An Empirical Study of Meta- and Hyper-Heuristic Search for Multi-Objective Release Planning

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Abstract—A variety of meta-heuristic search algorithms have been introduced for optimising software release planning. However, there has been no comprehensive empirical study of different search algorithms across multiple different real world datasets. In this paper we present an empirical study of global, local and hybrid meta- and hyper-heuristic search based algorithms on 10 real world datasets. We find that the hyper-heuristics are particularly effective. For example, the hyper-heuristic genetic algorithm significantly outperformed the other six approaches (and with high effect size) for solution quality 85% of the time, and was also faster than all others 70% of the time. Furthermore, correlation analysis reveals that it scales well as the number of requirements increases.

I. Introduction

Release planning is the problem of determining the sets of requirements that should be included in each release of a software system. In order to plan the software release cycle, a number of different conflicting objectives need to be considered. For example, the estimated cost of implementing a requirement has to be balanced against the perceived value to the customer of that requirement. There may be multiple stakeholders, and their different interpretations of cost and value may lead to complex solution spaces for release planners.

In order to help decision-makers navigate these complex solution spaces, meta-heuristic search has been widely studied as a candidate solution technique [1], [2], [3], [4]. This work has placed release planning within the general area of Search Based Software Engineering (SBSE) [5]. Table I summarises the literature on search based release planning, listing the meta-heuristic algorithms proposed and the datasets on which they have been evaluated.

The different cost and value objectives for each stakeholder are typically measured along incomparable dimensions. To avoid the familiar problem of 'comparing apples with oranges', much of the previous work on multi-objective release planning has used Pareto optimal search, whose results are Pareto fronts. Each element on this front is a candidate solution to the release planning problem. All solutions on the Pareto front are non-denominated: no other solution on the front is better according to all objectives. The Pareto front thus represents a set of 'best compromises' between the objectives that can be found by the search based algorithm.

The overview in Table I reveals that previous work has evaluated meta-heuristic algorithms on very few real world datasets. Much of the previous work presents results for only a single real world dataset. In the absence of real world datasets, many authors have relied upon synthetically generated data.

While studies on synthetic data can answer experimental research questions [6], they cannot address the critical concern of the practitioners: "how well can I expect these techniques to behave on real world data?"

As a result, the state-of-the-art is currently poorly understood: though a variety of different algorithms has been proposed, there has been no empirical study across multiple different algorithms and multiple different datasets. We wish to address this issue by providing a thorough empirical study of optimised release planning. We believe that this may help to understand the different strengths and weaknesses of algorithms for release planning and their performance on real world datasets. We hope that our study will also provide results against which future work can compare¹.

We report the results of an empirical study using 10 real world datasets. We investigate multi-objective release planning with respect to these datasets, which we optimise using Hill Climbing (HC), Genetic Algorithms (GA) and Simulated Annealing (SA).

As a sanity check, recommended for SBSE work [7], [8], we also report results for purely random search. Random search provides a baseline against which to benchmark more 'intelligent' search techniques. Our study also includes hyperheuristic versions of HC, GA and SA. Hyper-heuristics [9] are a more recent trend in search methodologies, not previously been used in any SBSE research. The findings we report here indicate that they are promising for release planning problems.

Overall, our study thus involves 7 different algorithms. We assess solutions found using these algorithms according to 4 different measures of solution quality, over each of the 10 real world datasets. We include standard, widely used, measures of multi-objective solution quality: convergence, hypervolume and two different assessments of each algorithm's contribution to the Pareto front. We also measure diversity and speed. For algorithms that produce good quality solutions, these are important additional algorithmic properties for decision-makers, because they need quick answers that enable them to base their decisions on the full diversity of candidate solutions.

The primary contributions of the paper are as follows:

1. **Comprehensive study**: We provide a comprehensive study of the performance of global, local and hybrid meta-heuristic algorithms for release planning problems on 10 real world datasets. The results facilitate detailed algorithm comparison

¹Note to referee: if the paper is accepted then we will make all our implementations and data publicly available on the web to support replication and further study.

and reveal that dataset specifics can lead to important a differences in study findings.

- 2. **Hyper-heuristic search**: We introduce and evaluate hyper-heuristic algorithms for release planning. We present evidence that they provide good solution quality, diversity, and speed.
- 3. **Scalability assessment**: We investigate the scalability of meta- and hyper-heuristic algorithms on real world datasets for the first time. The results provide evidence that hyper-heuristics have attractive scalability and that random search is surprisingly *un*scalable for release planning.

The rest of the paper is organised as follows: Section II sets our work in the context of related work. Section III introduces our three hyper-heuristic release planning algorithms. Section IV explains our experimental methodology. Section V presents the results and discusses the findings and threats to validity. Section VI concludes the paper.

TABLE I
PREVIOUS META-HEURISTIC ALGORITHM STUDIES

Paper Authors (Year)	Algorithm(s)	Dataset				
Bagnall et al. [1] (2001)	Greedy, HC, SA	Synthetic				
Feather & Menzies [10] (2002),	SA	NASA				
Feather et al. [11] (2004),						
Feather et al. [12] (2006)						
Ruhe & Greer [2] (2003),	GA	Synthetic				
Ruhe & Ngo-The[13] (2004),						
Amandeep et al.[14] (2004),						
Ruhe [15] (2010)						
Zhang et al. [4] (2007)	NSGA2, Pareto GA	Synthetic				
Durillo et al. [16] (2009),	NSGA2, MOCell,	Synthetic &				
Durillo et al. [17] (2011)	PAES	Motorola				
Colares et al. [18] (2009)	NSGA2	Synthetic				
Finkelstein et al. [19] (2008),	NSGA2, Two-	Motorola				
Zhang et al. [20] (2011)	Archive					
Zhang & Harman [21] (2010),	NSGA2, Archive-	Synthetic &				
Zhang et al. [22] (2013)	based NSGA2	RALIC				
Sagrado & Aguila [23] (2009),	ACO, SA, GA	Synthetic				
Sagrado et al.[24] (2010)						
Jinag et al. [25] (2010),	Backbone	Eclipse,				
Xuan et al. [26] (2012)	Multilevel	Mozilla,				
		Gnome				
Zhang et al. [27] (2010)	NSGA2	Ericsson				
Jinag et al. [28] (2010)	ACO, GRASP, SA	Synthetic				
Kumari et al. [29] (2012)	QEMEA	Synthetic				
Souza et al. [30] (2011),	ACO	Synthetic				
Ferrreira & Souza [31] (2012)						
Cai et al. [32] (2012)	NSGA2, SPEA2	Synthetic				
Brasil et al. [33] (2012)	NSGA2, MOCell	Synthetic				
Cai & Wei [34] (2013)	MOEAs	Synthetic				

II. THE CONTEXT OF OUR STUDY

Bagnall *et al.* [1] first suggested the term *Next Release Problem* and described various meta-heuristic optimisation algorithms for solving it. Feather and Menzies [10] were the first to use a real world dataset, but this dataset is no longer publicly available. Ruhe *et al.* [2], [13], [14], [35] introduced the software release planning process together with exact optimisation algorithms [36], [37], [15], [38] and meta-heuristics, such as genetic algorithms. Van den Akker *et al.* [39], [40], [3], [41] also studied exact approaches to single objective constrained requirements selection problems.

Zhang *et al.* [4] introduced the Multi-Objective Next Release Problem formulation as a Pareto optimal problem with a set of objectives. However, Feather *et al.* [12], [11] had previously used Simulated Annealing to construct a form of Pareto front for visualisation of choices. Also, at the same

time, Saliu and Ruhe [42] introduced a multi-objective searchbased optimisation to balance the tension between user-level and system-level requirements.

Subsequently, Finkelstein *et al.* [43] used multi-objective formulations to characterise different notions of fairness in the satisfaction of multiple stakeholders with different views on the priorities for requirement choices.

The multi-objective formulation subsumes previous single objective formulations: any single objective formulation that has a single objective and no constraints is clearly a special case of a multi-objective formulation for n objectives where n=1. Furthermore, a constrained single objective formulation, in which there is a single optimisation objective and a set of constraints to be satisfied, can be transformed into a multi-objective formulation in which the constraints become additional objectives to be met.

The technical details of the various approaches used for constrained single objective formulations and their multiobjective counterparts are, of course, different. However, in this paper we want to study the most general setting in which requirements optimisation choices might be cast. Therefore, we adopt the multi-objective paradigm.

The Next Release Problem (NRP) considers only a single release, while Release Planning (RP) considers a series of releases. NRP is thus a special case of the RP; we shall focus on RP as the more general case.

