphs can be stored with an **adjacency list** (better for sparse) or an **adjacency matrix** (better for dense)

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19

next:

breadth-first search

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20
