My Template

Billy Inn

August 7, 2014

Contents

1	Dat	a Structures	3
	1.1	Fenwick Tree	3
		1.1.1 unidimension	3
		1.1.2 two-dimension	3
	1.2	RMQ	4
	1.3	Splay Tree	4
			4
		1.3.2 Implemented by array	7
2	Mat	th 1	1
	2.1	High Precision Class	1
	2.2	Number Theory	3
		2.2.1 Sieve of Eratosthenes	3
		2.2.2 Sieve of Euler	4
		2.2.3 Euclid Algorithm	4
		2.2.4 Euler Function	4
		2.2.5 Calculate the Inversion	5
		2.2.6 Chinese Remainder Theorem	5
		2.2.7 Bezout Theorem	5
	2.3	Combinatorics	6
		2.3.1 Calculate Combinations by Recurrence	6
	2.4	Matrix and System of Linear Equations	6
		2.4.1 Matrix Fast Power	6
	2.5	Numerical Method	7
		2.5.1 Adaptive Simpson's Rule	7
3	Gra	ph Theory 1	8
	3.1	Fundamentals	8
		3.1.1 Topo Sort	8
		3.1.2 Eulerian trail and circuit	9
	3.2	Shortest Paths	0
	3.3	Minimal Spanning Tree	0
		3.3.1 Fundamentals	0
		3.3.2 Kruskal algorithm	0

		3.3.3 Second-best Minimum Spanning Tree 21			
		3.3.4 Minimal Bottleneck Spanning Tree			
		3.3.5 Minimal Directed Spanning Tree			
	3.4	Bipartite Graph Matching			
		3.4.1 Augmenting Path Algorithm			
		3.4.2 Kuhn-Munkers Algorithm			
	3.5	Network Flows			
	3.6	Depth First Search			
		3.6.1 Calculate DFS Sequence			
		3.6.2 Multiplication LCA Algorithm(online) 27			
		3.6.3 Targan's Off-line LCA Algorithm 28			
4	Stri	ng 30			
	4.1	KMP			
	4.2	AC Automation			
	4.3	Suffix Array			
	4.4	Suffix Automation			
	4.5	String Hash			
	4.6	Longest Palindromic Substring			
5	Geometry 38				
	5.1	Basic			
	5.2	Problems About Circle and Sphere			
		5.2.1 Definition of circle			
		5.2.2 Calculate the intersection of line and circle 41			
		5.2.3 Calculate the intersection of circles 41			
	5.3	2D Algorithm			
		5.3.1 Judge whether point in polygon			
	5.4	3D Algorithm			
	5.5	Stimulated Annealing Algorithm			
		5.5.1 Minimum Sphere Coverage			
6	Dvr	namic Programming 45			
	6.1°	Longest Common Increasing Sequence			
7	Oth	er Topics 47			
•	7.1	Dichotomy			
		7.1.1 condition 1			
		7.1.2 condition 2			
		7.1.3 condition 3			
	7.2	Hash List			
	7.3	Construct a proper answer for N Queens Problem			

Chapter 1

Data Structures

1.1 Fenwick Tree

Test: CF 383C

1.1.1 unidimension

```
int C[maxn];
    int lowbit(int x){return x\&-x;}
    void add(int x, int v)
        while(x <= n)
            C[x]+=v;x+=lowbit(x);
    int getsum(int x)
10
11
        int ans=0;
12
        while(x>0)
13
14
            ans+=C[x];x-=lowbit(x);
15
16
17
        return ans;
18
```

1.1.2 two-dimension

```
int C[maxn][maxn];
int n,m;
int lowbit(int x){return x&-x;}
void add(int x,int y,int v)
```

```
5
       for(int i=x; i<=n; i+=lowbit(i))
6
            for(int j=y;j<=m;j+=lowbit(j))
7
8
9
               C[i][j]+=v;
10
11
    int getsum(int x, int y)
12
13
       int ans=0;
14
       for(int i=x;i>0;i==lowbit(i))
15
            for(int j=y;j>0;j-=lowbit(j))
16
17
               ans+=C[i][j];
18
19
20
       return ans;
21
            RMQ
    1.2
    Test: CF 6E
   int d[maxn][power];
    void RMQ_init(int* A,int n)
3
            for(int i=0; i< n; i++) d[i][0]=A[i];
4
            for(int j=1;(1<< j)<=n;j++)
                   for (int i=0; i+(1<< j)-1< n; i++)
                           d[i][j]=min(d[i][j-1],d[i+(1<<(j-1))][j-1]);
8
   int RMQ(int L,int R)
10
11
           int k=0;
12
           while((1 << (k+1)) <= R-L+1) k++;
13
           return min(d[L][k], d[R-(1 << k)+1][k]);
14
15
   1.3
           Splay Tree
    1.3.1
            Implemented by List
    Test: hdu1754,1890,3436,3487,poj3468,UVa11922
   const int maxn=200010;
   struct Node
```

```
{
 3
        Node *ch[2];
                        //children
 4
        Node *pre:
                        //ancestor
 5
                        //the number of nodes that are rooted in this node
 6
        int sz;
                        //the value to be maintained
        int val:
        inline void push_up()
                                    //like the function of segment tree
 8
 9
            sz=1+ch[0]->sz+ch[1]->sz;
10
11
        inline void push_down()
                                    //like the function of segment tree
12
13
    };
    Node *null=new Node();
                                    //virtual node
14
    struct SplaySequence
15
16
        Node seq[maxn];
17
        Node *root;
18
        int a[maxn];
19
        Node* build(int I, int r, Node *f)
20
21
            if(|>r) return null;
22
            int m=(1+r)>>1;
23
            Node *o=\&seq[m];
24
            o->pre=f;
25
            o->val=a[m];
26
            o->ch[0]=build(l,m-1,o);
27
28
            o->ch[1]=build(m+1,r,o);
            o->push_up();
29
            return o;
30
31
        void init (int sz)
32
33
            null -> sz=0;
34
            root=build(1,sz, null );
35
36
        //d=0 rotate toward the leftd =1 rotate toward the right
37
        inline void rotate(Node *x,int d)
38
39
            Node *y=x->pre;
40
            y->push_down();
41
            x->push_down();
42
            y->ch[d^1]=x->ch[d];
43
            if (x->ch[d]!=null) x->ch[d]->pre=y;
44
            x->pre=y->pre;
45
            if (y->pre!=null) y->pre->ch[y->pre->ch[1]==y]=x;
46
            x->ch[d]=y;
47
            y->pre=x;
48
```

```
y->push_up();
49
             if (y==root) root=x;
50
51
        //rotate x to the position below f and return x
52
53
         inline Node* splay(Node *x,Node *f)
54
            while(x->pre!=f)
55
56
                 if (x->pre->pre==f) rotate(x,x->pre->ch[0]==x);
57
                else
58
59
                     Node y=x->pre, z=y->pre;
60
                     int d=(z->ch[0]==y);
61
                     if (y->ch[d]==x) rotate(x,d^1), rotate(x,d);
62
                     else rotate(y,d), rotate(x,d);
63
64
65
            x->push_up();
66
            return x;
67
68
        //rotate the node that is rooted in r
69
        //and ranks k to the position below f, and return it
70
        inline Node* select(int k, Node *f, Node *r)
71
72
             Node *t=r;
73
74
            while(1)
75
                int s=t->r-t->l+1;
76
                int tmp=t->ch[0]->sz;
77
                 if (k \le tmp+s \&\& k \ge tmp) break;
78
                 if(k \le tmp) t = t - sch[0];
79
                else k=tmp+s,t=t->ch[1];
80
81
            return splay(t,f);
82
83
        //split o to two trees
84
        //the first k nodes are rooted in l, the other are rooted in r
85
        inline void split (Node *o,int k, Node* &I, Node* &r)
86
87
             l=select(k, null, o);
88
             r=l->ch[1];
89
             l -> ch[1] = null;
90
             r->pre=null;
91
             l -> push_up();
92
93
        //merge l and r, and return the root
94
```

```
inline Node* merge(Node *I,Node *r)
95
96
             Node *t=select(I->sz,null,I);
97
             t->ch[1]=r;
98
99
             r->pre=t;
             t->push_up();
100
             return t;
101
102
103
     } ss;
    void debug(Node *o)
                             //debug function
104
105
         if(o!=null)
106
107
             debug(o->ch[0]);
108
             printf ("%d\n",o->val);
109
             debug(o->ch[1]);
110
111
    }
112
```