A. Representation of Release Planning

Requirements management in RP is formulated as a combinatorial optimisation problem. Approaches to the NRP represent the solution as a bitset of requirements for the next release. The RP formulation, being more general, is typically represented as a sequence of integers that denote release sequence numbers.

Let us assume a set of stakeholders, $C = \{c_1, \ldots, c_j, \ldots, c_m\}$ whose requirements are to be considered in the development of the next release of the software. Each stakeholder may have a degree of importance for the company that can be reflected by a weight factor. The set of relative weights associated with each stakeholder c_j $(1 \le j \le m)$ is denoted by a weight set: $Weight = \{w_1, \ldots, w_j, \ldots, w_m\}$ where $w_j \in [0,1]$ and $\sum_{j=1}^m w_j = 1$.

The set of requirements is denoted by: $\Re = \{r_1, \dots, r_i, \dots, r_n\}$. The study is based on the assumption that the requirements represented in the paper are at a similar level of abstraction. They are not be stated in too much detail nor on too high a level of abstraction. If the level differs among the requirements, there might be difficulty in selecting the correct subset of requirements for RP.

There are S (S>1) possible releases considered in the model, which allow us to look further rather than just the next release. In each release s $(1 \le s \le S)$, the decision vector $\overrightarrow{x_i} = \{x_1, \dots, x_n\} \in \{0, \dots, S\}$ determines the requirements that are to be satisfied in release s. In this vector, $x_i > 0$ if requirement r_i is delivered in release s and s if requirement s is not selected in the first s releases. In this paper we use

the three-release formulation, with weights 5, 3 and 1 for the first, second and third releases respectively.

B. Objectives and Fitness

In any approach to SBSE, it is necessary to choose the objectives, which define the fitness function(s) used to guide the search [8], [44].

The choice of the objectives to be considered in any multiobjective NRP/RP instance is governed by the specifics of the dataset and scenario for which the search based optimiser seeks requirement sets. In the 10 datasets used in this paper the objectives are to be maximised our Revenue, Sales Value and Importance, while those to be minimised are Cost and effect on *Urgency*. Figure 1 provides descriptive statistics that characterise the datasets, their sources and the requirement optimisation objectives pertinent to each dataset. In all but one case, the problem is a bi-objective problem in which there is a single objective to be maximised (such as revenue) and a single objective to be minimised (such as cost). The exception is the Ericsson dataset. It has two 'importance' objectives to be maximised: one for the present and one for the future.

The fitness functions are defined as follows:

Maximize
$$f_1(\overrightarrow{x}) = \sum_{i=1}^n Revenue_i \cdot x_i$$
 (1)

Minimize $f_2(\overrightarrow{x}) = \sum_{i=1}^n Cost_i \cdot x_i$ (2)

Minimize
$$f_2(\overrightarrow{x}) = \sum_{i=1}^n Cost_i \cdot x_i$$
 (2)

The objectives of Revenue and Cost in the fitness functions can also be repleced by other objectives. The specific choice of objectives to be maximised or minimised are parameters to the search based optimisation algorithm used to search for requirement sets. The algorithms use these objectives as fitness functions that guide the search. We can compare different algorithms across different datasets, because the algorithm itself does not change, merely the fitness functions used to guide the search.

III. HYPER-HEURISTIC SEARCH

Hyper-heuristics [9] are search methodologies, more recently introduced to the optimisation literature than the metaheuristics that have been enthusiastically adopted by the SBSE community [44], [45]. Hyper-heuristics differ from metaheuristics because they search the space of heuristics rather than the space of solutions [46].

Hyper-heuristics encapsulate problem specific information using a pool of low-level heuristics (known as search operators). In the wider literature on engineering and design in general (rather than software engineering in particular), hyperheuristics have been widely applied. However, even in this wider context, the design and engineering problems attacked using hyper-heuristics have tended to be single objective problems, with only a very few previous attempts to apply hyper-heuristics to multi-objective problems [47], [48].

Though hyper-heuristics have been suggested as a future direction for SBSE [6], there has been no previous investigation of the potential of hyper-heuristics for any SBSE problem,

TABLE II THE DESCRIPTION OF 10 SEARCH OPERATORS

IHEI	DESCRIPTION OF TO SEARCH OPERATORS
Operator	Description
Random	With uniform probability, select a requirement and change its release number to another release version (uniformly selected)
Swap	Swap the release numbers of two requirements
Delete Add	With uniform probability, exclude a requirement from the
Delett_: Iuu	current release, and add it to another release (also selected with uniform probability)
Delete Add	With uniform probability, select a requirement, r , and
Best	an objective o , replacing r with the best performing
	requirement according to objective o
Delete_Worst_	With uniform probability, select a requirement, r , and
Add	an objective o , replacing r with the worst performing
	requirement according to objective o
Delete_Worst_	With uniform probability, select an objective o, and
Add_Best	a release number, n , replacing the worst performing
	requirement according to objective o at release n with
	the best performing requirement (according to o)
Delay_Ahead	With uniform probability, select two requirements r_1
	(for a release other than the first) and r_2 (for a release
	other than the last). Move the release date of r_1 to a
	later release number than its current position, and move
	the release number of r_2 forward to an earlier number
	than its current position (both sampled with uniform
	probability). That is, r_1 is 'delayed' and r_2 is 'advanced'
Delay_Ahead_	With uniform probability, select a requirement, r , and an
Best	objective o . Delay r and advance the best requirement
	according to objective o
Delay_Worst_	With uniform probability, select a requirement, r , and an
Ahead	objective o. Delay the release of the worst requirement
D. I. W.	(according to objective o) and advance the release of r
Delay_Worst_	With uniform probability, select an objective o. Delay the
Ahead_Best	release of the worst requirement and advance the release
	of the best ('worst' and 'best' according to objective o)

so their performance on release planning has, hitherto, been unknown. This paper is thus the first paper in the requirements engineering literature to explore the use of hyper-heuristics and one of the few papers in the optimisation literature to use multi-objective hyper-heuristics.

We implemented 10 search operators for hyper-heuristic release planning. These are explained in Figure II. 10 search operators are used to generate the neighbourhoods of solution for Hyper Hill Climbing and Hyper Simulated Annealing. We also deploy these 10 operators to mutate the offsprings of solutions in Hyper NSGA-II. The first two (Random and Swap) represent standard mutation operators (i.e. they perform a small change on the solution, by swapping or changing solution components), while the remaining 8 follow the so-called 'ruin-recreate' (also known as the 'destructionconstruction') principle [49]. Ruin-recreate operators partly decompose (ruin) the solution and subsequently recreate it, incorporating problem-specific reconstruction heuristics to rebuild the solutions from their decomposed fragments.

A. Adaptive Operator Selection

An adaptive operator selection scheme consists of two components: the 'credit assignment' mechanism and the 'selection' mechanism [50]. Credit assignment involves the attribution of credit (or reward) to the hyper-heuristic's operators, determined by their performance during the search process.

Our hyper-heuristic release planners use the scheme proposed by Fialho et al. [51], known as 'extreme value credit assignment', which is based on the principle that infrequent, yet large, improvements in the objective score are likely to be

more effective than frequent, small improvements. Therefore, it rewards operators that have had a recent large positive impact on the objective score, whilst consistent operators that only yield small improvements receive less credit, and therefore have lower chances of selection.

The credit assignment mechanism needs to be coupled with a selection mechanism that uses the accumulated credits to select the operator to apply in the current iteration. We used an operator selection rule called 'Probability Matching' [52] to select the operators according their probabilites, which corresponds to the well-known roulette wheel selection used by meta-heuristic SBSE work [44].

IV. EXPERIMENTAL SET UP

A. Algorithms

More than 15 different meta-heuristic algorithms have been used in SBSE research [44]. Many of these have also been used in research on the NRP/RP (as in Table I). Indeed, all of these 15 (and many more meta-heuristics [53]) could be used, in principle, since the formulation of the problem is sufficiently generic that any search based technique could be applied.

We chose to investigate the performance of six search techniques: three meta-heuristics, Hill Climbing, Simulated Annealing and the Nondominated Sorting Genetic algorithm (NSGA2), which we denote HC, SA and GA respectively, together with hyper-heuristic versions of each of the three meta-heuristics, which we denote HHC, HSA and HGA. The motivation for this choice derives from the way in which computational search algorithms can be classified as either 'local' or 'global'. Local search techniques, such as HC, tend to be fast, but they can become stuck in a local optimum, thereby producing sub-optimal solutions. By contrast, global search techniques, such as genetic algorithms, may be computationally more expensive, but they incorporate mechanisms to avoid local optima confinement.

