

CSCI 311

Algorithms and Data Structures

CSU Chico

Week 11.01

Today

■ Graphs

- Notation
- Types of graphs
- Representations

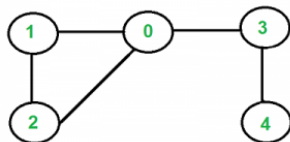
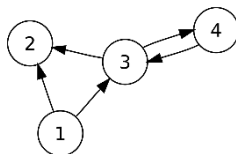
■ Breadth-first search (BFS)

■ Depth-first search (DFS)

■ Edge classification

Graphs

Notation



- $G = (V, E)$
 - › V - a set of nodes or vertices, often labelled $1, \dots, n$
 - › E - a set of edges
- Edges can be undirected, $\{u, v\}$, or directed, (u, v)
- Vertices and edges can have attributes
 - › $v.id, v.cost, v.dist$ where $v \in V$
 - › $(u, v).weight, \{u, v\}.dist$ where $(u, v) \in E$ or $\{u, v\} \in E$
- Run time often measured as a function of $|V|$ and $|E|$

Graphs

Types

- G is a **simple** graph if
 - › Edges are undirected and unweighted
 - › There are no multi-edges (no same edge)
 - › There are no self loops (no edge starts and ends at the same vertex)

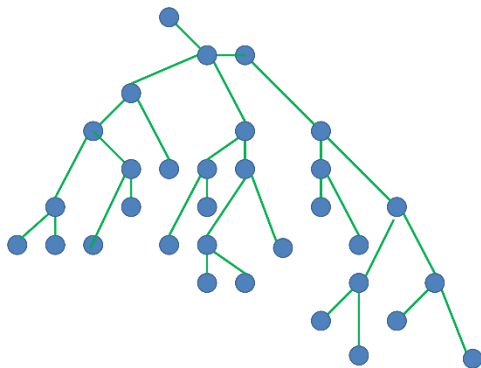


Graphs

Types

■ G is a **tree** if

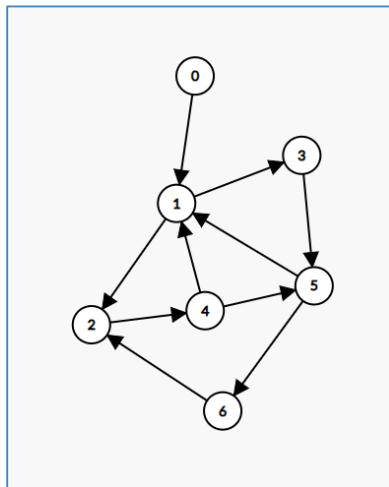
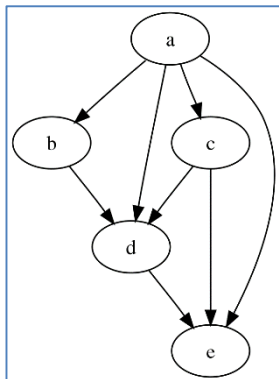
For every pair $u, v \in V$ is connected by exactly one path



Graphs

Types

- G is a **directed acyclic graph (DAG)** if directed but with no cycles



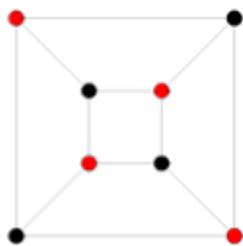
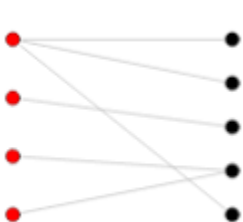
Graphs

Types

■ G is a **bipartite graph** if

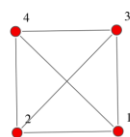
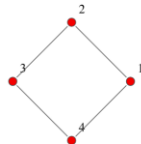
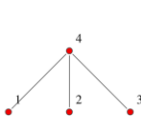
Two disjoint sets of vertices U and V