1.3.2 Implemented by array

```
Test: POJ 3468
```

```
#include <cstdio>
    #define keyTree (ch[ ch[root ][1] ][0])
    const int maxn = 222222;
    struct SplayTree{
4
            int sz[maxn];
5
            int ch[maxn][2];
6
            int pre[maxn];
 7
            int root , top1 , top2;
 8
            int ss[maxn], que[maxn];
 9
10
            inline void Rotate(int x, int f) {
11
                    int y = pre[x];
12
13
                    push_down(y);
                    push_down(x);
14
                    ch[y][!\ f]\ = ch[x][f];
15
                     pre[ch[x][f]] = y;
16
                     pre[x] = pre[y];
17
                     if (pre[x]) ch [pre[y]] [ch[pre[y]] == y] = x;
18
19
                    ch[x][f] = y;
                     pre[y] = x;
20
21
                    push_up(y);
22
            inline void Splay(int x, int goal) {
23
```

```
push_down(x);
24
                   while(pre[x] != goal) {
25
                           if(pre[pre[x]] == goal) {
26
                                   Rotate(x , ch[pre[x]][0] == x);
27
28
                           } else {
                                   int y = pre[x], z = pre[y];
29
                                   int f = (ch[z][0] == y);
30
                                   if(ch[y][f] == x) \{
31
                                          Rotate(x, !f), Rotate(x, f);
32
                                   } else {
33
                                          Rotate(y , f) , Rotate(x , f);
34
                                   }
35
                           }
36
37
                   push_up(x);
38
                    if (goal == 0) root = x;
39
40
            inline void RotateTo(int k, int goal) {
41
                   int x = root;
42
                   push_down(x);
43
                   \textbf{while}(sz[\ ch[x][0]\ ]\ !=k)\ \{
44
                           if (k < sz[ch[x][0]]) {
45
                                   x = ch[x][0];
46
                           } else {
47
                                   k = (sz[ch[x][0]] + 1);
48
                                   x = ch[x][1];
49
50
                           push_down(x);
51
52
                   Splay(x, goal);
53
54
            inline void erase(int x) {
55
                   int father = pre[x];
56
                   int head = 0, tail = 0;
57
                   for (que[tail++] = x; head < tail; head ++) {
58
                           ss[top2 ++] = que[head];
59
                           if (ch[que[head]][0]) que[tail++]=ch[que[head]][0];
60
                           if (ch[ que[head] ][1]) que[tail ++] = ch[ que[head] ][1];
61
62
                   ch[ father ][ ch[father ][1] == x] = 0;
63
                   pushup(father);
64
           }
65
    66
67
            inline void NewNode(int &x,int c) {
68
                    if (top2) \times = ss[--top2];
69
```

```
else x = ++top1;
70
                      ch[x][0] = ch[x][1] = pre[x] = 0;
71
72
                      sz[x] = 1;
73
74
                      val[x] = sum[x] = c; /*specialized function*/
                      add[x] = 0;
75
             }
76
77
             inline void push_down(int x) {/* specialized function*/
78
                      if (add[x]) {
79
                              \mathsf{val}\,[\mathsf{x}] \ += \mathsf{add}[\mathsf{x}];
80
                              add[ch[x][0]] += add[x];
                              add[ch[x][1]] += add[x];
82
                              sum[ ch[x][0] ] += (long long)sz[ ch[x][0] ] * add[x];
83
                              sum[ ch[x][1] ] += (long long)sz[ ch[x][1] ] * add[x];
84
85
                              add[x] = 0;
                      }
86
87
             inline void push_up(int x) {
88
                      sz[x] = 1 + sz[ch[x][0]] + sz[ch[x][1]];
89
                      /* specialized function*/
90
                      sum[x] = add[x] + val[x] + sum[ch[x][0]] + sum[ch[x][1]];
91
             }
92
93
94
             /* initialize */
95
             inline void makeTree(int &x,int I, int r, int f) {
                      if (1 > r) return;
96
                      int m = (l + r) >> 1;
97
                      NewNode(x, num[m]);
                                                       /*varies according to the problem*/
98
                      makeTree(ch[x][0], I, m-1, x);
99
100
                      makeTree(ch[x][1], m + 1, r, x);
                      pre[x] = f;
101
                      push_up(x);
102
103
             inline void init (int n) {/* specialized function*/
104
                      ch [0][0] = ch [0][1] = pre[0] = sz[0] = 0;
105
                      add[0] = sum[0] = 0;
106
107
                      root = top1 = 0;
108
                      //for convience to add two nodes
109
                      NewNode(root, -1);
110
                      NewNode(ch[root][1], -1);
111
                      pre[top1] = root;
112
                      sz[root] = 2;
113
114
115
```

```
for (int i = 0; i < n; i ++) scanf("%d",&num[i]);
116
                     makeTree(keyTree , 0 , n-1 , ch[root ][1]);
117
                     push_up(ch[root ][1]);
118
                     push_up(root);
119
120
             inline void update( ) {/* specialized function*/
121
                     int | , r , c;
122
                      scanf("%d%d%d",&l,&r,&c);
123
                     RotateTo(I-1,0);
124
                     RotateTo(r+1,root);
125
                     add[ keyTree ] += c;
126
                     sum[keyTree] += (long long)c * sz[keyTree];
127
128
             inline void query() {/* specialized function*/
129
                     int | , r;
130
                      scanf("%d%d",&l,&r);
131
                     RotateTo(I-1, 0);
132
                     RotateTo(r+1, root);
133
                      printf ("%lld\n",sum[keyTree]);
134
             }
135
136
             int num[maxn];
137
             int val[maxn];
138
             int add[maxn];
139
             long long sum[maxn];
140
141
     }spt;
142
143
     int main() {
144
             int n , m;
145
             scanf("%d%d",&n,&m);
146
             spt. init (n);
147
             while(m --) {
148
                     char op [2];
149
                      scanf("%s",op);
150
                      if(op[0] == 'Q')
151
                              spt.query();
152
                      } else {
153
                              spt.update();
154
                      }
155
156
             return 0;
157
158
```

Chapter 2

Math

2.1 High Precision Class

```
const int maxn = 200;
    struct bign{
 2
      int len, s[maxn];
      bign() {
        memset(s, 0, sizeof(s));
        len = 1;
 6
 7
      bign(int num) {
 8
        *this = num;
 9
10
      bign(const char* num) {
11
        *this = num;
12
13
      bign operator = (int num) {
14
15
        char s[maxn];
        sprintf (s, "%d", num);
16
17
        *this = s;
        return *this;
18
19
      bign operator = (const char* num) {
20
21
        len = strlen(num);
        for(int i = 0; i < len; i++) s[i] = num[len-i-1] - '0';
22
        return *this;
23
24
      string str() const {
25
        string res = "";
26
        for(int i = 0; i < len; i++) res = (char)(s[i] + '0') + res;
27
        if (res == "") res = "0";
```

```
29
        return res;
30
      bign operator + (const bign& b) const{
31
32
        bign c;
33
        c.len = 0;
        for(int i = 0, g = 0; g || i < max(len, b.len); i++) {
34
           int x = g;
35
           if (i < len) \times += s[i];
36
           if (i < b.len) x += b.s[i];
37
          c.s[c.len++] = x \% 10;
38
          g = x / 10;
39
40
        return c;
41
42
      void clean() {
43
        while(len > 1 \&\& !s[len-1]) len--;
44
45
      bign operator * (const bign& b) {
46
        bign c; c.len = len + b.len;
47
        for(int i = 0; i < len; i++)
48
           for(int j = 0; j < b.len; j++)
49
            c.s[i+j] += s[i] * b.s[j];
50
        for(int i = 0; i < c.len-1; i++){
51
          c.s[i+1] += c.s[i] / 10;
52
          c.s[i] \% = 10;
53
54
        c.clean();
55
        return c;
56
57
      bign operator - (const bign& b) {
58
59
        bign c; c.len = 0;
        for(int i = 0, g = 0; i < len; i++) {
60
61
           int x = s[i] - g;
           if (i < b.len) \times -= b.s[i];
62
           if (x >= 0) g = 0;
63
           else {
64
65
            g = 1;
            x += 10;
66
67
          c.s[c.len++] = x;
68
69
        c.clean();
70
        return c;
71
72
      bool operator < (const bign& b) const{
73
         if(len != b.len) return len < b.len;</pre>
74
```

```
for(int i = len-1; i >= 0; i--)
75
           if (s[i] != b.s[i]) return s[i] < b.s[i];
76
         return false:
77
78
79
       bool operator > (const bign& b) const{
         return b < *this;
80
81
       bool operator <= (const bign& b) {
82
         return !(b > *this);
83
84
       bool operator == (const bign& b) {
85
         return !(b < *this) && !(*this < b);
86
87
       bign operator += (const bign& b) {
88
         *this = *this + b;
89
         return *this;
90
91
92
     };
     istream & operator >> (istream &in, bign& x) {
93
94
       string s;
       in >> s;
95
       x = s. c_str();
96
       return in;
97
98
     ostream& operator << (ostream &out, const bign& x) {
99
100
       out << x.str();
       return out;
101
102
```

