…I have to make a confession: when it comes to my sense of orientation I am a total failure… sometimes it feels like GPS and Google maps were actually invented for me!

Well, nowadays anybody uses those practical little helpers. But how do they actually manage to find the shortest path from A to B?

If you want to understand the father of all routing algorithms, *Dijkstra’s algorithm*, and want to know how to program it in R read on!

This post is partly based on this essay Python Patterns – Implementing Graphs, the example is from the German book “Das Geheimnis des kürzesten Weges” (“The secret of the shortest path”) by my colleague Professor Gritzmann and Dr. Brandenberg.

Dijkstra’s algorithm is a *recursive algorithm*. If you are not familiar with *recursion* you might want to read my post To understand Recursion you have to understand Recursion… first.

First, we are going to define the *graph* in which we want to navigate and we attach *weights* for the time it takes to cover it. We use the excellent igraph package (on CRAN) for visualizing the graph:

library(igraph)

## Attaching package: 'igraph'

## The following objects are masked from 'package:stats': ##

## decompose, spectrum

## The following object is masked from 'package:base': ##

## union

graph <- list(s = c("a", "b"),

a = c("s", "b", "c", "d"),

b = c("s", "a", "c", "d"),

c = c("a", "b", "d", "e", "f"),

d = c("a", "b", "c", "e", "f"),

e = c("c", "d", "f", "z"),

f = c("c", "d", "e", "z"),

z = c("e", "f"))

weights <- list(s = c(3, 5),

a = c(3, 1, 10, 11),

b = c(5, 3, 2, 3),

c = c(10, 2, 3, 7, 12),

d = c(15, 7, 2, 11, 2),

e = c(7, 11, 3, 2),

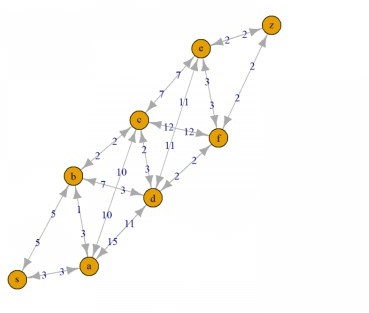
f = c(12, 2, 3, 2), z = c(2, 2))

# create edgelist with weights

G <- data.frame(stack(graph), weights = stack(weights)[[1]]) set.seed(500)

el <- as.matrix(stack(graph)) g <- graph\_from\_edgelist(el)

oldpar <- par(mar = c(1, 1, 1, 1)) plot(g, edge.label = stack(weights)[[1]]) par(oldpar)



Next, we create a helper function to calculate the path length:

path\_length <- function(path) {

# if path is NULL return infinite length if (is.null(path)) return(Inf)

# get all consecutive nodes

pairs <- cbind(values = path[-length(path)], ind = path[-1]) # join with G and sum over weights

sum(merge(pairs, G)[ , "weights"])

}

And now for the core of the matter, Dijkstra’s algorithm: the general idea of the algorithm is very simple and elegant: start at the starting node and call the algorithm recursively for all nodes linked from there as new starting nodes and thereby build your path step by step. Only keep the shortest path and stop when reaching the end node (base case of the recursion). In case you reach a dead-end in between assign infinity as length (by the path\_length function above).

I added a lot of documentation to the code so it is hopefully possible to understand how it works:

find\_shortest\_path <- function(graph, start, end, path = c()) {

# if there are no nodes linked from current node (= dead end) return NULL

if (is.null(graph[[start]])) return(NULL) # add next node to path so far

path <- c(path, start)

# base case of recursion: if end is reached return path if (start == end) return(path)

# initialize shortest path as NULL shortest <- NULL

# loop through all nodes linked from the current node (given in start)

for (node in graph[[start]]) {

# proceed only if linked node is not already in path if (!(node %in% path)) {

# recursively call function for finding shortest path with node as start and assign it to newpath

newpath <- find\_shortest\_path(graph, node, end, path)

# if newpath is shorter than shortest so far assign newpath to shortest

if (path\_length(newpath) < path\_length(shortest)) shortest <- newpath

}

}

# return shortest path shortest

}

Now, we can finally test the algorithm by calculating the shortest path from **s** to **z** and back:

find\_shortest\_path(graph, "s", "z") # via b ## [1] "s" "b" "c" "d" "f" "z"

find\_shortest\_path(graph, "z", "s") # back via a ## [1] "z" "f" "d" "b" "a" "s"

Note that the two routes are actually different because of the different weights in both directions (e.g. think of some construction work in one direction but not the other).