# Humuhumunukunukuapua'a UFMG

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## 31 de agosto de 2023

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### 1 Primitivas

#### 1.1 Aritmetica Modular

```
// O mod tem q ser primo
429 template <int p> struct mod_int {
        ll expo(ll b, ll e) {
c68
c85
            ll ret = 1;
c87
            while (e) {
                if (e % 2) ret = ret * b % p;
cad
9d2
                e /= 2, b = b * b % p;
            }
c42
edf
            return ret;
734
        11 inv(11 b) { return expo(b, p-2); }
1f6
4d7
        using m = mod_int;
d93
        int v;
fe0
        mod int() : v(0) {}
        mod int(ll v ) {
e12
019
            if (v_ >= p or v_ <= -p) v_ %= p;</pre>
bc6
            if (v_{-} < 0) v_{-} += p;
2e7
            v = v_{-};
7f3
        }
74d
        m& operator +=(const m& a) {
2fd
            v += a.v;
ba5
            if (v >= p) v -= p;
357
            return *this;
c8b
eff
        m& operator -=(const m& a) {
8b4
            v -= a.v;
cc8
            if (v < 0) v += p;
357
            return *this;
f8d
        m& operator *=(const m& a) {
4c4
8a5
            v = v * 11(a.v) \% p;
357
            return *this;
d4c
        m& operator /=(const m& a) {
3f9
5d6
            v = v * inv(a.v) % p;
357
            return *this;
62d
d65
        m operator -(){ return m(-v); }
b3e
        m& operator ^=(ll e) {
06d
            if (e < 0) {
```

```
6e2
                v = inv(v):
00c
                e = -e:
275
94a
            v = \exp(v, e\%(p-1));
357
            return *this;
ba3
423
        bool operator ==(const m& a) { return v == a.v; }
        bool operator !=(const m& a) { return v != a.v; }
69f
1 c 6
        friend istream& operator >>(istream& in, m& a) {
d1c
            11 val; in >> val;
d48
            a = m(val):
091
            return in:
870
44f
        friend ostream& operator <<(ostream& out, m a) {</pre>
5a0
            return out << a.v;</pre>
214
        }
        friend m operator +(m a, m b) { return a += b; }
399
f9e
        friend m operator -(m a, m b) { return a -= b; }
        friend m operator *(m a, m b) { return a *= b; }
9c1
        friend m operator /(m a, m b) { return a /= b; }
51b
08f
        friend m operator ^(m a, ll e) { return a ^= e; }
1af };
055 typedef mod_int < (int) 1e9+7 > mint;
1.2 Big Integer
// Complexidades: (para n digitos)
// Soma, subtracao, comparacao - O(n)
// Multiplicacao - O(n log(n))
// Divisao, resto - O(n^2)
864 struct bint {
669
        static const int BASE = 1e9;
990
        vector < int > v;
3bd
        bool neg;
609
        bint() : neg(0) {}
d53
        bint(int val) : bint() { *this = val; }
e8f
        bint(long long val) : bint() { *this = val; }
        void trim() {
a0f
f42
            while (v.size() and v.back() == 0) v.pop_back();
df8
            if (!v.size()) neg = 0;
8e3
        }
```

```
// converter de/para string | cin/cout
294
        bint(const char* s) : bint() { from_string(string(s)); }
548
        bint(const string& s) : bint() { from_string(s); }
4ab
        void from_string(const string& s) {
            v.clear(), neg = 0;
0a6
d72
            int ini = 0:
            while (ini < s.size() and (s[ini] == '-' or s[ini] == '+'
8e2
   or s[ini] == '0'))
                if (s[ini++] == '-') neg = 1;
71d
883
            for (int i = s.size()-1; i >= ini; i -= 9) {
05e
                int at = 0:
5b1
                for (int j = max(ini, i - 8); j \le i; j++) at = 10*at
   + (s[j]-'0');
1fd
                v.push_back(at);
a5a
df8
            if (!v.size()) neg = 0;
e9a
2ff
        string to_string() const {
            if (!v.size()) return "0";
8be
793
            string ret;
            if (neg) ret += '-';
73e
            for (int i = v.size()-1; i >= 0; i--) {
3e9
582
                string at = ::to_string(v[i]);
ced
                int add = 9 - at.size();
75e
                if (i+1 < v.size()) for (int j = 0; j < add; j++) ret</pre>
   += '0':
f9f
                ret += at:
f64
edf
            return ret;
770
d2f
        friend istream& operator >> (istream& in, bint& val) {
eb6
            string s: in >> s:
966
            val = s:
091
            return in;
328
99d
        friend ostream& operator << (ostream& out, const bint& val) {
8b9
            string s = val.to_string();
396
            out << s:
fe8
            return out;
ce1
       }
        // operators
        friend bint abs(bint val) {
60a
c5f
            val.neg = 0;
d94
            return val:
44b
        }
```

```
bee
        friend bint operator-(bint val) {
815
            if (val != 0) val.neg ^= 1;
d94
            return val;
326
        }
        bint& operator=(const bint& val) { v = val.v, neg = val.neg;
41f
   return *this; }
249
        bint& operator=(long long val) {
            v.clear(), neg = 0;
0a6
3a6
            if (val < 0) neg = 1, val *= -1;
            for (; val; val /= BASE) v.push_back(val % BASE);
fdc
357
            return *this;
220
        }
3bd
        int cmp(const bint& r) const { // menor: -1 | igual: 0 |
   maior: 1
b14
            if (neg != r.neg) return neg ? -1 : 1;
0bb
            if (v.size() != r.v.size()) {
ff7
                int ret = v.size() < r.v.size() ? -1 : 1;</pre>
                return neg ? -ret : ret;
91b
1f6
478
            for (int i = int(v.size())-1; i >= 0; i--) {
405
                if (v[i] != r.v[i]) {
2e5
                    int ret = v[i] < r.v[i] ? -1 : 1;</pre>
91b
                    return neg ? -ret : ret;
9a9
                }
c32
            }
bb3
            return 0;
07d
152
        friend bool operator < (const bint& 1, const bint& r) { return
   1.cmp(r) == -1;}
        friend bool operator > (const bint& 1, const bint& r) { return
c7a
   1.cmp(r) == 1; }
        friend bool operator <= (const bint& 1, const bint& r) { return
   1.cmp(r) <= 0; }
        friend bool operator >= (const bint& 1. const bint& r) { return
   1.cmp(r) >= 0; }
        friend bool operator == (const bint& 1, const bint& r) { return
   1.cmp(r) == 0; }
        friend bool operator!=(const bint& 1, const bint& r) { return
   1.cmp(r) != 0; }
38e
        bint& operator +=(const bint& r) {
            if (!r.v.size()) return *this;
6bf
a93
            if (neg != r.neg) return *this -= -r;
256
            for (int i = 0, c = 0; i < r.v.size() or c; i++) {</pre>
e28
                if (i == v.size()) v.push_back(0);
08f
                v[i] += c + (i < r.v.size() ? r.v[i] : 0);
                if ((c = v[i] >= BASE)) v[i] -= BASE:
baa
```

```
8bb
357
            return *this:
ab1
54c
        friend bint operator+(bint a, const bint& b) { return a += b; }
9c8
        bint& operator -=(const bint& r) {
            if (!r.v.size()) return *this;
6bf
524
            if (neg != r.neg) return *this += -r;
358
            if ((!neg and *this < r) or (neg and r < *this)) {
b10
                *this = r - *this;
a10
                neg ^= 1;
357
                return *this;
807
            }
256
            for (int i = 0, c = 0; i < r.v.size() or c; i++) {
                v[i] -= c + (i < r.v.size() ? r.v[i] : 0);
9ef
                if ((c = v[i] < 0)) v[i] += BASE;</pre>
c8c
687
            }
0eb
            trim();
357
            return *this;
f72
        friend bint operator-(bint a, const bint& b) { return a -= b: }
f44
        // operators de * / %
        bint& operator *=(int val) {
6b0
bca
