



# Topic 1

## Introduction to Computer

# The Computer Revolution

- Progress in computer technology
  - Moore's Law
- Makes novel applications feasible
  - Computers in automobiles
  - Cell phones
  - Biomedical project
  - Internet+
  - Security
- Computers are pervasive



# Classes of Computers

- Desktop computers
  - General purpose, variety of software
  - Subject to cost/performance tradeoff
- Embedded computers
  - Hidden as components of systems
  - Subject to power/performance/cost constraints
- Server computers
  - Network based
  - High capacity, performance, reliability
  - Range from small servers to building sized

# Where can we find Computers

- Desktop, Laptop, hand held PC, ...
- Automotive
  - Automatic Ignition Systems, Cruise, ABS, traction control, airbag release system...
- Consumer Electronics
  - TV, PDA, appliances, toys, cell phones, camera ...
- Industrial Control
  - robotics, control systems ...
- Medical
  - Infusion Pumps, Dialysis Machines, Prosthetic Devices, Cardiac Monitors, ...
- Networking
  - wired and wireless routers, hubs, ...
- Office Automation
  - fax, photocopiers, printers, scanners, ...
- Aerospace applications
  - Flight-control systems, engine controllers, auto-pilots and passenger in-flight entertainment systems...
- Defense systems
  - Radar systems, fighter aircraft flight-control systems, radio systems, and missile guidance systems...