For all edges $\{u, v\} \in E$, u and v are in different vertex sets



Graphs

Representations



■ Two sets V and E

■ Adjacency matrix

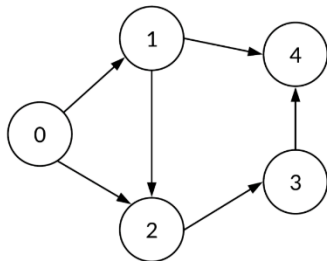
A is an $n \times n$ matrix

$A[i][j] = 1$ if $(i, j) \in E$ and 0 otherwise

$$\begin{pmatrix} 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix}$$

$$\begin{pmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \end{pmatrix}$$

$$\begin{pmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix}$$

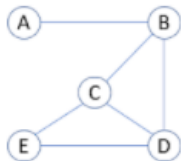


Adjacency Matrix

	0	1	2	3	4
0	0	1	1	0	0
1	0	0	1	0	1
2	0	0	0	1	0
3	0	0	0	0	1
4	0	0	0	0	0

Graphs

Representations

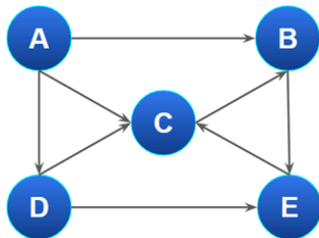
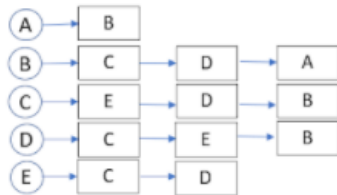


- Two sets V and E

- Adjacency list

A is a list of n lists

$j \in A[i]$ if and only if $(i, j) \in E$



A	B, C, D
B	E
C	B
D	C, E
E	C

Breadth-First Search

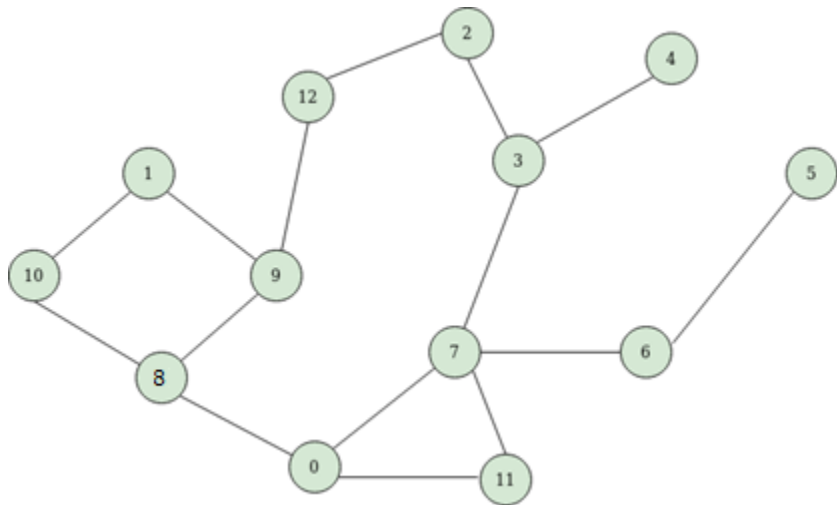
Intuition

Given a graph $G = (V, E)$ and a source node $s \in V$

- 1 Add s to the (initially empty) *work* queue
- 2 Mark s as visited
- 3 While this queue is not empty
 - 1 Remove the first node
 - 2 Mark this node as visited
 - 3 Add all of its unvisited neighbors to the queue

Breadth-First Search

Example



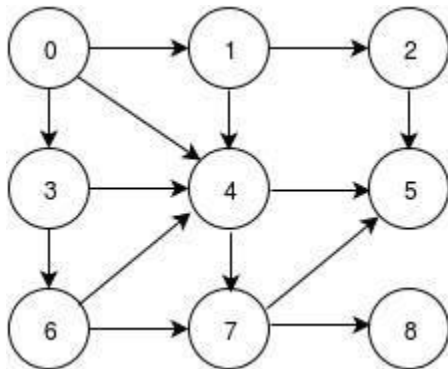
Breadth-First Search

Pseudocode (also calculating the distances)

```
function BFS(G, s)
  for u in V
    u.dist =  $\infty$ 
    u.visited = false
    u.predecessor = null
  s.dist = 0
  s.visited = true
  Q = queue()
  Q.push(s)
  while length(Q) > 0
    u = Q.pop()
    for v in neighbors(G, u)
      if v.visited == false
        v.dist = u.dist + 1
        v.visited = true
        v.predecessor = u
        Q.push(v)
```