            if (val < 0) val *= -1, neg ^= 1;</pre>
566
            for (int i = 0, c = 0; i < v.size() or c; i++) {</pre>
e28
                if (i == v.size()) v.push_back(0);
352
                long long at = (long long) v[i] * val + c;
6a3
                v[i] = at % BASE:
b3d
                c = at / BASE:
cb1
            }
            trim():
0eb
357
            return *this;
a57
        }
        friend bint operator *(bint a. int b) { return a *= b: }
480
        friend bint operator *(int a, bint b) { return b *= a; }
d5c
13b
        using cplx = complex <double >;
        void fft(vector<cplx>& a, bool f, int N, vector<int>& rev)
bfb
   const {
bc7
            for (int i = 0; i < N; i++) if (i < rev[i]) swap(a[i],
   a[rev[i]]);
bad
            vector < cplx > roots(N);
            for (int n = 2; n <= N; n *= 2) {</pre>
192
4e9
                const static double PI = acos(-1);
                for (int i = 0; i < n/2; i++) {
71a
40d
                    double alpha = (2*PI*i)/n;
                    if (f) alpha = -alpha;
1a1
3f6
                    roots[i] = cplx(cos(alpha), sin(alpha));
```

```
f16
3e9
                for (int pos = 0; pos < N; pos += n)
898
                     for (int 1 = pos, r = pos+n/2, m = 0; m < n/2;
   1++, r++, m++) {
297
                         auto t = roots[m]*a[r];
                         a[r] = a[1] - t;
254
b8f
                         a[1] = a[1] + t:
                     }
b0d
e07
3f1
            if (!f) return;
08Ъ
            auto invN = cplx(1)/cplx(N);
873
             for (int i = 0; i < N; i++) a[i] *= invN;</pre>
c75
0e0
        vector<long long> convolution(const vector<int>& a, const
   vector<int>& b) const {
            vector < cplx > l(a.begin(), a.end()), r(b.begin(), b.end());
ff9
996
            int ln = 1.size(), rn = r.size(), N = ln+rn+1, n = 1,
   log_n = 0;
821
            while (n \le N) n \le 1, \log_n + 1;
            vector < int > rev(n);
808
            for (int i = 0; i < n; i++) {</pre>
603
434
                 rev[i] = 0:
f44
                for (int j = 0; j < log_n; j++) if (i >> j & 1)
4ff
                     rev[i] = 1 << (log_n-1-i);
256
230
            l.resize(n), r.resize(n);
            fft(1, false, n, rev), fft(r, false, n, rev);
a89
917
            for (int i = 0; i < n; i++) l[i] *= r[i];
88ъ
            fft(1, true, n, rev);
7ae
            vector < long long > ret;
            for (auto& i : 1) ret.push_back(round(i.real()));
c14
edf
            return ret;
917
        }
633
        vector < int > convert base (const vector < int > & a. int from. int
   to) const {
498
            static vector<long long> pot(10, 1);
             if (pot[1] == 1) for (int i = 1; i < 10; i++) pot[i] =</pre>
671
   10*pot[i-1];
4b8
            vector < int > ret:
156
             long long at = 0;
            int digits = 0;
608
941
            for (int i : a) {
412
                 at += i * pot[digits];
035
                 digits += from;
684
                 while (digits >= to) {
0.c8
                     ret.push_back(at % pot[to]);
cf9
                     at /= pot[to];
```

```
fd4
                    digits -= to;
                }
122
87b
            }
944
            ret.push_back(at);
            while (ret.size() and ret.back() == 0) ret.pop back();
384
edf
            return ret;
090
       }
        bint operator*(const bint& r) const { // O(n log(n))
edb
2af
            bint ret;
            ret.neg = neg ^ r.neg;
968
d5d
            auto conv = convolution(convert_base(v, 9, 4),
   convert_base(r.v, 9, 4));
            long long c = 0:
a0e
            for (auto i : conv) {
a74
f6d
                long long at = i+c;
4cb
                ret.v.push_back(at % 10000);
a25
                c = at / 10000;
773
            for (; c; c /= 10000) ret.v.push_back(c%10000);
3cb
            ret.v = convert base(ret.v, 4, 9);
0e2
            if (!ret.v.size()) ret.neg = 0;
25 c
            return ret:
edf
c6b
359
        bint& operator*=(const bint& r) { return *this = *this * r; };
9a3
        bint& operator/=(int val) {
d9a
            if (val < 0) neg ^= 1, val *= -1;</pre>
            for (int i = int(v.size())-1, c = 0; i >= 0; i--) {
f18
                long long at = v[i] + c * (long long) BASE;
2a7
                v[i] = at / val;
e02
fb1
                c = at % val;
            }
fdb
0eb
            trim();
357
            return *this:
db6
        friend bint operator/(bint a, int b) { return a /= b; }
e74
        int operator %=(int val) {
4a9
23b
            if (val < 0) val *= -1;</pre>
156
            long long at = 0;
f31
            for (int i = int(v.size())-1: i >= 0: i--)
1b3
                at = (BASE * at + v[i]) \% val;
            if (neg) at *= -1;
d22
ce6
            return at;
4b4
        friend int operator%(bint a, int b) { return a %= b; }
2fb
        friend pair < bint, bint > divmod(const bint& a_, const bint& b_)
   { // O(n^2)}
611
            if (a == 0) return {0, 0}:
```

```
d8a
            int norm = BASE / (b_.v.back() + 1);
b4e
            bint a = abs(a_) * norm;
027
            bint b = abs(b_) * norm;
14d
            bint q, r;
            for (int i = a.v.size() - 1: i >= 0: i--) {
c91
                r *= BASE, r += a.v[i];
b71
4ff
                long long upper = b.v.size() < r.v.size() ?</pre>
   r.v[b.v.size()] : 0;
86d
                int lower = b.v.size() - 1 < r.v.size() ?</pre>
   r.v[b.v.size() - 1] : 0:
431
                int d = (upper * BASE + lower) / b.v.back();
5d4
30f
                while (r < 0) r += b, d--: // roda O(1) vezes
738
                q.v.push_back(d);
c6a
a48
            reverse(q.v.begin(), q.v.end());
            q.neg = a_.neg ^ b_.neg;
ae2
88b
            r.neg = a_.neg;
8e5
            q.trim(), r.trim();
            return {q, r / norm};
0ef
4fd
1d8
        bint operator/(const bint& val) { return divmod(*this,
   val).first: }
        bint& operator/=(const bint& val) { return *this = *this /
7f9
   val: }
1f9
        bint operator%(const bint& val) { return divmod(*this.
   val).second: }
        bint& operator%=(const bint& val) { return *this = *this %
   val; }
6c3 };
1.3 Fracao
// Funciona com o Big Int
```

```
a4e template < typename T = int > struct frac {
        T num, den;
a40
e3f
        template < class U, class V>
        frac(U num_ = 0, V den_ = 1) : num(num_), den(den_) {
61d
             assert(den != 0);
bad
583
             if (den < 0) num *= -1, den *= -1:
             T g = gcd(abs(num), den);
a51
572
             num /= g, den /= g;
fbf
        }
51f
        friend bool operator < (const frac& 1, const frac& r) {</pre>
```

```
fa0
            return l.num * r.den < r.num * l.den;</pre>
a4e
4b5
        friend frac operator+(const frac& 1, const frac& r) {
            return {1.num*r.den + 1.den*r.num, 1.den*r.den};
b61
25 f
74d
        friend frac operator-(const frac& 1, const frac& r) {
            return {1.num*r.den - 1.den*r.num, 1.den*r.den};
2cd
8a7
c80
        friend frac operator*(const frac& 1, const frac& r) {
510
            return {1.num*r.num, 1.den*r.den};
14b
       }
a1b
        friend frac operator/(const frac& 1, const frac& r) {
8f3
            return {1.num*r.den, 1.den*r.num};
b2c
       friend ostream& operator << (ostream& out, frac f) {</pre>
012
37a
            out << f.num << '/' << f.den:
fe8
            return out;
b49
       }
cdb };
```