# Product: Hunter Programmable Digital Thermostat.

## Microprocessor: 4-bit

by Daniel W. Lewis



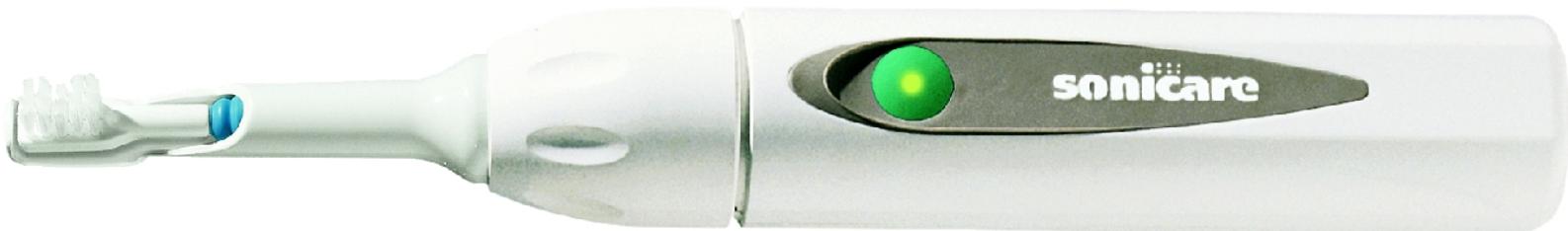


**Product: Vendo V-MAX 720 vending machine.**

**Microprocessor:  
8-bit Motorola  
68HC11.**

by Daniel W. Lewis

**Product: Sonicare Plus toothbrush.  
Microprocessor: 8-bit Zilog Z8.**



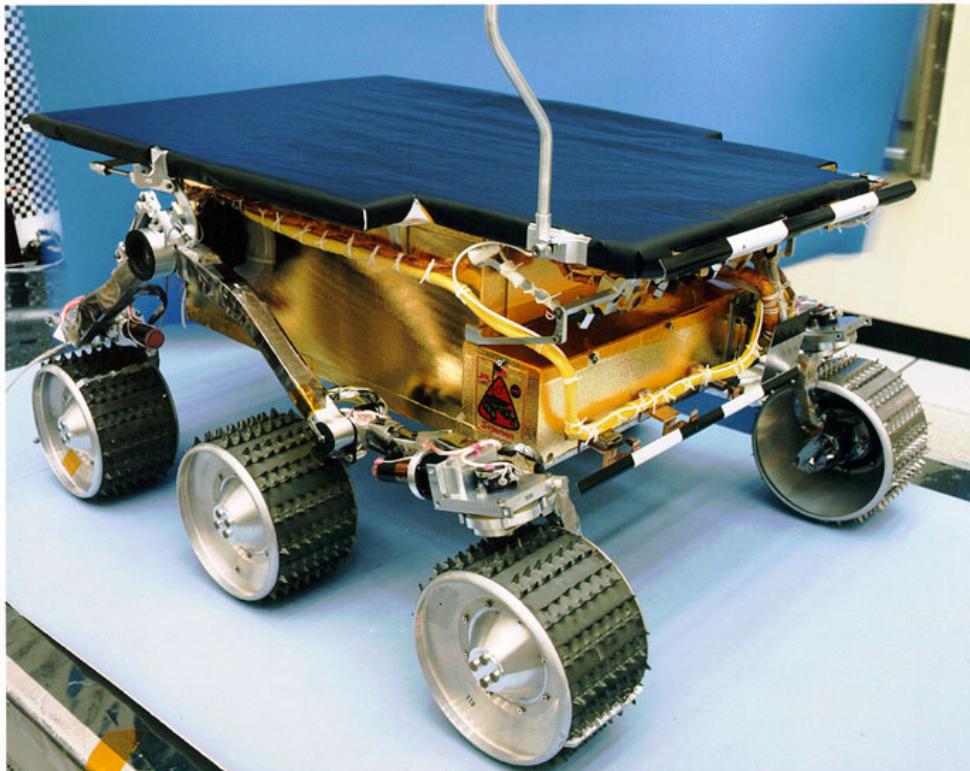
by Daniel W. Lewis



## Product: Miele dishwashers.

**Microprocessor:  
8-bit Motorola  
68HC05.**

by Daniel W. Lewis



# Product: NASA's Mars Sojourner Rover.

## Microprocessor: 8-bit Intel 80C85.

by Daniel W. Lewis



**Product: CoinCo  
USQ-712 coin  
changer.**

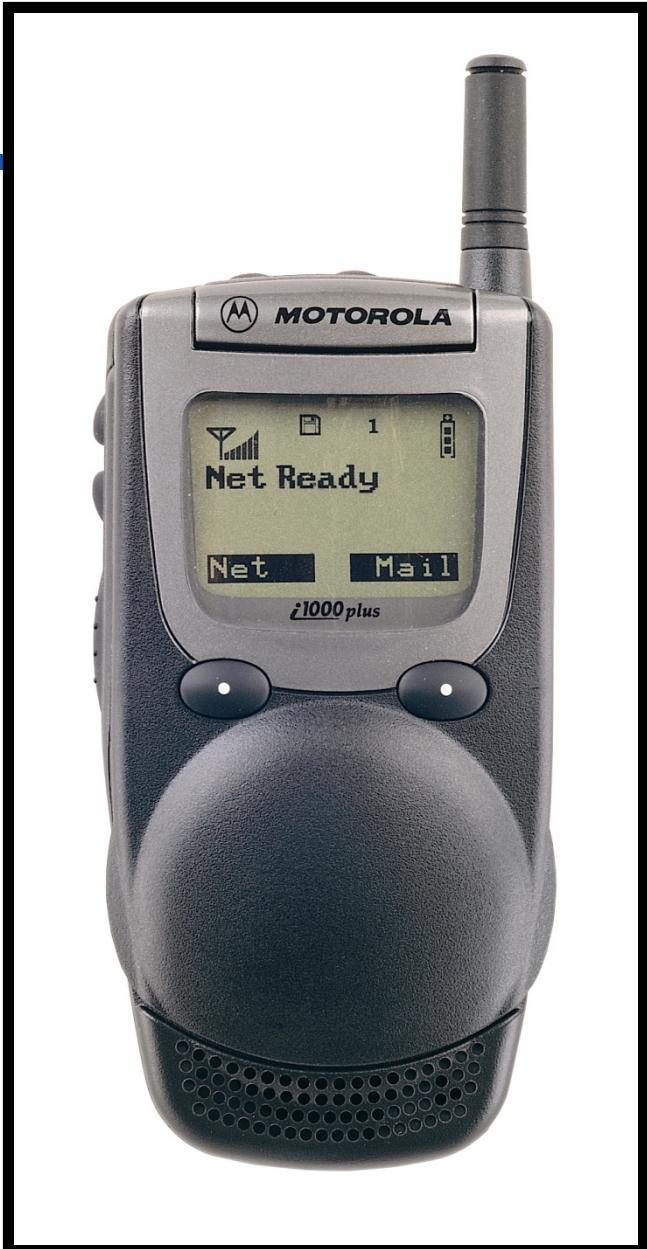
**Microprocessor:  
8-bit Motorola  
68HC912.**