2.2 Number Theory

2.2.1 Sieve of Eratosthenes

```
const int maxn=10000010;
    const int maxp=700000;
    int vis [maxn],prime[maxp];
    void sieve(int n)
                              //sieve the primes no larger than n
    {
 5
             int m=(int) \operatorname{sqrt}(n+0..5);
 6
             memset(vis,0, sizeof( vis ));
 7
             for(int i=2;i<=m;i++) if(!vis[i])
 8
                      for(int j=i*i; j<=n; j+=i) vis[j]=1;
 9
             int c=0;
10
             for(int i=2;i \le n;i++) if(!vis[i])
11
                      prime[c++]=i;
12
13
```

2.2.2 Sieve of Euler

The Sieve can be used to solve problem of multiplicative function. Test: CF 10C

```
int vis [maxn],prime[maxp];
    void sieve(int n)
2
3
            int tot=0;
 4
            for(int i=2;i \le n;i++)
5
                     if (! vis [i]) prime[tot++]=i;
                     for(int j=0; j<tot; j++)
 8
 9
                              if(i*prime[j]>n) break;
10
                              vis[i*prime[j]]=1;
11
                              if (i%prime[j]==0) break;
12
                     }
13
    }
14
```

2.2.3 Euclid Algorithm

```
typedef long long LL;
    LL gcd(LL a,LL b)
    {
 3
            return b==0?a:gcd(b,a%b);
 4
    }
5
    void exgcd(LL a,LL b,LL& d,LL& x,LL& y)
6
 7
            if (!b){d=a;x=1;y=0;}
 8
            else
9
10
                    gcd(b,a\%b,d,y,x);
11
                    y=x*(a/b);
12
            }
13
14
```

2.2.4 Euler Function

```
int euler_phi (int n)

int m=(int)sqrt(n+0.5);

int ans=n;

for(int i=2;i<=m;i++) if(n%i==0)

{
    ans=ans/i*(i-1);
}</pre>
```

```
while(n\%i==0) n/=i;
8
9
            if (n>1) ans=ans/n*(n-1);
10
11
12
    int phi[maxn];
    void phi_table(int n)
13
14
            for(int i=2; i<=n; i++) phi[i]=0;
15
            phi[1]=1;
16
            for(int i=2;i<=n;i++) if(!phi[i])
17
18
                    for(int j=i; j \le n; j+=i)
19
20
                            if (!phi[j]) phi[j]=j;
21
                            phi[j]=phi[j]/i*(i-1);
22
23
24
25
    }
    2.2.5
             Calculate the Inversion
   LL inv(LL a, LL n)
1
 2
 3
            LL d,x,y;
 4
            exgcd(a,n,d,x,y);
            return d==1?(x+n)\%n:-1;
 5
    }
 6
    2.2.6
             Chinese Remainder Theorem
    //n equations:x=a[i] \pmod{m[i]} (0<=i<n)
    LL china(int n, int *a, int *m)
 2
 3
    {
            LL M=1,d,y,x=0;
 4
            for(int i=0; i< n; i++) M*=m[i];
 5
            for(int i=0; i< n; i++)
 6
 7
            {
                    LL w=M/m[i];
 8
 9
                    exgcd(m[i], w,d,d,y);
                    x=(x+y*w*a[i])%M;
10
11
            return (x+M)%M;
12
13
```

2.2.7 Bezout Theorem

For any two integers a,b, assume that d is their gcd. The equation ax + by = m has integer solutions (x,y) if and only if m is the multiple of d. When the

equation has solutions, the set of solutions are $\{(\frac{mx_0+kb}{d}, \frac{my_0-ka}{d})|k \in \mathbb{Z}\}$, where (x_0, y_0) is one solution of this equation.

2.3 Combinatorics

2.3.1 Calculate Combinations by Recurrence

```
1  void init ()
2  {
3     memset(C,0,sizeof(C));
4     C[0][0]=1;
5     for(int i=1;i<n;i++)
6     {
7         C[i][0]=C[i][i]=0;
8         for(int j=1;j<i;j++)
9         C[i][j]=C[i-1][j]+C[i-1][j-1];
10     }
11  }</pre>
```

2.4 Matrix and System of Linear Equations

2.4.1 Matrix Fast Power

```
typedef matrix int [maxn][maxn];
    void mat_mul(matrix A,matrix B,matrix res)
 3
        matrix C;
 4
        memset(C,0,sizeof(C));
 5
        for(int i=0; i< n; i++)
             for(int j=0; j< n; j++)
                 for(int k=0;k< n;k++)
                     C[i][j]+=A[i][k]*B[k][j];
 9
        memcpy(res,C,sizeof(C));
10
    }
11
    void mat_pow(matrix A,int m,matrix res)
12
13
        matrix a,r;
14
        memcpy(a,A,sizeof(a));
15
        memset(r,0,sizeof(r));
16
        for(int i=0; i< n; i++) r[i][i]=1;
17
        while(m)
18
19
             if (m&1) mat_mul(r,a,r);
20
             m >> = 1;
21
            mat_mul(a,a,a);
22
```

```
23 }
24 memcpy(res,r,sizeof(r));
25 }
```

2.5 Numerical Method

2.5.1 Adaptive Simpson's Rule

```
//F is a global function
    double simpson(double a,double b)
3
            double c=a+(b-a)/2;
 4
            return (F(a)+4*F(c)+F(b))*(b-a)/6;
5
    }
6
    //recursive part
    double asr(double a,double b,double eps,double A)
 8
9
            double c=a+(b-a)/2;
10
            double L=simpson(a,c),R=simpson(c,b);
11
            if (fabs(L+R-A)<=15*eps) return L+R+(L+R-A)/15.0;
12
            return asr(a,c,eps/2,L)+asr(c,b,eps/2,R);
13
    }
14
    //main part
15
    double asr(double a,double b,double eps)
16
17
            return asr(a,b,eps,simpson(a,b));
18
19
```

Chapter 3

Graph Theory

3.1 Fundamentals

3.1.1 Topo Sort

Test:UVaLive 4255

```
const int maxn=20;
    int g[maxn][maxn]; //adjacency matrix
    int in[maxn];
                         //indegree
    int topo[maxn];
                         //result
 5
    int toposort(int n)
 6
        int top=-1;
        for(int i=0;i< n;i++)
 8
             if (in[i]==0)
 9
10
                 in\,[\,i\,]{=}top;top{=}i;
11
12
13
        for(int i=0; i< n; i++)
             if (top==-1) return 0; //exist cycle
14
15
16
            {
17
                 int j=top;
                 top=in[j];
18
                 topo[i]=j;
19
                 for(int k=0;k< n;k++)
20
                     if (g[j][k] \&\& (--in[k]) == 0)
21
22
23
                         in[k] = top;
24
                         top=k;
                     }
25
```

```
26 }
27 return 1; //succeed
28 }
```

