

Interactive Theorem Proving in HOL4

Course 06: Basic Tactics

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Acknowledgement of Country

We acknowledge and celebrate the First Australians on whose traditional lands we meet, and pay our respect to the elders past and present.

More information about Acknowledgement of Country can be found [here](#) and [here](#)



Goal-Directed Proofs and Tactics

Goal-Directed Proofs

- ▶ User starts by setting a proof goal (as statements of the targeting theorem);
- ▶ *Tactics* (like `CONJ_TAC`) to reduce the current goal to zero or more subgoals;
- ▶ *Tacticals* (like `THEN` and `THEN1`) to organize tactics in tree-like structure.
- ▶ When no subgoal is left, the forward proof is automatically constructed from bottom up, to generate the final theorem.

Category of Built-in Tactics

- ▶ Goal deconstruction;
- ▶ Quantifiers;
- ▶ Assumption management;
- ▶ Subgoal management;
- ▶ Rewriting;
- ▶ Case analysis (case splits);
- ▶ Induction;
- ▶ Renaming and abbreviation;
- ▶ Automatic Provers.



Goal Deconstruction: Conjunctive Goals (1)

Prove two conjunctives separately

CONJ_TAC : tactic

Before

Initial goal:

$P \wedge Q$: proof

After

Q

P (* first subgoal *)

2 subgoals
: proof



Goal Deconstruction: Conjunctive Goals (2)

Prove the first conjunctive, then use the first to prove the second

CONJ_ASM1_TAC : tactic

Before

Initial goal:

$P \wedge Q$: proof

After

Q

0. P

P

2 subgoals

Why it works?

$\vdash P \wedge Q \iff P \wedge (P \Rightarrow Q).$



Goal Deconstruction: Conjunctive Goals (3)

Use the 2nd conjunctive to prove the 1st one (then prove the 2nd)

```
CONJ_ASM2_TAC : tactic
```

Before

Initial goal:

$P \wedge Q$: proof

After

Q

P

0. Q

2 subgoals

Alternative Approach

Use `ONCE_REWRITE_TAC [CONJ_SYM]` to rewrite the goal to $Q \wedge P$.



Goal Deconstruction: Conjunctive Goals (4)

```
hurdUtils.STRONG_CONJ_TAC : tactic
```

Before

Initial goal:

$P \wedge Q$: proof

After

val it =

$P \implies Q$

P

2 subgoals

For the other direction

Use `ONCE_REWRITE_TAC [CONJ_SYM]` and then `STRONG_CONJ_TAC` to get $Q \implies P$.



Goal Deconstruction: Disjunctive Goals (1)

Prove the first disjunctive subgoal only

```
DISJ1_TAC : tactic
```

Before

Initial goal:

$P \vee Q$: proof

After

1 subgoal:

val it =

P

: proof

See also

DISJ2_TAC for proving the second disjunctive only.



Goal Deconstruction: Disjunctive Goals (2)

Use the first (negated) disjunctive to prove the second

```
hurdUtils.STRONG_DISJ_TAC : tactic
```

Before

Initial goal:

$\sim P \vee Q$: proof

After

1 subgoal:

> val it =

Q

0. P

: proof



Quantifiers: Eliminate Universal Quantifier

```
qx_gen_tac   : term quotation -> tactic
qx_genl_tac  : term quotation list -> tactic
```

Before

Initial goal:

`!x y. P x y: proof`

Tactic

`qx_genl_tac ['a', 'b']`

After

1 subgoal:

`val it =`

`P a b`

`: proof`

Alternatives

`X_GEN_TAC ``x`` >> X_GEN_TAC ``y```

“rpt GEN_TAC” or “NTAC 2 STRIP_TAC”.



Quantifiers: Introduce Universal Quantifier

```
qspec_tac      : term quotation * term quotation -> tactic  
qid_spec_tac   : term quotation list -> tactic
```

Before

Initial goal:

P a b: proof

Tactic

```
qspec_tac ('b', 'y')
```

After

1 subgoal:

val it =

!y. P a y

: proof

Alternatives

qid_spec_tac 'a' is equivalent to qspec_tac ('a', 'a').



Quantifiers: Eliminate Existential Quantifier

```
qexists_tac : term quotation -> tactic
```

Before

Initial goal:

```
?n. SUC n = 1: proof
```

Tactic

```
qexists_tac '0'
```

After

1 subgoal:

```
val it =
```

```
    SUC 0 = 1
```

```
    : proof
```



Quantifiers: Choose for Existential Quantifier

Choose a variable for existential quantifier and push to assumptions

```
Q.X_CHOOSE_TAC : term quotation -> thm_tactic  
STRIP_TAC : tactic
```

Before

Initial goal:

$(\exists x. P\ x) \implies Q$: proof

Tactic

```
DISCH_THEN (Q.X_CHOOSE_TAC 'y')
```

After

1 subgoal:
val it =

Q

0. $P\ y$

: proof



Assumption Management (1)

Adding a theorem as a new assumption

```
ASSUME_TAC : thm_tactic (= thm -> tactic)
```

Before

Initial goal:

$\sim P \implies \sim Q$: proof

Tactic

```
ASSUME_TAC (Q.SPECL ['P', 'Q']  
            CONTRAPOS_THM)
```

After

1 subgoal:

val it =

$\sim P \implies \sim Q$

0. $\sim P \implies \sim Q \iff Q \implies P$

: proof



Assumption Management (2)

