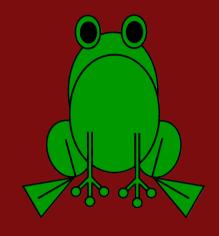


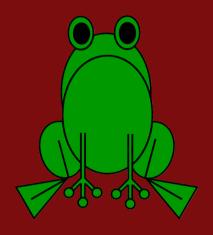
# Late Data Layout: Unifying Data Representation Transformations

Vlad Ureche Eugene Burmako Martin Odersky École polytechnique fédérale de Lausanne, Switzerland {first.last}@epfl.ch

> 23<sup>rd</sup> of October 2014 **OOPSLA '14** Portland, OR

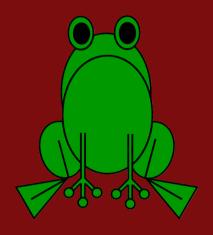


Unifying Data Representation Transformations



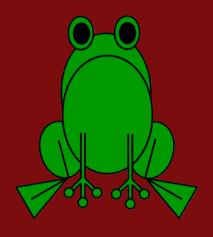
Unifying Data Representation Transformations

- compiler transformations
- separate compilation
- global scope



Unifying Data Representation Transformations

- unboxing, value classes
- how data is represented



Unifying Data Representation Transformations

- what is there to unify?
- why bother?



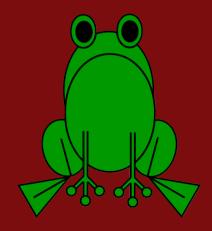
Motivation

Transformation

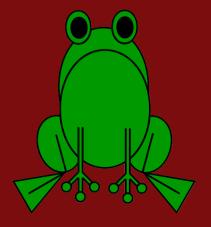
**Properties** 

Benchmarks

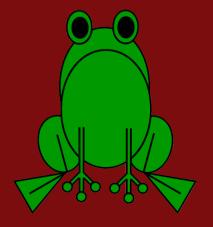
Conclusion



## Representation Transformations









- value
- no garbage collection
- locality



- value
- no garbage collection
- locality

## java.lang.Integer (\*\*)



- indirect access
- object allocation
  - and thus garbage collection
- no locality guarantees
- compatible with erased generics



- value
- no garbage collection
- locality

in Java, programmers are responsible for the choice of representation

## java.lang.Integer



- indirect access
- object allocation
  - and thus garbage collection
- no locality guarantees
- compatible with erased generics





- value
- no garbage collection
- locality

in Java, programmers are responsible for the choice of representation

### java.lang.Integer



- indirect access
- object allocation
  - and thus garbage collection
- no locality guarantees
- compatible with erased generics



What about Scala?



- value
- no garbage collection
- locality

## java.lang.Integer (\*\*)



- indirect access
- object allocation
  - and thus garbage collection
- no locality guarantees
- compatible with erased generics



#### int

- value
- no garbage collection
- locality

### java.lang.Integer



- indirect access
- object allocation
  - and thus garbage collection
- no locality guarantees
- compatible with erased generics

#### scala.Int



- value
- no garbage collection
- locality

### java.lang.Integer



- indirect access
- object allocation
  - and thus garbage collection
- no locality guarantees
- compatible with erased generics

#### scala.Int



- value
- no garbage collection
- locality





- indirect access
- object allocation
  - and thus garbage collection
- no locality guarantees
- compatible with erased generics

scalac {

#### scala.Int



- value
- no garbage collection
- locality





- indirect access
- object allocation
  - and thus garbage collection
- no locality guarantees
- compatible with erased generics



1 Choice of representation



#### scala.Int



- value
- no garbage collection
- locality





- indirect access
- object allocation
  - and thus garbage collection
- no locality guarantees
- compatible with erased generics



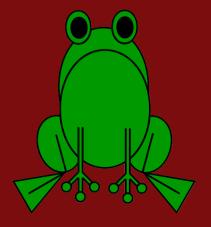


Choice of representationCoercions between representations





Value Classes



value class

#### value class



### struct (by-val)

- preferred encoding
- fields are inlined
- no heap allocations

#### value class



#### Struct (by-val)

- preferred encoding
- fields are inlined
- no heap allocations

### object (by-ref)



- fallback encoding
- compatible with
  - subtyping
  - erased generics

#### value class



#### Struct (by-val)

object (by-ref)