Breadth-First Search

Example



Breath-First Search

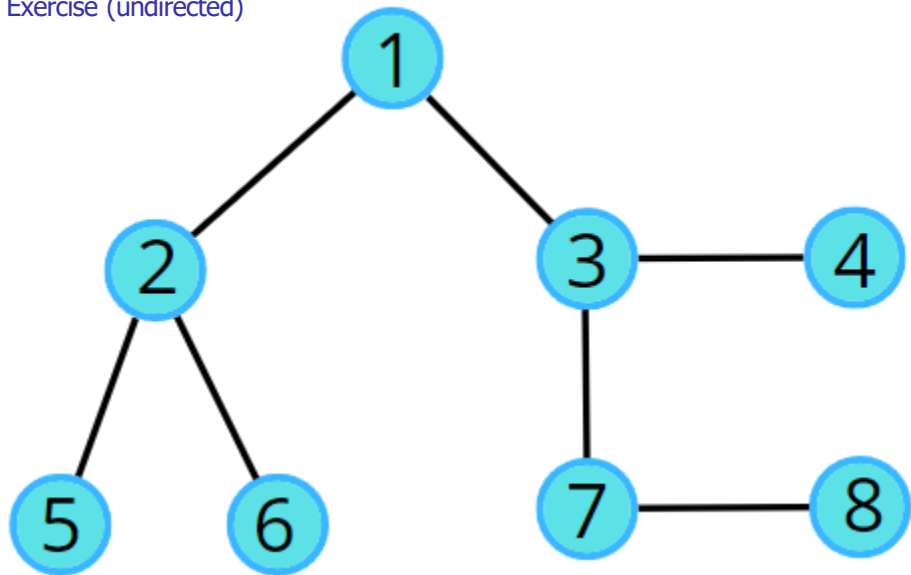
Keep Track of Distance

```
Node{  
    dist  
    visited  
    predecessor  
}
```

```
function BFS(G, s)  
    for u in V  
        u.dist =  $\infty$   
        u.visited = false  
        u.predecessor = null  
    s.dist = 0  
    s.visited = true  
    Q = queue()  
    Q.push(s)  
    while length(Q) > 0  
        u = Q.pop()  
        for v in neighbors(G, u)  
            if v.visited == false  
                v.dist = u.dist + 1  
                v.visited = true  
                v.predecessor = u  
                Q.push(v)
```

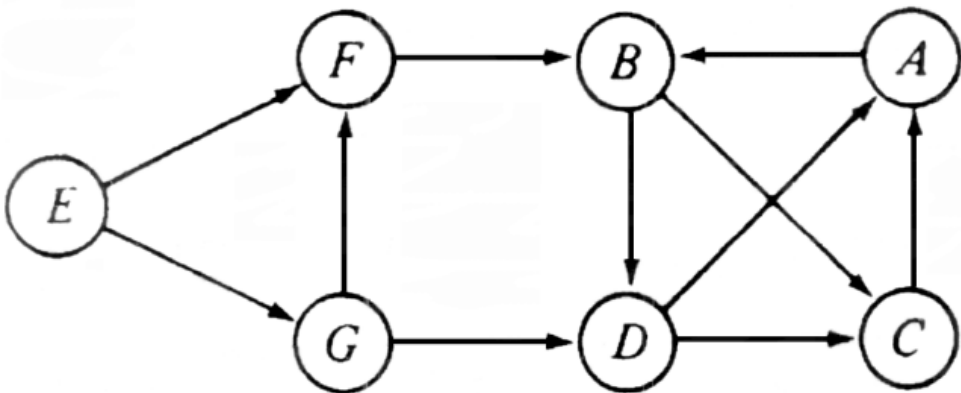
BFS-Undirected-Keep-Distance

Exercise (undirected)

