# Programming Language and Compilers - Assignment 1

# **Calculator Project**

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# The parser that was chose was a bottom up parser (SLR-1)

The grammar rules are as follows:

- R0:  $S' \rightarrow S$
- R1:  $S \rightarrow S * T$
- R2:  $\mathbf{S} \to \mathbf{T}$
- R3:  $T \rightarrow T + F$
- R4:  $\mathbf{T} \rightarrow \mathbf{F}$
- R5:  $\mathbf{F} \rightarrow \mathbf{N}$

The canonical LR(0) items are:

#### State 0

- $S' \rightarrow . S$
- $S \rightarrow . S * T$
- $S \rightarrow . T$
- $T \rightarrow . T + F$
- $T \rightarrow . F$
- $F \rightarrow . N$

#### State 1

- $S' \rightarrow S..$
- $S \rightarrow S \cdot * T$

## State 2

- $S \rightarrow T$ .
- $T \rightarrow T \cdot + F$

#### State 3

 $T \rightarrow F$ .

## State 4

 $F \rightarrow N$ .

# State 5

 $S \rightarrow S * . T$ 

 $T \rightarrow . T + F$ 

 $T \rightarrow . F$ 

 $F \rightarrow . N$ 

# State 6

 $T \rightarrow T + . F$ 

 $F \rightarrow . N$ 

# State 7

 $E \rightarrow E * T$ .

 $T \rightarrow T . + F$ 

# State 8

 $T \rightarrow T + F$ .

The parsing table is constructed as below:

	Action				Go To		
State	+	*	N	\$	S	Т	F
0			s4		1	2	3
1		s5		acc			
2	s6	r2		r2			
3	r4	r4		r4			
4	r5	r5		r5			
5			s4			7	3
6			s4				8
7	s6	r1		r1			
8	r3	r3		r3			

# **Rules of Translation**

#### 1. Value Calculation

Production	Semantic Rules
$S \rightarrow S1 * T$	S.val := S1.val * T.val
$S \rightarrow T$	S.val := T.val
$T \rightarrow T1 + F$	T.val := T1.val + F.val
$T \rightarrow F$	T.val := F.val
$F \rightarrow N$	F.val := N.lexval

## 2. Prefix Notation

Production	Semantic Rules
$S \rightarrow S1 * T$	S.pf := '*'    S1.pf    T.pf
$S \rightarrow T$	S.pf := T.pf
$T \rightarrow T1 + F$	$T.pf := '+' \parallel T1.pf \parallel F.pf$
$T \rightarrow F$	T.pf := F.pf
$F \rightarrow N$	F.pf := N.pf

## 3. Postfix Notation

Production	Semantic Rules
$S \rightarrow S * T$	S.pf := S.pf    T.pf    '*'
$S \to T$	S.pf := T.pf
$T \rightarrow T + F$	$T.pf := T1.pf \parallel F.pf \parallel '+'$
$T \rightarrow F$	T.pf := F.pf
$F \rightarrow N$	F.pf := N.pf