VE281

Data Structures and Algorithms

Trees

Learning Objectives:

- Know some basic terminology of trees and binary trees
- Know some basic properties of binary trees
- Know how to represent a binary tree by an array and a linked list

Outline

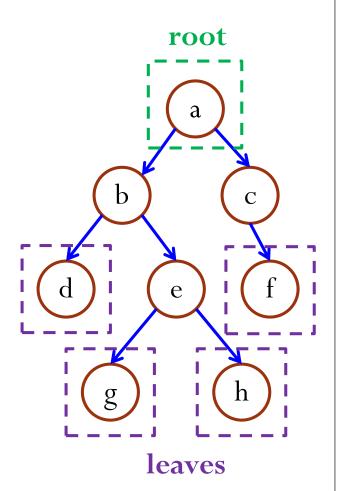
- Trees
- Binary Trees

Trees

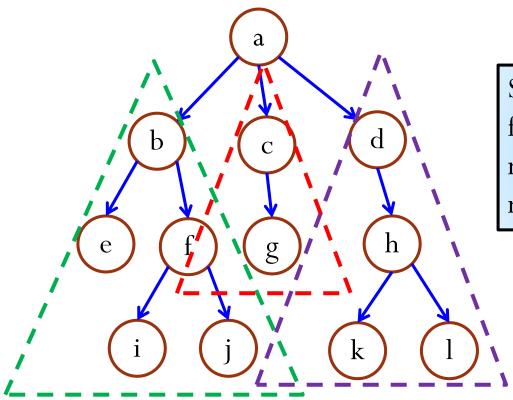
- Tree is an extension of linked list data structure:
 - Each node connects to **multiple** nodes.
- A tree is a "natural" way to represent hierarchical structure and organization.
- Many problems in computer science can be solved by breaking it down into smaller pieces and arranging the pieces in some form of hierarchical structure.
 - For example: merge sort.

Tree Terminology

- Just like lists, trees are collections of nodes.
- The node at the top of the hierarchy is the **root**.
- Nodes are connected by edges.
- Edges define **parent-child** relationship.
 - Root has no parent.
 - All other nodes have <u>exactly one</u> parent.
- A node with no children is called a **leaf**.



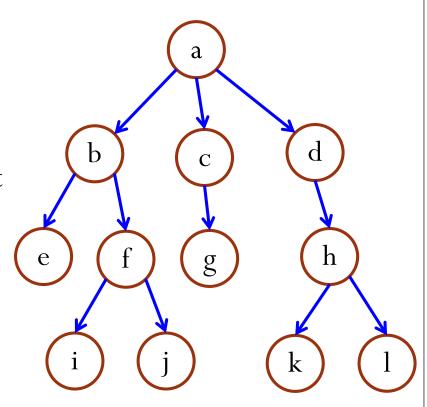
Subtrees



Subtree can be defined for any node in general, not just for the root node.

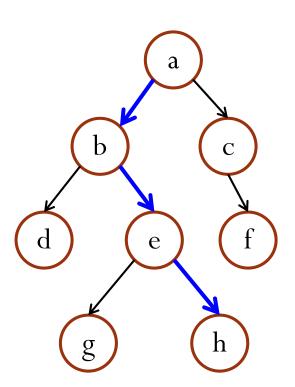
More Tree Terminology

- f is the **child** of b.
- b is the **parent** of f.
- Nodes that share the same parent are **siblings**.
 - b and c are the **siblings** of d.
 - e is the **sibling** of f.



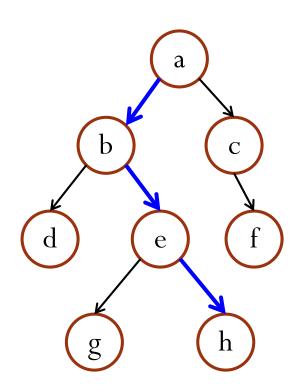
Path

- A **path** is a sequence of nodes such that the next node in the sequence is a child of the previous.
 - E.g., $a \rightarrow b \rightarrow e \rightarrow h$ is a path.
 - The path length is 3.
- Path length may be 0, e.g., b going to itself is a path and its length is 0.
- Claim: If there exists a path between two nodes, then this path is the unique path between these two nodes.



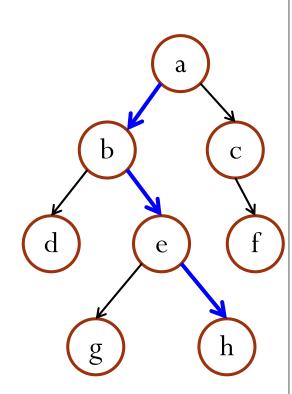
Ancestors and Descendants

- If there exists a path from a node A to a node B, then A is an **ancestor** of B and B is a **descendant** of A.
 - E.g., a is an ancestor of h and h is a descendant of a.