- preferred encoding
- fields are inlined
- no heap allocations

- fallback encoding
- compatible with
  - subtyping
  - erased generics





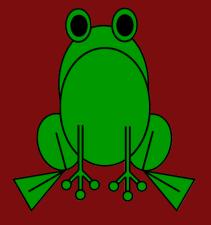
Choice of representationCoercions between representations





Value Classes

Miniboxing



## Miniboxing





#### Miniboxing: Improving the Speed to Code Size Tradeoff in Parametric Polymorphism Translations

Vlad Ureche Cristian Talau Martin Odersky

EPFL, Switzerland {first.last}@epfl.ch



#### **Abstract**

Parametric polymorphism enables code reuse and type safety. Underneath the uniform interface exposed to programmers, however, its low level implementation has to cope with inherently non-uniform data: value types of different sizes and semantics (bytes, integers, floating point numbers) and reference types (pointers to heap objects). On the Java Virtual Machine, parametric polymorphism is currently translated to bytecode using two competing approaches: homogeneous and heterogeneous. Homogeneous translation requires boxing, and thus introduces indirect access delays. Heterogeneous translation duplicates and adapts

#### 1. Introduction

Parametric polymorphism allows programmers to describe algorithms and data structures irrespective of the data they operate on. This enables code reuse and type safety. For the programmer, generic code, which uses parametric polymorphism, exposes a uniform and type safe interface that can be reused in different contexts, while offering the same behavior and guarantees. This increases productivity and improves code quality. Modern programming languages offer generic collections, such as linked lists, array buffers or maps as part of their standard libraries.







T (primitive)





#### T (primitive)



### long integer

- preferred encoding
- for all primitive types





#### T (primitive)



### long integer

- preferred encoding
- for all primitive types

### T (erased to Object)



- fallback encoding
- compatible with
  - virtual dispatch
  - subtyping
  - erased generics

## Miniboxing



#### T (primitive)



### long integer





- preferred encoding
- for all primitive types

- fallback encoding
- compatible with
  - virtual dispatch
  - subtyping
  - erased generics





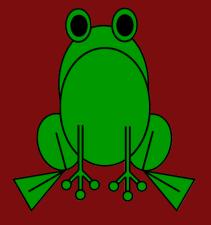
Choice of representation
Coercions between representations





Value Classes

Miniboxing

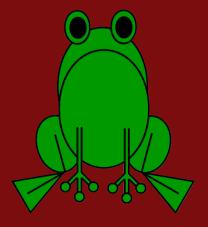




Value Classes

Miniboxing

motivated by erased generics



## Representation Transformations

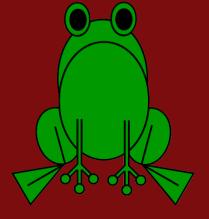
Unboxing Primitive Types

Value Classes

Miniboxing

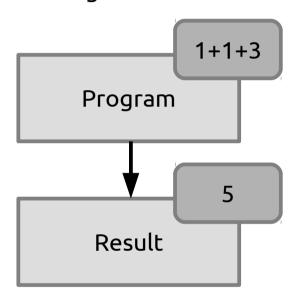
Staging (Multi-Stage Programming)

