

The Sound of Sorting Algorithm Cheat Sheet

Function selectionSort(A : Array of Element; n : \mathbb{N})

```

for  $i$  := 1 to  $n$  do
   $min$  :=  $i$ 
  for  $j$  :=  $i + 1$  to  $n$  do
    if  $A[j] < A[min]$  then
       $min$  :=  $j$ 
  endfor
  swap( $A[i]$ ,  $A[min]$ )
  invariant  $A[1] \leq \dots \leq A[i]$ 
endfor

```

Function insertionSort(A : Array of Element; n : \mathbb{N})

```

for  $i$  := 2 to  $n$  do
   $j$  :=  $i$ 
  while ( $j > 0$ ) & ( $A[j - 1] < A[j]$ )
    swap( $A[j - 1]$ ,  $A[j]$ )
     $j$  :=  $j - 1$ 
  endwhile
  invariant  $A[1] \leq \dots \leq A[i]$ 
endfor

```

Function mergeSort(A : Array of Element; lo, hi : \mathbb{N})

```

if  $hi - lo \leq 1$  then return
 $mid$  := ( $lo + hi$ ) / 2
mergeSort( $lo, mid$ ), mergeSort( $mid, hi$ )
 $B$  := allocate (Array of Element size  $hi - lo$ )
 $i$  :=  $lo$ ,  $j$  :=  $mid$ ,  $k$  := 1
while ( $i < mid$ ) & ( $j < hi$ )
  if  $A[i] < A[j]$   $B[k++]$  :=  $A[i++]$ 
  else  $B[k++]$  :=  $A[j++]$ 
endwhile
while  $i < mid$  do  $B[k++]$  :=  $A[i++]$ 
while  $j < hi$  do  $B[k++]$  :=  $A[j++]$ 
 $A[lo, \dots, hi - 1]$  :=  $B[1, \dots, (hi - lo)]$ 
dispose ( $B$ )

```

Procedure bubbleSort(A : Array $[1 \dots n]$ of Element)

```

for  $i$  := 1 to  $n$  do
  for  $j$  := 1 to  $n - i$  do
    if  $A[j] > A[j + 1]$  then
      swap( $A[j]$ ,  $A[j + 1]$ )

```

Procedure heapSort(A : Array $[1 \dots n]$ of Element)

```

buildHeap( $A$ )
while  $n > 1$ 
  swap( $A[1]$ ,  $A[n]$ )
   $n$  :=  $n - 1$ 
  siftDown( $A, 1$ )

```

Procedure buildHeap(A : Array $[1 \dots n]$ of Element)

```

for  $i$  :=  $\lfloor n/2 \rfloor$  downto 1 do
  siftDown( $i$ )

```

Procedure siftDown(A : Array $[1 \dots n]$ of Element; i : \mathbb{N})

```

if  $2i > n$  then return
 $k$  :=  $2i$ 
if ( $2i + 1 \leq n$ ) & ( $A[2i] \leq A[2i + 1]$ ) then
   $k$  :=  $k + 1$ 
if  $A[i] < A[k]$  then
  swap( $A[i]$ ,  $A[k]$ )
  siftDown( $A, k$ )

```

Procedure cocktailShakerSort(A : Array $[1 \dots n]$ of Element)

```

 $lo$  := 1,  $hi$  :=  $n$ ,  $mov$  :=  $lo$ 
while  $lo < hi$  do
  for  $i$  :=  $hi$  downto  $lo + 1$  do
    if  $A[i - 1] > A[i]$  then
      swap( $A[i - 1]$ ,  $A[i]$ ),  $mov$  :=  $i$ 
  endfor
   $lo$  :=  $mov$ 
  for  $i$  :=  $lo$  to  $hi - 1$  do
    if  $A[i] > A[i + 1]$  then
      swap( $A[i]$ ,  $A[i + 1]$ ),  $mov$  :=  $i$ 
  endfor
   $hi$  :=  $mov$ 

```

Procedure gnomeSort(A : Array $[1 \dots n]$ of Element)

```

 $i$  := 2
while  $i \leq n$  do
  if  $A[i] \geq A[i - 1]$  then
     $i$  ++
  else
    swap( $A[i]$ ,  $A[i - 1]$ )
    if  $i > 2$  then  $i$  --
  endwhile

```

Procedure quickSort(A : **Array of** Element; $\ell, r : \mathbb{N}$)

```

if  $\ell \geq r$  then return
 $q := \text{pickPivotPos}(A, \ell, r)$ 
 $m := \text{partition}(A, \ell, r, q)$ 
quickSort( $A, \ell, m - 1$ ), quickSort( $A, m + 1, r$ )

```

Function partition(A : **Array of** Element; $\ell, r : \mathbb{N}$; $q : \mathbb{N}$)

```

 $p := A[q]$  // pivot element
swap( $A[q], A[r]$ ) // swap to the end
 $i := \ell$ 
invariant

|          |       |     |     |
|----------|-------|-----|-----|
| $\ell$   | $i$   | $j$ | $r$ |
| $\leq p$ | $> p$ | $?$ | $p$ |

for  $j := \ell$  to  $r - 1$  do
    if  $A[j] \leq p$  then
        swap( $A[i], A[j]$ ),  $i++$  // move smaller to the front
assert

|          |       |     |
|----------|-------|-----|
| $\ell$   | $i$   | $r$ |
| $\leq p$ | $> p$ | $p$ |


swap( $A[i], A[r]$ ) // move pivot into the middle
assert

|          |     |       |
|----------|-----|-------|
| $\ell$   | $i$ | $r$   |
| $\leq p$ | $p$ | $> p$ |

return  $i$ 

