

## Functional Dependencies and Boyce-Codd Normal Form (BCNF) Decomposition

Let  $X$  and  $Y$  be two sets of attributes. A functional dependency  $X \rightarrow Y$  holds if rows that have the same values of  $X$  also have the same values of  $Y$ .

Example: Consider a table with addresses and suppose that  $\text{postal\_code} \rightarrow \text{city}$ . What does that mean? What about  $\text{city} \rightarrow \text{postal\_code}$ ?

Note:  $X$  is a key if  $X \rightarrow$  all the attributes of a table

By the decomposition rule, if we know that  $A \rightarrow BC$ , then  $A \rightarrow B$  and  $A \rightarrow C$  are both true. Also,  $A \rightarrow A$  (by the reflexivity rule), so instead of saying  $A \rightarrow ABC$ , we can say  $A \rightarrow BC$ .

Recall: an FD  $X \rightarrow Y$  is trivial if  $Y$  is a subset of  $X$ . For example,  $A \rightarrow A$  or  $AB \rightarrow A$  are trivial.

### Boyce-Codd Normal Form (BCNF)

A table  $T$  is in BCNF if for all \*nontrivial\* functional dependencies  $X \rightarrow A$  that hold on  $T$ ,  $X$  contains a key of  $T$

Algorithm for BCNF decomposition

Input: Table  $T$ , Set  $F$  of all non-trivial functional dependencies that hold on  $T$

While there remains at least one FD in  $F$ , call it  $X \rightarrow Y$ , where  $X$  does not contain a key

Let  $B$  be the table containing  $X$  and  $Y$

Decompose  $B$  into two tables:

one with just  $X$  and  $Y$

one with all the columns of  $B$  except  $Y$

### Example

$T(ABCDE)$

$F = \{ AE \rightarrow BCD, A \rightarrow B, B \rightarrow A, DE \rightarrow C \}$

Decompose  $T$  into BCNF. Identify the candidate keys of the decomposed tables and prove that your decomposed schema is in BCNF.

Solution:

First, note that the  $AE$  is a key of  $T$ . So,  $AE \rightarrow BCD$  does not break BCNF.

Next, note that  $A \rightarrow B$  breaks BCNF.

Decompose  $T$  into  $T_1(AB)$ , with keys  $A$  and  $B$ , and  $T_2(ACDE)$  with key  $AE$

In the second iteration of the algorithm, note that  $DE \rightarrow C$  breaks BCNF

Decompose  $T_2$  into  $T_{2a}(CDE)$ , with key  $DE$ , and  $T_{2b}(ADE)$ , with key  $AE$

Final answer: T(AB), keys A or B, T2A(CDE) key DE, T2B(ADE) key AE

In BCNF because all the dependencies in F now contain a key on the left-hand side.

### Practice questions

1. Let T be a table with four attributes: ABCD. Assume  $F = \{ B \rightarrow ACD, C \rightarrow D, C \rightarrow A \}$ . Decompose T into BCNF, identify the keys of the decomposed tables, and prove that your decomposed schema is in BCNF.

2. Now assume  $F = \{ BD \rightarrow AC, B \rightarrow C, D \rightarrow A \}$ . Decompose T into BCNF, identify the keys of the decomposed tables, and prove that your decomposed schema is in BCNF.