#### 1.4 Geometria

```
c83 typedef double ld;
e3b const ld DINF = 1e18;
43a const ld pi = acos(-1.0);
107 const ld eps = 1e-9;
b32 #define sq(x) ((x)*(x))
d97 bool eq(ld a, ld b) {
ba0
        return abs(a - b) <= eps;</pre>
bfc }
b2a struct pt { // ponto
c1e
        ld x, v;
        pt(1d x_{-} = 0, 1d y_{-} = 0) : x(x_{-}), y(y_{-}) {}
3dd
5bc
        bool operator < (const pt p) const {</pre>
059
            if (!eq(x, p.x)) return x < p.x;
f98
            if (!eq(y, p.y)) return y < p.y;
            return 0;
bb3
f61
        }
        bool operator == (const pt p) const {
ed0
            return eq(x, p.x) and eq(y, p.y);
589
cb9
        pt operator + (const pt p) const { return pt(x+p.x, y+p.y); }
        pt operator - (const pt p) const { return pt(x-p.x, y-p.y); }
a24
```

```
4a8
        pt operator * (const ld c) const { return pt(x*c , y*c ); }
        pt operator / (const ld c) const { return pt(x/c , y/c ); }
a60
3b6
        ld operator * (const pt p) const { return x*p.x + y*p.y; }
        ld operator ^ (const pt p) const { return x*p.y - y*p.x; }
6df
        friend istream& operator >> (istream& in, pt& p) {
5ed
e37
            return in >> p.x >> p.y;
e45
        }
a8b };
b3a struct line { // reta
730
        pt p, q;
0d6
        line() {}
4b8
        line(pt p_, pt q_) : p(p_), q(q_) {}
8d7
        friend istream& operator >> (istream& in, line& r) {
4cb
            return in >> r.p >> r.q;
858
7ab };
// PONTO & VETOR
364 ld dist(pt p, pt q) { // distancia
5f3
        return hypot(p.y - q.y, p.x - q.x);
c68 }
9d7 ld dist2(pt p, pt q) { // quadrado da distancia
        return sq(p.x - q.x) + sq(p.y - q.y);
80f }
483 ld norm(pt v) { // norma do vetor
490
        return dist(pt(0, 0), v);
cf7 }
589 ld angle(pt v) { // angulo do vetor com o eixo x
        ld ang = atan2(v.v. v.x):
587
        if (ang < 0) ang += 2*pi;</pre>
6f8
19c
        return ang;
404 }
298 ld sarea(pt p, pt q, pt r) { // area com sinal
606
        return ((q-p)^(r-q))/2;
1b1 }
e32 bool col(pt p, pt q, pt r) { // se p, q e r sao colin.
        return eq(sarea(p, q, r), 0);
e7d
98c }
Ocd bool ccw(pt p, pt q, pt r) { // se p, q, r sao ccw
```

```
fa7
        return sarea(p, q, r) > eps;
85d }
1ef pt rotate(pt p, ld th) { // rotaciona o ponto th radianos
e5c
        return pt(p.x * cos(th) - p.y * sin(th),
ff1
                p.x * sin(th) + p.y * cos(th));
41a }
ab1 pt rotate90(pt p) { // rotaciona 90 graus
        return pt(-p.y, p.x);
e4a }
// RETA
edc bool isvert(line r) { // se r eh vertical
87d
        return eq(r.p.x, r.q.x);
Ofb }
099 bool isinseg(pt p, line r) { // se p pertence ao seg de r
        pt a = r.p - p, b = r.q - p;
f65
        return eq((a \hat{b}), 0) and (a * b) < eps;
b04
726 }
98d ld get_t(pt v, line r) { // retorna t tal que t*v pertence a reta r
        return (r.p^r.q) / ((r.p-r.q)^v);
a0a }
256 pt proj(pt p, line r) { // projecao do ponto p na reta r
       if (r.p == r.q) return r.p;
97a
       r.q = r.q - r.p; p = p - r.p;
9f8
        pt proj = r.q * ((p*r.q) / (r.q*r.q));
2cd
        return proj + r.p;
232 }
d5c pt inter(line r, line s) { // r inter s
       if (eq((r.p - r.q) ^ (s.p - s.q), 0)) return pt(DINF, DINF);
146
205
       r.q = r.q - r.p, s.p = s.p - r.p, s.q = s.q - r.p;
543
        return r.q * get_t(r.q, s) + r.p;
111 }
676 bool interseg(line r, line s) { // se o seg de r intersecta o seg
        if (isinseg(r.p, s) or isinseg(r.q, s)
19b
            or isinseg(s.p, r) or isinseg(s.q, r)) return 1;
c21
        return ccw(r.p, r.q, s.p) != ccw(r.p, r.q, s.q) and
9fa
413
                ccw(s.p, s.q, r.p) != ccw(s.p, s.q, r.q);
```

```
359 }
fcb ld disttoline(pt p, line r) { // distancia do ponto a reta
        return 2 * abs(sarea(p, r.p, r.q)) / dist(r.p, r.q);
1b7 }
bcc ld disttoseg(pt p, line r) { // distancia do ponto ao seg
        if ((r.q - r.p)*(p - r.p) < 0) return dist(r.p, p);
951
        if ((r.p - r.q)*(p - r.q) < 0) return dist(r.q, p);
        return disttoline(p, r):
367 }
11d ld distseg(line a, line b) { // distancia entre seg
        if (interseg(a, b)) return 0;
4df
        ld ret = DINF;
349
341
        ret = min(ret, disttoseg(a.p, b));
        ret = min(ret, disttoseg(a.q, b));
ceb
        ret = min(ret, disttoseg(b.p, a));
093
        ret = min(ret, disttoseg(b.q, a));
448
edf
        return ret;
222 }
// POLIGONO
// corta poligono com a reta r deixando os pontos p tal que
// ccw(r.p. r.q. p)
1a9 vector<pt> cut_polygon(vector<pt> v, line r) { // O(n)
8af
        vector<pt> ret;
        for (int j = 0; j < v.size(); j++) {</pre>
8a4
dac
            if (ccw(r.p, r.q, v[j])) ret.push_back(v[j]);
            if (v.size() == 1) continue;
dce
            line s(v[j], v[(j+1)%v.size()]);
0.30
            pt p = inter(r, s);
ae3
a3d
            if (isinseg(p, s)) ret.push_back(p);
d44
        }
8a1
        ret.erase(unique(ret.begin(), ret.end()), ret.end());
24d
        if (ret.size() > 1 and ret.back() == ret[0]) ret.pop_back();
edf
        return ret;
253 }
// distancia entre os retangulos a e b (lados paralelos aos eixos)
// assume que ta representado (inferior esquerdo, superior direito)
5f5 ld dist_rect(pair<pt, pt> a, pair<pt, pt> b) {
080
        ld hor = 0, vert = 0:
        if (a.second.x < b.first.x) hor = b.first.x - a.second.x;</pre>
34b
```

```
f5f
        else if (b.second.x < a.first.x) hor = a.first.x - b.second.x;</pre>
        if (a.second.y < b.first.y) vert = b.first.y - a.second.y;</pre>
4fd
80a
        else if (b.second.y < a.first.y) vert = a.first.y - b.second.y;</pre>
96f
        return dist(pt(0, 0), pt(hor, vert));
630 }
13d ld polarea(vector<pt> v) { // area do poligono
9c5
        1d ret = 0;
сбе
        for (int i = 0; i < v.size(); i++)</pre>
            ret += sarea(pt(0, 0), v[i], v[(i + 1) % v.size()]);
80f
d03
        return abs(ret);
5df }
// se o ponto ta dentro do poligono: retorna O se ta fora,
// 1 se ta no interior e 2 se ta na borda
8e7 int inpol(vector\langle pt \rangle \& v, pt p) { // O(n)
8de
        int qt = 0;
        for (int i = 0; i < v.size(); i++) {</pre>
f14
            if (p == v[i]) return 2;
bda
            int j = (i+1)%v.size();
6af
e38
            if (eq(p.y, v[i].y) and eq(p.y, v[j].y)) {
97f
                 if ((v[i]-p)*(v[j]-p) < eps) return 2;
                 continue;
5e2
48b
            }
388
            bool baixo = v[i].y+eps < p.y;</pre>
464
            if (baixo == (v[j].y+eps < p.y)) continue;</pre>