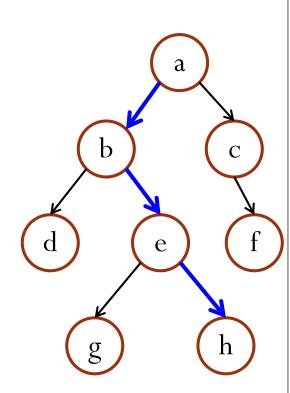
Depth, Level, and Height of a Node

- The **depth** or **level of a node** is the length of the unique path from the **root** to the node.
 - E.g., depth(b)=1, depth(a)=0.
- The height of a node is the length of the longest path from the node to a leaf.
 - E.g., height(b)=2, height(a)=3.
 - All leaves have height zero.



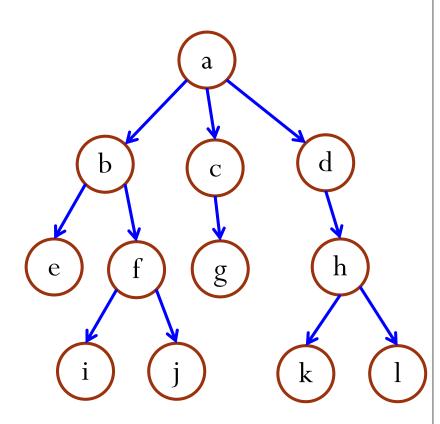
Depth, Level, and Height of a Tree

- The **height of a tree** is the height of its root.
 - This is also known as the **depth of a tree**.
 - The depth of the tree on the right is 3.
- The **number of levels of a tree** is the height of the tree **plus one**.
 - The number of levels of the tree on the right is 4.



Degree

- The **degree of a node** is the number of children of a node.
 - E.g., degree(a) = 3, degree(c) = 1.
- The degree of a tree is the maximum degree of a node in the tree.
 - The degree of the tree on the right is 3.



A Simple Implementation of Tree

- Each node is part of a **linked list** of **siblings**.
- Additionally, each node stores a pointer to its **first child**.

```
struct node {
  Item item;
  node *firstChild;
  node *nextSibling;
};
```

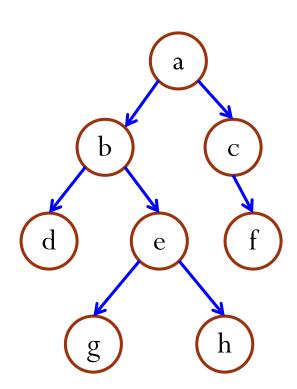
Outline

- Trees
- Binary Trees

Binary Tree

• Every node can only have **at most two** children.

• An empty tree is a special binary tree.



Binary Tree Properties

- What is the **minimum** number of nodes in a binary tree of height h (i.e., has h + 1 levels)?
 - Answer: At least one node at each level.
 - h + 1 levels means at least h + 1 nodes.
- What is the **maximum** number of nodes in a binary tree of height h (i.e., has h+1 levels)?
 - Answer: At most 2^k nodes at level k.
 - Maximum number of nodes is

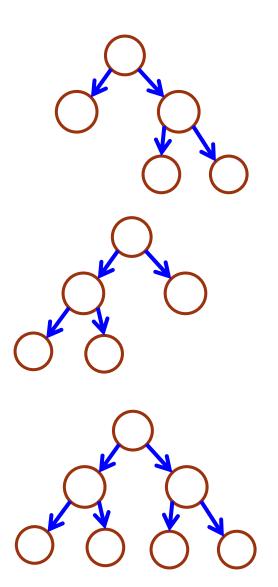
$$1 + 2 + 2^2 + \dots + 2^h = 2^{h+1} - 1$$

Number Of Nodes and Height

- Claim (from the previous slide): Let n be the number of nodes in a binary tree whose height is h (i.e., has h+1 levels).
 - We have $h + 1 \le n \le 2^{h+1} 1$.
- Question: given n nodes, what is the height h of the tree?
 - $\log_2(n+1) 1 \le h \le n-1$

Types of Binary Trees

- A binary tree is **proper** if every node has 0 or 2 children.
- A binary tree is **complete** if:
- 1. every level **except** the lowest is fully populated, and
- 2. the lowest level is populated from left to right.
- A binary tree is **perfect** if **every level** is fully populated.