motivated by erased generics



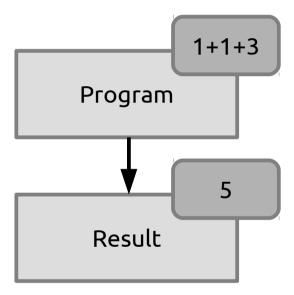
## Staging

#### 1-stage execution

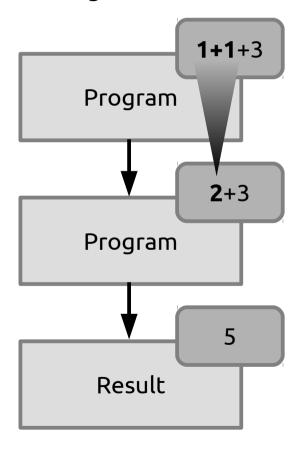


## Staging





#### 2-stage execution



#### value

#### value



#### direct value (5)

- is a computed value
- from an expression evaluated in the current stage

#### value



#### direct value (5)

- is a computed value
- from an expression evaluated in the current stage

#### lifted expression (2+3)



- executed in the next stage
- stores the expression that produces the value

#### value



#### direct value (5)





- is a computed value
- from an expression evaluated in the current stage

- executed in the next stage
- stores the expression that produces the value

scalac



Choice of representation – domain-specific



Coercions between representations



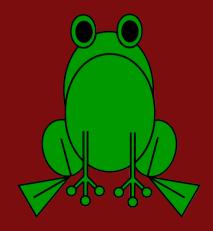
Motivation

Transformation

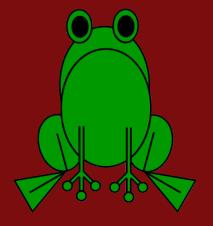
**Properties** 

Benchmarks

Conclusion

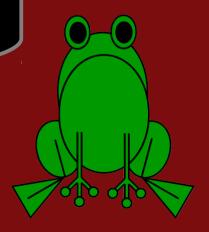


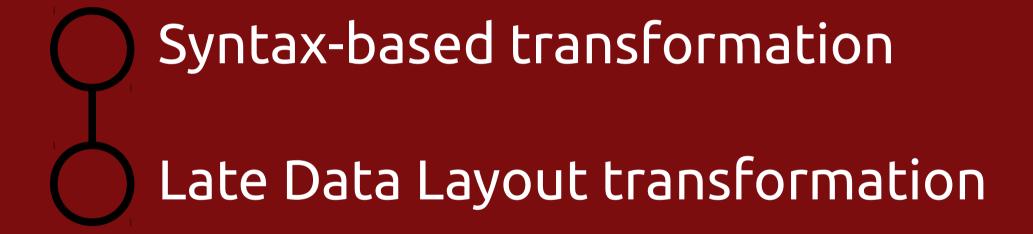
#### How to transform a program?

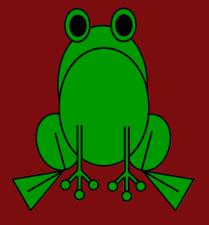


#### How to transform a program?

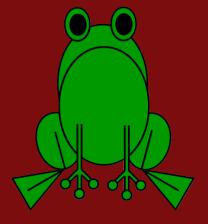
We'll use primitive unboxing as the running example, to keep things simple











- we need coercions between representations
- simple set of syntax-based rules
  - example

val x: Int = ...
val y: Int = x

```
<u>val x: Int = ...</u>
val y: Int = x
```

```
val x: Int = ...
val y: Int = x

val x: int = unbox(...)
val y: Int = box(x)
```

```
val x: Int = ...

val y: Int = x

val x: int = unbox(...)

val y: Int = box(x)
Coerce the definition right-hand side
```

```
val x: Int = ...
val y: Int = x

val x: int = unbox(...)
val y: Int = box(x)
Coerce all occurences of
```

the transformed value

```
val x: Int = ...
val y: Int = x

val x: int = unbox(...)
val y: Int = box(x)
```

```
val x: Int = ...
val y: Int = x
val x: int = unbox(...)
val y: Int = box(x)
val x: int = unbox(...)
val y: int = unbox(box(x))
```

```
val x: Int = ...
      val y: Int = x
      val x: int = unbox(...)
      val y: Int = box(x)
      val x: int = unbox(...)
      val y: int = unbox(box(x))
                                suboptimal
scala-ldl.org
```

# Peephole Optimization

val y: int = unbox(box(x))

peephole

val y: int = x

another example

```
def choice(t1: Int, t2: Int): Int =
  if (Random.nextBoolean())
   t1
  else
  t2
```

```
def choice(t1: Int, t2: Int): Int =
if (Random.nextBoolean())
t1
else
t2
```

```
def choice(t1: int, t2: Int): Int =
  if (Random.nextBoolean())
  box(t1)
  else
  t2
```

```
def choice(t1: int, t2: int): Int =
  if (Random.nextBoolean())
  box(t1)
  else
  box(t2)
```

```
def choice(t1: int, t2: int): Int =
  if (Random.nextBoolean())
  box(t1)
  else
  box(t2)
```