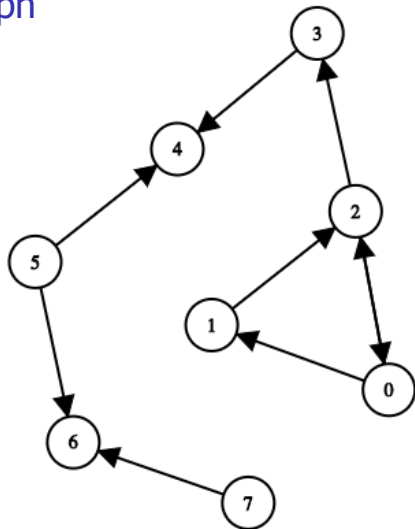


BFS-Directed-Keep-Distance

Exercise (directed)

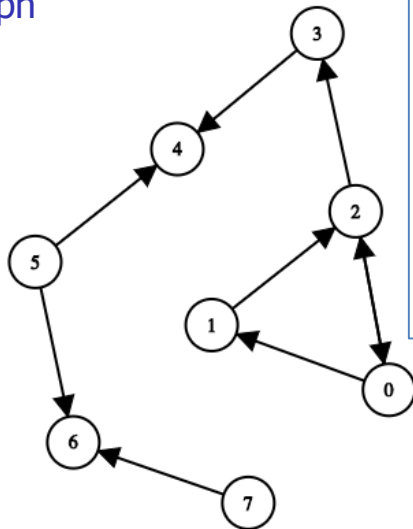


BFS on Disconnected Graph



```
function BFS(G, s)
  for u in V
    u.dist =  $\infty$ 
    u.visited = false
    u.predecessor = null
  s.dist = 0
  s.visited = true
  Q = queue()
  Q.push(s)
  while length(Q) > 0
    u = Q.pop()
    for v in neighbors(G, u)
      if v.visited == false
        v.dist = u.dist + 1
        v.visited = true
        v.predecessor = u
        Q.push(v)
```

BFSN on Disconnected Graph



Function **BFSN(G)**

```
for u in V
```

```
    u.dist =  $\infty$ 
```

```
    u.visited = false
```

```
    u.predecessor = null
```

While any s in V is not visited

```
    s.dist = 0
```

```
    s.visited = true
```

```
    Q = queue()
```

```
    Q.push(s)
```

```
    while length(Q) > 0
```

```
        u = Q.pop()
```

```
        for v in neighbors(G, u)
```

```
            if v.visited == false
```

```
                v.dist = u.dist+1
```

```
                v.visited = true
```

```
                v.predecessor = u Q.push(v)
```

Depth-First Search (DFSVisit)

Intuition

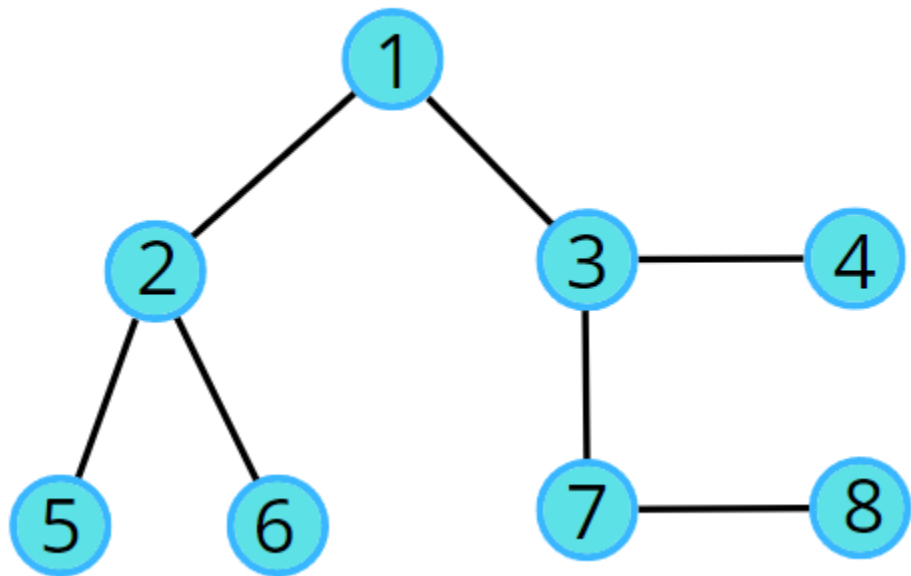
Given a graph $G = (V, E)$

- 1 Pick a source node s
- 2 Mark s as visited
- 3 for all neighbors u of s
 - 1 If u has not been visited, run depth first search on G with u as the source
- 4 Repeat until all vertices have been visited