3.1.2 Eulerian trail and circuit

Test:UVa 10054,10441

- undirected graph
 - 1. whether there exists an Eulerian circuit
 - (a) the graph is connected
 - (b) every vertex has even degree
 - 2. whether there exists an Eulerian trail
 - (a) the graph is connected
 - (b) there are only two vertex with odd degree which are the start point and the end point
- directed graph
 - 1. whether there exists an Eulerian circuit
 - (a) the graph is connected
 - (b) every vertex has equal indegree and outdegree
 - 2. whether there exists an Eulerian trail
 - (a) the graph is connected
 - (b) there is only one vertex with outdegree greater than indegree by one as the start point, and there is only one vertex with indegree greater than outdegree by one as end point

```
int g[maxn][maxn];
                             //adjacency matrix
                             //degree
    int cnt[maxn];
 2
    vector<pair<int,int> > ans;
                                     //edge list, stored in reverse order
                                     //n is the number of vertices, m is the number of edges
    int n,m;
 5
    void euler(int u)
 6
        for(int i=0; i< n; i++) if(g[u][i])
 7
 8
            g[u][i]--;g[i][u]--;
 9
            euler (i);
10
            ans.push_back(make_pair(u,i));
11
        }
12
13
    int solve(int start)
14
15
    {
```

```
int flag =1;
16
         for(int i=0;i< n;i++)
17
              if (cnt[i]%2)
18
19
20
                  flag =0; break;
21
         if (flag)
22
23
              ans. clear ();
24
              euler (start);
25
              if (ans. size ()!=m \parallel ans [0]. second!=ans[ans. size ()-1]. first ) flag =0;
26
27
                                //whether exists an Eulerian circuit
28
         return flag;
    }
29
```

3.2 Shortest Paths

3.3 Minimal Spanning Tree

3.3.1 Fundamentals

Theorem 1

For an *n*-vertex graph $G(\text{with } n \geq 1)$, the following are equivalent (and character the trees with *n* vertices).

- 1. G is connected and has no cycles
- 2. G is connected and has n-1 edges
- 3. G has n-1 edges and no cycles
- 4. G has no loops and has, for each $u, v \in V(G)$, exactly one u,v-path

Corollary 1

- 1. Every edge of a tree is a cut-edge
- 2. Adding one edge to a tree forms exactly one cycle
- 3. Every connected graph contains a spanning tree

3.3.2 Kruskal algorithm

```
int cmp(const int i, const int j){return w[i]<w[j];} //indirect sorting function
int find(int x){return p[x]==x?x:p[x]=find(p[x]);} //disjoint set
int Kruskal()
{
    int ans=0;</pre>
```

```
for (int i=0; i< n; i++) p[i]=i; i=1; i
  6
                            for(int i=0;i<m;i++) r[i]=i; // initial edge index
                            sort(r,r+m,cmp); //sort edges
                            for(int i=0;i< m;i++)
  9
10
                                               int e=r[i]; int x=find(u[e]); int y=find(v[e]);
11
                                                if (x!=y) ans +=w[e], p[x]=y;
12
13
                            return ans;
14
15
                              Second-best Minimum Spanning Tree
         Test:UVa 10600
         #include <iostream>
         #include <cstdio>
         #include <cstring>
         #include <algorithm>
         using namespace std;
         \textbf{const int} \  \, \text{maxn}{=}110;
         const int maxm=maxn*maxn/2;
         struct edges
  9
                            int u,v,c;
10
                            bool operator<(const edges& tmp) const
11
12
                                              return c<tmp.c;
13
14
         } edge[maxm];
15
16
         int n,m,e;
         int head[maxn],pnt[maxn*2],nxt[maxn*2],cost[maxn*2];
17
         int maxcost[maxn][maxn];
18
         int f[maxn];
19
         int vis [maxn], vis2 [maxm];
20
         void addedge(int u,int v,int c)
21
22
                            pnt[e]=v; cost[e]=c; nxt[e]=head[u]; head[u]=e++;
23
24
         int find (int u){return f[u]==u?u:f[u]=find(f[u]);}
25
         void dfs(int u,int fa,int c)
26
27
28
                            maxcost[u][fa]=maxcost[fa][u]=c;
                            for(int i=1; i <= n; i++) if(vis[i]) maxcost[u][i]=maxcost[i][u]=max(maxcost[i][fa], c);
29
                            vis[u]=1;
30
                            for(int i=head[u];i!=-1;i=nxt[i])
31
32
```

```
int v=pnt[i];
33
                     if (v!=fa) dfs(v,u,cost[i]);
34
             }
35
36
37
    int main()
    {
38
             //freopen("in. txt", "r", stdin );
39
             int T;
40
             scanf("%d",&T);
41
             while(T--)
42
43
                     scanf("%d%d",&n,&m);
44
                      for(int i=0; i< m; i++) scanf("%d%d%d", \&edge[i].u, \&edge[i].v, \&edge[i].c); 
45
                     for(int i=1; i <= n; i++) f[i]=i;
46
                     memset(head,-1,sizeof(head));
47
                     e=0;
48
                     sort (edge,edge+m);
49
                     int ans1=0;
50
                     memset(vis2,0, sizeof(vis2));
51
                     for(int i=0;i< m;i++)
52
53
                              int u=find(edge[i].u),v=find(edge[i].v);
54
                              if(u!=v)
55
                              {
56
                                      f[u]=v;
57
58
                                      addedge(u,v,edge[i].c);
                                      addedge(v,u,edge[i].c);
59
                                      ans1+=edge[i].c;
60
                                       vis2[i]=1;
61
62
63
                     memset(vis,0, sizeof(vis));
64
                     dfs(1,0,0);
65
                     int ans2=0 x3fffffff ;
66
                     for(int i=0;i< m;i++) if(!vis2[i])
67
                     {
68
69
                              int u=edge[i].u,v=edge[i].v,c=edge[i].c;
                              ans2=min(ans2,ans1+c-maxcost[u][v]);
70
71
                      printf ("%d∟%d\n",ans1,ans2);
72
73
             return 0;
74
75
```

3.3.4 Minimal Bottleneck Spanning Tree

Actually, the MBST is the MST. And we can get it by Kruskal Algorithm. If we want to query minimal bottleneck path between any two vertices, we can use dfs to calculate maxcost array just as we do in Second-best Minimal Spanning Tree. Then we can answer each query in O(1), and the total complexity is $O(n^2)$.

3.3.5 Minimal Directed Spanning Tree

```
const int INF = 10000000000:
 1
    const int maxn = 100 + 10;
 3
    struct MDST {
 4
      int n;
 5
      int w[maxn][maxn];
      int vis [maxn];
      int ans:
      int removed[maxn];
 9
      int cid[maxn];
10
      int pre[maxn];
11
      int iw[maxn];
12
      int max_cid;
13
14
      void init (int n) {
15
         this -> n = n;
16
         for(int i = 0; i < n; i++)
17
           for(int j = 0; j < n; j++) w[i][j] = INF;
18
19
20
      void AddEdge(int u, int v, int cost) {
21
22
         w[u][v] = min(w[u][v], cost);
23
24
      int dfs(int s) {
25
         vis[s] = 1;
26
27
         int ans = 1;
         for(int i = 0; i < n; i++)
28
29
           if (! \text{ vis } [i] \&\& \text{ w}[s][i] < \text{INF}) \text{ ans } += \text{dfs}(i);
         return ans;
30
31
32
      bool cycle(int u) {
33
         max\_cid++;
34
35
         int v = u;
         while(cid[v] != max_cid) { cid[v] = max_cid; v = pre[v]; }
36
         return v == u;
37
```

```
}
38
39
      void update(int u) {
40
        iw[u] = INF;
41
42
        for(int i = 0; i < n; i++)
           if (!removed[i] \&\& w[i][u] < iw[u]) {
43
            iw[u] = w[i][u];
44
            pre[u] = i;
45
46
      }
47
48
      bool solve(int s) {
49
        memset(vis, 0, sizeof(vis));
50
        if (dfs(s) != n) return false;
51
52
        memset(removed, 0, sizeof(removed));
53
        memset(cid, 0, sizeof(cid));
54
        for(int u = 0; u < n; u++) update(u);
55
        pre[s] = s; iw[s] = 0;
56
        ans = \max_{\text{cid}} = 0;
57
        for (;;) {
58
          bool have_cycle = false;
59
           for(int u = 0; u < n; u++) if(u != s \&\& !removed[u] \&\& cycle(u)){
60
             have_cycle = true;
61
             int v = u;
62
            do {
63
               if (v != u) removed [v] = 1;
64
              ans += iw[v];
65
               for(int i = 0; i < n; i++) if(cid[i] != cid[u] && !removed[i]) {
66
                 if(w[i][v] < INF) w[i][u] = min(w[i][u], w[i][v]-iw[v]);
67
68
                 w[u][i] = min(w[u][i], w[v][i]);
                 if(pre[i] == v) pre[i] = u;
69
70
              }
71
              v = pre[v];
             } while(v != u);
72
            update(u);
73
74
            break;
75
           if (! have_cycle ) break;
76
77
        for(int i = 0; i < n; i++)
78
           if (!removed[i]) ans += iw[i];
79
80
        return true;
81
    };
82
```