Anything missing?

```
Yes, unboxing the returned value
def choice(t1: int, t2: int): Int =
 if (Random.nextBoolean())
  box(t1)
 else
  box(t2)
              Anything missing?
```

```
def choice(t1: int, t2: int): Int =
  if (Random.nextBoolean())
  box(t1)
  else
  box(t2)
```

```
def choice(t1: int, t2: int): Int
if (Random.nextBoolean())
  box(t1)
  else
  box(t2)
```

```
def choice(t1: int, t2: int): int =
  if (Random.nextBoolean())
  box(t1)
  else
  box(t2)
```

```
def choice(t1: int, t2: int): int =
  unbox(
  if (Random.nextBoolean())
    box(t1)
  else
  box(t2)
)
```

```
def choice(t1: int, t2: int): int =
  unbox(
  if (Random.nextBoolean())
    box(t1)
  else
  box(t2)
)
```

```
def choice(t1: int, t2: int): int =
   unbox(
   if (Random.nextBoolean())
      box(t1)
   else
      box(t2)
)
```

```
def choice(t1: int, t2: int): int =
   unbox(
   if (Random.nextBoolean())
      box(t1)
   else
      box(t2)
)
```

```
def choice(t1: int, t2: int): int =
 unbox(
  if (Random.nextBoolean())
   <u>box</u>(t1)
  else
   box(t2)
                           new peephole rule
                                 sink outside coercions
                                  into the if branches
```

```
def choice(t1: int, t2: int): int =
  if (Random.nextBoolean())
    unbox(box(t1))
  else
    unbox(box(t2))
```

```
def choice(t1: int, t2: int): int =
  if (Random.nextBoolean())
    unbox(box(t1))
  else
  unbox(box(t2))
```

## Syntax-based

```
def choice(t1: int, t2: int): int =
  if (Random.nextBoolean())
   t1
  else
  t2
```

## Syntax-based

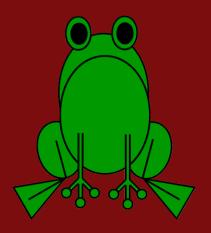
```
def choice(t1: int, t2: int): int =
  if (Random.nextBoolean())
   t1
  else
  t2
complicated
```

## Syntax-based

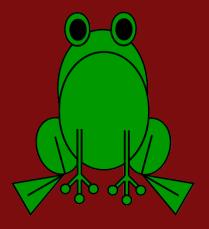
- peephole transformation does not scale
  - needs **multiple rewrite rules** for each node
  - needs stateful rewrite rules
  - leads to an explosion of rules x states

Details in the paper

### Coercions are fixed in the tree

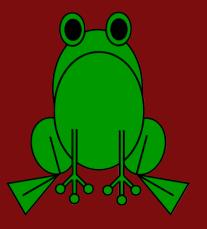


## Coercions are **fixed in the tree** and moving them around is difficult.



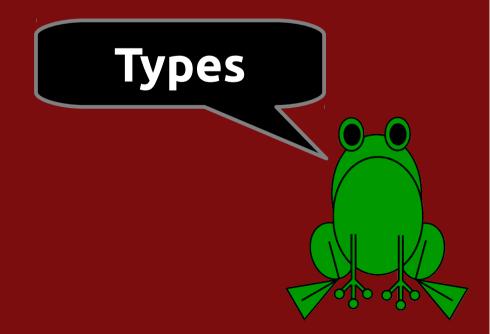
Coercions are **fixed in the tree** and moving them around is difficult.

