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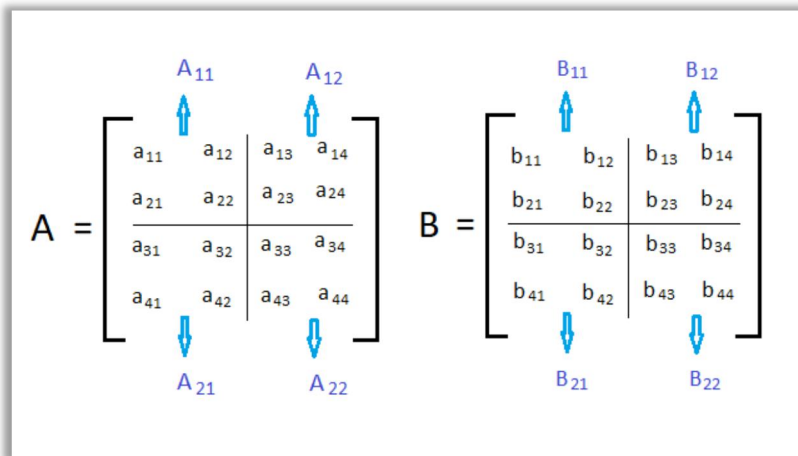
Homework#1: Matrix Multiplication

- the time complexity of a naïve approach is :

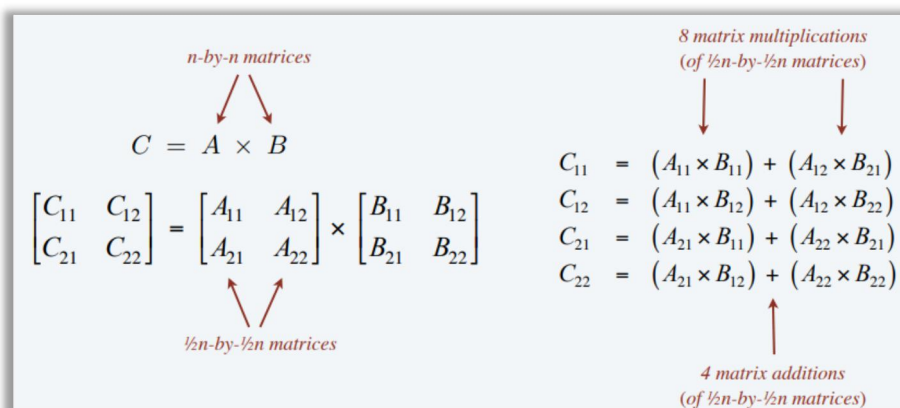
- ✓ there are 3 for loops.
- ✓ It needs three for loops; For accessing all the elements of any matrix we need two for loops. But for finding the product, it requires one additional for loop. That's how it is taking 3 for loops.
- ✓ So, Time complexity: $O(n^3)$

- the time complexity of the used divide-and-conquer approach:

- ✓ To multiply two n -by- n matrices A and B :
 >>Divide: partition A and B into $\frac{1}{2}n$ -by- $\frac{1}{2}n$ blocks.



- >>Conquer: multiply 8 pairs of $\frac{1}{2}n$ -by- $\frac{1}{2}n$ matrices, recursively.
- >>Combine: add appropriate products using 4 matrix additions.



- ✓ So, the time complexity:

#	work
1	kn
8	$8K(n/2) = K4n$
\vdots	\vdots
8^i	$8^i K(n/2^i) = K4^i n$
\vdots	\vdots
$8^{\log_2 n}$	$K8^{\log_2 n} = Kn^3$

So, total=

$$\sum_{i=0}^{\log_2 n} 8^i K(n/2^i) = \Theta(n^3)$$

- ✓ Thus, a naïve divide and conquer algorithm runs in $O(N^3)$.
- ✓ The passed matrices must be $n \times n$ matrices (square matrices), where n is a power of 2
- ✓ parameter n is the size of matrices.

● Optimized divide and conquer:

- ✓ Strassen's trick:

>>Key idea: Can multiply two n -by- n matrices via 7 ($\frac{1}{2}n$ -by- $\frac{1}{2}n$ matrix) multiplications (plus 11 additions and 7 subtractions).

$\frac{1}{2}n$ -by- $\frac{1}{2}n$ matrices

$$\begin{bmatrix} C_{11} & C_{12} \\ C_{21} & C_{22} \end{bmatrix} = \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix} \times \begin{bmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{bmatrix}$$

$$\begin{aligned} C_{11} &= P_5 + P_4 - P_2 + P_6 \\ C_{12} &= P_1 + P_2 \\ C_{21} &= P_3 + P_4 \\ C_{22} &= P_1 + P_5 - P_3 - P_7 \end{aligned}$$

$$\begin{aligned} P_1 &\leftarrow A_{11} \times (B_{12} - B_{22}) \\ P_2 &\leftarrow (A_{11} + A_{12}) \times B_{22} \\ P_3 &\leftarrow (A_{21} + A_{22}) \times B_{11} \\ P_4 &\leftarrow A_{22} \times (B_{21} - B_{11}) \\ P_5 &\leftarrow (A_{11} + A_{22}) \times (B_{11} + B_{22}) \\ P_6 &\leftarrow (A_{12} - A_{22}) \times (B_{21} + B_{22}) \\ P_7 &\leftarrow (A_{11} - A_{21}) \times (B_{11} + B_{12}) \end{aligned}$$

7 matrix multiplications (of $\frac{1}{2}n$ -by- $\frac{1}{2}n$ matrices)

Pf. $C_{12} = P_1 + P_2$
 $= A_{11} \times (B_{12} - B_{22}) + (A_{11} + A_{12}) \times B_{22}$
 $= A_{11} \times B_{12} + A_{12} \times B_{22}. \quad \checkmark$

✓ So, the time complexity:

#	work
1	kn
7	$7K(n/2)=K(7/2)n$
⋮	⋮
7^i	$7^i K(n/2^i)=K(7/2)^i n$
⋮	⋮
$7^{\log_2 n}$	$K 7^{\log_2 n}=Kn^{\log_2 7}$

So, total=

$$\sum_{i=0}^{\log_2 n} 7^i K(n/2^i) = O(n^{\log_2 7}) = O(n^{2.8074})$$

✓ Thus, a divide and conquer(Strassen Algorithm) which runs in $O(N^{\log_2 7})$, that is almost equal to $O(N^{2.8074})$.

- **Output of naïve.cpp ,divide_conquer.cpp & strassen_algorithm.cpp :** are the same which showing in the following figure.

```

Enter n: dimension of matrices n x n (n is a power of 2)
4
Enter elements of first matrix
1 1 1 1
2 2 2 2
3 3 3 3
4 4 4 4
Enter elements of second matrix
1 1 1 1
2 2 2 2
3 3 3 3
4 4 4 4
Result matrix is
10 10 10 10
20 20 20 20
30 30 30 30
40 40 40 40

Process returned 0 (0x0)   execution time : 44.813 s
Press any key to continue.

```

- **Note:** I used pointers in **strassen_algorithm.cpp** and vectors in **divide_conquer.cpp** to make functions accept any dimension of square matrices($n \times n$).
- why I didn't use pointers again in **divide_conquer.cpp**? Just for remembering C++ and learning more about it.