3.4 Bipartite Graph Matching

3.4.1 Augmenting Path Algorithm

```
Test:POJ 3020
```

```
int n,m; //size of bipartite sets
     int G[maxn][maxn]; //adjacency matrix
     int linker [maxn], vis [maxn];
     int dfs(int u)
 5
               for(int v=0;v< m;v++)
 6
                         \textbf{if}\left(\mathsf{G}[\mathsf{u}][\mathsf{v}] \ \&\& \ !\mathsf{vis}[\mathsf{v}]\right)
 8
                                   vis[v]=1;
 9
                                   if ( linker [v] = -1 \parallel dfs( linker [v]))
10
11
                                             linker [v]=u;
12
13
                                             return 1;
14
15
               return 0;
16
17
     int hungary()
18
19
               int ans=0;
20
21
               memset(linker,-1,sizeof( linker ));
               for(int u=0;u< n;u++)
22
23
                         memset(vis,0, sizeof( vis ));
24
                         if (dfs(u)) ans++;
25
26
               return ans;
27
28
     }
```

3.4.2 Kuhn-Munkers Algorithm

Test:UVaLive 4043

```
int W[maxn][maxn],n;
int Lx[maxn],Ly[maxn]; //label on X and Y
int linker [maxn],slack[maxn];
int S[maxn],T[maxn]; //mark whether visited
int dfs(int u)
{
    S[u]=1;
    for(int v=0;v<n;v++)</pre>
```

```
{
 9
                      if (T[v]) continue;
10
                      int tmp=Lx[u]+Ly[v]-W[u][v];
11
                      if(tmp==0)
12
13
                      {
                              T[v]=1;
14
                              if ( linker [v] = -1 \parallel dfs( linker [v]))
15
16
                                       linker [v]=u;
17
                                      return 1;
18
19
20
                     else if (slack [v]>tmp)
21
                              slack [v]=tmp;
22
23
             return 0;
24
25
    void update()
26
27
             int a=inf;
28
             for(int i=0;i< n;i++)
29
                      if (!T[i] \&\& a-slack[i]>0)
30
                              a=slack[i];
31
             for(int i=0;i< n;i++)
32
33
                      if (S[i]) Lx[i]-=a;
34
                      if (T[i]) Ly[i]+=a;
35
                      else slack [i]—a;
36
             }
37
    }
38
    void KM()
39
    {
40
             for(int i=0; i< n; i++)
41
42
                      linker [i]=-1;
43
                     Lx[i]=Ly[i]=0;
44
45
                      for (int j=0; j< n; j++)
                              Lx[i]=max(Lx[i],W[i][j]);
46
47
             for(int i=0;i< n;i++)
48
49
                     for(int j=0; j< n; j++) slack[j]=inf;
50
                     while(1)
51
52
                              memset(S,0,sizeof(S));
53
                              memset(T,0,sizeof(T));
54
```

3.5 Network Flows

3.6 Depth First Search

3.6.1 Calculate DFS Sequence

```
Test:CF 383C
```

```
//tree is stored in adjacency lists
    //dfs\_clock is initialed as 0 vis array is initialed as 0
    //dfn is the dfs sequence, ldfs is the time the vertex enters and rdfs the time leaves
    int Idfs [maxn],rdfs [maxn],dfn[maxn],vis [maxn];
    int dfs_clock;
    void dfs(int u)
 7
        vis[u]=1;
        Idfs[u]=++dfs\_clock;
 9
        dfn[dfs\_clock]=u;
10
        for(int i=head[u];i!=-1;i=nxt[i])
11
12
            int v=pnt[i];
13
            if (! vis [u]) dfs(v);
14
15
        rdfs[u]=++dfs\_clock;
16
        dfn[dfs\_clock]=u;
17
18
    }
```

3.6.2 Multiplication LCA Algorithm(online)

Test:POJ 1330,3728

```
for(int i=1;i < pow;i++) p[u][i]=p[p[u][i-1]][i-1];
9
             for(int i=head[u]; i!=-1; i=nxt[i])
10
11
                     int v=pnt[i];
12
13
                     dfs(v,u);
14
    }
15
                                               //get the lca of a and b in O(logn)
    int lca(int a, int b)
16
17
             if (d[a]>d[b]) a^=b,b^=a,a^=b;
18
             if (d[a] < d[b])
19
20
             {
                     int del=d[b]-d[a];
21
                     for(int i=0; i < pow; i++) if(del&(1<<i)) b=p[b][i];
22
23
             if(a!=b)
24
25
                     for(int i=pow-1; i>=0; i--)
26
                              if (p[a][i]!=p[b][i])
27
                                      a=p[a][i], b=p[b][i];
28
                     a=p[a][0], b=p[b][0];
29
30
             return a;
31
    }
32
```

3.6.3 Targan's Off-line LCA Algorithm

Test:POJ 1330,1470

```
vector < int > g[maxn],q[maxn];//adjacency lists and query lists
    int f[maxn];
    int a [maxn]; //record the current lca of the vertex
    int vis [maxn];
    //disjoint set
    int find (int x) {return x==f[x]?x:f[x]=find(f[x]);}
    void uni(int x, int y)
 7
8
             int a = find(x);
 9
             int b=find(y);
10
             if (x!=y) f[a]=b;
11
12
    //use:lca(root);
13
    //usually insert (u,v) and (v,u) into q
15
    void lca(int u)
    {
16
            a[u]=u;
17
```

```
for(int i=0;i < g[u]. size(); i++)
18
19
                          lca(g[u][i]);
uni(u,g[u][i]);
20
21
                          a[find(u)]=u;
22
                } vis [u]=1;
23
24
                for(int i=0; i < q[u]. size(); i++)
25
26
                           \textbf{if}\,(\,\mathsf{vis}\,[\mathsf{q}[\mathsf{u}\,][\,\mathsf{i}\,]])
27
28
                                int LCA=a[find(q[u][i])];
                                                                     //answer the query
29
                                operation ();
30
                                                                     //operate corresponding operations
31
                }
32
    }
33
```

Chapter 4

String

4.1 KMP

```
//fail function
    int f[maxn];
 2
    void getfail (char* P)
 4
 5
        int m=strlen(P);
        f[0]=0; f[1]=0;
 6
        for(int i=1;i< m;i++)
 8
            int j=f[i];
 9
10
            while(j && P[i]!=P[j]) j=f[j];
            f[i+1]=P[i]==P[j]?j+1:0;
11
12
13
14
    //find function
15
    void find(char* T,char* P)
16
17
        int n=strlen(T),m=strlen(P);
18
19
        int i=0;
        for(int i=0;i< n;i++)
20
21
            while(j && P[j]!=T[i]) j=f[j];
22
            if (P[j]==T[i]) j++;
23
            if(j==m)
24
25
                operate();
26
                            //note when find a match, remember to get backward along fail edge once
27
                j=f[j];
```

```
29 }
30 }
```

