

Homework 7

CS386D Database Systems

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1 Part A**15.3.2**

$B(S) + (B(S)B(R))/(M - 1) = 10000 + \lceil (10000/(1000 - 1)) \rceil \times 10000 = 120000$ disk I/O.

15.3.3 a

$$100000 = 10000 + \lceil 10000/(M - 1) \rceil (10000)$$

$$M = 1112.11$$

$$M \geq 1113$$

15.4.2 a, b, c

- a) $3(B(S) + B(R)) = 3(10000 + 10000) = 60000$ I/Os.
 b) $5(B(S) + B(R)) = 5(10000 + 10000) = 100000$ I/Os.
 c) $3(B(S) + B(R)) = 3(10000 + 10000) = 60000$ I/Os.

15.4.3

For 2PPMS, we can decrease the size of the sublists such that we do have M sublists, eg. we only use pM buffers to sort sublists and sublists are pM long where $0 < p < 1$. Then when we sort a sublist, we can place the first block of the sublist in our unused blocks, the $(1 - p)M$ blocks not used for sorting, until these are full. Then the algorithm would be more or less the same, where we must write out entire lists and reread them to merge sort. Thus we save $2((1 - p)M)$ disk I/Os.

2 Part B**16.4.1 a, d, i**

a) $\frac{100 \times 200 \times 300 \times 400}{60 \times 100 \times 50} = 8000$ tuples

d) Selection gives us $T(\sigma_{c=20}(Y)) = 6$.
 $\frac{6 \times 400}{50} = 48$ tuples.

i) $\frac{200 \times 300}{3} = 20000$ tuples.

3 Part C

1. Text problem 16.6.1a, (which concerns precisely part a of 16.4.1). In addition to answering the text question, show the entries in the dynamic programming matrix.

		W	X	Y	Z
		W	X	Y	Z
Size		100	200	300	400
Cost		0	0	0	0
Best Plan		W	X	Y	Z

	WX	WY	WZ	XY	XZ	YZ
	W	X	Y	Z		
	(W, X)	(W, Y)	(W, Z)	(X, Y)	(X, Z)	(Y, Z)
Size	333	30000	40000	600	80000	2400
Cost	0	0	0	0	0	0
Best Plan	$W \bowtie X$	$W \bowtie Y$	$W \bowtie Z$	$X \bowtie Y$	$X \bowtie Z$	$Y \bowtie Z$

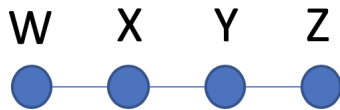
	WXY	WXZ	WYZ	XYZ		
	WX	WY	WZ	XY	XZ	YZ
	W	X	Y	Z		
	(W, X, Y)	(W, X, Z)	(W, Y, Z)	(X, Y, Z)		
Size	1000	133333	240000	4800		
Cost	333	333	2400	600		
Best Plan	$Y \bowtie (W \bowtie X)$	$(W \bowtie X) \bowtie Z$	$W \bowtie (Y \bowtie Z)$	$Z \bowtie (X \bowtie Y)$		

	WXYZ					
	WXY	WXZ	WYZ	XYZ		
	WX	WY	WZ	XY	XZ	YZ
	W	X	Y	Z		
	Plan				Cost	
	$Z \bowtie (Y \bowtie (W \bowtie X))$				1333	
	$Y \bowtie ((W \bowtie X) \bowtie Z)$				133666	
	$X \bowtie (W \bowtie (Y \bowtie Z))$				242400	
	$W \bowtie (Z \bowtie (X \bowtie Y))$				5400	
	$(W \bowtie X) \bowtie (Y \bowtie Z)$				2733	
	$(W \bowtie Y) \bowtie (X \bowtie Z)$				110000	
	$(X \bowtie Y) \bowtie (W \bowtie Z)$				40600	

$Z \bowtie (Y \bowtie (W \bowtie X))$ is the best plan

2. Repeat 16.6.1a, this time.

a) draw the query graph.



b) leveraging the query graph to avoid Cartesian products, limit the amount of work and number of entries you create in the dynamic programming matrix.

```

graph LR
    W
  
```

	W	X	Y	Z
Size	100	200	300	400
Cost	0	0	0	0
Best Plan	W	X	Y	Z

```

graph LR
    W --- X
    X --- Y
    Y --- Z
  
```

	WX	WY	WZ	XY	XZ	YZ
	W	X	Y	Z		
	(W, X)		(X, Y)		(Y, Z)	
Size	333		600		2400	
Cost	0		0		0	
Best Plan	W ⋈ X		X ⋈ Y		Y ⋈ Z	

```

graph LR
    WX --- Y
    Y --- Z
    W --- XY
    XY --- Z
    W --- X
    X --- YZ
  
```

	WXY	WXZ	WYZ	XYZ		
	WX	WY	WZ	XY	XZ	YZ
	W	X	Y	Z		
	(W, X, Y)			(X, Y, Z)		
Size	1000			4800		
Cost	333			600		
Best Plan	Y ⋈ (W ⋈ X)			Z ⋈ (X ⋈ Y)		

WXYZ						Plan	Cost
WXY	WXZ	WYZ	XYZ			$Z \bowtie (Y \bowtie (W \bowtie X))$	1333
WX	WY	WZ	XY	XZ	YZ	$W \bowtie (Z \bowtie (X \bowtie Y))$	5400
W	X	Y	Z			$(W \bowtie X) \bowtie (Y \bowtie Z)$	2733