by Daniel W. Lewis



**Product: Garmin  
Nuvi GPS  
Receiver**

**Microprocessor:  
32-bit.**



# Product: Motorola i1000plus iDEN Multi- Service Digital Phone.

## Microprocessor: Motorola 32-bit MCORE.

by Daniel W. Lewis



**Product: Nintendo Wii Controller**  
**Microprocessor: IBM 32-bit Power RISC**



## Product: Apple iPad

**Microprocessor:  
Apple A4,  
a 32-bit ARM RISC.**



**Product: Apple iWatch**

**Microprocessor:  
32-bit Apple A6 and  
M7 Coprocessor**



## Product: HUAWEI P20 Pro

**Microprocessor:  
8 64-bit ARMv8-A, the  
same architecture as  
iPhone X**

by Daniel W. Lewis

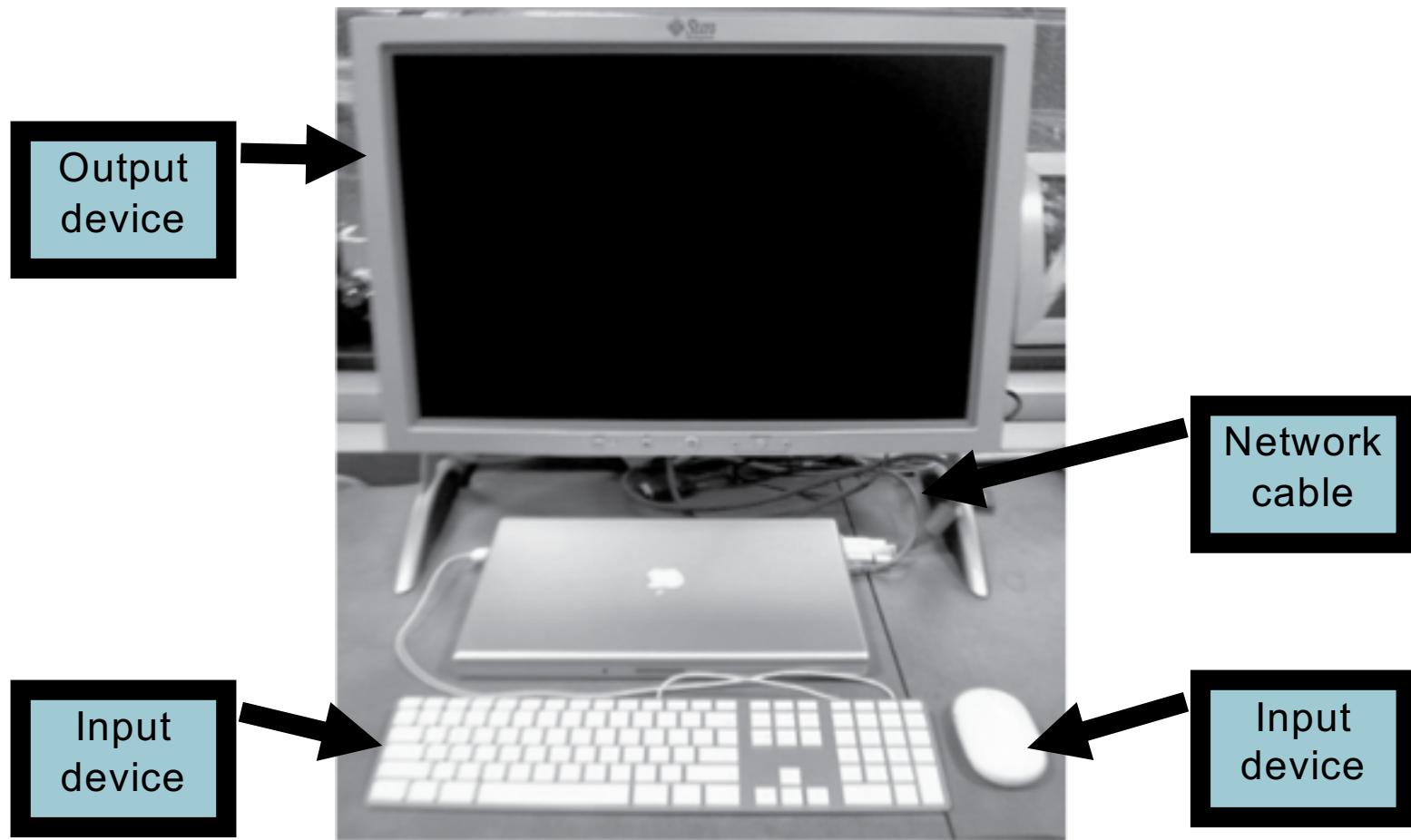


**Product: Sony Aibo  
ERS-110 Robotic  
Dog.**