            auto t = (p-v[i])^(v[j]-v[i]);
366
1b4
            if (eq(t, 0)) return 2:
839
            if (baixo == (t > eps)) qt += baixo ? 1 : -1;
d13
b84
        return qt != 0;
a64 }
6ff bool interpol(vector<pt> v1, vector<pt> v2) { // se dois poligonos
   se intersectam - O(n*m)
7d1
        int n = v1.size(), m = v2.size();
        for (int i = 0; i < n; i++) if (inpol(v2, v1[i])) return 1;</pre>
c36
        for (int i = 0; i < n; i++) if (inpol(v1, v2[i])) return 1;
ab8
523
        for (int i = 0; i < n; i++) for (int j = 0; j < m; j++)
0c8
            if (interseg(line(v1[i], v1[(i+1)%n]), line(v2[j],
   v2[(j+1)%m]))) return 1;
        return 0;
bb3
c58 }
494 ld distpol(vector<pt> v1, vector<pt> v2) { // distancia entre
   poligonos
        if (interpol(v1, v2)) return 0;
f6b
```

```
349
        ld ret = DINF;
1c8
        for (int i = 0; i < v1.size(); i++) for (int j = 0; j < v1.size()
   v2.size(); j++)
            ret = min(ret, distseg(line(v1[i], v1[(i + 1) %
6c2
   v1.size()]).
                         line(v2[j], v2[(j + 1) % v2.size()])));
9d9
edf
        return ret;
125 }
138 vector <pt> convex_hull(vector <pt> v) { // convex hull - 0(n log(n))
        sort(v.begin(), v.end());
        v.erase(unique(v.begin(), v.end()), v.end());
d76
52d
        if (v.size() <= 1) return v;</pre>
526
        vector<pt> 1, u;
f14
        for (int i = 0; i < v.size(); i++) {</pre>
            while (l.size() > 1 and !ccw(l.end()[-2], l.end()[-1],
fb2
   v[i]))
364
                l.pop_back();
c35
            l.push_back(v[i]);
58e
        for (int i = v.size() - 1; i >= 0; i--) {
3e9
f19
            while (u.size() > 1 \text{ and } !ccw(u.end()[-2], u.end()[-1],
   v[i]))
7a8
                u.pop_back();
a95
            u.push_back(v[i]);
0b8
        }
cfc
        1.pop_back(); u.pop_back();
82b
        for (pt i : u) l.push_back(i);
792
        return 1:
10d }
483 struct convex_pol {
f50
        vector<pt> pol;
        // nao pode ter ponto colinear no convex hull
d98
        convex_pol() {}
a04
        convex_pol(vector < pt > v) : pol(convex_hull(v)) {}
        // se o ponto ta dentro do hull - O(\log(n))
        bool is_inside(pt p) {
8af
            if (pol.size() == 0) return false;
b6e
            if (pol.size() == 1) return p == pol[0];
eae
67f
            int 1 = 1, r = pol.size();
            while (1 < r) {
40 c
                int m = (1+r)/2:
ee4
```

```
48f
                if (ccw(p, pol[0], pol[m])) 1 = m+1;
ef3
                else r = m:
91c
            }
            if (1 == 1) return isinseg(p, line(pol[0], pol[1]));
00a
9e7
            if (l == pol.size()) return false;
            return !ccw(p, pol[1], pol[1-1]);
1c0
6b0
        }
        // ponto extremo em relacao a cmp(p, q) = p mais extremo q
        // (copiado de https://github.com/gustavoM32/caderno-zika)
        int extreme(const function < bool(pt, pt) > & cmp) {
719
b1c
            int n = pol.size();
4a2
            auto extr = [&](int i, bool& cur_dir) {
22a
                cur dir = cmp(pol[(i+1)%n], pol[i]):
                return !cur_dir and !cmp(pol[(i+n-1)%n], pol[i]);
61a
364
            };
63d
            bool last_dir, cur_dir;
            if (extr(0, last_dir)) return 0;
a0d
            int 1 = 0, r = n;
993
            while (1+1 < r) {
ead
                int m = (1+r)/2:
ee4
                if (extr(m, cur_dir)) return m;
f29
                bool rel_dir = cmp(pol[m], pol[l]);
44a
                if ((!last_dir and cur_dir) or
b18
261
                        (last_dir == cur_dir and rel_dir == cur_dir)) {
8a6
                    1 = m:
1 f 1
                    last dir = cur dir:
94a
                } else r = m;
            }
606
792
            return 1;
56c
316
        int max_dot(pt v) {
ec1
            return extreme([&](pt p, pt q) { return p*v > q*v; });
3b7
        }
a54
        pair < int , int > tangents(pt p) {
            auto L = [\&](pt q, pt r) \{ return ccw(p, r, q); \};
ffb
            auto R = [\&](pt q, pt r) \{ return ccw(p, q, r); \};
8fd
fa8
            return {extreme(L), extreme(R)};
736
        }
3ec };
// CIRCUNFERENCIA
911 pt getcenter(pt a, pt b, pt c) { // centro da circunf dado 3 pontos
174
       b = (a + b) / 2:
2ae
        c = (a + c) / 2;
        return inter(line(b, b + rotate90(a - b)),
98b
3f8
                line(c, c + rotate90(a - c))):
```

```
a12 }
4b3 vector <pt> circ_line_inter(pt a, pt b, pt c, ld r) { // intersecao
   da circunf (c, r) e reta ab
        vector<pt> ret:
8af
        b = b-a, a = a-c;
f2b
4b1
        1d A = b*b:
20a
        1d B = a*b:
2e9
        1d C = a*a - r*r;
        1d D = B*B - A*C:
1fa
818
        if (D < -eps) return ret;</pre>
dc5
        ret.push_back(c+a+b*(-B+sqrt(D+eps))/A);
20e
        if (D > eps) ret.push back(c+a+b*(-B-sqrt(D))/A):
edf
        return ret:
cd8 }
ad2 vector <pt> circ_inter(pt a, pt b, ld r, ld R) { // intersecao da
   circunf (a, r) e (b, R)
        vector<pt> ret;
8af
        1d d = dist(a, b);
b7e
        if (d > r+R \text{ or } d+min(r, R) < max(r, R)) return ret;
5ce
398
        1d x = (d*d-R*R+r*r)/(2*d);
183
        1d y = sqrt(r*r-x*x);
325
        pt v = (b-a)/d;
76e
        ret.push_back(a+v*x + rotate90(v)*y);
2cb
        if (y > 0) ret.push_back(a+v*x - rotate90(v)*y);
edf
        return ret:
fb1 }
6e0 bool operator <(const line& a, const line& b) { // comparador pra
   reta
        // assume que as retas tem p < q</pre>
        pt v1 = a.q - a.p, v2 = b.q - b.p;
a13
        if (!eq(angle(v1), angle(v2))) return angle(v1) < angle(v2);
f82
780
        return ccw(a.p, a.q, b.p); // mesmo angulo
27e }
b14 bool operator ==(const line& a, const line& b) {
76c
        return !(a < b) and !(b < a);
449 }
// comparador pro set pra fazer sweep line com segmentos
2c4 struct cmp_sweepline {
        bool operator () (const line& a, const line& b) const {
d80
            // assume que os segmentos tem p < q
191
            if (a.p == b.p) return ccw(a.p, a.q, b.q);
            if (!eq(a.p.x, a.q.x)) and (eq(b.p.x, b.q.x)) or a.p.x+eps <
231
   b.p.x))
```

```
780
                return ccw(a.p, a.q, b.p);
            return ccw(a.p, b.q, b.p);
dc0
243
367 };
// comparador pro set pra fazer sweep angle com segmentos
bef pt dir:
5b0 struct cmp_sweepangle {
d80
       bool operator () (const line& a, const line& b) const {
522
            return get_t(dir, a) + eps < get_t(dir, b);</pre>
653
       }
97f }:
```