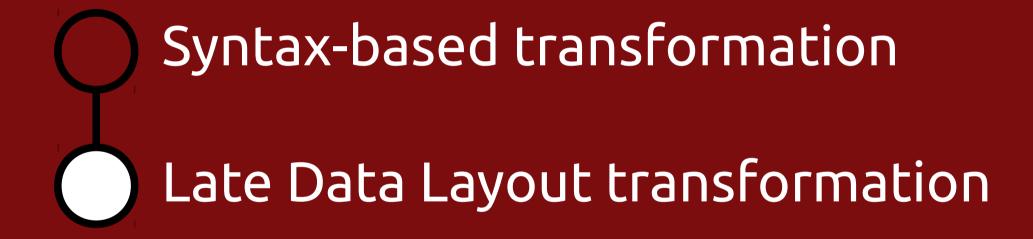
We need a more fluid abstraction.

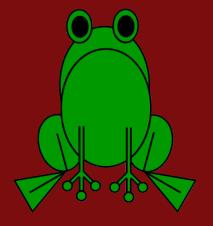


Coercions are **fixed in the tree** and moving them around is difficult.

We need a more fluid abstraction.





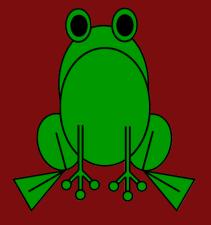




## Late Data Layout transformation

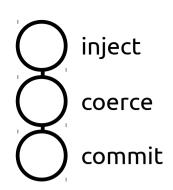
Phases Inject
Coerce

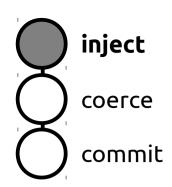
Commit

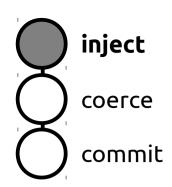


# LDL Transformation The Inject Phase

- propagates representation information
  - into the type system
    - based on annotated types
    - e.g. an @unboxed annotation added to integers







operations can be performed as well

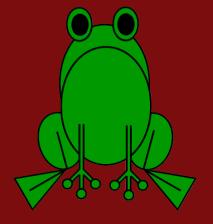
(e.g. miniboxing duplicates methods)



## Late Data Layout transformation

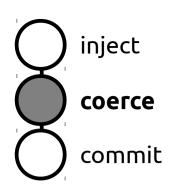
Phases Inject
Coerce

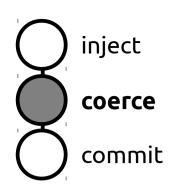
Commit

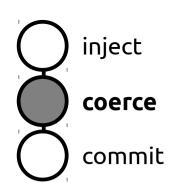


## LDL Transformation The Coerce Phase

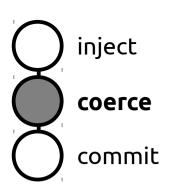
- introduces coercions
  - re-type-checks the tree
  - exposes representation mismatches
    - as annotation mismatches (Int vs @unboxed Int)
    - leading to coercions







the return type of choice is **@unboxed Int** 



the return type of choice is **@unboxed Int** 

def choice(t1: @unboxed Int,

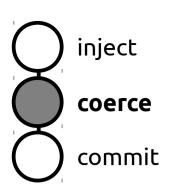
t2: @unboxed Int): @unboxed Int =

if (Random.nextBoolean())

t1

else

t2



the return type of choice is **@unboxed Int** 

def choice(t1: @unboxed Int,

t2: @unboxed Int): @unboxed Int =

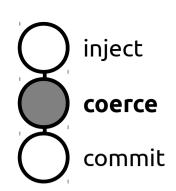
if (Random.nextBoolean())

t1

else

**t2** 

: @unboxed Int



the return type of choice is **@unboxed Int** 

def choice(t1: @unboxed Int,

t2: @unboxed Int): @unboxed Int =

if (Random.nextBoolean())