Move proposition from goal to assumption

```
DISCH_TAC : tactic
```

```
STRIP_TAC : tactic
```

Before

Initial goal:

$P \implies Q$

Tactic

```
DISCH_TAC
```

After

```
val it =
```

Q

0. P

: proof



Assumption Management (3)

Moving last assumption to the goal

```
POP_ASSUM : thm_tactic -> tactic
MP_TAC      : thm_tactic
```

Before

```
      R
-----
0.  P
1.  Q

: proof
```

Tactic

```
POP_ASSUM MP_TAC
```

After

```
val it =

      Q ==> R
-----
0.  P

: proof
```



Assumption Management (4)

Moving matched assumption to the goal

```
Q.PAT_X_ASSUM : term quotation -> thm_tactic -> tactic
```

Before

```
      R
-----
0.  P
1.  Q

: proof
```

After

```
val it =

      P ==> R
-----
0.  Q

: proof
```

Tactic

```
Q.PAT_X_ASSUM 'P' MP_TAC
```



Subgoal Management (1)

Using subgoals

- ▶ Goal-directed proofs do not always revert the informal proof;
- ▶ Good formal proofs are in forwarding direction, aligned with the informal proofs;
- ▶ Subgoals are stage work of the proof stored into assumptions;
- ▶ Subgoals make proofs read easier;
- ▶ Subgoals can be either in forward or backward styles.



Subgoal Management (2)

Prove and place a theorem on the assumptions of the goal

```
op by : term quotation * tactic -> tactic
```

Before

Initial goal:

$\forall x. P\ x : \text{proof}$

Tactic

'Q' by cheat

After

val it =

$\forall x. P\ x$

0. Q

: proof

The subgoal inherits all assumptions of the current goal.



Subgoal Management (3)

Prove and place a theorem into the goal

```
hurdUtils.Know : term quotation -> tactic
```

Before

Initial goal:

```
!x. P x: proof
```

Tactic

```
Know 'Q :bool'
```

After

2 subgoals:

```
val it =
```

```
Q ==> !x. P x
```

```
Q
```

2 subgoals



Subgoal Management (4)

Replace the goal's conclusion with a sufficient alternative.

```
hurdUtils.Suff : term quotation -> tactic
```

Before

Initial goal:

$\text{!}x. P\ x : \text{proof}$

Tactic

`Suff 'Q :bool'`

After

Q

$Q \implies \text{!}x. P\ x$

2 subgoals



Rewriting Tactics (1)

Rewriting goal using theorems

```
REWRITE_TAC : (thm list -> tactic)
```

Before

```
val it =  
  Proof manager status: 1 proof.  
  1. Incomplete goalstack:  
      Initial goal:  
      SUC n = n + 1  
      : proofs  
> ADD1;  
val it = |- !m. SUC m = m + 1: thm
```

Tactic

```
REWRITE_TAC [ADD1]
```

After

```
OK..  
val it =  
  Initial goal proved.  
  |- SUC n = n + 1: proof
```



Rewriting Tactics (2)

Rewriting goal using theorems and assumptions

`ASM_REWRITE_TAC : (thm list -> tactic)`

Before

```
      n + 1 = 2
-----
0.  n = 1

: proof
```

Tactic

`ASM_REWRITE_TAC []`

After

```
1 subgoal:
val it =

      1 + 1 = 2
-----
0.  n = 1

: proof
```



Rewriting Tactics (3)

Rewriting goal and assumptions using simpset, theorems and *newer* assumptions

```
FULL_SIMP_TAC : simpset -> thm list -> tactic
```

Before

```
      m + 1 = 2
-----
0.   m = n
1.   n = 1

: proof
```

Tactic

```
FULL_SIMP_TAC std_ss []
```

After

```
1 subgoal:
val it =

      1 + 1 = 2
-----
0.   m = 1
1.   n = 1

: proof
```



Rewriting Tactics (4)

Rewriting goal and assumptions using simpset, theorems and *older* assumptions

```
REV_FULL_SIMP_TAC : simpset -> thm list -> tactic
```

Before

```
      m + 1 = 2
-----
0.   n = 1
1.   m = n

: proof
```

Tactic

```
FULL_SIMP_TAC std_ss []
```

After

```
1 subgoal:
val it =

      1 + 1 = 2
-----
0.   n = 1
1.   m = 1

: proof
```



Induction (1)

How induction works

1. There must be an induction theorem to apply, e.g.:

`[numTheory.INDUCTION]`

$\vdash P\ 0 \wedge (\forall n. P\ n \Rightarrow P\ (SUC\ n)) \Rightarrow \forall n. P\ n$

2. The tactic `HO_MATCH_MP_TAC` is used for applying the induction theorem.

Before

Initial goal:

`!n. n + n = 2 * n: proof`

Tactic

`HO_MATCH_MP_TAC numTheory.INDUCTION`

After

`val it =`

`0 + 0 = 2 * 0 /\`

`!n. n + n = 2 * n ==>`

`SUC n + SUC n = 2 * SUC n`

`: proof`



Induction (2)

Induction-related tactics

`Induct` : tactic
`Induct_on` : term quotation \rightarrow tactic

Before

Initial goal:

`!n. n + n = 2 * n: proof`

Tactics

`Induct`, or `Induct_on 'n'`

After

$\text{SUC } n + \text{SUC } n = 2 * \text{SUC } n$

0. $n + n = 2 * n$

$0 + 0 = 2 * 0$

2 subgoals
: proof