Many techniques embody elements of both local and global search, with the local search facilitating 'exploitation' in the search landscape, while the global search facilitates 'exploration' of the landscape [54]. One widely studied algorithm that does this is SA, which augments the basic HC approach with a 'cooling' coefficient that mimics cooling in annealing processes [55]. This process can enable the algorithm to escape local optima. In the SBSE literature, SA is the most widely used compromise between global and local search [44].

Our choice of the three meta-heuristic algorithms reflects our desire to sample from the set of possible algorithm choices, three that, in some sense, 'cover' the spectrum of algorithmic behaviours from local to global search. As Table I shows, all three of these meta-heuristic search techniques have been proposed and studied for release planning problems. We also wish to study the effect of hyper-heuristics as well as meta-heuristics, motivating our choice of the three hyper-heuristic variants of the three meta-heuristics we selected for study.

In the hyper-heuristic algorithms (HGA, HSA and HHC), the mutation operator of the meta-heuristic version (GA, SA and HC respectively) is replaced by the adaptive operator selection mechanism outlined in Section III-A. For SA, HC, HSA and HHC, a random seeding is used for generating the initial population. For the GA and HGA, we set the population size to 100 and the number of generations to 50. To ensure a fair comparison, SA, HSA, HC and HHC and the random search were all given the same budget of fitness evaluations as GA and HGA.

B. Datasets

We used 10 datasets, of which 7 are drawn from real world requirements selection problems in a variety of different organisations. Of these, 5 have been used in separate previous studies in the literature and 2 (StoneGate and MS Word) are newly introduced for the first time here. The other 3 datasets contain bug fixes requested for Eclipse, Mozilla and Gnome. They might be regarded as 'pseudo real world'; they are taken from real world applications but it is debatable whether they truly denote 'requirements'.

We include these three in the study since they have previously been used to act as a surrogates for real world datasets. Their use in previous work was motivated by the need to overcome the difficulty of finding sufficiently many real world datasets on which to evaluate [26]. A summary of each dataset studied in this paper can be found in Figure 1.

Previous studies have included at most two real world requirements datasets (or all three of the bug fixing pseudoreal world datasets), often augmenting these with synthetic data to compensate for the lack of real world data. Our study is therefore the most comprehensive study of metaheuristics for release planning so far reported in the literature. Our use of these 10 datasets is sufficient to allow us to ask an important research question that has, hitherto, eluded the research community: 'how well do the algorithms scale with respect to the size of the real world requirements problem to which they are applied?'.

C. Performance Metrics

We use 4 quality metrics to compare the performance of each of the 6 search-based optimisation algorithms (and random search). In most multi-objective optimisation problems the globally optimal Pareto front is unobtainable. Release planning is no exception to this. In such situations it is customary to construct a 'reference' front. The reference front is defined to be the largest nondominated subset of the union of solutions from all algorithms studied. As such, the reference front represents the best current approximation available to the true location of the globally optimal Pareto front. Three of the quality metrics we use (Contribution, Unique Contribution and Convergence) are computed in terms of each algorithm's distance from or contribution to this reference front:

Contribution (denoted 'Contrib' in our results tables) for algorithm A is the number of solutions produced by A that lie on the reference front. This is the simplest (and most intuitive) quality metric. It assesses how many of the best solutions found overall are found by algorithm A.

Name and Source Number of		Object	ives	Summary Description of Dataset and Software System				
of Dataset	R	SH	Max	Min	Summary Description of Dataset and Software System			
Baan [41]	100	17	Revenue	Cost	ERP product developed by 600 engineers spread over four countries			
StoneGate	100	91	Sales Value Impact		Industrial software security release planning project (confidential source)			
Motorola [20]	35	4	Revenue Cost		UK service provider requirements for range of handheld communication device			
RalicP [56], [22]	143	77	Revenue	Cost	Library and ID Card System in current use at University College London (UCI			
RalicR [56], [22]	143	79	Revenue	Cost	Library and ID Card System in current use at UCL (a variant of RalicP)			
Ericsson [27]	124	14	Importance	Cost	Requirements for a software testing tool for now and into the future			
MS Word	50	4	Revenue	Urgency	Text editing system for use with Microsoft Word			
Eclipse [26]	3502	536	Importance	Cost	The Eclipse environment with bug fix requests treated as requirements			
Mozilla [26]	4060	768	Importance	Cost	The Mozilla system with bug fix requests treated as requirements			
Gnome [26]	2690	445	Importance	Cost	The Gnome desktop system with bug fix requests treated as requirements			

Fig. 1. The 10 datasets used and their numbers of Requirements (R), stakeholders (SH) and objectives to be Maximised (Max) and Minimised (Min). Those datasets with accompanying citations are taken from previous studies; those without citations are used in this paper for the first time.

Unique Contribution (denoted 'UContrib') for algorithm A is the number of solutions produced by A that lie on the reference front and are not produced by any algorithm under study except A. This is a variant of the 'Contribution' metric that takes account of the fact that an algorithm may contribute relatively few of the best solutions found, but may still be valuable if it contributes a set of unique best solutions that no other algorithm finds.

Convergence (denoted 'Conv') for algorithm A is the Euclidean distance between the Pareto front produced by A and the reference front.

Hypervolume (denoted 'HVol') is the volume covered by the solutions in the objective space. HVolis the union of hypercubes of solutions on the Pareto front [57]. By using a volume HVolis less susceptible to bias than Unique Contribution when the numbers of points on the two compared fronts differ significantly.

Quality of solutions is clearly important, but diversity is also an important secondary criterion for algorithms that exhibit acceptable solution quality. We measured the diversity using a standard metric introduced by Deb [58]:

Diversity measures the extent of distribution in the obtained solutions and spread achieved between approximated solutions and the reference front [58].

Finally, in order to assess the compensation or effort required to produce the quality and diversity of solutions observed, we measure the computational effort:

Speed is measured in terms of the wall clock time required to produce the solutions reported, averaged over 30 executions. All experiments were carried out on a desktop computer with a 6 core 3.2GHz Intel CPU and 8GB memory.

In order to facilitate a more easy comparison of the six overall metrics used in our study, we normalise all of them to lie between 0.0 and 1.0 and convert all of them to 'maximising metrics' (such that higher values denote superior performance). For example, 'speed' (to give it a name that captures it 'maximising form') is computed as 1-T, where T is the normalised wall clock time. Thus, in all tables of data presented in this paper (including the correlation analyses) the reader can safely assume 'higher means better'. We normalise a value x, drawn from a set of observed values, ranging from x_{min} to x_{max} , using the standard normalising equation: $\frac{x-x_{min}}{x_{max}-x_{min}}$.

Our algorithms are executed 30 times each to cater for their stochastic natures, so the normalised metric values reported are averaged (using mean and median) over these 30 runs.

Averaging means that there is often no value reported in our results that happens to be exactly 1.0 or exactly 0.0, despite normalisation using maxima and minima.

D. Statistical Testing

We need to take account of the stochastic nature of each algorithm due to their partial reliance on randomisation. This is a well-known phenomenon for which it is widely advised [7], [8] that inferential statistical testing should be used as an appropriate way to compare algorithm performance. The pseudo random number sequence used by the algorithms is the cause of uncertainty. We are therefore sampling over the population of pseudo random number sequences [8].

We use inferential statistical testing techniques to draw inferences about the population of all possible executions of the algorithm on a particular instance, based on a sample of these executions. In our experiments we set our sample size to 30. That is, each of the 7 algorithms is executed 30 times on each of the 10 datasets.

We had no knowledge of the distribution of the population from which we sample executions so we use nonparametric statistical techniques, thereby increasing the robustness of our statistical inferences [7], [59], [8]. However, many widely used nonparametric statistical techniques are not as robust as the researcher might hope.

For example, though the widely used Mann-Whitney [60] (and closely-related Wilcoxon [61]) test and the Kruskal-Wallis test [62], make fewer assumptions than parametric tests they do, nevertheless assume that that *variance is consistent* across all populations [63]. Unfortunately, we can make no such assumption about our data in these experiments.

Therefore, we use Cliff's method [64] for assessing statistical significance. Cliff's method is not only nonparametric, but it is also specifically designed for ordinal data. Our research questions are ordinal and we have no reason to believe that our measurement scales are anything but ordinal [65]. Furthermore, unlike other popular nonparametric inferential statistical techniques such as the Kruskal-Wallis and Wilcoxon-Mann-Whitney tests, Cliff's method makes no assumptions about the variance of the data, thereby making it more robust.