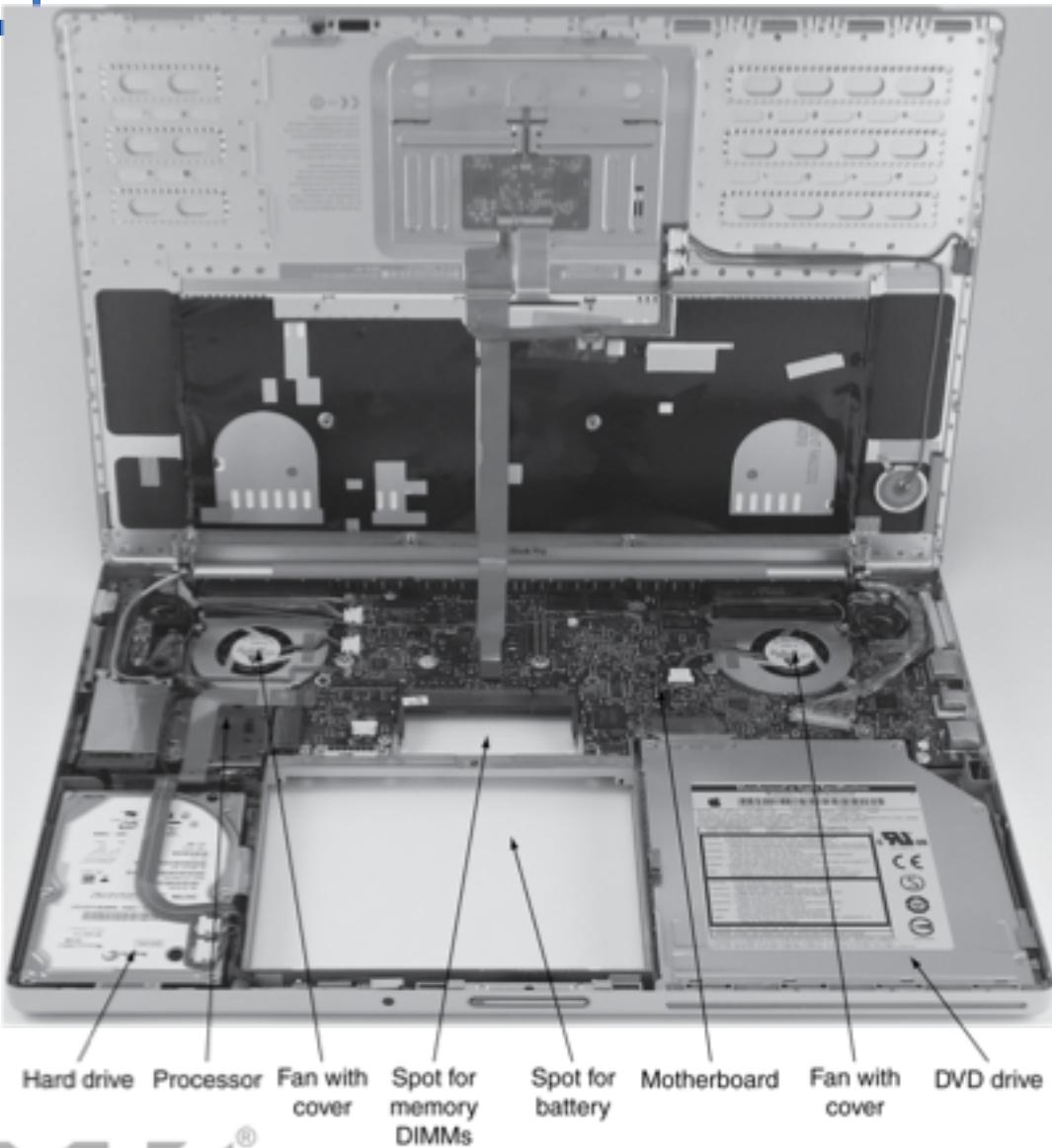
**Microprocessor:  
64-bit MIPS RISC.**

by Daniel W. Lewis

# Anatomy of a Computer



# Opening the Box

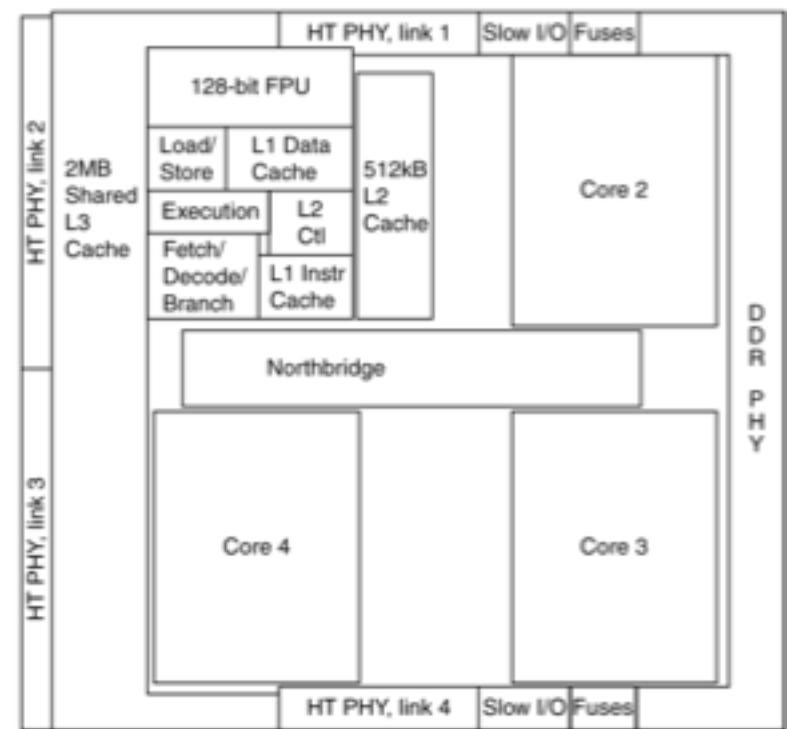
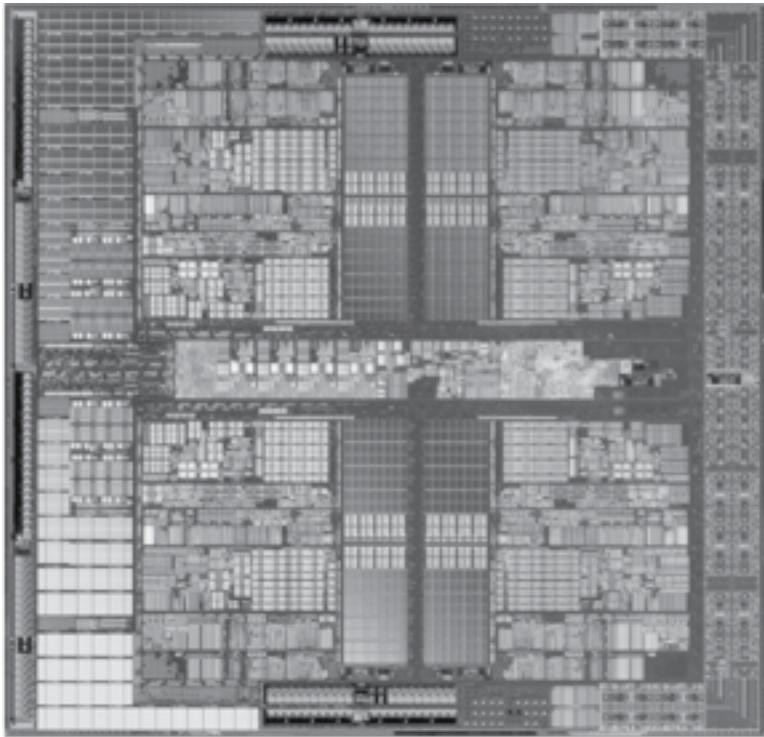


# Inside the Processor (CPU)

- Datapath: performs operations on data
- Control: controls how data flows
- Cache memory
  - Small fast SRAM memory for immediate access to data

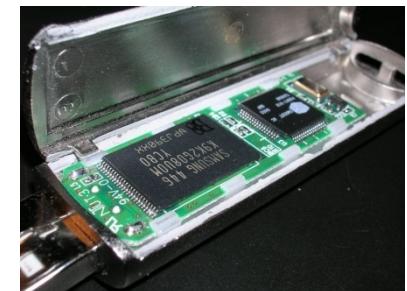
# Inside the Processor

- AMD Barcelona: 4 processor cores



# Peripheral – Memory

- Volatile main memory
  - Loses instructions and data when power off
- Non-volatile secondary memory
  - Magnetic disk
  - Flash memory
  - Optical disk (CDROM, DVD)