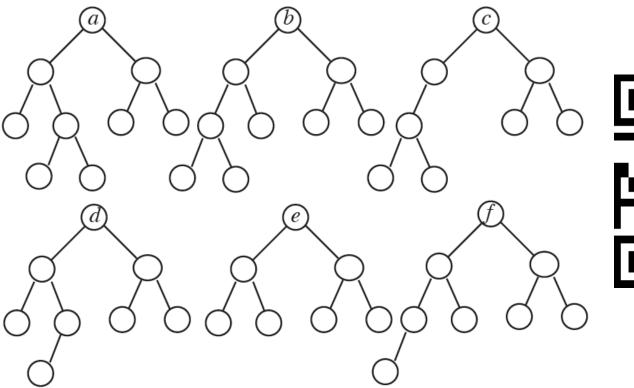




Which Statements Are Correct?

A. Trees a and d are proper. **B.** Tree c is complete.

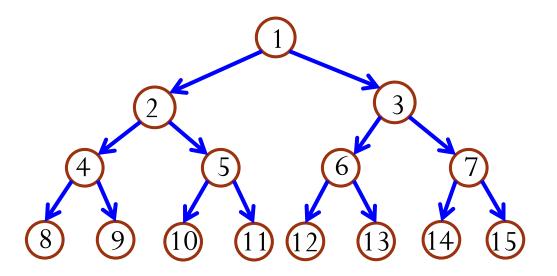
 ${\bf C.}$ Trees b and f are complete. ${\bf D.}$ Tree e is perfect.



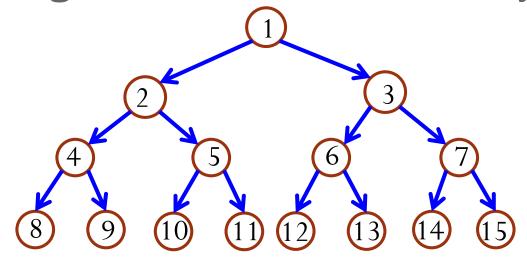


Numbering Nodes In a Perfect Binary Tree

- Numbering nodes from 1 to $2^{h+1} 1$.
- Numbering from top to bottom level.
- Within a level, numbering from left to right.



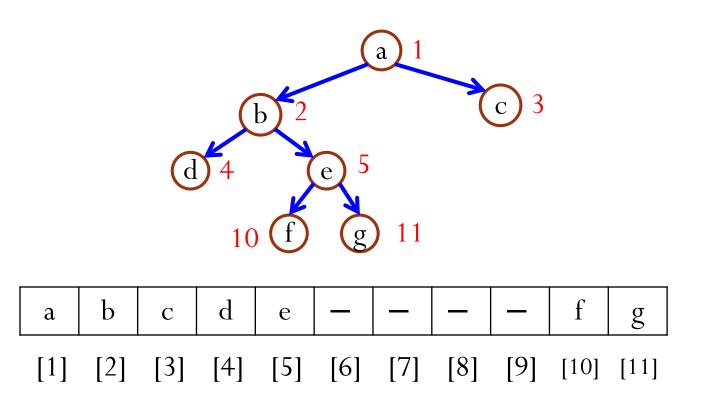
Numbering Nodes In a Perfect Binary Tree



- What is the parent of node i?
 - For $i \neq 1$, it is $\lfloor i/2 \rfloor$. For node 1, it has no parent.
- What is the left child of node i? Let *n* be the number of nodes.
 - If $2i \le n$, it is 2i; If 2i > n, no left child.
- What is the right child of node i?
 - If $2i + 1 \le n$, it is 2i + 1; If 2i + 1 > n, no right child.

Representing Binary Tree Using Array

- Based on the numbering scheme for a perfect binary tree.
- If the number of the node in a perfect binary tree is i, then the node is put at index i of the array.



How Would You Represent a Rightskewed Binary Tree?

• Assume array index starts from 1.

[1] [3] [5] [7] [9] [11] [13] [15]

a
b
c
-

C. | a | - | b | - | - | c | - | - | - | - | d

D. a - b - c - d - - - - - - [1] [3] [5] [7] [9] [11] [13] [15]

An n node binary tree needs an array whose length is between n + 1 and 2^n .



Representing Binary Tree Using Linked Structure

```
struct node {
  Item item;
  node *left;
  node *right;
};
```

- left/right points to a left/right subtree.
 - If the subtree is an empty one, the pointer points to **NULL**.
- For a leaf node, both its left and right pointers are NULL.