4.2 AC Automation

```
//tire tree
            const int maxn=100010;//maximal node number
             const int sigma=26;
             //the size of character size, the value of characters are 0 ignsize in size in the size of characters are <math>0 ignsize in size in the size in the size of characters are <math>0 ignsize in size in the size in t
             int ch[maxn][sigma]; //ch[i][j] means the node going from i along j to
             int val[maxn];
                                                                                  //additive information
   6
                                                                                  //the total number of node
             int sz;
             void init ()
                                                                                  // initial
   9
                         sz=1;
10
                         memset(ch[0],0,sizeof(ch[0]));
11
12
             int idx(char c)
                                                                                 //index the character
13
14
                                                                                  //vary depending on specific condition
15
                         return c-'a':
16
             //insert string s, whose additive information is v
17
             //note that v cannot be \theta
18
             void insert (char *s, int v)
19
20
21
                         int u=0, n=strlen(s);
                         for(int i=0; i< n; i++)
22
23
                                      int c=idx(s[i]);
24
                                      if (!ch[u][c])
25
26
                                                  memset(ch[sz],0, sizeof(ch[sz]));
27
28
                                                  val[sz]=0;
                                                  ch[u][c]=sz++;
29
30
                                      u=ch[u][c];
31
32
                         val[u]=v;
33
                         //val[u] = (1 << v);
34
                          //operating like this when it comes to state compressing
35
             }
36
37
             //get fail function
38
            int f[maxn];
                                                                                        //fail function
39
             int last [maxn];
                                                                                        //suffix link
             void getfail ()
```

```
42
         queue<int> q;
43
         f[0]=0;
44
         for(int c=0;c<sigma;c++)</pre>
45
46
             int u=ch[0][c];
47
             if (u) \{f[u]=0;q.push(u); last [u]=0;\}
48
49
         while(!q.empty())
50
51
             int r=q.front(); q.pop();
52
             for(int c=0;c<sigma;c++)</pre>
53
54
                  int u=ch[r][c];
55
                  if (!u)
56
57
                      //ch[r][c]=ch[f[r]][c];
58
                      //add all nonexist edges, applying on dp
59
                      continue;
60
61
                  q.push(u);
62
                  int v=f[r];
63
                  \textbf{while}(v \&\& !ch[v][c]) \ v = f[v]; // \textit{when all nonexist edges are added, code here can be deleted}
64
                  f[u]=ch[v][c];
65
                  last [u]=val[f[u]]?f[u]: last [f[u]];
66
67
             }
         }
68
    }
69
70
    int find(char *T)
71
72
         int n=strlen(T);
73
         int i=0;
74
         for(int i=0; i< n; i++)
75
76
             int c=idx(T[i]);
77
78
             while(j \&\& !ch[j][c]) j=f[j];
                                                //when all nonexist edges are added, code here can be deleted
             j=ch[j][c];
79
             if (val[j]) process(j);
80
             else if ( last [j]) process( last [j]);
81
         }
82
    }
83
84
    void process(int j)
85
    {
86
         if(j)
87
```

4.3 Suffix Array

```
int s[maxn];
                                     //string to be constructed
 1
    int sa[maxn];
                                     //suffix array
    int t[maxn],t2[maxn],c[maxn];
    //every character's value is 0-m-1, n is usually the length of string plus 1
     void build_sa (int m,int n)
                                            //construct suffix array
 5
 6
 7
        int i,*x=t,*y=t2;
        for(i=0;i< m;i++) c[i]=0;
 8
        for(i=0;i< n;i++) c[x[i]=s[i]]++;
 9
        for(i=1;i< m;i++) c[i]+=c[i-1];
10
        for(i=n-1;i>=0;i--) sa[--c[x[i]]]=i;
11
        for(int k=1; k <= n; k << =1)
12
13
             int p=0;
14
             for(i=n-k;i< n;i++) y[p++]=i;
15
             for (i=0;i< n;i++) if (sa[i]>=k) y [p++]=sa[i]-k;
16
             for (i=0;i< m;i++) c[i]=0;
17
18
             for (i=0;i< n;i++) c[x[y[i]]]++;
             for (i=0;i< m;i++) c[i]+=c[i-1];
19
             for (i=n-1;i>=0;i--) sa[--c[x[y[i]]]]=y[i];
20
             swap(x,y);
21
            p=1;x[sa[0]]=0;
22
23
             for (i=1; i< n; i++)
                 x[sa[i]]=y[sa[i-1]]==y[sa[i]]\&\&y[sa[i-1]+k]==y[sa[i]+k]?p-1:p++;
24
25
             if (p>=n) break;
26
             m=p;
        }
27
    }
28
29
    //calculate rank and height array
30
    int rank[maxn];
                             //suffix i's index in sa
31
    int height[maxn];
                             //the\ LCP\ of\ sa[i-1]\ and\ sa[i]
32
    void getheight (int n) //n is the length of string
33
34
        int i, j, k=0;
35
        for (i=0; i \le n; i++) rank [sa[i]]=i;
36
        for(i=0;i< n;i++)
37
38
```

```
if (k) k--;
39
            int j=sa[rank[i]-1];
40
            while(s[i+k]==s[j+k]) k++;
41
            height [rank[i]]=k;
42
43
        }
    }
44
45
    //RMQ
46
    int d[maxn][50];
47
    void RMQ_init(int n)
48
49
        for(int i=0; i< n; i++) d[i][0]=height[i];
50
        for(int j=1;(1<< j)<=n;j++)
51
            for(int i=0;i+(1<< j)-1< n;i++)
52
                d[i][j]=min(d[i][j-1],d[i+(1<<(j-1))][j-1]);
53
54
    int RMQ(int L,int R)
55
56
        if (L>R) swap(L,R);
57
        if (L==R) return L-sa[L];
58
        L++;
59
        int k=0;
60
        while((1 << (k+1)) <= R-L+1) k++;
61
        return min(d[L][k], d[R-(1 << k)+1][k]);
62
    }
63
64
    //inital
65
    build_sa (m,n+1);
66
    getheight(n);
    RMQ_init(n+1);
68
```

4.4 Suffix Automation

```
//string to be constructed
    char s[maxn];
    int sz;
                                 //the number of node
    struct state
        state *pre;
 5
        state *go[sigma];
 6
        int val;
                                 //the length of the longest path
        void clear ()
                                 // initialize node
 8
9
10
            pre=0;val=0;
            memset(go,0,sizeof(go));
11
12
    };
13
```

```
state *root,* last;
                                 //root node and last node added
14
    state st [maxn];
15
    void init ()
                                 // initialize SAM
16
17
18
        sz=0:
        root=last=\&st[sz++];
19
        root->clear();
20
21
                                 //extend node
    void extend(int w)
22
23
        state *p=last;
                                 //last node added
24
        state *np=&st[sz++];
                                 //create a new node
25
        np->clear();
26
        np->val=p->val+1;
27
        while(p && p->go[w]==0) //add edge to np when p->go[w]=0
28
29
             p->go[w]=np;
30
31
             p=p->pre;
32
        if (p==0) np->pre=root; //if w occurs first, add edge to the root
33
34
35
             state *q=p->go[w];
36
             if (q->val==p->val+1) //q is the proper state
37
38
                 np->pre=q;
39
            else
40
             {
                 state *nq=&st[sz++];//create a new node and copy the information
41
                 nq->clear();
42
                 \mathsf{nq}{-}{>}\mathsf{val}{=}\mathsf{p}{-}{>}\mathsf{val}{+}1;
43
                 memcpy(nq->go,q->go,sizeof(q->go));
                 nq->pre=q->pre;
45
46
                 q->pre=nq;
                 np->pre=nq;
47
                 while(p && p->go[w]==q)
48
49
                     p->go[w]=nq;
50
                     p=p->pre;
51
52
53
54
         last = np;
55
56
57
    //sort the nodes
58
    int t[maxn];
59
```

```
state *r[maxn];
//l is the length of string, sz is the number of node
//note here contains root node, root node is indexed as 1 in r
for(int i=0;i<|;i++) t[i]=0;
for(int i=0;i<|sz;i++) t[st[i]. val]++;
for(int i=1;i<|;i++) t[i]+=t[i-1];
for(int i=0;i<|sz;i++) r[t[st[i]. val]--]=&st[i];</pre>
```

4.5 String Hash

```
typedef unsigned long long ULL;
    const int maxn=10000;
    const int x=123;
                                 //hash's parameter
    ULL H[maxn];
                                 //the hash value of string
                                 //x^n
    ULL xp[maxn];
    void hash_init (char *s) // initialize H(n),x^n
 7
         int n=strlen(s);
 8
         H[n] = 0;
 9
         \label{eq:for_int} \begin{tabular}{ll} \begin{tabular}{ll} for (int & i=n-1; i>=0; i--) & H[i]=H[i+1]*x+(s[i]-'a'); & //calculate & H(n) \\ \end{tabular}
10
11
12
         for(int i=1; i<=n; i++) xp[i]=xp[i-1]*x;
                                                                       //calculate x^n
13
     //hash function, return the hash value of string starts from i with length of L
14
    ULL hash(int i, int L)
15
16
         return H[i]-H[i+L]*xp[L];
17
    }
18
```