We use the Vargha-Delaney \hat{A}_{12} metric for effect size (as recommended by Arcuri and Briand [7]). Vargha-Delaney \hat{A}_{12} also makes few assumptions and is highly intuitive: $\hat{A}_{12}(A,B)$ is simply the probability that algorithm A will outperform algorithm B in a head-to-head comparison.

We have set our α level to the widely used 'standard' value of 0.05. Each comparison of a p value with the chosen α

level is essentially a claim about probability; the probability of committing Type I error (the error of incorrectly rejecting the Null Hypothesis). However, if the experimenter conducts a series of tests, then the chances of committing a Type I error increase, potentially quite dramatically, unless some adjustment (or 'correction') is made.

One popular (but not necessarily ideal) adjustment is the Bonferroni correction, which was first used to control for multiple statistical inferences by Dunn [66]. Unfortunately, this is been shown to be highly conservative; while it avoids Type I errors, it does so at the risk of introducing Type II errors (the error of incorrectly accepting the Null Hypothesis), thereby reducing statistical power.

Fortunately, more recent techniques have been developed that retain the property of the Bonferroni correction (avoiding Type I errors), while simultaneously reducing its tendency to increase Type II errors. In our work we use just such a technique, the Hochberg's method [67] for controlling multiple hypothesis testing. This method ranks the statistical tests applied, adjusting the α level for each successive test. It is a less conservative adjustment in the Bonferroni correction, while retaining its ability to avoid Type I errors [68].

As well as investigating the quality of solutions produced by each algorithm, we also want to investigate the correlation between the size of the problem instance and the behaviour of each algorithm. Though we believe there may be correlations of interest, we have no reason to believe that they will be linear. Furthermore, since our data is measured on an ordinal scale, the use of the Pearson correlation [69], [70], [71] is inappropriate; we need to choose an ordinal correlation method. In order to ensure robustness of our conclusions we chose to use both Kendall's τ [72] and Spearman rank correlation [71], [73], both of which are nonparametric, rank-based assessments of correlation.

E. Research questions

This section explains and motivates the four research questions we ask in our study. When comparing different algorithms for release planning problems, a natural question to ask is the quality of the solutions produced, according to the standard measures of multi-objective solution quality. Our first research question therefore investigates solution quality:

RQ1: **Quality**: According to each of the 4 quality measures, and on each of the 10 datasets, which algorithm performs best? To answer this question we use the Cliff's inferential statistical comparison, as explained in Section IV-D to determine which algorithms significantly outperform others and the \hat{A}_{12} measure of effect size in each case.

Quality of solutions is important, but from the release planner's point of view a wide diversity of candidate solutions may also be important. In the most extreme case, a degenerate Pareto front (containing only a single solution) may have maximum quality but will have no diversity and will thus offer the release planner no choice at all. We therefore also investigate the diversity of solutions produced by each algorithm: **RQ2**: **Diversity**: What is the diversity of the solutions produced by each algorithm? We used Cliff's method to report on statistically significant differences in Diversity and \hat{A}_{12} to assess the effect size of any such differences observed.

Naturally, the computational effort required to produce these solutions is also important. An algorithm that produces slightly lower quality solutions, but which does so almost instantaneously will have different applications to one that produces better quality solutions, but takes several minutes to produce them. The former can be used in a situation where the release planner wants to repeatedly investigate 'what if' questions, rebalancing estimates of cost and value in real time. In this scenario, speed trumps quality, provided quality is sufficient to be acceptable for 'what if' analysis. The latter will be more useful in scenarios where requirements optimisation provides decision support for an important overall choice about release planning. In this situation, quality trumps speed, provided a solution can be found in reasonable time (which might be hours or even days).

RQ3: **Speed**: how fast can the algorithms produce the solutions? An algorithm that produces good solutions with acceptable diversity in reasonable time for small problems may scale less well to larger problems. In RP problems, scalability is not merely a question of the computational effort required for a larger problem; it is to be expected that computational effort will be directly proportional to problem size.

However, perhaps more importantly, there may also be some degradation solution quality and/or diversity as the size of the problem scales up. We therefore investigate scalability from the point of view of all of the metrics we collected in our answers to the foregoing three research questions.

RQ4: **Scalability**: What is the scalability of each of the algorithms with regard to solution quality, solution diversity and speed? In order to answer this question we report the rank of correlation between the size of the problem (measured in terms of the number of requirements in the dataset) and each of the metrics for quality, diversity and speed.

Since each algorithm is executed 30 times to facilitate statistical comparisons, we report correlations between the number of requirements and each of the mean, median and best case for quality, diversity and speed. As explained in Section IV-D, we use Kendall's τ and Spearman rank correlation to assess the degree of correlation between the quality, diversity and speed metrics and the size of the problem (measured as the number of requirements).

V. RESULTS AND ANALYSIS

This section presents the results of the empirical study of meta- and hyper-heuristic approaches for multi-objective RP. **RQ1: Quality**: Table III presents the mean and median values of the metrics for quality, diversity and speed for each of the 7 algorithms on each of the 10 datasets. Table IV presents the results of the inferential statistical tests. Since we need to compare 7 different algorithms with each other, this yields 21 pairwise comparisons (and thus 21 columns of data). Each column contains the Vargha-Delaney \hat{A}_{12} effect size measure

where the result is significant (at the $0.05~\alpha$ level), and is left blank otherwise. The column heading indicates which algorithm is being compared with which others, in groups of 6, 5, 4, 3, 2, and 1 pairwise comparisons (6+5+4+3+2+1 = 21 pairwise comparisons in total).

For example, in the pair of columns headed by the title $\frac{\text{HHC}}{\text{HC} \mid \text{R}}$, the HHC algorithm is compared against each of the HC and R algorithms. The value 1.00 in the first row under the first of these two columns indicates the following: HHC outperforms HC significantly for its contribution to the Baan dataset's reference front (because the entry is not blank) and the probability of this observation being made is 1.00. That is, HHC always beats HC for its contribution to the Baan dataset reference front in our sample of 30 runs.

The fifth row of data in this same column contains the effect size measure 0.06, indicating a significant result. However, this time the probability of HHC beating HC is 0.06, so the HC algorithm significantly outperforms the HHC algorithm on the metric assessed (Diversity for the Baan dataset).

From these two entries in the table of results we can see that, for the Baan dataset, HHC contributes far more strongly to the reference front than HC, but HC is far more diverse. As can be seen, the other three quality metrics for the comparison of HHC and HC on the Baan dataset also strongly favour HHC. Based on these observations we would prefer HHC to HC, since diversity is only interesting if the algorithm's quality is strong; a highly diverse set of sub-optimal solutions is easy to achieve and is of little value to the release planner.

The forgoing discussion indicates the density of information contained in Table IV. Space does not permit a further individual discussion of each of the 21 pairwise comparisons, but some general observations do emerge: From Table III, we can see that the random search (R) and also HC and SA make few contributions to the reference front:

Random search contributes in two cases, while the other two algorithms only manage a contribution in a single case: the Ericsson dataset. Even when these three search strategies do make a contribution to the reference front they contribute only tiny proportion of solutions (no more than 2%).

Looking at the results for the three meta-heuristic algorithms (GA, SA and HC), we see that GA performs best overall for quality on smaller datasets, while SA performs noticeably better on the three larger datasets (Eclipse, Mozilla and Gnome). This highlights the risk of drawing conclusions based on too narrow a selection of real world datasets.

The results for the three hyper-heuristic algorithms are less equivocal; they outperform their meta-heuristic counterparts according to all 4 quality measures. That is, HSA and HHC each significantly outperform both SA and HC on all 10 datasets with high effect size in every case, while HGA significantly outperforms GA on 9 out of the 10 datasets with high effect size. HGA is beaten by its meta-heuristic counterpart only on the Ericsson datasset.

HGA clearly offers the best performance over all datasets, algorithms and quality metrics: It significantly outperforms the other algorithms in 85% (205/240) of the pairwise algo-

rithm quality comparisons. However, though HGA performs strongest in terms of solution quality, it would be a mistake to conclude that is the *only* algorithm that should be used. HSA outperforms the HGA for 'contribution' in 4 of the 10 datasets and, perhaps more importantly, for '*unique* contribution' in three cases. Even HHC significantly outperforms HGA in terms of 'contribution' in one case.

We also observe further evidence that suggests that results from one dataset may not generalise to others. The most extreme example of this is the Ericsson dataset, for which all of the algorithms behaved very differently (when compared to each other) than they did for the other datasets.