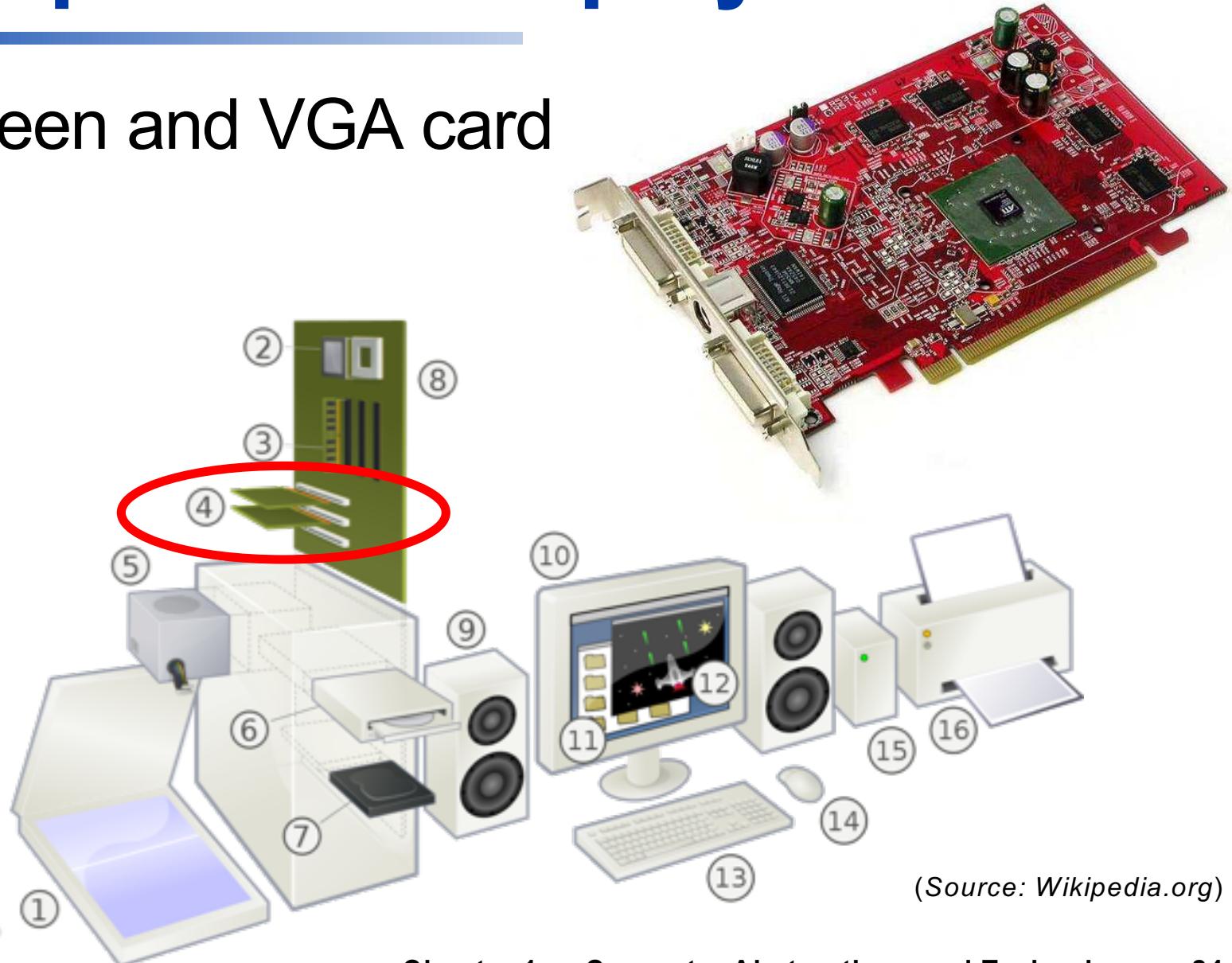
# Peripheral – Networks

- Communication and resource sharing
- Local area network (LAN): Ethernet
  - Within a building
- Wide area network (WAN): the Internet
- Wireless network: WiFi, Bluetooth



# Peripheral – Display

## Screen and VGA card



# Together with Software



- Application software
  - Written in high-level language (HLL)
- System software
  - Operating System: service code
    - Handling input/output
    - Managing memory and storage
    - Scheduling tasks & sharing resources
  - Written in C and assembly
- Hardware
  - Processor, memory, I/O controllers

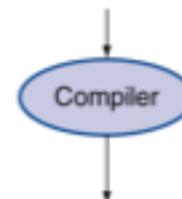
# Levels of Program Code

## ■ High-level language

- ## ■ What we use

High-level  
language  
program  
(in C)

```
swap(int v[], int k)
{int temp;
    temp = v[k];
    v[k] = v[k+1];
    v[k+1] = temp;
}
```

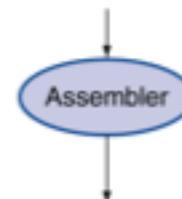


## ■ Assembly language

- What both we and computers can use

## Assembly language program (for MIPS)

```
swap:  
    mul $2. $5.4  
    add $2. $4,$2  
    lw   $15. 0($2)  
    lw   $16. 4($2)  
    sw   $16. 0($2)  
    sw   $15. 4($2)  
    jr   $31
```



## ■ Machine instruction

- ## ■ What computers use

Binary machine  
language  
program  
(for MIPS)

```
00000000101000010000000000000011000  
00000000000110000001100000100001  
1000110001100010000000000000000000  
1000110011110010000000000000000100  
1010110011110010000000000000000000  
1010110001100010000000000000000100  
000000111110000000000000000001000
```

# Levels of Program Code

## ■ High-level language

- Syntax is similar to English
  - A translator is required to translate the program – compile
  - Allows the user to work on the program logic at higher level

High-level  
language  
program  
(in C)

```

swap(int v[], int k)
{int temp;
    temp = v[k];
    v[k] = v[k+1];
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```

## Assembly language program (for MIPS)

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swap:  
    mul $2, $5,4  
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```

Binary machine  
language  
program  
(for MIPS)



# Levels of Program Code

## ■ Assembly language

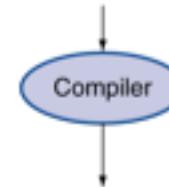
- Composed of assembly **instructions**
  - An assembly instruction is a mnemonic representation of a machine instruction
  - Assembly instruction must be translated before it can be executed – assembler
  - Programmers need to work on the program logic at a very low level, hard to achieve high productivity.

High-level  
language  
program  
(in C)

```

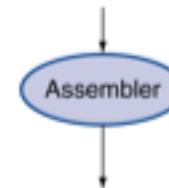
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```



## Assembly language program (for MIPS)

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swap:  
    mul $2, $5,4  
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    lw   $15, 0($2)  
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    sw   $15, 4($2)  
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```



Binary machine language program (for MIPS)

```
00000000101000010000000000000011000  
000000000000110000001100000100001  
1000110001100010000000000000000000  
1000110011110010000000000000000100  
1010110011110010000000000000000000  
1010110001100010000000000000000100  
0000001111100000000000000000000000001000
```

# Levels of Program Code

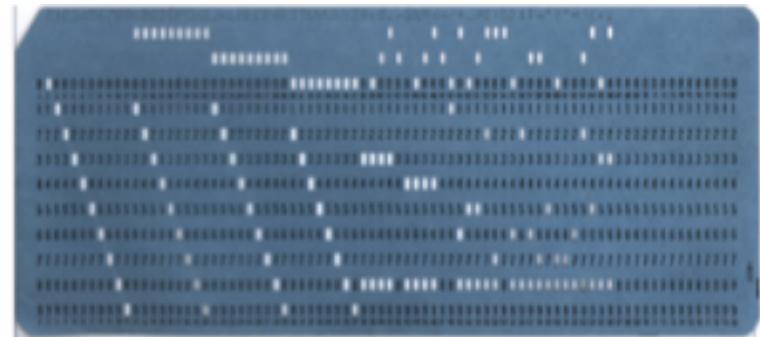
## Machine instruction

- A sequence of binary digits which can be executed by the processor
- Hard to understand, program, and debug for human being