4.6 Longest Palindromic Substring

Test:LeetCode, Longest Palindromic Substring

```
int p[2010];
    void get( string str )
 2
 3
              int mx=0,id;
 4
              for(int i=1; i < str. size(); i++)
 5
 6
                        if (mx>i) p[i]=min(p[2*id-i],mx-i);
 7
                       else p[i]=1;
                       for (; str[i+p[i]] = str[i-p[i]]; p[i]++)
 9
10
                       \textbf{if} \, (p[i] + i > mx)
11
12
                                mx=p[i]+i;
13
```

```
14
                                id=i;
                       }
15
16
17
     string longestPalindrome(string s)
18
    {
19
              string t="$";
20
              for(int i=0; i < s. size(); i++)
21
                       t+='#',t+=s[i];
22
              t+='#';
23
24
              get(t);
              int ans=0,pos;
25
              string str="";
26
              for(int i=1; i < t. size(); i++)
27
                       if(p[i]>ans)
28
29
                                ans=p[i];
30
                                pos=i;
31
32
              \quad \textbf{for(int} \ \ \overset{.}{i} = & pos-ans+1; \\ i < & pos+ans; \\ i++)
33
                       if(t[i]!='#') str+=t[i];
34
              return str;
35
36
```

Chapter 5

Geometry

5.1 Basic

comparsion between real numbers

```
const double eps=1e-8;
 1
    int dcmp(double x)
 3
        if (fabs(x)<eps) return 0;</pre>
 4
        else return \times < 0? -1:1;
 5
   }
 6
       definition of point and vector
    struct point
 1
    {
 2
        double x,y;
 3
        point(double x=0,double y=0):x(x),y(y){}
 5
    point operator+(point a,point b){return point(a.x+b.x,a.y+b.y);}
    point operator-(point a,point b){return point(a.x-b.x,a.y-b.y);}
    point operator*(point a, double p){return point(a.x*p,a.y*p);}
    point operator/(point a,double p){return point(a.x/p,a.y/p);}
    bool operator <(const point& a,const point& b)
10
11
        return dcmp(a.x-b.x)<0 || (dcmp(a.x-b.x)==0&&dcmp(a.y-b.y)<0);
12
13
    bool operator==(point a,point b)
14
15
        return dcmp(a.x-b.x)==0 \&\& dcmp(a.y-b.y)==0;
16
17
```

```
Dot(\vec{u}, \vec{v}) = |\vec{u}| * |\vec{v}| * cos < \vec{u}, \vec{v} >
   double dot(point a, point b){return a.x*b.x+a.y*b.y;}
1
   double length(point a){return sqrt(dot(a,a));}
   double angle(point a, point b){return acos(dot(a,b)/length(a)/length(b));}
       Cross(\vec{u}, \vec{v}) is the twice of the directed area of the triangle formed by \vec{u} \& \vec{v}
   double cross(point a, point b){return a.x*b.y-a.y*b.x;}
   double area2(point a, point b, point c){return cross(b-a,c-a);}
       rotate the vector
   point rotate (point a, double rad)
1
2
       return point(a.x*cos(rad)-a.y*sin(rad), a.x*sin(rad)+a.y*cos(rad));
3
   }
4
       calculate normal of the vector
   point normal(point a)
1
2
       double L=length(a);
3
       return point(-a.y/L,a.x/L);
4
5 }
       get intersection of two lines which are not parallel to each other
   //p,q are points on two lines
   //v, w are direction vector of two lines
   point getlineinter (point p, point v, point q, point w)
3
4
5
        point u=p-q;
       double t=cross(w,u)/cross(v,w);
       return p+v*t;
8
  }
       distance from point to line
   double distoline (point p, point a, point b)
2
        point v1=b-a, v2=p-a;
       return fabs(cross(v1,v2))/length(v1);
4
   }
       distance from point to segment
   double distoseg(point p,point a,point b)
1
2
        if (a==b) return length(p-a);
3
        point v1=b-a, v2=p-a, v3=p-b;
4
        if (dcmp(dot(v1,v2))<0) return length(v2);
        else if (dcmp(dot(v1,v3))>0) return length(v3);
```

```
7
       else return fabs(cross(v1,v2))/length(v1);
  }
8
      projection of point on line
   point getpro(point p,point a,point b)
1
2
       point v=b-a;
3
       return a+v*(dot(v,p-a)/dot(v,v));
4
5
      judge whether two lines are properly intersected
   bool segproperint (point a1, point a2, point b1, point b2)
1
2
       double c1=cross(a2-a1,b1-a1),c2=cross(a2-a1,b2-a1),
3
              c3=cross(b2-b1,a1-b1),c4=cross(b2-b1,a2-b1);
       return dcmp(c1)*dcmp(c2)<0 \&\& dcmp(c3)*dcmp(c4)<0;
5
  }
6
      judge whether point on segment, endpoints exclusive
   bool onseg(point p, point a1, point a2)
1
2
       return dcmp(cross(a1-p,a2-p))==0 \&\& dcmp(dot(a1-p,a2-p))<0;
3
  }
4
      directed area of polygon
   double polyarea(point* p, int n)
1
2
       double area=0;
3
       for(int i=1; i< n-1; i++)
4
           area += cross(p[i]-p[0],p[i+1]-p[0]);
5
       return area /2;
7
  }
```

5.2 Problems About Circle and Sphere

5.2.1 Definition of circle

```
1 struct circle
2 {
3          point c;
4          double r;
5          circle (point c=point(0,0), double r=0):c(c),r(r){}
6          point getpoint(double a)
7          {
8                return point(c.x+cos(a)*r,c.y+sin(a)*r);
}
```

```
}
10
           };
                                          Calculate the intersection of line and circle
             5.2.2
             \\return the number of intersections
             \\the intersections are stored in sol
                              getlinecircleinter (point p, point v, circle C, double & t1,
   3
                                                                                                                                             double& t2,vector<point>& sol)
             {
   5
                                       double a=v.x,b=p.x-C.c.x,c=v.y,d=p.y-C.c.y;
   6
                                       double e=a*a+c*c, f=2*(a*b+c*d), g=b*b+d*d-C.r*C.r;
                                       double delta=f*f-4*e*g:
                                       if (dcmp(delta)<0) return 0;</pre>
10
                                       if (dcmp(delta) = = 0)
11
                                                                t1=t2=-f/(2*e);
12
                                                                 sol . push_back(p+v*t1);
13
                                                                return 1;
14
15
16
                                       t1=(-f-sqrt(delta))/(2*e); sol.push_back(p+v*t1);
                                       t2=(-f+sqrt(delta))/(2*e); sol.push_back(p+v*t2);
17
                                       return 2;
18
19
             5.2.3
                                          Calculate the intersection of circles
             \label{eq:cont_v} \textbf{double} \  \, \text{angle(point v)} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \hspace{0.5cm} \setminus \  \, \text{calculate the polar angle} \\ \  \, \text{angle(point v)} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \hspace{0.5cm} \setminus \  \, \text{calculate the polar angle} \\ \  \, \text{angle(point v)} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \hspace{0.5cm} \setminus \  \, \text{calculate the polar angle} \\ \  \, \text{angle(point v)} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \  \, \text{atan2(v.y,v.x);} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \  \, \text{atan2(v.y,v.x);} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \  \, \text{atan2(v.y,v.x);} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \  \, \text{atan2(v.y,v.x);} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \  \, \text{atan2(v.y,v.x);} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \  \, \text{atan2(v.y,v.x);} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \  \, \text{atan2(v.y,v.x);} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \  \, \text{atan2(v.y,v.x);} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \  \, \text{atan2(v.y,v.x);} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \  \, \text{atan2(v.y,v.x);} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \  \, \text{atan2(v.y,v.x);} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \  \, \text{atan2(v.y,v.x);} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \  \, \text{atan2(v.y,v.x);} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \  \, \text{atan2(v.y,v.x);} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \  \, \text{atan2(v.y,v.x);} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \  \, \text{atan2(v.y,v.x);} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \  \, \text{atan2(v.y,v.x);} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \  \, \text{atan2(v.y,v.x);} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \  \, \text{atan2(v.y,v.x);} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \  \, \text{atan2(v.y,v.x);} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \  \, \text{atan2(v.y,v.x);} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \  \, \text{atan2(v.y,v.x);} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \  \, \text{atan2(v.y,v.x);} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \  \, \text{atan2(v.y,v.x);} \\ \{ \textbf{return } \  \, \text{atan2(v.y,v.x);} \} \\ \  
   1
             \\return the number of intersections
   3
             \\the intersections are stored in sol
   4
                            getcircleinter (circle C1, circle C2, vector < point > & sol)
   5
   6
                                       double d=length(C1.c-C2.c);
   7
                                       if (dcmp(d)==0)
   8
   9
                                       {
                                                                 if (dcmp(C1.r-C2.r)==0) return -1;
10
                                                                return 0;
11
12
                                       if (dcmp(C1.r+C2.r-d)<0) return 0;
13
                                        if (dcmp(fabs(C1.r-C2.r)-d)>0) return 0;
14
15
                                       double a=angle(C2.c-C1.c);
16
                                       double da=acos((C1.r*C1.r+d*d-C2.r*C2.r)/(2*C1.r*d));
17
                                       point p1=C1.getpoint(a-da),p2=C1.getpoint(a+da);
18
```