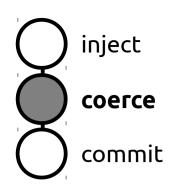
t1

else

**t2** 

: @unboxed Int

**expected type** (part of local type inference)



```
def choice(t1: @unboxed Int,
```

t2: @unboxed Int): @unboxed Int =

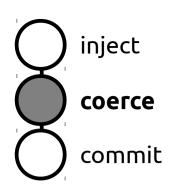
```
if (Random.nextBoolean())
```

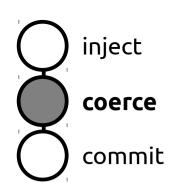
t1

else

t2

: @unboxed Int





```
def choice(t1: @unboxed Int,
t2: @unboxed Int): @unboxed Int =
```

if (Random.nextBoolean()): Boolean

t1

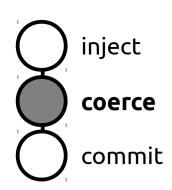
else

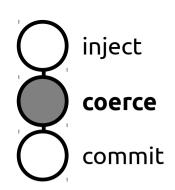
**t2** 

matches:

expected: Boolean

found: Boolean





def choice(t1: @unboxed Int,

t2: @unboxed Int): @unboxed Int =

if (Random.nextBoolean())

t1:@unboxed Int

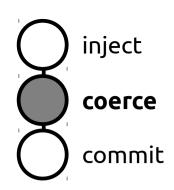
else

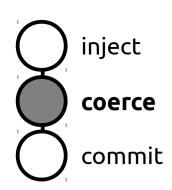
**t2** 

#### matches:

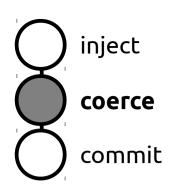
expected: @unboxed Int

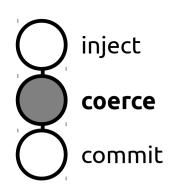
found: @unboxed Int

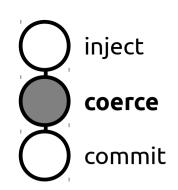




```
def choice(t1: @unboxed Int,
                  t2: @unboxed Int): @unboxed Int =
       if (Random.nextBoolean())
         t1
       else
        t2: @unboxed Int
             matches:
scala-ldl.d
```







No peephole transformation

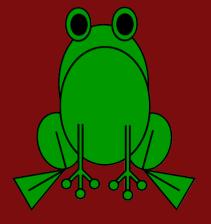


## Late Data Layout transformation

Phases Inject

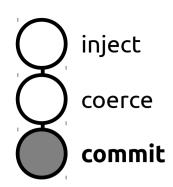
Coerce

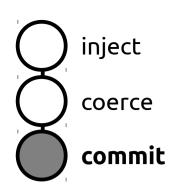
Commit

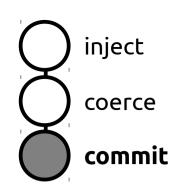


## LDL Transformation The Commit Phase

- converts annotations to representations
  - @unboxed Int → int
  - Int → java.lang.Integer
- coercion markers are also transformed
  - box(...) → new Integer(...)
  - unbox(...) → ....intValue





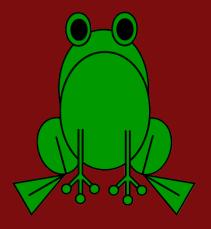


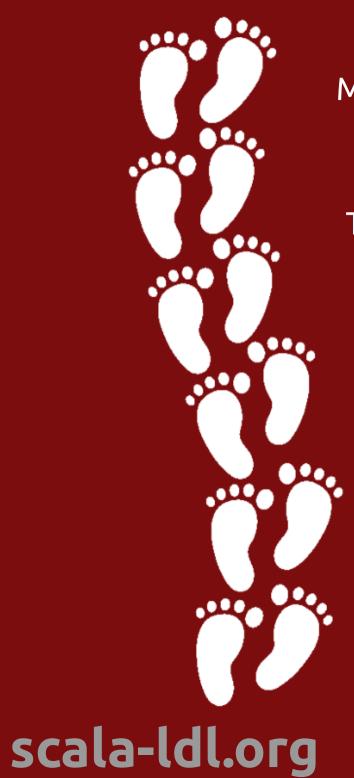


Phases Inject

Coerce

Commit





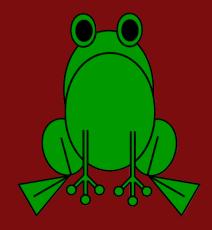
Motivation

Transformation

Properties

Benchmarks

Conclusion





## Late Data Layout transformation

Properties



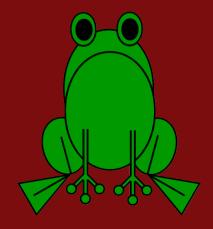
Selectivity

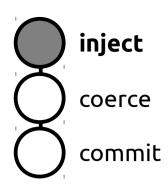


Consistency

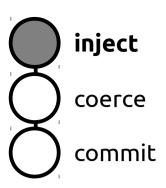


Optimality (not formally proven yet)

