

1

(A)

input: array a of n elements and array b of m elements  
output: Number of distinct elements

Algorithm: distinct(a, n, b, m)  
result <- 0

```
for i <- 1 to n
  if (NOT(BinarySearch(b, m, a[i])))
    result <- result + 1;
```

```
for j <- 1 to m
  if (NOT(BinarySearch(a, n, b[j])))
    result <- result + 1;
```

return result

input: array a of n elements and array b of m elements  
output: Number of distinct elements in array a from b

Algorithm: distinct\_in\_first(a, n, b, m)  
result <- 0

```
for i <- 1 to n
  if (NOT(BinarySearch(b, m, a[i])))
    result <- result + 1
```

return result;

(B)

input: array a of n elements and array b of m elements  
output: Number of unique elements in the intersection

Algorithm: fixed\_intersection(a, n, b, m)  
intersect <- new int[1]  
count <- 0  
for i <- 1 to n  
 if search(b,m,a[i])  
 //intersect is of size count  
 if NOT search(intersect, count, a[i] )  
 if count == intersect.length  
 intersect2 <- new int[intersect.length\*2]  
 for j <- 1 to intersect.length  
 intersect2[j] <- intersect[j]  
 intersect <- intersect2  
 count <- count + 1  
 intersect[count] <- a[i]  
return count

2

(a)

input: array a of n elements and natural number k

output: array of elements occurring k times in array a

Algorithm: k\_times\_elements(a, n, k)

```
result <- new int[1]
```

```
count <- 0
```

```
for i <- 1 to n
```

```
  if occurrences(a, n, a[i]) <= k AND occurrences(result, count, a[i]) > 0
```

```
    //This is to mimic a list by increasing the size of the array
```

every time it is full

```
    if count == result.length
```

```
      result2 <- new int[result.length+1]
```

```
      for j <- 1 to result.length
```

```
        result2[j] <- result[j]
```

```
      result <- result2
```

```
      count <- count + 1
```

```
      result[count] <- a[i]
```

```
    return result
```

subroutine:

input: array a of n elements and value b

output: number of occurrences of b in array a

Algorithm: occurrences(a, n, b)

```
result <- 0
```

```
for i <- 1 to n
```

```
  if a[i] == b
```

```
    result <- result + 1
```

```
return result
```

(b)

The occurrences subroutine has a running time of  $O(n)$

The k\_times\_elements algorithm has a running time of  $O(n^2)$

In the worst case running time: each iteration is bounded by  $O(2n) = O(n)$ ,

because the occurrences subroutine is called twice,

and i is at most n.

3

(e)

For the multiplyConstant method, the worst case running time is  $O(n)$ . This is because only the inner loop will iterate at most n times and the outer loop iterates 1 time only inside the multiplyPolys method.

For the multiplyPolys method, the worst case running time is  $O(n^2)$

for the clean method, the worst case running time is  $O(n^2)$

For the derive method, the worst case running time is  $O(n)$

4

(c)

The factorization method has a worst case running time of  $O(1)$ . This is

because there is no loop and no matter the input the running time will remain constant and does not grow.