

#### SQL - Part 2

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# **Learning Outcomes**

- develop working knowledge of the SQL data definition language (DDL)
  - SQL data types
  - creating tables and constraints
- learn additional SQL data manipulation language (DML) concepts
  - insert, update, and delete rows
  - string pattern matching
  - inner and outer joins
  - views
  - auto-increment attributes
  - stored procedures
  - cursors
  - triggers
- Textbook sections (6<sup>th</sup> ed.): 3.2, 3.6, 3.9, 4.1, 4.2, 4.4, <del>5.1, 5.3</del>



# **Schema for Running Example**

- instructor = (<u>ID</u>, name, dept\_name, salary)
- teaches = (ID, course\_id, sec\_id, semester, year)
- section = (<u>course\_id</u>, <u>sec\_id</u>, <u>semester</u>, <u>year</u>, building, room\_no, time\_slot\_id)
- course = (<u>course\_id</u>, title, dept\_name, credits)
- department = (<u>dept\_name</u>, building, budget)
- ...

#### instructor

#### ID dept name salary name 10101 Srinivasan Comp. Sci. 65000 Finance 12121 90000 Wu 15151 Mozart Music 40000 22222 **Physics** 95000 Einstein El Said History 60000

#### teaches

	ID	course_id	sec_id	semester	year
	10101	CS-101	1	Fall	2009
ı	10101	CS-315	1	Spring	2010
ı	10101	CS-347	1	Fall	2009
ı	12121	FIN-201	1	Spring	2010
	15151	MU-199	1	Spring	2010
	22222	PHY-101	1	Fall	2009



# SQL DDL



### **Data Definition Language**

The SQL data-definition language (DDL) is used to specify of information about relations, including:

- the schema for each relation
- the domain of values associated with each attribute
- integrity constraints
- the set of indexes to be maintained for each relation.
- the physical storage structure of each relation (e.g., choice of InnoDB vs. MyISAM storage engine in MySQL)



# **Domain Types in SQL**

- **char(n).** Fixed length character string, with user-specified length *n*.
- varchar(n). Variable length character string, with user-specified maximum length n.
- int. Integer (a machine-dependent finite subset of the integers).
- **smallint.** Small integer (a machine-dependent subset of int).
- numeric(p,d). Fixed point number, with user-specified precision of p significant digits, with d digits to the right of decimal point.
- real, double precision. Floating point and double-precision floating point numbers, with machine-dependent precision.
- float(n). Floating point number, with user-specified precision of at least n digits.
- date, time. Calendar date (YYYY-MM-DD format), and time of day (hh:mm:ss format).



#### **Create Table Construct**

An SQL relation is defined using the create table command:

```
create table r(A_1 D_1, \dots, A_n D_n, \dots, A_n D_n, \dots, (integrity-constraint_1), \dots, (integrity-constraint_k)
[ Note: r is a relation, A_i is an attribute, and D_i is the domain of A_i]
```

Example:



### **Drop and Alter Table Constructs**

- drop table student
  - Deletes the table and its contents. Assumes table exists.
- drop table if exists student
  - Deletes the table and its contents if table exists.
- **delete from** student
  - Deletes all tuples from table, but retains table.
- alter table
  - alter table r add A D
    - Adds attribute with name A and domain D to relation r. All existing tuples in the relation are assigned null as the value for the new attribute.
  - alter table r drop A
    - Drops attribute A from relation r. Not supported by many databases.



## **Declaring Integrity Constraints**

- not null: disallows null values
- **primary key**  $(A_1, ..., A_n)$ : ensures uniqueness
- **unique**  $(A_1, ..., A_n)$ : ensures uniqueness (think superkey)
- **foreign key**  $(A_m, ..., A_n)$  **references**  $r(A'_m, ..., A'_n)$ : defines a foreign key in the child (referencing) table that points to a referenced key in a parent (referenced) table r
- default V: makes V the default value for an attribute
- Example:



# **Primary Keys and Superkeys**

Example of primary key and unique constraints in SQL:

```
create table customer (
    customer_id int,
    social_insurance_num numeric(9,0),
    first_name varchar(20),
    last_name varchar(20),
    primary key (customer_id),
    unique (social_insurance_num)
)
```

- Note 1: primary key and unique both identify superkeys. You can use primary key at most once per table but you can use unique more than once.
- Note 2: A primary key attribute cannot be null. The primary key constraint on an attribute implies the not null constraint.
- Note 3: For a **unique** attribute, the number of tuples that may have a *null* value is system-dependent (zero, one, or many).



## Foreign Keys

Example of a foreign key (a.k.a. referential integrity) constraint:

```
create table instructor (
    ID
                   char(5) primary key,
                   varchar(20) not null,
    name
                   varchar(20),
    dept_name
                   numeric(8,2),
    salary
    foreign key (dept_name) references department(dept_name)
create table department (
                   varchar(20) primary key,
    dept_name
                   varchar(20) not null,
    building
    budget
                   int
```

Note: In MySQL the referenced key (in this case dept\_name) must be a superkey of the referenced table, or a prefix of a multi-attribute primary key of the referenced table.