RQ2: Diversity: As might be expected, random search performs very well in terms of diversity. Table IV shows that it outperforms almost every other algorithm for almost every dataset, often significantly and with a large effect size.

However, we know from the answer to RQ1 that random search only contributes to the reference front for 2 of the 10 datasets, and even then it only contributes at most 1% of the unique solutions. We therefore conclude that the diversity exhibited by the random search is largely suboptimal; any solutions it offers (diverse or otherwise) are likely to be dominated by solutions found by one or more of the others.

Of the three hyper-heuristic algorithms (which were competitive for the quality metrics), HGA exhibits the best diversity. It significantly outperforms HSA in 9, and HHC in 8, of the 10 datasets. As with the quality metrics, we observe that the Ericsson dataset also produces very different behaviour in terms of Diversity. The NSGA2 algorithm (on which both GA and HGA algorithms are based) was designed to promote diversity and so we might expect that it should perform best, both its meta- and hyper-heuristic versions.

Even the meta-heuristic version (GA) outperforms the hyper-heuristic versions of simulated annealing (HSA) and hill climbing (HHC) with respect to Diversity for 9 of the 10 datasets. However, there is no evidence that it outperforms its own hyper-heuristic version (HGA) with respect to Diversity. That is, GA significantly outperforms HGA for one dataset (Ericsson), while HGA significantly outperforms GA on 2 datasets (Mozilla and Gnome). In all other cases neither significantly outperforms the other.

RQ3: Speed: One might expect that a random search would be fast, since it is such a simple algorithm. However, we find (quite surprisingly) that the speed of random search is worse than all other algorithms studied for the larger datasets. We studied these results further and found that the explanation lies in the cost of invalid solutions: As the problem scale increases, a randomly constructed solution to the release planning problem is increasingly likely to be invalid. For example, it is increasingly likely to contain gaps in the release plan. The low cost of random search is nullified by the cost of such repair, more so as the problem size increases. By contrast, meta-heuristic and hyper-heuristic algorithms spend most of their time adapting existing release plans, making them more scalable than random search, despite being more sophisticated. Interestingly, HGA is fastest overall: it significantly outper-

		HGA		GA	Н	SA	S	A	Н	НС	Н	IC]	R
Data Sets	Metrics	Mean	Mean	Median	Mean	Median	Mean	Median	Mean	Median	Mean	Median	Mean	Median
Baan	Contrib UContrib Conv	0.83 0.8 0.88 0.9 0.98 0.9	6 0.02 0 0.03 8 0.79	0.00 0.00 0.79	0.20 0.20 0.78	0.20 0.18 0.79	0.00 0.00 0.26	0.00 0.00 0.26	0.14 0.13 0.71	0.13 0.12 0.71	0.00 0.00 0.19	0.00 0.00 0.19	0.00 0.00 0.19	0.00 0.00 0.20
Div	HVol Diversity Speed	0.98 0.9 0.54 0.5 0.94 0.9	4 0.49	0.75 0.48 0.96	0.69 0.17 0.37	0.69 0.18 0.36	0.26 0.73 0.78	0.27 0.75 0.78	0.62 0.14 0.15	0.62 0.13 0.16	0.20 0.57 0.67	0.21 0.61 0.70	0.15 0.88 0.98	0.15 0.88 0.98
Stone- Gate	Contrib UContrib Conv HVol Diversity	0.83 0.8 0.86 0.8 0.96 0.9 0.96 0.9 0.55 0.5	5 0.19 6 0.67 7 0.60 3 0.48	0.15 0.15 0.68 0.58	0.47 0.45 0.71 0.70 0.15	0.46 0.44 0.70 0.71 0.15	0.00 0.00 0.16 0.23 0.65	0.00 0.00 0.19 0.24 0.66	0.37 0.36 0.68 0.71 0.14	0.35 0.32 0.68 0.72 0.13	0.00 0.00 0.14 0.20 0.55	0.00 0.00 0.14 0.21 0.59	0.00 0.00 0.21 0.29 0.87	0.00 0.00 0.21 0.28 0.88
Motorola	Speed Contrib UContrib Conv HVol Diversity Speed	0.99 0.9 0.81 0.8 0.85 0.8 0.98 0.9 0.97 0.9 0.26 0.2 0.90 0.9	0 0.46 6 0.49 8 0.94 7 0.91 3 0.29	0.99 0.47 0.48 0.94 0.92 0.29 0.88	0.39 0.39 0.32 0.80 0.75 0.28	0.39 0.39 0.31 0.81 0.74 0.29 0.28	0.86 0.00 0.00 0.37 0.29 0.53 0.75	0.89 0.00 0.00 0.37 0.29 0.53 0.78	0.16 0.38 0.31 0.79 0.74 0.31 0.16	0.16 0.39 0.30 0.80 0.74 0.30 0.16	0.85 0.00 0.00 0.39 0.34 0.45 0.71	0.85 0.00 0.00 0.40 0.33 0.48 0.70	0.99 0.00 0.00 0.20 0.08 0.70 0.97	0.99 0.00 0.00 0.21 0.07 0.69 0.97
RalicP	Contrib UContrib Conv HVol Diversity Speed	0.75 0.7 0.81 0.7 0.96 0.9 0.96 0.9 0.56 0.5 0.95 0.9	8 0.26 6 0.85 6 0.75 7 0.55	0.21 0.25 0.84 0.76 0.55	0.16 0.16 0.60 0.53 0.50	0.15 0.15 0.61 0.52 0.48 0.27	0.00 0.00 0.24 0.27 0.70 0.79	0.00 0.00 0.23 0.27 0.77 0.82	0.13 0.12 0.55 0.51 0.48 0.17	0.12 0.11 0.55 0.50 0.47 0.17	0.00 0.00 0.19 0.20 0.66 0.76	0.00 0.00 0.18 0.19 0.65 0.79	0.00 0.00 0.19 0.13 0.88 0.96	0.00 0.00 0.18 0.13 0.89
RalicR	Contrib UContrib Conv HVol Diversity Speed	0.84 0.8 0.78 0.7 0.97 0.9 0.97 0.9 0.39 0.4 0.96 0.9	5 0.10 7 0.80 7 0.71 2 0.39	0.07 0.07 0.81 0.73 0.40 0.98	0.12 0.10 0.74 0.68 0.18	0.12 0.10 0.74 0.68 0.18	0.00 0.00 0.21 0.24 0.50 0.84	0.00 0.00 0.20 0.25 0.54 0.84	0.08 0.07 0.67 0.61 0.17 0.18	0.08 0.07 0.67 0.62 0.18 0.17	0.00 0.00 0.13 0.17 0.47 0.71	0.00 0.00 0.13 0.18 0.49	0.00 0.00 0.17 0.11 0.78 0.97	0.00 0.00 0.17 0.11 0.76 0.97
Ericsson	Contrib UContrib Conv HVol Diversity Speed	0.00 0.0 0.00 0.0 0.51 0.5 0.30 0.4 0.44 0.4 0.98 0.9	0 0.01 0 0.01 0 0.56 1 0.33 5 0.87	0.01 0.01 0.42 0.40 0.88 0.99	0.96 0.94 0.99 0.96 0.77 0.10	0.96 0.94 0.99 0.96 0.75 0.10	0.02 0.02 0.16 0.75 0.64 0.72	0.02 0.02 0.15 0.74 0.84 0.74	0.97 0.95 0.99 0.95 0.75 0.11	0.97 0.97 0.99 0.93 0.77 0.10	0.02 0.02 0.19 0.76 0.64 0.72	0.02 0.02 0.18 0.76 0.74	0.01 0.01 0.15 0.79 0.85 0.88	0.01 0.01 0.15 0.80 0.81
MS Word	Contrib UContrib Conv HVol Diversity Speed	0.62 0.6 0.73 0.7 0.96 0.9 0.93 0.9 0.30 0.3 0.98 0.9	3 0.49 7 0.88 2 0.81 2 0.31	0.41 0.49 0.89 0.80 0.29 0.97	0.83 0.81 0.82 0.83 0.14 0.27	0.83 0.80 0.82 0.84 0.14 0.26	0.00 0.00 0.22 0.26 0.55 0.86	0.00 0.00 0.20 0.25 0.62 0.90	0.76 0.73 0.82 0.84 0.21 0.12	0.76 0.73 0.82 0.85 0.20 0.13	0.00 0.00 0.27 0.32 0.51 0.84	0.00 0.00 0.28 0.30 0.56 0.85	0.00 0.00 0.20 0.23 0.79 0.99	0.00 0.00 0.21 0.24 0.79 0.99
Eclipse	Contrib UContrib Conv HVol Diversity Speed	0.60 0.5 0.53 0.5 0.98 0.9 0.98 0.9 0.69 0.6 0.95 0.9	3 0.00 9 0.59 8 0.62 9 0.65	0.00 0.00 0.61 0.62 0.65 0.98	0.81 0.73 0.92 0.92 0.20 0.79	0.82 0.77 0.92 0.92 0.19 0.80	0.00 0.00 0.68 0.75 0.45 0.96	0.00 0.00 0.68 0.75 0.48 0.96	0.58 0.52 0.92 0.90 0.14 0.79	0.58 0.52 0.92 0.90 0.14 0.80	0.00 0.00 0.63 0.59 0.27 0.93	0.00 0.00 0.68 0.66 0.30	0.00 0.00 0.75 0.80 0.93 0.05	0.00 0.00 0.75 0.80 0.94 0.05
Mozilla	Contrib UContrib Conv HVol Diversity Speed	0.73 0.7 0.63 0.6 0.97 0.9 0.98 0.9 0.72 0.7 0.96 0.9	3 0.00 5 0.00 7 0.70 8 0.66 2 0.64	0.00 0.00 0.72 0.68 0.65	0.85 0.75 0.92 0.92 0.20 0.83	0.84 0.75 0.93 0.92 0.20 0.83	0.00 0.00 0.72 0.75 0.41 0.97	0.00 0.00 0.72 0.75 0.43 0.97	0.61 0.54 0.92 0.90 0.14 0.83	0.62 0.53 0.92 0.90 0.14 0.83	0.00 0.00 0.64 0.57 0.34 0.95	0.00 0.00 0.69 0.66 0.36	0.00 0.00 0.77 0.80 0.91 0.06	0.00 0.00 0.77 0.80 0.91 0.06
Gnome	Contrib UContrib Conv HVol Diversity Speed	0.61 0.6 0.48 0.4 0.98 0.9 0.97 0.9 0.72 0.7 0.94 0.9	9 0.00 8 0.59 7 0.61 0 0.68	0.00 0.00 0.58 0.61 0.68	0.84 0.68 0.92 0.92 0.21 0.71	0.84 0.67 0.92 0.92 0.22 0.72 LE III	0.00 0.00 0.67 0.74 0.42 0.95	0.00 0.00 0.67 0.74 0.47 0.95	0.56 0.46 0.91 0.89 0.15 0.71	0.55 0.45 0.91 0.90 0.15 0.71	0.00 0.00 0.62 0.65 0.42 0.92	0.00 0.00 0.64 0.70 0.45 0.93	0.01 0.01 0.74 0.80 0.95 0.06	0.01 0.01 0.74 0.80 0.94 0.07