#### 1.5 Geometria - inteiro

```
2de #define sq(x) ((x)*(11)(x))
b2a struct pt { // ponto
        int x, y;
        pt(int x_{-} = 0, int y_{-} = 0) : x(x_{-}), y(y_{-}) {}
df1
        bool operator < (const pt p) const {</pre>
95a
           if (x != p.x) return x < p.x;
89 c
            return y < p.y;</pre>
dcd
       }
a83
        bool operator == (const pt p) const {
d74
            return x == p.x and y == p.y;
7b4
        pt operator + (const pt p) const { return pt(x+p.x, y+p.y); }
cb9
a 24
        pt operator - (const pt p) const { return pt(x-p.x, y-p.y); }
0ef
        pt operator * (const int c) const { return pt(x*c, y*c); }
        11 operator * (const pt p) const { return x*(11)p.x +
   v*(ll)p.v; }
       11 operator ^ (const pt p) const { return x*(11)p.y -
   y*(11)p.x; }
        friend istream& operator >> (istream& in, pt& p) {
5ed
            return in >> p.x >> p.y;
e37
e45
       }
840 };
b3a struct line { // reta
730
       pt p, q;
       line() {}
4b8
       line(pt p_, pt q_) : p(p_), q(q_) {}
8d7
        friend istream& operator >> (istream& in, line& r) {
4cb
            return in >> r.p >> r.q;
858
        }
```

```
7ab }:
// PONTO & VETOR
ea8 ll dist2(pt p, pt q) { // quadrado da distancia
        return sq(p.x - q.x) + sq(p.y - q.y);
515 }
5a2 11 sarea2(pt p, pt q, pt r) { // 2 * area com sinal}
        return (q-p)^(r-q);
bf4 }
e32 bool col(pt p, pt q, pt r) { // se p, q e r sao colin.
        return sarea2(p, q, r) == 0;
034
a08 }
Ocd bool ccw(pt p, pt q, pt r) { // se p, q, r sao ccw
276
        return sarea2(p, q, r) > 0;
42b }
c31 int quad(pt p) { // quadrante de um ponto
        return (p.x<0)^3*(p.y<0);
dbb
fcf }
2df bool compare_angle(pt p, pt q) { // retorna se ang(p) < ang(q)
        if (quad(p) != quad(q)) return quad(p) < quad(q);</pre>
        return ccw(q, pt(0, 0), p);
ea1
771 }
ab1 pt rotate90(pt p) { // rotaciona 90 graus
        return pt(-p.y, p.x);
a0d
e4a }
// RETA
099 bool isinseg(pt p, line r) { // se p pertence ao seg de r
f65
        pt a = r.p - p, b = r.q - p;
2ac
        return (a ^ b) == 0 and (a * b) <= 0;
c9f }
676 bool interseg(line r, line s) { // se o seg de r intersecta o seg
19b
        if (isinseg(r.p, s) or isinseg(r.q, s)
            or isinseg(s.p, r) or isinseg(s.q, r)) return 1;
c21
        return ccw(r.p, r.q, s.p) != ccw(r.p, r.q, s.q) and
9fa
413
                ccw(s.p, s.q, r.p) != ccw(s.p, s.q, r.q);
```

```
359 }
9e0 int segpoints(line r) { // numero de pontos inteiros no segmento
        return 1 + \_gcd(abs(r.p.x - r.q.x), abs(r.p.y - r.q.y));
9ce
dd8 }
88a double get_t(pt v, line r) { // retorna t tal que t*v pertence a
   reta r
        return (r.p^r.q) / (double) ((r.p-r.q)^v);
1ad
d27 }
// POLIGONO
// quadrado da distancia entre os retangulos a e b (lados paralelos
   aos eixos)
// assume que ta representado (inferior esquerdo, superior direito)
485 ll dist2_rect(pair<pt, pt> a, pair<pt, pt> b) {
        int hor = 0, vert = 0;
c59
        if (a.second.x < b.first.x) hor = b.first.x - a.second.x;</pre>
34b
        else if (b.second.x < a.first.x) hor = a.first.x - b.second.x:</pre>
f5f
        if (a.second.y < b.first.y) vert = b.first.y - a.second.y;</pre>
4fd
80a
        else if (b.second.y < a.first.y) vert = a.first.y - b.second.y;</pre>
        return sq(hor) + sq(vert);
869
e13 }
9c3 ll polarea2(vector<pt> v) { // 2 * area do poligono
b73
        ll ret = 0;
сбе
        for (int i = 0: i < v.size(): i++)</pre>
            ret += sarea2(pt(0, 0), v[i], v[(i + 1) % v.size()]);
532
d03
        return abs(ret);
d5f }
// se o ponto ta dentro do poligono: retorna O se ta fora,
// 1 se ta no interior e 2 se ta na borda
8e7 int inpol(vector<pt>&v, pt p) { // O(n)
8de
        int qt = 0;
        for (int i = 0; i < v.size(); i++) {</pre>
f14
bda
            if (p == v[i]) return 2;
6af
            int j = (i+1)%v.size();
cc6
            if (p.v == v[i].v \text{ and } p.v == v[i].v) {
                if ((v[i]-p)*(v[j]-p) <= 0) return 2;</pre>
547
5e2
                 continue;
b47
            }
78c
            bool baixo = v[i].y < p.y;</pre>
057
            if (baixo == (v[j].y < p.y)) continue;</pre>
            auto t = (p-v[i])^(v[j]-v[i]);
366
            if (!t) return 2:
2ad
```

```
0bb
            if (baixo == (t > 0)) qt += baixo ? 1 : -1;
9cf
b84
        return qt != 0;
afd }
138 vector <pt > convex_hull(vector <pt > v) { // convex hull - O(n log(n))
fca
        sort(v.begin(), v.end());
d76
        v.erase(unique(v.begin(), v.end()), v.end());