Lab 7 Challenge DFS Discover & Finish Times

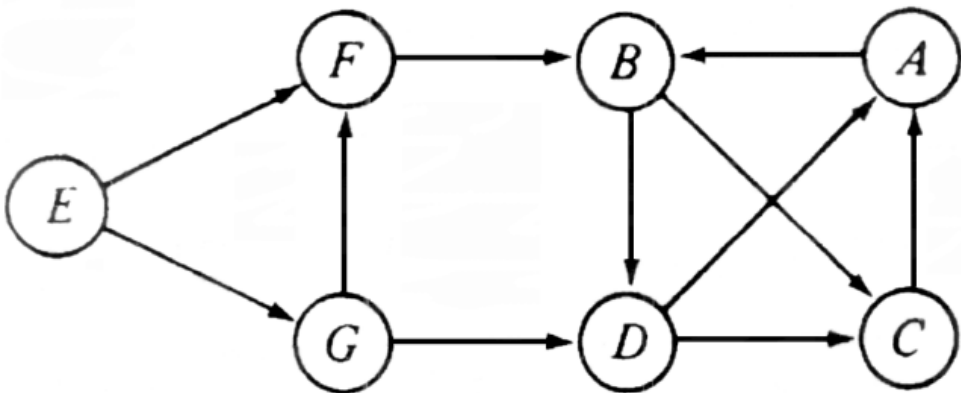
DFS-Undirected

Example (Discovery & Finish Time)

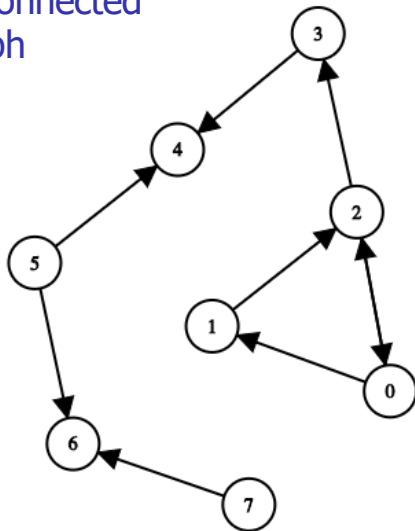


DFS-Directed

Example (Discovery & Finish Time)



DFS on Disconnected Graph



```
function DFS(G)
  for u in V
    u.visited = false
    u.predecessor = null
    u.discovered = -1
    u.finished = -1
  time = 0
  for u in V
    if u.visited == false
      time = DFSVisit(G, u, time)
```

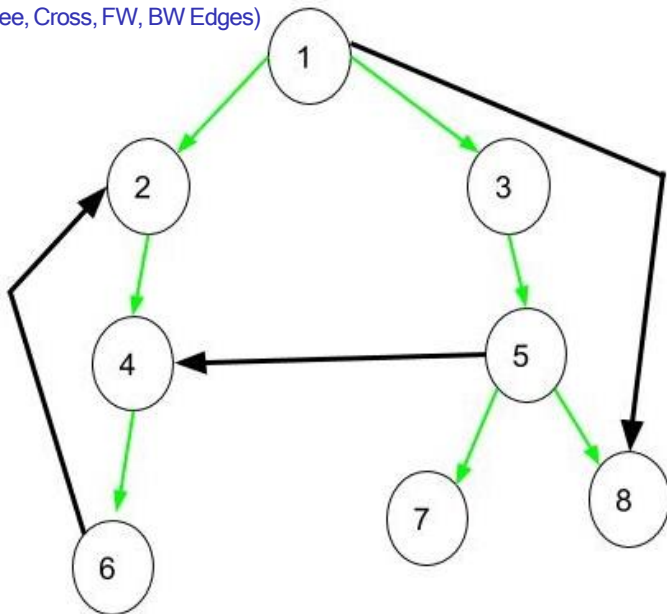
```
function DFSVisit(G, u, time)
  time = time + 1
  u.discovered = time
  u.visited = true
  for v in neighbors(G, u)
    if v.visited == false
      v.predecessor = u
      time = DFSVisit(G, v, time)
  time = time + 1
  u.finished = time
  return time
```