19

sol .push_back(p1);

5.3 2D Algorithm

5.3.1 Judge whether point in polygon

```
//p is the point, poly is the polygon
    int isinpoly (point p, vector < point > & poly)
 3
        int wn=0;
 4
 5
        int n=poly.size();
        for(int i=0; i< n; i++)
             if (onseg(p, poly[i], poly[(i+1)\%n])) return -1; //on the boarder
 8
            int k=dcmp(cross(poly[(i+1)%n]-poly[i],p-poly[i]));
            int d1=dcmp(poly[i].y-p.y);
10
11
            int d2=dcmp(poly[(i+1)\%n].y-p.y);
            if (k>0 \&\& d1<=0 \&\& d2>0) wn++;
12
             if (k<0 \&\& d2<=0 \&\& d1>0) wn--;
13
14
                                 //inside
        if (wn!=0) return 1;
15
        return 0;
                                 //outside
16
17
    }
```

5.4 3D Algorithm

5.5 Stimulated Annealing Algorithm

5.5.1 Minimum Sphere Coverage

```
Test:UVa 10095

#include <iostream>
#include <cstdio>
#include <cstring>
#include <cmath>

makes in the state of the state
```

```
else return \times < 0? -1:1;
12
    }
13
    struct point
14
15
16
         double x,y,z;
    } p[maxn];
17
    double length(point a, point b)
18
19
20
         double dx=a.x-b.x;
         double dy=a.y-b.y;
21
        double dz=a.z-b.z;
22
         return sqrt (dx*dx+dy*dy+dz*dz);
23
24
    int max_dis(int n,point q)
25
26
    {
         int i;
27
         double res=0;
28
         for(int i=0; i< n; i++)
29
30
             double tmp=length(q,p[i]);
31
             if (dcmp(tmp-res)>0)
32
33
                 res=tmp;
34
                 j=i;
35
36
37
        return j;
38
39
    int main()
40
41
42
         int n;
         while(scanf("%d",&n)&&n)
43
44
             for(int i=0; i< n; i++) scanf("%lf%lf",&p[i].x,&p[i].y,&p[i].z);
45
             if (n==1) printf("0.0000_{\square}%.4f_{\square}%.4f_{\square}%.4f_{\square}m",p[0].x,p[0].y,p[0].z);
46
             else
47
48
                 double r=inf;
49
                 double step=100000;
                                           //maximal coordinate range
50
                 point q;
51
                 q.x=q.y=q.z=0;
                                           //select a node randomly
52
                 while(step>eps)
53
54
                      int j=max_dis(n,q);
55
                      double tmp=length(p[j],q);
56
                      if (dcmp(r-tmp)>0) r=tmp;
57
```

```
double dx=(p[j].x-q.x)/tmp;
58
                         \begin{array}{l} \textbf{double} \ dy = & (p[j].y - q.y)/tmp; \\ \textbf{double} \ dz = & (p[j].z - q.z)/tmp; \end{array} 
59
60
                        q.x+=dx*step;
61
62
                        q.y+=dy*step;
                        q.z+=dz*step;
63
                        step*=0.993;
                                                //decided\ by\ precision\,,\,better\ no\ less\ than\ 0.99
64
65
                   66
              }
67
68
         return 0;
69
70
    }
```

Chapter 6

Dynamic Programming

6.1 Longest Common Increasing Sequence

```
Time Complexity: O(n^2)
    Test: CF 10D
   #include <iostream>
 2 #include <cstdio>
    #include <cstring>
    using namespace std;
    const int maxn = 510;
    int n,m;
    int a[maxn],b[maxn];
    int dp[maxn][maxn];
    int path[maxn][maxn];
    void print(int i,int j,int pre)
10
11
             if (path[i][j]==0)
12
13
                     if (j!=pre) printf ("%d<sub>\(\pi\)</sub>",b[j]);
14
                     return;
15
16
             print (i-1,path[i][j], j);
17
             if (j!=pre) printf ("d_{\sqcup}",b[j]);
18
19
    int main()
20
    {
21
             scanf("%d",&n);
22
             for(int i=1; i <= n; i++) scanf("%d",&a[i]);
23
             scanf("%d",&m);
24
             for(int i=1;i \le m;i++) scanf("%d",&b[i]);
25
             memset(dp,0,sizeof(dp));
26
```

```
27
              memset(path,0,sizeof(path));
              for(int i=1; i <= n; i++)
28
29
                       int tmp=0,k=0;
30
                       for(int j=1; j \le m; j++)
31
32
                                dp[i][j]=dp[i-1][j];
33
                                path[i][j]=j;
34
                                if(a[i]>b[j] && tmp<dp[i-1][j])
35
36
37
                                          tmp=dp[i-1][j];
                                          k=j;
38
39
                                if(a[i]==b[j])
40
41
                                          \begin{array}{l} dp[i][j] = tmp + 1; \\ path[i][j] = k; \end{array}
42
43
                                }
44
                       }
45
46
              int ans=0,p=0;
47
              for(int i=1;i \le m;i++)
48
                       if (ans < dp[n][i])
49
50
                                ans=dp[n][i];
51
52
                                p=i;
53
              printf ("%d\n",ans);
54
              print (n, p, 0);
55
              puts("");
56
              return 0;
57
58
    }
```

Chapter 7

Other Topics

7.1 Dichotomy

7.1.1 condition 1

```
\textbf{while}(L{<}R)
1
2
          int M=L+(R-L+1)/2;
3
          if (judge(M)) L=M;
4
          else R=M-1;
6
   7.1.2 condition 2
  while(L<R)
1
2
          int M=L+(R-L)/2;
3
          if (judge(M)) R=M;
          else L=M+1;
   7.1.3 condition 3
  while(R-L>eps)
1
2
3
          double M=(L+R)/2;
          if (judge(M)) L=M;
          else R=M;
  }
```

7.2 Hash List

Test:UVa 10085

```
const int maxn=1000000;
    const int mod=1000003;
 2
    state st [maxn];
    int head[mod],nxt[mod];
    void init (){memset(head,0,sizeof(head));}
                                                   // initial, 0 represents nonexist
    int try_to_insert (int s)
                                                  //insert function
 7
        int tmp=hash(s);
                                              //calculate the hash value
 8
        int u=head[tmp];
 9
        while(u)
10
11
             if (memcmp(st[u],st[s], sizeof(st[s]))==0) return 0;
12
            u=n\times t[u];
13
14
        nxt[s]=head[tmp];
15
        head[tmp]=s;
16
        return 1;
17
    }
18
```

7.3 Construct a proper answer for N Queens Problem

Test:UVa 10094

```
1. n\%2 \neq 0 \land n\%3 \neq 0

n is even: 2,4,6,8,\ldots,n,1,3,5,7,\ldots,n-1

n is odd: 2,4,6,8,\ldots,n-1,1,3,5,7,\ldots,n

2. n\%2 = 0 \lor n\%3 = 0

when n is even, k is n/2; when n is odd k is (n-1)/2.

n is even and k is even: k,k+2,\ldots,n,2,4,6,\ldots,k-2,k+3,k+5,\ldots,n-1,1,3,5,\ldots,k+1

n is odd and k is even: k,k+2,\ldots,n-1,2,4,6,\ldots,k-2,k+3,k+5,\ldots,n-2,1,3,5,\ldots,k+1,n

n is even and k is odd: k,k+2,\ldots,n-1,1,3,5,\ldots,k-2,k+3,k+5,\ldots,n,2,4,6,\ldots,k+1

n is odd and k is odd: k,k+2,\ldots,n-2,1,3,5,\ldots,k-2,k+3,k+5,\ldots,n-1,2,4,6,\ldots,k+1,n
```