52d
        if (v.size() <= 1) return v;</pre>
        vector<pt> 1, u;
526
f14
        for (int i = 0; i < v.size(); i++) {</pre>
fb2
            while (1.size() > 1 \text{ and } !ccw(1.end()[-2], 1.end()[-1],
   v[i]))
364
                1.pop_back();
c35
            l.push_back(v[i]);
58e
3e9
        for (int i = v.size() - 1; i >= 0; i--) {
            while (u.size() > 1 \text{ and } !ccw(u.end()[-2], u.end()[-1],
f19
   v[i]))
7a8
                 u.pop_back();
a95
            u.push_back(v[i]);
0b8
cfc
        1.pop_back(); u.pop_back();
82b
        for (pt i : u) l.push_back(i);
792
        return 1:
10d }
786 ll interior_points(vector<pt> v) { // pontos inteiros dentro de um
   poligono simples
        11 b = 0;
c4e
        for (int i = 0; i < v.size(); i++)</pre>
c6e
Осе
            b += segpoints(line(v[i], v[(i+1)\%v.size()])) - 1;
        return (polarea2(v) - b) / 2 + 1;
a1c
af2 }
483 struct convex_pol {
        vector<pt> pol;
f50
        // nao pode ter ponto colinear no convex hull
d98
        convex_pol() {}
        convex_pol(vector < pt > v) : pol(convex_hull(v)) {}
a 04
        // se o ponto ta dentro do hull - O(\log(n))
8af
        bool is_inside(pt p) {
b6e
            if (pol.size() == 0) return false;
            if (pol.size() == 1) return p == pol[0];
eae
            int 1 = 1, r = pol.size();
67f
```

```
40c
             while (1 < r) {
ee4
                 int m = (1+r)/2;
48f
                 if (ccw(p, pol[0], pol[m])) 1 = m+1;
ef3
                 else r = m;
91c
            if (1 == 1) return isinseg(p, line(pol[0], pol[1]));
00a
9e7
            if (l == pol.size()) return false;
1c0
            return !ccw(p, pol[1], pol[1-1]);
6b0
        // ponto extremo em relacao a cmp(p, q) = p mais extremo q
        // (copiado de https://github.com/gustavoM32/caderno-zika)
719
        int extreme(const function < bool(pt, pt) > & cmp) {
b1c
            int n = pol.size();
             auto extr = [&](int i, bool& cur_dir) {
4a2
22a
                 \operatorname{cur\_dir} = \operatorname{cmp}(\operatorname{pol}[(i+1)\%n], \operatorname{pol}[i]);
                 return !cur_dir and !cmp(pol[(i+n-1)%n], pol[i]);
61a
364
            };
            bool last_dir, cur_dir;
63d
            if (extr(0, last_dir)) return 0;
a0d
            int 1 = 0, r = n:
993
             while (1+1 < r) {
                 int m = (1+r)/2:
ee4
                 if (extr(m, cur_dir)) return m;
f29
                 bool rel_dir = cmp(pol[m], pol[l]);
44a
b18
                 if ((!last dir and cur dir) or
261
                         (last dir == cur dir and rel dir == cur dir)) {
8a6
                     1 = m:
1 f 1
                     last dir = cur dir:
                 } else r = m;
94a
606
792
            return 1:
56c
316
        int max_dot(pt v) {
ec1
             return extreme([&](pt p, pt q) { return p*v > q*v; });
3b7
a54
        pair < int , int > tangents(pt p) {
ffb
             auto L = [\&](pt q, pt r) \{ return ccw(p, r, q); \};
8fd
             auto R = [k](pt q, pt r) \{ return ccw(p, q, r); \};
fa8
            return {extreme(L), extreme(R)};
736
        }
3ec }:
6e0 bool operator <(const line& a, const line& b) { // comparador pra
   reta
        // assume que as retas tem p < q</pre>
        pt v1 = a.q - a.p, v2 = b.q - b.p;
a13
036
        bool b1 = compare_angle(v1, v2), b2 = compare_angle(v2, v1);
```

```
73c
        if (b1 or b2) return b1;
780
        return ccw(a.p, a.q, b.p); // mesmo angulo
b61 }
b14 bool operator ==(const line& a, const line& b) {
76c
        return !(a < b) and !(b < a):
449 }
// comparador pro set pra fazer sweep line com segmentos
2c4 struct cmp_sweepline {
        bool operator () (const line& a, const line& b) const {
            // assume que os segmentos tem p < q
191
            if (a.p == b.p) return ccw(a.p, a.q, b.q);
614
            if (a.p.x != a.g.x and (b.p.x == b.g.x or a.p.x < b.p.x))
780
                return ccw(a.p, a.q, b.p);
dc0
            return ccw(a.p, b.q, b.p);
baf
677 };
// comparador pro set pra fazer sweep angle com segmentos
bef pt dir:
5b0 struct cmp_sweepangle {
d80
        bool operator () (const line& a, const line& b) const {
            return get_t(dir, a) < get_t(dir, b);</pre>
261
dc5
        }
f6d }:
1.6 Geometria 3D
c83 typedef double ld;
e3b const ld DINF = 1e18;
107 const ld eps = 1e-9;
b32 #define sq(x) ((x)*(x))
d97 bool eq(ld a, ld b) {
ba0
            return abs(a - b) <= eps;</pre>
bfc }
b2a struct pt { // ponto
```