Classification of Edges (Directed)

- Let $G = (V, E)$, $s \in V$, and T_s be a D(B)FS tree
- Tree edge
 - › $(u, v) \in E$ is a tree edge if $\{u, v\}$ is in T_s
 - › v was visited following $(u, v) \in E$
- Back edge
 - › $(u, v) \in E$ is a back edge if it is not a tree edge and v is an ancestor of u in T_s
 - › u is visited as a result of the call $\text{DFS}(G, v)$
- Forward edge
 - › $(u, v) \in E$ is a forward edge if it is not a tree edge and v is a descendent of u in T_s
 - › v is visited as a result of the call $\text{DFS}(G, u)$
- Cross edge
 - › All other edges

DFS Tree & Edges

Example (Tree, Cross, FW, BW Edges)

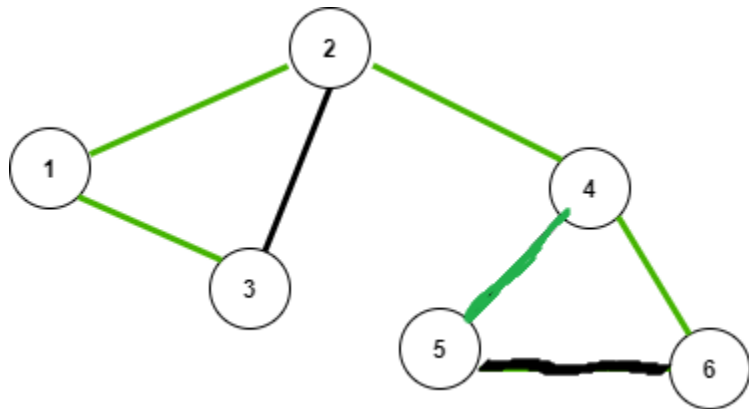


Classification of Edges (Undirected)

- Let $G = (V, E)$, $s \in V$, and T_s be a B(D)FS
- tree Tree edge
 -) $(u, v) \in E$ is a tree edge if $\{u, v\}$ is in T_s
 -) v was visited following $(u, v) \in E$
- Cross edge
 -) All other edges

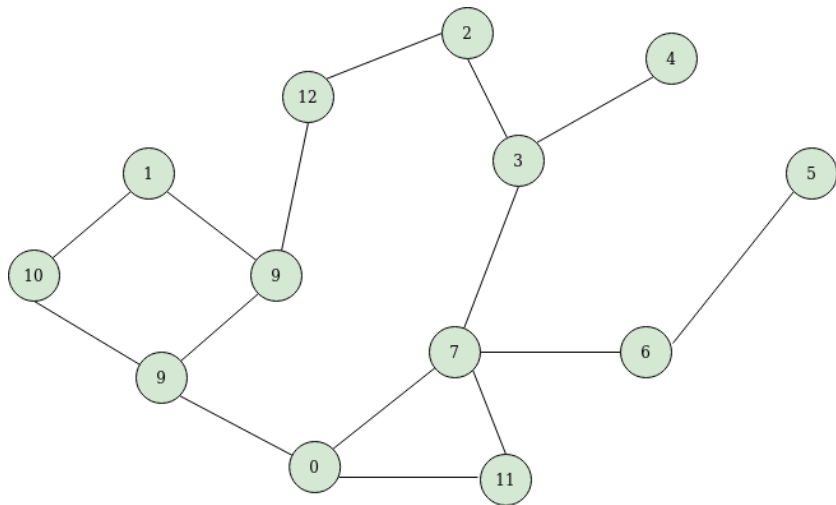
BFS Tree & Edges

Example (Tree & Cross Edges)



Depth-First Search (DFSVisit)

Example (Undirected) - Classification of Edges



Depth-First Search (DFSVisit)

Example (Directed) - Classification of Edges