THE PERFORMANCE (MEAN AND MEDIAN) OF THE 7 ALGORITHMS FOR THE 10 DATASETS. ALL METRICS REPORTED IN THIS TABLE ARE NORMALISED AND MAXIMISING SO THE READER CAN ASSUME THAT 'HIGHER NUMBERS MEAN BETTER PERFORMANCE' IN ALL CASES. UNSURPRISINGLY, THE RESULTS SHOW THAT RANDOM SEARCH TENDS TO PRODUCE LOW QUALITY SOLUTIONS. A LITTLE MORE SURPRISINGLY, THE META-HEURISTIC ALGORITHMS (HC AND SA) ALSO CONTRIBUTE LITTLE TO THE BEST SOLUTIONS FOUND (AS ASSESSED BY THE METRICS 'CONTRIB' AND 'UCONTRIB' IN THE TABLE). THE RESULTS ALSO SHOW THAT THE ERICSSON DATASET OCCASIONS VERY DIFFERENT BEHAVIOUR FROM THE ALGORITHMS COMPARED TO THE OTHER DATASETS, INDICATING THAT THE DATASET STUDIED REALLY DOES MATTER IN EMPIRICAL STUDIES OF RELEASE PLANNING.

Data					F	IGA					GA			1	H	SA			SA		I	IHC	НС
Set	M.	etric	GA	HSA	SA	HHC	HC	R	HSA	SA	HHC	HC	R	SA	HHC	HC	R	HHC	HC	R	HC	R	R
	l	Contrib	1.00	1.00	1.00	1.00	1.00	1.00	0.03	0.43	0.05	0.43	0.43	1.00	0.81	1.00	1.00	0.00	i	i	1.00	1.00	Ŧ
		UContrib	1.00	1.00	1.00	1.00	1.00	1.00	0.03	0.43	0.06	0.43	0.43	1.00	0.80	1.00	1.00	0.00			1.00	1.00	
D	Quality	Conv	1.00	1.00	1.00	1.00	1.00	1.00		1.00	0.93	1.00	1.00	1.00	0.90	1.00	1.00	0.00	0.75	0.81	1.00	1.00	
Baan		HVol	1.00	1.00	1.00	1.00	1.00	1.00	0.85	1.00	0.98	1.00	1.00	1.00	0.94	1.00	1.00	0.00	0.84	1.00	1.00	1.00	0.83
	Diversity			1.00	0.16	1.00		0.00	1.00	0.10	1.00	0.28	0.00	0.00		0.07	0.00	1.00	0.72	0.20	0.06	0.00	0.04
	Speed		0.12	1.00	1.00	1.00	1.00	0.00	1.00	1.00	1.00	1.00	0.02	0.00	0.99	0.06	0.00	1.00	0.77	0.00	0.00	0.00	0.00
		Contrib	1.00	1.00	1.00	1.00	1.00	1.00	0.01	0.97	0.10	0.96	0.97	1.00	0.80	1.00	1.00	0.00			1.00	1.00	Т
	Quality	UContrib	1.00	1.00	1.00	1.00	1.00	1.00	0.06	0.96	0.16	0.96	0.97	1.00	0.75	1.00	1.00	0.00		İ	1.00	1.00	
Stone-	Quanty	Conv	1.00	1.00	1.00	1.00	1.00	1.00		1.00		1.00	1.00	1.00	0.77	1.00	1.00	0.00			1.00	1.00	0.20
Gate		HVol	1.00	1.00	1.00	1.00	1.00	1.00	0.12	1.00	0.11	1.00	1.00	1.00		1.00	1.00	0.00		0.25	1.00	1.00	0.12
	Diversity			1.00	0.27	1.00	1.00	0.02	1.00	0.11	1.00	1.00	0.00	0.01	1.00	0.03	0.00	0.99		0.05	0.02	0.00	0.02
	Speed		<u> </u>	1.00	1.00	1.00	1.00	0.75	1.00	1.00	1.00	1.00		0.00	1.00	0.00	0.00	1.00		0.00	0.00	0.00	0.00
		Contrib	0.98	1.00	1.00	1.00	1.00	1.00	0.71	1.00	0.73	1.00	1.00	1.00		1.00	1.00	0.00			1.00	1.00	
	Quality	UContrib	0.97	1.00	1.00	1.00	1.00	1.00	0.85	1.00	0.88	1.00	1.00	1.00		1.00	1.00	0.00		0.00	1.00	1.00	0.05
Motorola	` '	Conv	0.98 1.00	1.00	1.00	1.00	1.00	1.00	1.00	1.00	1.00	1.00	1.00	1.00		1.00	1.00	0.00	0.28	0.93	1.00	1.00	0.95
	Diversity	H VOI	1.00	1.00	0.07	1.00	0.14	0.02	1.00	0.05	1.00	0.15	0.01	0.05		0.13	0.00	0.00	0.28	0.99	0.19	0.01	0.10
	Speed	1	0.91	1.00	0.98	1.00	0.96	0.02	1.00	0.05	1.00	0.94	0.00	0.00	0.78	0.00	0.00	1.00	+	0.00	0.00	0.00	0.00
	. speed	1 0	1					1	100		1	1	1		10.70	1			+	15.00		-	15.00
		Contrib	0.99	1.00	1.00	1.00	1.00	1.00		1.00		1.00	1.00	1.00		1.00	1.00	0.00			1.00	1.00	
	Quality	UContrib Conv	0.99 1.00	1.00	1.00	1.00	1.00	1.00	1.00	1.00	1.00	1.00	1.00	1.00	0.78	1.00	1.00	0.00			1.00	1.00	
RalicP		HVol	1.00	1.00	1.00	1.00	1.00	1.00	1.00	1.00	1.00	1.00	1.00	1.00	0.78	1.00	1.00	0.00	0.79	0.99	1.00	1.00	0.79
	Diversity		1	0.74	0.15	0.77	0.22	0.00	0.74	0.12	0.79	0.19	0.00	0.12	1	0.13	0.00	0.89	1	0.12	0.12	0.00	0.02
	Speed		0.07	1.00	1.00	1.00	1.00	0.31	1.00	1.00	1.00	1.00	0.81	0.00	0.78	0.00	0.00	1.00		0.00	0.00	0.00	0.00
		Contrib	1.00	1.00	1.00	1.00	1.00	1.00		0.83	i i	0.83	0.83	1.00	0.78	1.00	1.00	0.00	i 	† 	1.00	1.00	Ť
		UContrib	1.00	1.00	1.00	1.00	1.00	1.00		0.83		0.83	0.83	1.00	0.74	1.00	1.00	0.00			1.00	1.00	
	Quality	Conv	1.00	1.00	1.00	1.00	1.00	1.00	0.88	1.00	0.99	1.00	1.00	1.00	0.91	1.00	1.00	0.00	0.78	0.71	1.00	1.00	
RalicR		HVol	1.00	1.00	1.00	1.00	1.00	1.00	0.71	1.00	0.93	1.00	1.00	1.00	0.97	1.00	1.00	0.00	0.81	0.97	1.00	1.00	0.80
	Diversity			0.92		0.93		0.00	0.96		0.98		0.00	0.11		0.12	0.00	0.89		0.09	0.12	0.00	0.07
	Speed		0.09	1.00	1.00	1.00	0.97	0.18	1.00	1.00	1.00	0.97	0.77	0.00	0.99	0.04	0.00	1.00	0.81	0.00	0.00	0.00	0.03
		Contrib	0.12	0.00	0.00	0.00	0.04		0.00	0.00	0.00	0.12		1.00		1.00	1.00	0.00		1.00	1.00	1.00	0.88
	Quality	UContrib	0.08	0.00	0.00	0.00	0.04	İ	0.00	0.00	0.00	0.12	İ	1.00		1.00	1.00	0.00		1.00	1.00	1.00	0.88