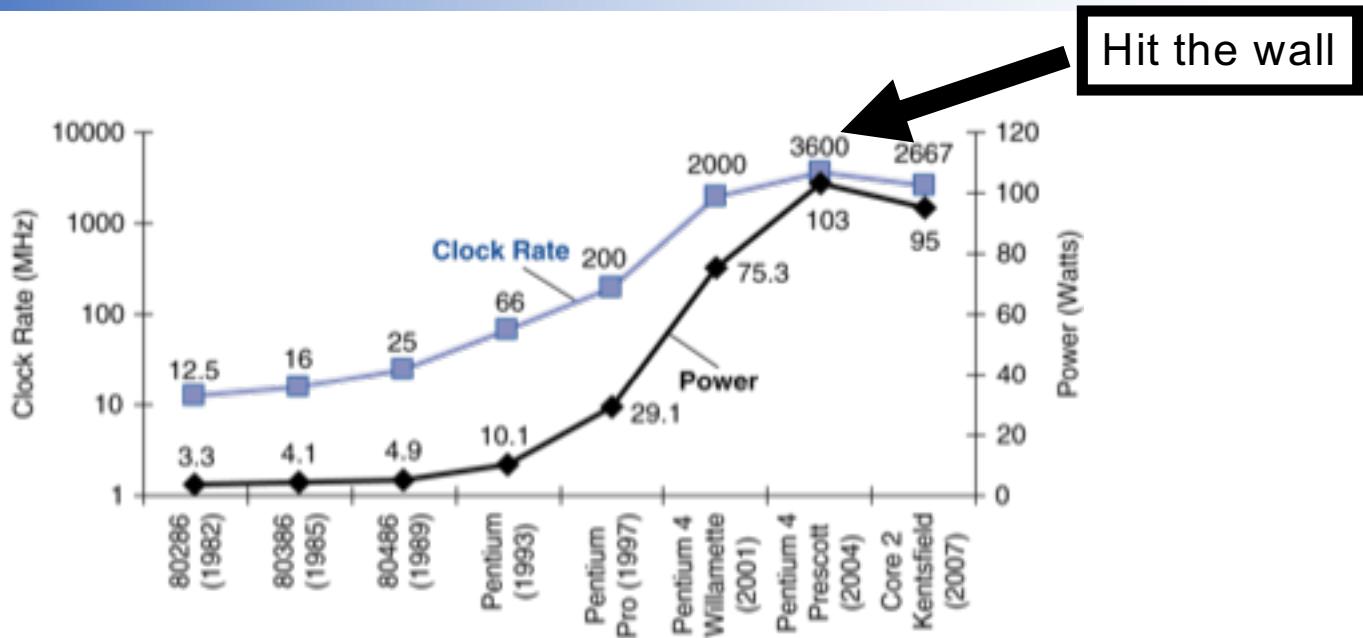
Binary machine language program (for MIPS)

```
000000000101000010000000000000110000  
00000000000011000000011000000100001  
100011000110001000000000000000000000  
100011001111001000000000000000000000  
101011001111001000000000000000000000  
101011000110001000000000000000000000  
00000001111100000000000000000000000000001000
```



*From Wikipedia*

# Power Trends



- In CMOS IC technology

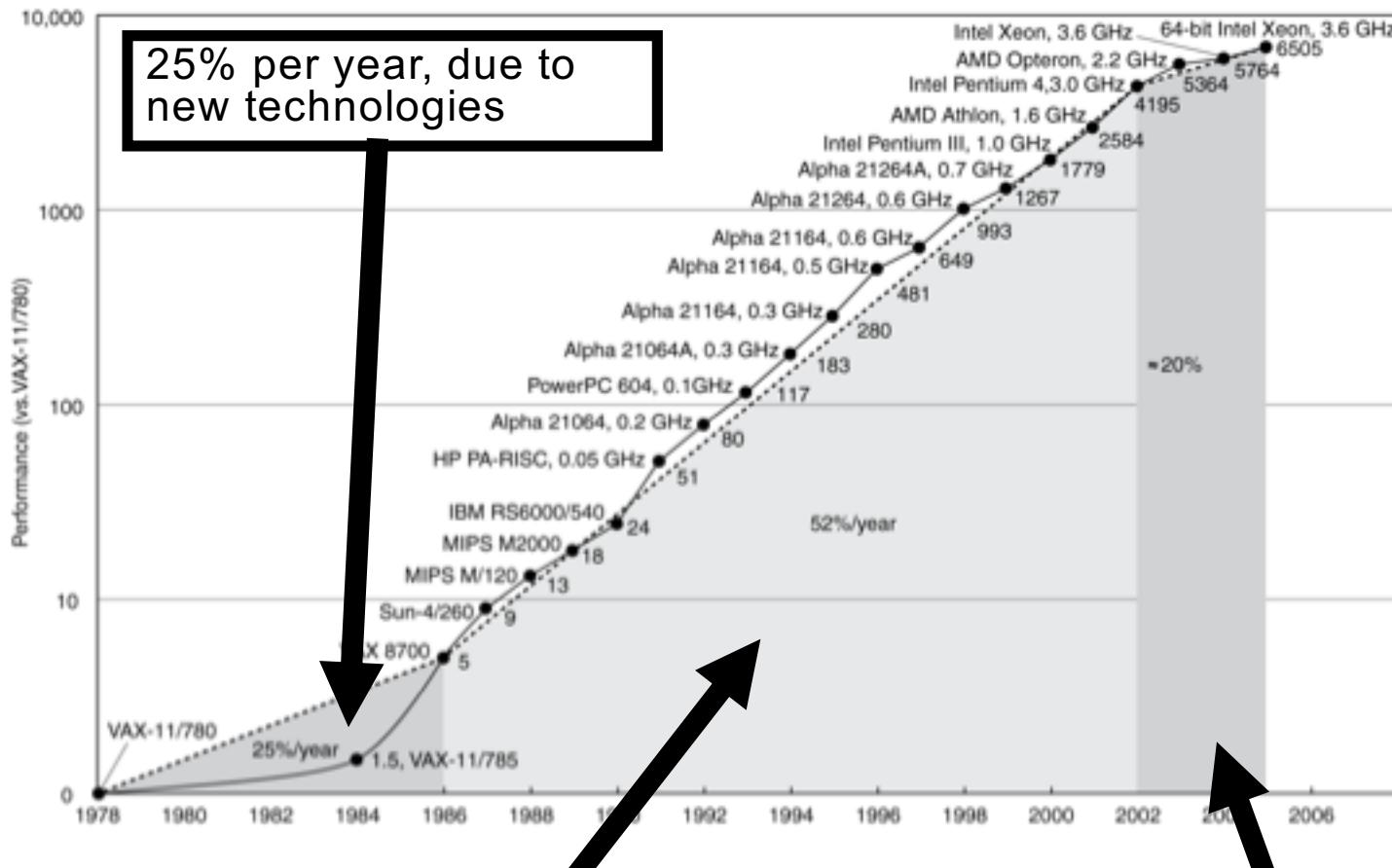
$$\text{Power} = \text{Capacitive load} \times \text{Voltage}^2 \times \text{Frequency}$$