 $pt(1d x_{-} = 0, 1d y_{-} = 0, 1d z_{-} = 0) : x(x_{-}), y(y_{-}), z(z_{-})$ 

if (!eq(x, p.x)) return x < p.x;

if (!eq(y, p.y)) return y < p.y;

if (!eq(z, p.z)) return z < p.z;

bool operator < (const pt p) const {</pre>

a50

5bc

059

f98

44c

{}

ld x, y, z;

```
bb3
                    return 0;
6cd
a83
            bool operator == (const pt p) const {
41c
                    return eq(x, p.x) and eq(y, p.y) and eq(z, p.z);
fb5
            pt operator + (const pt p) const { return pt(x+p.x, y+p.y,
   z+p.z); }
            pt operator - (const pt p) const { return pt(x-p.x, y-p.y,
   z-p.z); }
            pt operator * (const ld c) const { return pt(x*c , y*c ,
   z*c ); }
           pt operator / (const ld c) const { return pt(x/c , y/c ,
   z/c ): }
a65
            ld operator * (const pt p) const { return x*p.x + y*p.y +
   z*p.z; }
            pt operator ^ (const pt p) const { return pt(y*p.z -
   z*p.v, z*p.x - x*p.z, x*p.v - v*p.x); }
            friend istream& operator >> (istream& in, pt& p) {
5ed
                    return in >> p.x >> p.y >> p.z;
9bf
5e8
           }
3ee }:
b3a struct line { // reta
730
           pt p, q;
046
           line() {}
4b8
           line(pt p_, pt q_) : p(p_), q(q_) {}
            friend istream& operator >> (istream& in, line& r) {
8d7
4cb
                   return in >> r.p >> r.q:
858
            }
7ab };
79b struct plane { // plano
            array<pt, 3> p; // pontos que definem o plano
7e1
            arrav<ld. 4> eq: // equação do plano
29b
            plane() {}
bb7
fb0
            plane(pt p_, pt q_, pt r_) : p({p_, q_, r_}) { build(); }
            friend istream& operator >> (istream& in, plane& P) {
ca9
2ab
                    return in >> P.p[0] >> P.p[1] >> P.p[2];
70e
                    P.build();
544
            }
0a8
            void build() {
da2
                    pt dir = (p[1] - p[0]) ^ (p[2] - p[0]);
7d5
                    eq = \{dir.x, dir.y, dir.z, dir*p[0]*(-1)\};
41a
            }
d5d }:
```

```
// converte de coordenadas polares para cartesianas
// (angulos devem estar em radianos)
// phi eh o angulo com o eixo z (cima) theta eh o angulo de rotacao ao
 redor de z
2fb pt convert(ld rho, ld th, ld phi) {
            return pt(sin(phi) * cos(th), sin(phi) * sin(th),
   cos(phi)) * rho:
// projecao do ponto p na reta r
256 pt proj(pt p, line r) {
            if (r.p == r.q) return r.p;
bea
97a
            r.q = r.q - r.p; p = p - r.p;
9f8
            pt proj = r.q * ((p*r.q) / (r.q*r.q));
2cd
            return proj + r.p;
232 }
// projecao do ponto p no plano P
b1a pt proj(pt p, plane P) {
            p = p - P.p[0], P.p[1] = P.p[1] - P.p[0], P.p[2] = P.p[2]
   - P.p[0];
b69
            pt norm = P.p[1] ^ P.p[2];
            pt proj = p - (norm * (norm * p) / (norm*norm));
6ab
467
            return proj + P.p[0];
4a0 }
// distancia
a45 ld dist(pt a, pt b) {
            return sqrt(sq(a.x-b.x) + sq(a.y-b.y) + sq(a.z-b.z));
2d0 }
// distancia ponto reta
137 ld distline(pt p, line r) {
            return dist(p, proj(p, r));
ce1
3c4 }
// distancia de ponto para segmento
d43 ld distseg(pt p, line r) {
73d
            if ((r.q - r.p)*(p - r.p) < 0) return dist(r.p, p);
951
            if ((r.p - r.q)*(p - r.q) < 0) return dist(r.q, p);
200
            return distline(p, r):
42c }
// distancia de ponto a plano com sinal
7cc ld sdist(pt p, plane P) {
150
            return P.eq[0]*p.x + P.eq[1]*p.y + P.eq[2]*p.z + P.eq[3];
d49 }
```

```
// distancia de ponto a plano
768 ld distplane(pt p, plane P) {
            return abs(sdist(p, P));
33d }
// se ponto pertence a reta
099 bool isinseg(pt p, line r) {
a32
            return eq(distseg(p, r), 0);
31a }
// se ponto pertence ao triangulo definido por P.p
cd2 bool isinpol(pt p. vector<pt> v) {
            assert(v.size() >= 3);
fad
            pt norm = (v[1]-v[0]) ^ (v[2]-v[1]);
bf4
            bool inside = true;
8a4
           int sign = -1;
cec
           for (int i = 0; i < v.size(); i++) {</pre>
f14
                    line r(v[(i+1)\%3], v[i]);
834
2a9
                    if (isinseg(p, r)) return true;
4ef
                    pt ar = v[(i+1)\%3] - v[i];
                    if (sign == -1) sign = ((ar^(p-v[i]))*norm > 0);
320
                    else if (((ar^(p-v[i]))*norm > 0) != sign) inside
82b
   = false:
15e
aca
            return inside;
c81 }
// distancia de ponto ate poligono
361 ld distpol(pt p, vector<pt> v) {
3e7
            pt p2 = proj(p, plane(v[0], v[1], v[2]);
            if (isinpol(p2, v)) return dist(p, p2);
61a
349
            ld ret = DINF:
            for (int i = 0; i < v.size(); i++) {</pre>
f14
6af
                   int j = (i+1)%v.size();
5ee
                   ret = min(ret, distseg(p, line(v[i], v[j])));
7b2
            }
edf
            return ret;
a8d }
// intersecao de plano e segmento
// BOTH = o segmento esta no plano
// ONE = um dos pontos do segmento esta no plano
// PARAL = segmento paralelo ao plano
// CONCOR = segmento concorrente ao plano
e51 enum RETCODE {BOTH, ONE, PARAL, CONCOR};
```

```
26b pair < RETCODE, pt > intersect(plane P, line r) {
        1d d1 = sdist(r.p, P);
f8f
        1d d2 = sdist(r.q, P);
53a
        if (eq(d1, 0) \text{ and } eq(d2, 0))
504
                     return pair(BOTH, r.p);
72c
         if (eq(d1, 0))
847
                     return pair(ONE, r.p);
485
        if (eq(d2, 0))
168
                     return pair(ONE, r.q);
3fb
         if ((d1 > 0 \text{ and } d2 > 0) \text{ or } (d1 < 0 \text{ and } d2 < 0)) {}
463
             if (eq(d1-d2, 0)) return pair(PARAL, pt());
406
             return pair(CONCOR, pt());
91c
c84
        1d frac = d1 / (d1 - d2);
3ff
        pt res = r.p + ((r.q - r.p) * frac);
394
        return pair(ONE, res);
b92 }
// rotaciona p ao redor do eixo u por um angulo a
787 pt rotate(pt p, pt u, ld a) {
             u = u / dist(u, pt());
773
             return u * (u * p) + (u ^ p ^ u) * cos(a) + (u ^ p) *
e6f
   sin(a):
7f0 }
1.7 Matriz
945 #define MODULAR false
5ed template < typename T > struct matrix : vector < vector < T >> {
        int n, m;
30f
        void print() {
603
             for (int i = 0; i < n; i++) {</pre>
70f
                 for (int j = 0; j < m; j++) cout << (*this)[i][j] << "</pre>
1fb
                 cout << endl:</pre>
d98
            }
        }
101
         matrix(int n_, int m_, bool ident = false) :
aa3
b14
                 vector < vector < T > (n_, vector < T > (m_, 0)), n(n_), m(m_)  {
94 e
            if (ident) {
df7
                 assert(n == m);
a89
                 for (int i = 0; i < n; i++) (*this)[i][i] = 1;
359
            }
527
        }
```

```
b83
        matrix(const vector<vector<T>>& c) : vector<vector<T>>(c),
            n(c.size()), m(c[0].size()) {}
a3d
efc
        matrix(const initializer_list<initializer_list<T>>& c) {
f7e
            vector < vector < T >> val;
212
            for (auto& i : c) val.push_back(i);
            *this = matrix(val);
303
c50
        }
        matrix<T> operator*(matrix<T>& r) {
388
            assert(m == r.n):
1e2
82c
            matrix <T> M(n, r.m);
d69
            for (int i = 0; i < n; i++) for (int k = 0; k < m; k++)
df4
                for (int j = 0; j < r.m; j++) {</pre>
e34
                    T \text{ add} = (*this)[i][k] * r[k][j];
f98 #if MODULAR
d41 #warning Usar matrix<ll> e soh colocar valores em [0, MOD) na
   matriz!
868
                    M[i][j] += add%MOD;
                    if (M[i][j] >= MOD) M[i][j] -= MOD;
983
8c1 #else
7bb
                    M[i][j] += add;
f2e #endif
                }
620
474
            return M;
394
528
        matrix<T> operator^(ll e){
f10
            matrix<T> M(n, n, true), at = *this;
c87
            while (e) {
2e2
                if (e&1) M = M*at;
cc2
                e >>= 1;
c80
                at = at*at:
            }
eb6
474
            return M;
ca3
        }
582
        void apply_transform(matrix M, ll e){
1c3
            auto& v = *this;
c87
            while (e) {
                if (e\&1) v = M*v;
9ba
                e >>= 1:
cc2
419
                M = M * M;
d86
            }
4e5
        }
70d };
```