Ericsson	Quanty	Conv		0.00	1.00	0.00	0.92	1.00	0.00	1.00	0.00	0.88	1.00	1.00		1.00	1.00	0.00			1.00	1.00	
Elicason		HVol		0.00	0.00	0.00	0.00	0.00	0.00	0.00	0.00	0.00	0.00	1.00		1.00	1.00	0.00		0.04	1.00	1.00	
	Diversity		0.00	0.00	0.20	0.00	0.24	0.00	0.92	0.84	0.96	0.76	1.00	0.00		0.00	0.20	1.00		0.28	0.00	0.20	0.28
	Speed		0.15	1.00	1.00	1.00	1.00	1.00	1.00	1.00	1.00	1.00	1.00	0.00		0.00	0.00	1.00		0.00	0.00	0.00	0.00
		Contrib	0.96	0.04	1.00	0.13	1.00	1.00	0.00	1.00	0.01	1.00	1.00	1.00	0.70	1.00	1.00	0.00			1.00	1.00	
	Quality	UContrib	0.96		1.00	1.00	1.00	1.00	0.01	1.00	0.07	1.00	1.00	1.00		1.00	1.00	0.00			1.00	1.00	0.75
MS Word	` '	Conv	0.97 0.99	1.00	1.00	1.00	1.00	1.00	0.84	1.00	0.86	1.00	1.00	1.00		1.00	1.00	0.00			1.00	1.00	0.75 0.83
word	Diversity	H VOI	0.99	0.78	0.20	1.00	0.26	0.01	0.81	0.22		0.27	0.01	0.12	0.23	0.13	0.00	0.83	_	0.16	0.18	0.00	0.83
	Speed		1.00	1.00	1.00	1.00	1.00	0.12	1.00	1.00	1.00	1.00	0.00	0.00	0.99	0.00	0.00	1.00	+	0.00	0.00	0.00	0.00
		Contrib	1.00	0.04	1.00	+	1.00	1.00	0.00	1	0.00	1	0.03	1.00	0.97	1.00	1.00	0.00	† 	0.08	1.00	1.00	0.06
		UContrib	1.00	0.04	1.00		1.00	1.00	0.00		0.00		0.03	1.00	0.97	1.00	1.00	0.00		0.08	1.00	1.00	0.05
	Quality	Conv	1.00	1.00	1.00	1.00	1.00	1.00	0.00	0.19	0.00	0.30	0.02	1.00	0.84	1.00	1.00	0.00		0.03	1.00	1.00	0.02
Eclipse		HVol	1.00	1.00	1.00	1.00	1.00	1.00	0.00	0.03	0.00	1	0.00	1.00	0.95	1.00	1.00	0.00	0.94	0.00	1.00	1.00	0.00
	Diversity			1.00	0.98	1.00	1.00	0.00	1.00	0.98	1.00	1.00	0.00	0.02	0.94		0.00	1.00	0.83	0.00		0.00	0.00
	Speed		0.00	1.00	0.19	1.00	0.76	1.00	1.00	0.99	1.00	1.00	1.00	0.00		0.02	1.00	1.00	0.92	1.00	0.02	1.00	1.00
		Contrib	1.00	0.18	1.00	0.79	1.00	1.00	0.00	T	0.00	I	0.00	1.00	0.98	1.00	1.00	0.00	T	T	1.00	1.00	0.06
	Olite.	UContrib	1.00	0.21	1.00	0.78	1.00	1.00	0.00		0.00		0.00	1.00	0.87	1.00	1.00	0.00			1.00	1.00	0.04
Mozilla	Quality	Conv	1.00	1.00	1.00	1.00	1.00	1.00	0.00		0.00		0.14	1.00	0.76	1.00	1.00	0.00	0.82	0.04	1.00	1.00	0.00
MOZIIIa		HVol	1.00	1.00	1.00	1.00	1.00	1.00	0.00	0.10	0.00		0.00	1.00	1.00	1.00	1.00	0.00	0.93	0.00	1.00	1.00	0.00
	Diversity		0.81	1.00	0.99	1.00	0.99	0.00	1.00	0.96	1.00	0.98	0.00	0.08	0.94	0.22	0.00	0.94	0.01	0.00	0.13	0.00	0.00
	Speed		0.00	1.00	0.06	1.00	1	1.00	1.00	0.99	1.00	1.00	1.00	0.00	1	0.00	1.00	1.00	0.91	1.00	0.00	1.00	1.00
			1.00	0.08	1.00		1.00	1.00	0.00		0.00		0.04	1.00	0.97	1.00	1.00	0.00		0.08	1.00	1.00	0.03
	Quality	UContrib	1.00	0.09	1.00		1.00	1.00	0.00		0.00		0.03	1.00	0.86	1.00	1.00	0.00		0.07	1.00	1.00	0.02
Gnome	~	Conv	1.00	1.00	1.00	1.00	1.00	1.00	0.00	0.25	0.00	l	0.04	1.00	0.82	1.00	1.00	0.00	L	0.08	1.00	1.00	0.00
	D: :	HVol	1.00	1.00	1.00	1.00	1.00	1.00	0.00	0.07	0.00	0.32	0.00	1.00	0.97	1.00	1.00	0.00	0.94	0.00	1.00	1.00	0.00
	Diversity	-	0.69	1.00	0.99	1.00	0.99	0.00	1.00	0.98	1.00	0.99	0.00	0.12	0.93	0.10	0.00	0.90 1.00	0.02	0.00	0.07	0.00	0.00
	Speed	1	μ.σσ	11.00	0.20	1.00	D.//	μ.00	11.00			μ.00	μ.00	0.00	1	0.00	μ.00	1.00	0.92	1.00	0.00	μ.00	11.00
										TABI													

Results of the statistical analysis. The numbers reported in this table are the Vargha-Delaney \hat{A}_{12} metric for effect size, where the result is significant at the $0.05~\alpha$ level according to Cliff's test. Where the entry is blank, the corresponding algorithmic comparison result is not significant. The value of \hat{A}_{12} is the probability that the algorithm labelled at the head of the column outperforms the algorithms labelled immediately below. For example, in the pair of columns headed $\frac{\text{HHC}}{\text{HC} \mid \text{R}}$, the HHC algorithm is compared to each of the HC and R algorithms. Consider the first of this pair of columns, which compares HHC to HC. Where the \hat{A}_{12} is greater than 0.5, this means that HHC (significantly) outperforms HC and where it is below 0.5, HC outperforms HHC. The closer \hat{A}_{12} is to 1.0, the greater the probability that HHC would be expected to outperform HC; the closer to 0.0, the greater the probability that HC would be expected to outperform HHC.

forms its rival in 70% (42/60) of the pairwise comparisons.