× 30

5V → 1V in 15y,  
15% off/generation

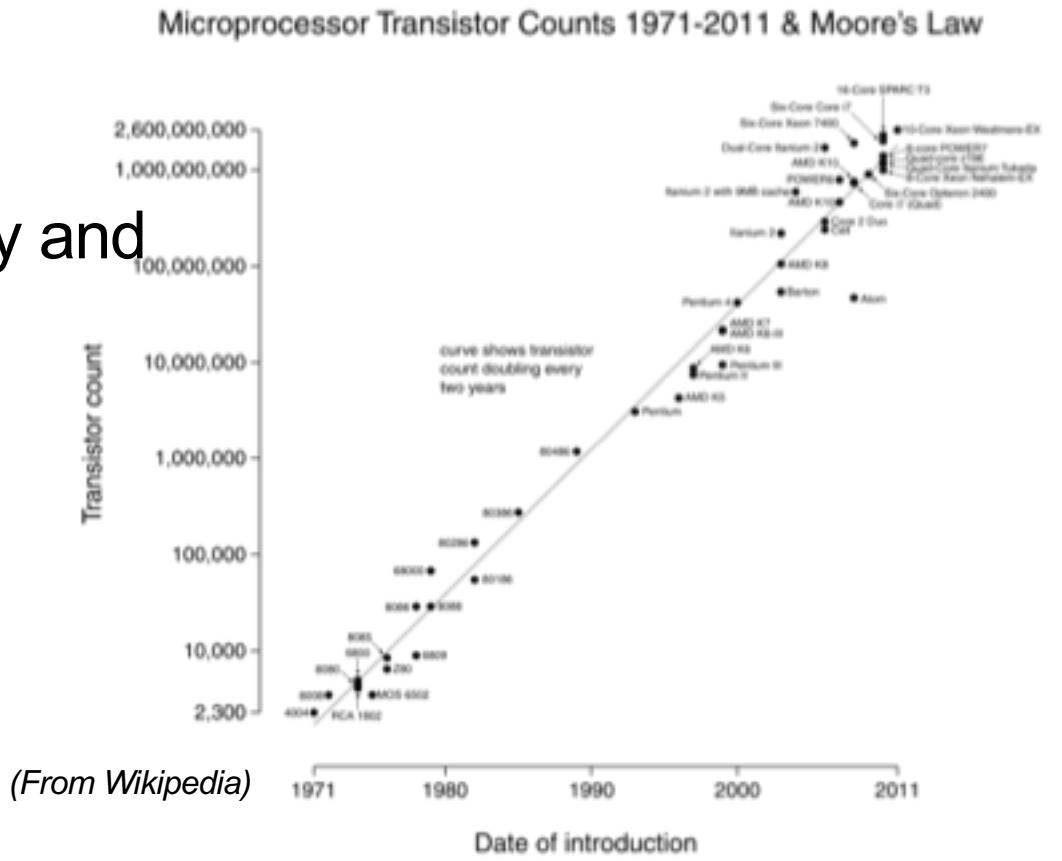
× 1000

# Uniprocessor Performance



# Technology Trends

- Electronics technology continues to evolve
  - Reduced cost
  - Parallelism
  - Low power
  - Increased capacity and performance



# How to Measure Computer Performance?

- Execution (response) time
  - How long it takes to do a task
  - Measure for PCs and embedded computers
- Throughput
  - Total work done per unit time
  - Servers
- We'll focus on execution time for now

# Measuring Performance - Time

- Elapsed time – for System Performance
  - Total execution time to complete a task, including
    - Processing, I/O, OS overhead, idle time, everything
  - But doesn't completely reflect computer's performance if it focuses on better throughput
- CPU time – for CPU Performance
  - CPU execution time processing a task
    - Exclude I/O time, time spent for other tasks
  - Comprises **user CPU time** and **system CPU time**
    - Hard to differentiate
  - Different programs run with different CPU performance and system performance
    - User must know what's the desired performance metric