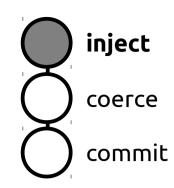


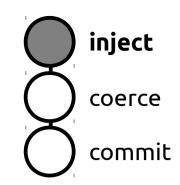


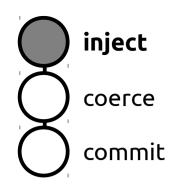
- annotated types
  - **selectively** pick the representation for each value

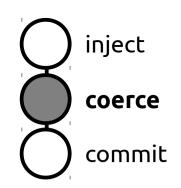


- annotated types
  - **selectively** pick the representation for each value
- selectivity is used for
  - bridge methods (some args boxed, others unboxed)
  - value classes (JVM: no multi-value returns)
  - staging (representation: domain-specific knowledge)
    - List[Int] vs List[@staged Int] vs @staged List[Int]









```
def choice(t1: Int,
```

t2: @unboxed Int): @unboxed Int =

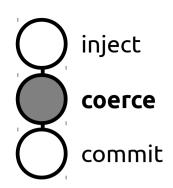
if (Random.nextBoolean())

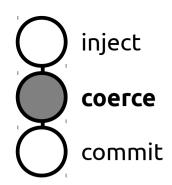
t1

else

t2

: @unboxed Int





```
def choice(t1: Int,
```

t2: @unboxed Int): @unboxed Int =

if (Random.nextBoolean())

t1: @unboxed Int

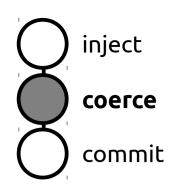
else

**t2** 

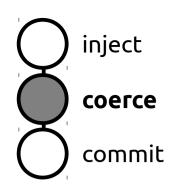
#### mismatch:

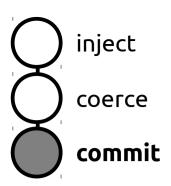
expected: @unboxed Int

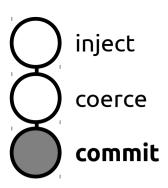
found: Int



coercion









## Late Data Layout transformation

**Properties** 



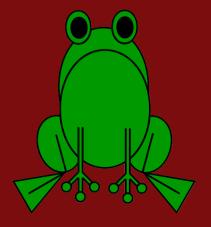
Selectivity



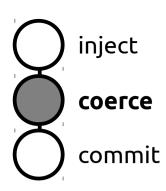
Consistency



Optimality (not formally proven yet)



## Consistency



- representations become part of types
- re-type-checking the program
  - proves type correctness
  - proves representation consistency



## Late Data Layout transformation

**Properties** 



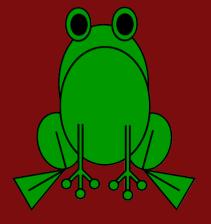
Selectivity

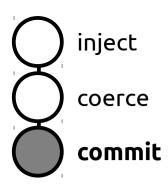


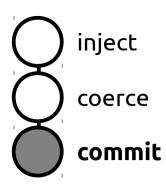
Consistency



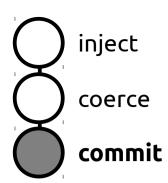
Optimality (not formally proven yet)



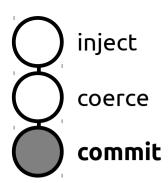




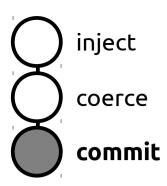
execution



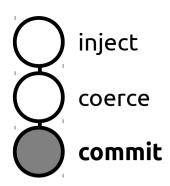
```
execution
```



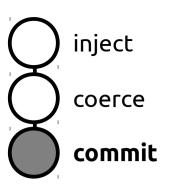
```
execution
```



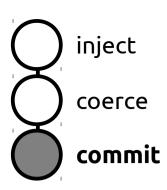
- on any execution trace through the program
  - the number of coercions executed is minimum
  - assuming the program terminates