#### **Referential Actions**

- Referential actions define the behavior of the DB in cases when an update or deletion on the parent (i.e., referenced) table SQL statement affects a value referenced by a child (i.e., referencing) table.
- SQL 92 defines four actions:
  - cascade: automatically update or delete foreign key in matching rows of child table.
  - set null: set foreign key columns to null in matching rows of child table. (Assumes foreign key is nullable.)
  - set default: set foreign key columns to the default value in matching rows of child table.
  - no action: reject operation and generate error.
     (Also known as restrict in MySQL.)



# **Referential Actions (Cont.)**

Example: cascade on update, set null on delete



# **SQL DML**



## **Modifying Relations – Insertion**

Example: Add a new tuple to course.

insert into course values ('ECE-356', 'Databases', 'ECE', 0.5)

Example: Add a new tuple to course with credits set to null.

insert into course values ('ECE-356', 'Databases', 'ECE', null)



## **Modifying Relations – Updates**

Example: Give a 3% salary increase to all instructors whose salary is below \$80,000.

update instructor
set salary = salary \* 1.03
where salary < 80000</pre>



### **Modifying Relations – Deletion**

Example: Delete all instructors.

delete from instructor

- Example: Delete all instructors from the Math department.
  delete from instructor
  where dept\_name= 'Math'
- Example: Delete all instructors whose departments are in the EIT building.

delete from instructor
where dept\_name in (select dept\_name
from department
where building = 'EIT')



# **String Pattern Matching**

- SQL includes a string matching operator for comparisons on character strings. The operator like uses patterns that are described using two special characters (wildcards):
  - percent (%): matches any substring
  - underscore (\_): matches any one character
- Example: Find the names of all instructors whose name includes the substring "dar".

**select** *name* **from** *instructor* **where** *name* **like** '%dar%'

Example: Match the string "100 %".

like '100 \%' escape '\'

Note: Patterns are case-sensitive.



#### **More Joins**

- Join operations take two relations and return another relation as a result.
- **Join condition:** defines which tuples in the two relations match, and what attributes are present in the result of the join.
- Join type: defines how tuples in each relation that do not match any tuple in the other relation (based on the join condition) are treated.

Join types
inner join
left outer join
right outer join
full outer join

Join Conditions

natural

on < predicate>
using  $(A_1, A_1, ..., A_n)$ 



# More Joins (Cont.)

#### Relation course

course_id	title	dept_name	credits
BIO-301	Genetics	Biology	4
CS-190	Game Design	Comp. Sci.	4
CS-315	Robotics	Comp. Sci.	3

#### Relation prereq

course_id	prereg_id
BIO-301	BIO-101
CS-190	CS-101
CS-347	CS-101

#### Observe that

prereq information is missing for CS-315 and course information is missing for CS-437



#### **Inner and Outer Joins**

select \* from course inner join prereq on course.course\_id = prereq.course\_id

course_id	title	dept_name	credits	prereg_id	course_id
BIO-301	Genetics	Biology	4	BIO-101	BIO-301
CS-190	Game Design	Comp. Sci.	4	CS-101	CS-190

Note: some databases allow you to omit the keyword "inner".

select \* from course left outer join prereq on course.course\_id = prereq.course\_id

course_id	title	dept_name	credits	prereq_id	course_id
BIO-301	Genetics	Biology	4	BIO-101	BIO-301
CS-190	Game Design	Comp. Sci.	4	CS-101	CS-190
CS-315	Robotics	Comp. Sci.	3	null	null



#### **Outer Joins**

select \* from course natural right outer join prereq

course_id	title	dept_name	credits	prereq_id
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-347	null	null	null	CS-101

select \* from course full outer join prereq using (course\_id)

course_id	title	dept_name	credits	prereg_id
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-315	Robotics	Comp. Sci.	3	null
CS-347	null	null	null	CS-101

Note: full outer join is not supported in MySQL 5.0.



#### **Views**

- In some cases, it is not desirable for all users to see all the relations stored in a database instance.
- Consider a person who needs to know an instructor's name and department, but not the salary. This person should see a relation described in SQL by the following query:

**select** *ID*, *name*, *dept\_name* **from** *instructor* 

- A view provides a mechanism to hide certain data from the view of certain users.
- Any relation that is not part of the conceptual model but is made visible to a user as a "virtual relation" is called a view.



#### **View Definition**

A view is defined using the create view statement, which follows the form

create view view\_name as < query expression >

- where <query expression> is any legal SQL query.
- Once a view is defined, the view name can be used to refer to the virtual relation.
- A view need not be defined over a single table. It can be defined over the result set of a join or group by query.
- Note: Views are <u>dynamic</u>. In other words, changing the data in the relations referenced by a view causes analogous changes in the virtual relation corresponding to the view.



### **Example Views**

- A view of instructors without their salary
   create view faculty as
   select ID, name, dept\_name
   from instructor
- Find all instructors in the Biology department select name from faculty where dept\_name = 'Biology'
- Create a view of department salary totals create view departments\_total\_salary(dept\_name, total\_salary) as select dept\_name, sum(salary) from instructor group by dept\_name

**select** \* **from** *departments\_total\_salary* 



#### **Auto-increment Attributes**

- Auto-increment attributes can be used to automatically generate primary key values.
- Example:

- Note 1: ID = 1 should be assigned automatically to the first tuple.
- Note 2: The auto-increment type does not exist in the ER model and therefore it does not eliminate the need for weak entity sets.