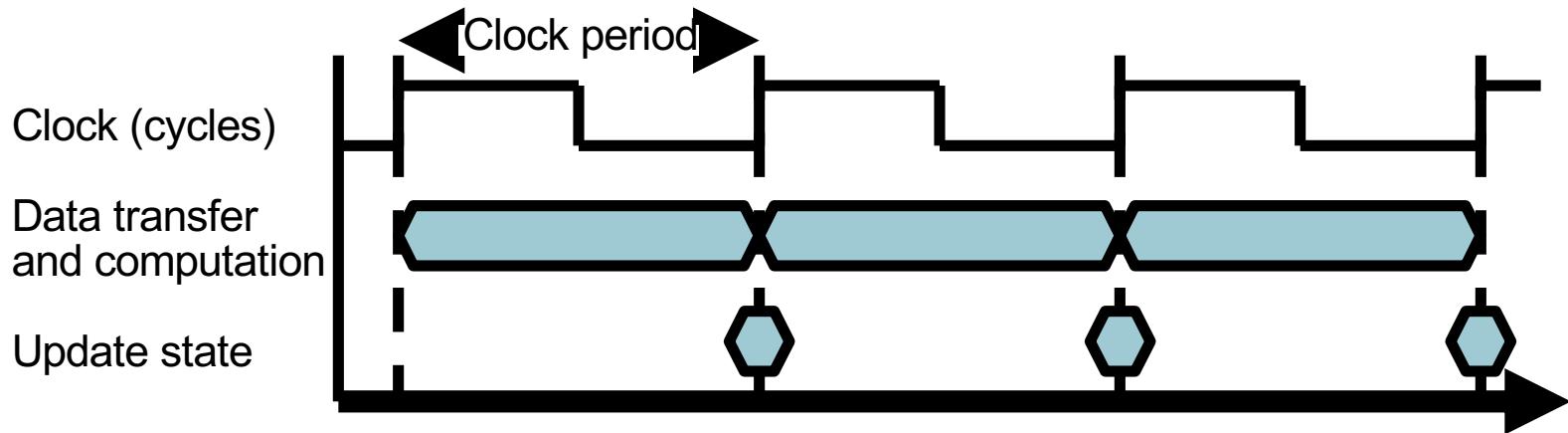
# Understanding Performance

Performance-related factors:

- Algorithm
  - Determines number of operations executed
- Programming language, compiler, architecture
  - Determine number of machine instructions executed per operation
- Processor and memory system
  - Determine how fast instructions are executed
- I/O system (including OS)
  - Determines how fast I/O operations are executed

# CPU Clocking

- Operation of digital hardware governed by a constant-rate clock



# CPU Time

$$\begin{aligned}\text{CPU Time} &= \text{CPU Clock Cycles per program} \times \text{Clock Cycle Time} \\ &= \frac{\text{CPU Clock Cycles}}{\text{Clock Rate}}\end{aligned}$$

- Performance improved by
    - Reducing number of clock cycles
    - Increasing clock rate
- Tradeoff we will face when designing the CPU



# CPU Time Example

- Computer A: 2GHz clock, 10s CPU time
- Designing Computer B
  - Aim for 6s CPU time
  - Can do faster clock, but causes  $1.2 \times$  clock cycles<sub>A</sub>
- How fast must Computer B clock be?

$$\text{Clock Rate}_B = \frac{\text{Clock Cycles}_B}{\text{CPU Time}_B} = \frac{1.2 \times \text{Clock Cycles}_A}{6s}$$

$$\begin{aligned}\text{Clock Cycles}_A &= \text{CPU Time}_A \times \text{Clock Rate}_A \\ &= 10s \times 2\text{GHz} = 20 \times 10^9\end{aligned}$$

$$\text{Clock Rate}_B = \frac{1.2 \times 20 \times 10^9}{6s} = \frac{24 \times 10^9}{6s} = 4\text{GHz}$$

# Instruction Count (IC) and Clock cycle Per Instruction (CPI)

- How instructions related to performance?

Clock Cycles = Instruction Count  $\times$  Clock Cycles per Instruction(CPI)

CPU Time = Instruction Count  $\times$  CPI  $\times$  Clock Cycle Time

$$= \frac{\text{Instruction Count} \times \text{CPI}}{\text{Clock Rate}}$$

Classic CPU  
performance equation

- Instruction Count for a program
  - Determined by program, Instruction Set Architecture (ISA), and compiler
- Cycles per instruction
  - Determined by CPU hardware
  - If different instructions have different CPI
    - Average CPI, affected by instruction mix

# CPI Example

- Computer A: Cycle Time = 250ps, CPI = 2.0
- Computer B: Cycle Time = 500ps, CPI = 1.2
- Same ISA
- Which is faster, and by how much?

$$\begin{aligned}\text{CPU Time}_A &= \text{Instruction Count} \times \text{CPI}_A \times \text{Cycle Time}_A \\ &= I \times 2.0 \times 250\text{ps} = I \times 500\text{ps}\end{aligned}$$

A is faster...

$$\begin{aligned}\text{CPU Time}_B &= \text{Instruction Count} \times \text{CPI}_B \times \text{Cycle Time}_B \\ &= I \times 1.2 \times 500\text{ps} = I \times 600\text{ps}\end{aligned}$$

$$\frac{\text{CPU Time}_B}{\text{CPU Time}_A} = \frac{I \times 600\text{ps}}{I \times 500\text{ps}} = 1.2$$

...by this much

# Performance Summary

## The BIG Picture

$$\text{CPU Time} = \frac{\text{Instructions}}{\text{Program}} \times \frac{\text{Clock cycles}}{\text{Instruction}} \times \frac{\text{Seconds}}{\text{Clock cycle}}$$

- Performance depends on
  - Algorithm: affects IC, possibly CPI
  - Programming language: affects IC, CPI
  - Compiler: affects IC, CPI
  - Instruction set architecture: affects IC, CPI, T<sub>c</sub>

# Concluding Remarks

- Classes and applications of computer
- Cost/performance is improving
  - Due to underlying technology development
- Hierarchical layers of abstraction
  - In both hardware and software
- Execution time: the best performance measure
- Power is a limiting factor
  - Use parallelism to improve performance