

Joins (1)

Physical Join Operators
Block Nested Loops
Merge Join(s)

Objectives:

- Introduce three primary join families
 - Block nested loops
 - Merge Join(s)
 - Hash Join(s)
- Examples, first two

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Three Primary Classes of Join Algorithm

1. Block-Nested Loops (vs. Simple Nested Loops.)
2. Merge-Join
3. Hash-Joins

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Simple Nested Loops

- R join S

For each $r \in R$

For each $s \in S$

if $\text{pred}(r,s)$ then output result

// ignores that data is on disk, read in blocks

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Block-Nest Loops - Assumptions

- Dedicate Most Memory to Outer-Loop
- Consider that there are many tuples per page.

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Block Nested-Loops Detailed

```

Join(relation R, relation S, buffer_space M){ // M in units of disk pages
  Until done( table_scan(S,M-1)){ // Read M-1 pages of S
    if equijoin then build index on portion of S in memory
    Open(R);
    Until done(table_scan(R, 1)){
      for each tuple r of R in memory{
        "find joining tuples of S" // if equijoin use index
        "output them" // if thetajoins another loop
      }
    }
  }
}

```

What to Notice:

1. Outer-loop iterates over “inner” relation
2. If equi-join, loop structure nested 3 deep
3. If theta-join, loop structure nested 4 deep
4. Call structure
 - not parameterized by theta/equi
 - theta predicate not passed in

In Your Text, Notice:

- Implementation of operators
 - Use an abstract iterator method
 - Open()
 - Next()
 - Close()
- Hide I/O by abstracting iterators.
 - IT WORKS
- They too ignore a lot of the arguments in the call structure.

Cost Model - Notation

- $B(R)$, number of blocks in R
- $T(R)$, number of tuples in R
- $T(R)/B(R)$, an estimate of the number of tuples per block.

Number of I/O's (cost) of Block-Nested Loops

Run through the complete inner relation, S , for each chunk of R

If S fits in $M-1$ buffers $\rightarrow B(R) + B(S)$

If S does not fit

$\rightarrow B(S)/M-1$ iterations

$I/O \text{ Cost} = (B(S)/M-1) * (M-1 + B(R))$

What About Writing the Result

- Book convention –
 - no cost in the operator for the result
 - in some cases, results are immediate input to another operator.
 - no I/O cost for output
 - introduce an explicit new operator
 - “Save intermediate result to disk”