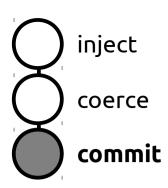
- on any execution trace through the program
  - the number of coercions executed is minimum
  - assuming the program terminates
- modulo
  - annotations introduced by the inject phase
    - unbox both parameters → no coercions at all



- on any execution trace through the program
  - the number of coercions executed is minimum
  - assuming the program terminates
- modulo
  - annotations introduced by the **inject phase** 
    - unbox both parameters → no coercions at all
  - post-transformations done by the commit phase
    - box(...) → new Integer(new Integer(...).intValue)

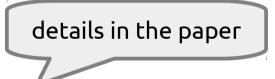


- peephole optimization
  - propagates coercions
- type system
  - propagates types



- peephole optimization
  - propagates coercions
- type system
  - propagates types

but types are fluid whereas coercions are not





## Late Data Layout transformation

**Properties** 



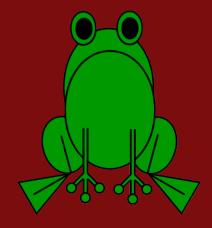
Selectivity



Consistency



Optimality (not formally proven yet)





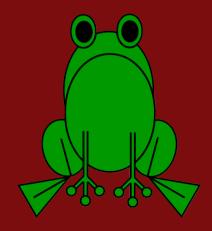
Motivation

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Properties

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Conclusion



## LDL is used in



- Scala compiler plugins
  - miniboxing (specialization)
  - value-class plugin
  - staging plugin

# Benchmarks ... in the paper



- implementation effort
  - Late Data Layout mechanism
    - developed as part of miniboxing
    - reused by the other compiler plugins
  - value class plugin → 2 developer-weeks
  - staging plugin → 1 developer-week

# Benchmarks ... in the paper



- performance
  - baseline vs transformed code
- numbers
  - up to 2x speedup when transforming value classes
  - up to 22x speedup when using miniboxing
  - up to 59x speedup when staging



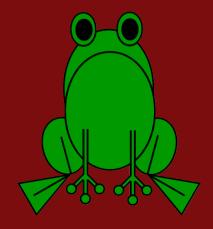
Motivation

**Transformation** 

**Properties** 

Benchmarks

Conclusion



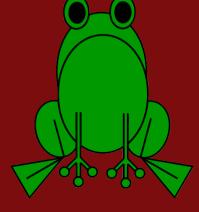
# Conclusion Insights

- use annotated types
  - to **selectively** mark values with the representation
- use expected type propagation
  - to provide **optimal** transformation
- use the type system
  - to provide representation consistency

## Credits and Thank you-s

- Cristian Talau developed the initial prototype, as a semester project
- Eugene Burmako the value class plugin based on the LDL transformation
- Aymeric Genet developing collection-like benchmarks for the miniboxing plugin
- Martin Odersky, for his patient guidance
- Eugene Burmako, for trusting the idea enough to develop the value-plugin based on the LDL transformation
- Iulian Dragos, for his work on specialization and many explanations
- · Miguel Garcia, for his original insights that spawned the miniboxing idea
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- Heather Miller for the eye-opening discussions we had
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- Denys Shabalin, Dmitry Petrashko for their patient reviews of the LDL paper
- Xiaoya Xiang and Philip Stutz for trusting miniboxing enough to try it out

Special thanks to the Scala Community for their support! (@StuHood, @vpatryshev and everyone else!)





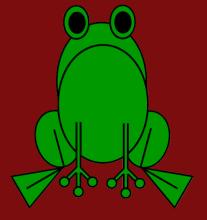
#### concept





герг. 2 ... герг. п





#### concept

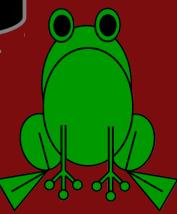




... repr. n



How would you use this?



#### concept





... repr. n