```

Procedure quickSortTernary(A : **Array of** Element; $\ell, r : \mathbb{N}$)

```

if  $\ell \geq r$  then return
 $q := \text{pickPivotPos}(A, \ell, r)$ 
 $(m, m') := \text{partitionTernary}(A, \ell, r, q)$ 
quickSortTernary( $A, \ell, m - 1$ ), quickSortTernary( $A, m' + 1, r$ )

```

Function partitionTernary(A : **Array of** Element; $\ell, r : \mathbb{N}$; $q : \mathbb{N}$)

```

 $p := A[q]$  // pivot element
 $i := \ell$ ,  $j := \ell$ ,  $k := r$ 
invariant

|        |       |     |       |     |
|--------|-------|-----|-------|-----|
| $\ell$ | $i$   | $j$ | $k$   | $r$ |
| $< p$  | $> p$ | $?$ | $= p$ |     |

while ( $j \leq k$ ) // three-way comparison
    if  $A[j] = p$  then swap( $A[j], A[k]$ ),  $k--$ ;
    else if  $A[j] < p$  then swap( $A[j], A[i]$ ),  $i++$ ,  $j++$ ;
    else  $j++$ ;
assert

|        |       |       |     |
|--------|-------|-------|-----|
| $\ell$ | $i$   | $k$   | $r$ |
| $< p$  | $> p$ | $= p$ |     |

 $i' := i + r - k + 1$ 
swap( $A[i \dots i'], A[k + 1 \dots r]$ ) // move = p area to the middle
assert

|        |       |       |     |
|--------|-------|-------|-----|
| $\ell$ | $i$   | $i'$  | $r$ |
| $< p$  | $= p$ | $> p$ |     |

return ( $i, i'$ )

```

Procedure LSDRadixSort(A : **Array** [$1 \dots n$] **of** Element)

```

 $K := 4$  // number of buckets per round
 $D := \lceil \log_K(\max\{A[i] + 1 \mid i = 1, \dots, n\}) \rceil$  // calculate number of rounds
 $B := \text{allocate}(\text{Array of Element size } n)$  // temporary array  $B$ 
for  $d := 0$  to  $D - 1$  do // sort by the  $d$ -th  $K$ -digit.
    redefine  $\text{key}(x) := (x \text{ div } K^d) \bmod K$ 
    KSortCopy( $A, B, n$ ), swap( $A, B$ ) // sort from  $A$  to  $B$ , and swap back
    invariant  $A$  ist nach den  $K$ -Ziffern  $d..0$  sortiert.
dispose ( $B$ )

```

Procedure KSortCopy(A, B : **Array** [$1 \dots n$] **of** Element; $K : \mathbb{N}$)

```

 $c = \langle 0, \dots, 0 \rangle$  : Array [ $0 \dots K - 1$ ] of  $\mathbb{N}$ 
for  $i := 1$  to  $n$  do  $c[\text{key}(A[i])]++$  // count occurrences
 $\text{sum} := 1$ 
for  $k := 0$  to  $K - 1$  do // exclusive prefix sum
     $\text{next} := \text{sum} + c[k]$ ,  $c[k] := \text{sum}$ ,  $\text{sum} := \text{next}$ 
for  $i := 1$  to  $n$  do
     $B[c[\text{key}(A[i])]++] := A[i]$  // move element  $A[i]$  into bucket of  $B$ 

```

Procedure MSDRadixSort(A : **Array** [$1 \dots n$] **of** Element)

```

 $K := 4$  // number of buckets per round
 $D := \lceil \log_K(\max\{A[i] + 1 \mid i = 1, \dots, n\}) \rceil$  // count number of round
MSDRadixSortRec( $A, D - 1, K$ )

```

Procedure MSDRadixSortRec(A : **Array** [$1 \dots n$] **of** Element; $d, K : \mathbb{N}$)

```

 $c = \langle 0, \dots, 0 \rangle$  : Array [ $0 \dots K - 1$ ] of  $\mathbb{N}$  // KSort with in-place permuting
redefine  $\text{key}(x) := (x \text{ div } K^d) \bmod K$ 
for  $i := 1$  to  $n$  do  $c[\text{key}(A[i])]++$  // count occurrences
 $b = \langle 0, \dots, 0 \rangle$  : Array [ $0 \dots K$ ] of  $\mathbb{N}$ 
 $\text{sum} := 1$ 
for  $k := 0$  to  $K$  do // inclusive prefix sum into  $b$ 
     $\text{sum} := \text{sum} + c[k]$ ,  $b[k] := \text{sum}$ 
assert  $b[K] = n$ 
for  $i := 1$  to  $n$  do
    while ( $j := -- b[\text{key}(A[i])] > i$ ) // walk on cycles until  $i$ 
        swap( $A[i], A[j]$ ) // move  $A[i]$  into right bucket
         $i := i + c[\text{key}(A[i])]$  // bucket of  $A[i]$  is finished
invariant  $A$  ist nach den  $K$ -Ziffern  $d..(D - 1)$  sortiert
if  $d = 0$  return // done?
for  $k := 0$  to  $K - 1$  do // recursion into each of the  $K$  buckets if
    if  $c[k] > 1$  // it contains two or more elements
        MSDRadixSortRec( $A[b[k] \dots b[k + 1] - 1], d - 1, K$ )
dispose ( $b$ ), dispose ( $c$ )

```