On the largest dataset, Mozilla, which has more than 4,000 requirements, each of the executions of random search took more than 13 minutes to complete, while each HGA execution took just over 3 minutes. Neither of these durations makes a big difference to the kind of release planning applications that could be undertaken. For the smaller datasets (with fewer than 200 requirements), each HGA executions completed in fewer than 10 seconds (sometimes merely 1 or 2 seconds). This puts HGA tantalisingly close to the threshold at which it could be used to investigate 'what if' scenarios; the release planner could modify the available requirements and/or their attributes and explore the impact of such changes in real time. **RQ4: Scalability**: The tables in Figure 2 highlight a scalability problem for meta-heuristic NSGA2 (denoted GA in the tables): as the number of requirements increase, GA's contributions

to the reference front decrease. This observation remains consistent whether we measure the mean, median or the best performance of each algorithm and also holds whether we use Kendall's or Spearman's correlation.

Figure 2 also reveals a negative correlation between the number of requirements and convergence of meta-heuristic NSGA2 and the best performance of meta-heuristic NSGA2 for hypervolume. Taken together, these negative correlation results for meta-heuristic NSGA2 quality metrics suggest that the quality of solutions it produces tend to decrease as the problem size increases.

We also observe a slightly less strong, but positive, correlation between the number of requirements and diversity of solutions produced by meta-heuristic NSGA2. This suggests meta-heuristic NSGA2 increases its diversity with scale. However, since its contribution and quality tend to decrease as the

Alg.	Medi	an value corr	elated wi	th numbe	r of requirem	ents					
Aig.		Quality			Diversity	Speed					
	Contrib	UContrib	Conv	HVol	Diversity	Specu					
HGA	-0.21	-0.52	0.21	0.30	0.66	-0.21					
GA	-0.61	-0.56	-0.52	-0.48	0.61	0.25					
HSA	0.11	0.02	0.11	0.21	0.25	0.43					
SA	-0.05	-0.05	0.30	0.43	-0.25	0.36					
HHC	0.07	-0.02	0.07	0.11	-0.16	0.57					
HC	-0.05	-0.05	0.11	0.16	-0.25	0.39					
R	0.26	0.26	0.07	0.52	0.71	-0.61					
Alg. Best value correlated with number of requirements											
Aig.		Quality			Diversity	Speed					
	Contrib	UContrib	Conv	HVol		Speed					
HGA	-0.17	-0.44	0.05	0.05	0.52	-0.25					
GA	-0.57	-0.57	-0.61	-0.57	0.43	0.75					
HSA	0.32	0.27	0.11	0.21	-0.07	0.43					
SA	0.17	0.12	0.30	0.25	-0.61	0.25					
HHC	0.02	-0.02	0.07	0.16	-0.39	0.57					
HC	-0.07	-0.07	0.21	0.11	-0.30	0.21					
R	0.71	0.66	0.34	0.57	0.00	-0.85					
Alg.	Mea	n value corre		h number	of requireme	nts					
rug.		Quality			Diversity	Speed					
	Contrib	UContrib	Conv	HVol	,						
HGA	-0.25	-0.48	0.21	0.25	0.71	-0.25					
GA	-0.61	-0.57	-0.52	-0.48	0.66	0.25					
HSA	0.11	0.05	0.11	0.21	0.25	0.43					
SA	0.02	-0.12	0.25	0.34	-0.39	0.43					
HHC	-0.07	-0.02	0.07	0.11	-0.23	0.57					
HC	-0.12	-0.02	0.07	0.16	-0.21	0.39					
R	0.61	0.61	0.25	0.52	0.66	-0.66					

Kendall correlation values

A1~	Medi	an value corr	elated wi	th numbe	r of requirem	ents	
Alg.		Quality	Diversity	Speed			
	Contrib	UContrib	Conv	HVol	Diversity	Speed	
HGA	-0.28	-0.63	0.21	0.41	0.83	-0.18	
GA	-0.73	-0.71	-0.61	-0.50	0.73	0.41	
HSA	0.23	0.06	0.31	0.32	0.27	0.63	
SA	-0.06	-0.06	0.52	0.62	-0.46	0.59	
HHC	0.18	0.05	0.27	0.29	-0.17	0.78	
HC	-0.06	-0.06	0.36	0.31	-0.46	0.62	
R	0.31	0.31	0.32	0.67	0.81	-0.78	
Alg.	Bes	t value correl	lated with	number	of requiremen	nts	
Aig.		Quality	Diversity	Speed			
	Contrib	UContrib	Conv	HVol	_		
HGA	-0.22	-0.53	0.05	0.06	0.65	-0.18	
GA	-0.70	-0.70	-0.78	-0.69	0.54	0.87	
HSA	0.47	0.36	0.31	0.32	-0.17	0.63	
SA	0.19	0.07	0.46	0.40	-0.70	0.40	
HHC	0.05	0.05	0.30	0.31	-0.51	0.79	
HC	-0.19	-0.19	0.43	0.34	-0.36	0.38	
R	0.82	0.79	0.54	0.66	0.00	-0.94	
Alg.	Mea	n value corre	lated wit	h number	of requireme	nts	
Aig.		Quality			Diversity	Speed	
	Contrib	UContrib	Conv	HVol			
HGA	-0.38	-0.60	0.21	0.35	0.85	-0.20	
GA	-0.80	-0.79	-0.61	-0.50	0.74	0.38	
HSA	0.23	-0.02	0.31	0.32	0.27	0.59	
SA	0.04	-0.22	0.47	0.51	-0.59	0.62	
HHC	0.02	0.05	0.27	0.29	-0.26	0.78	
HC	-0.21	-0.17	0.34	0.31	-0.45	0.57	
R	0.74	0.74	0.42	0.67	0.77	-0.81	

Spearman correlation values

Fig. 2. Correlations of metrics for quality, diversity and speed with size of problem. Perhaps surprisingly, the speed of Random search (labelled 'R' in the table), widely believed to be 'fast but low quality' does not scale well (indicated by a negative correlation). Also we observe that while the quality of meta-heuristic NSGA2 (labelled 'GA') tends to degrade with size, the quality of hyper-heuristic NSGA2 (labelled 'HGA') does not.

scale of the problem increases its diversity is of considerably lesser value; it is simply producing a wider range of increasing sub-optimal solutions as the problem scales.

Fortunately, the other algorithms found to perform well in

answer to RQ1 do not exhibit any such evidence for negative correlation between problem size and solution quality. In particular, hyper-heuristic NSGA2 exhibits no such correlation. Even more encouraging for this algorithm, we find consistent evidence, across all six correlation values, that it *increases its diversity* as the scale of the problem increases.

Actionable Findings and Threats to Validity: Our results are based on only 10 datasets. This is considerably larger than any previous empirical study of release planning. These datasets are obtained from open source as well as closed source, system tools as well as enterprise applications, and have sizes varying from 100 to 4,000 requirements. Nevertheless, it remains insufficient to generalise to every type of project in every scenario. Indeed, the results suggest that algorithms can behave very differently with respect to different datasets.

Therefore, as with other experimental/empirical SBSE work [6], this finding suggests that the use of synthetic datasets in experimental work on release planning should be augmented with the study of real world datasets most likely to share the characteristics of the problem domain to which the proposed algorithms are to be applied.

We have attempted to control potential threats to construct validity by the use of ordinal inferential statistical techniques which make no assumptions about distribution (including variance). For correlation analysis we use two different non-parametric correlations in order to increase confidence in our findings. These findings suggest that hyper-heuristics are an attractive new direction for release planning optimisation and that, in particular, hyper-heuristic NSGA2 (the algorithm labelled 'HGA') is highly attractive: it typically outperforms the other algorithms studied in terms of quality, diversity *and* speed. Furthermore, it appears that it scales well compared to its meta-heuristic counterpart and other algorithms studied.

VI. CONCLUSION

We have presented a comprehensive study of meta-heuristic and hyper-heuristic release planning on 10 real world datasets. Overall, hyper-heuristic NSGA2 performs the best in terms of quality, diversity, speed and scalability. However, our results also indicate that the hyper-heuristic versions of Simulated Annealing and Hill Climbing also make some contribution to the best solutions found and are also relatively scalable.

This finding suggests that if only a single algorithm is to be used then it should be hyper-heuristic NSGA2, but if resources allow, it may be advantageous to combine its results with those from other hyper-heuristic algorithms.

Furthermore, we found that algorithm behaviour can differ greatly from one dataset to another, indicating that research on synthetic datasets needs to be augmented with analysis of appropriate real world datasets.

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