#### 1.8 Matroid

```
// Matroids de Grafo e Particao
// De modo geral, toda Matroid contem um build() linear
// e uma funcao constante oracle()
// oracle(i) responde se o conjunto continua independente
// apos adicao do elemento i
// oracle(i, j) responde se o conjunto continua indepente
// apos trocar o elemento i pelo elemento j
//
// Intersecao sem peso O(r^2 n)
// em que n eh o tamanho do conjunto e r eh o tamanho da resposta
// Matroid Grafica
// Matroid das florestas de um grafo
// Um conjunto de arestas eh independente se formam uma floresta
//
// build() : O(n)
// oracle() : 0(1)
fda struct graphic_matroid {
5da
        int n, m, t;
        vector < array < int , 2>> edges;
32c
789
        vector < vector < int >> g;
62e
        vector < int > comp, in, out;
513
        graphic_matroid(int n_, vector<array<int, 2>> edges_)
            : n(n_), m(edges_.size()), edges(edges_), g(n), comp(n),
a1f
   in(n), out(n) {}
315
        void dfs(int u) {
ab8
            in[u] = t++:
17d
            for (auto v : g[u]) if (in[v] == -1)
863
                comp[v] = comp[u], dfs(v);
677
            out[u] = t:
d83
        }
        void build(vector<int> I) {
945
a34
            t = 0:
741
            for (int u = 0; u < n; u++) g[u].clear(), in[u] = -1;
667
            for (int e : I) {
d00
                auto [u, v] = edges[e];
125
                g[u].push_back(v), g[v].push_back(u);
a8a
809
            for (int u = 0; u < n; u++) if (in[u] == -1)
                comp[u] = u, dfs(u);
a7d
207
        bool is_ancestor(int u, int v) {
f31
a68
            return in[u] <= in[v] and in[v] < out[u];</pre>
0c2
        }
e6b
        bool oracle(int e) {
453
            return comp[edges[e][0]] != comp[edges[e][1]];
```

```
687
        bool oracle(int e, int f) {
f75
574
            if (oracle(f)) return true;
622
            int u = edges[e][in[edges[e][0]] < in[edges[e][1]]];</pre>
            return is_ancestor(u, edges[f][0]) != is_ancestor(u,
ff2
   edges[f][1]);
8a9
      }
691 }:
// Matroid de particao ou cores
// Um conjunto eh independente se a quantidade de elementos
// de cada cor nao excede a capacidade da cor
// Quando todas as capacidades sao 1, um conjunto eh independente
// se todas as suas cores sao distintas
// build() : O(n)
// oracle() : 0(1)
994 struct partition_matroid {
        vector < int > cap, color, d;
501
608
        partition_matroid(vector<int> cap_, vector<int> color_)
04d
            : cap(cap_), color(color_), d(cap.size()) {}
        void build(vector<int> I) {
945
def
            fill(d.begin(), d.end(), 0);
e9d
            for (int u : I) d[color[u]]++;
c58
514
        bool oracle(int u) {
            return d[color[u]] < cap[color[u]];</pre>
0a1
703
f7f
        bool oracle(int u, int v) {
2f7
            return color[u] == color[v] or oracle(v);
4b4
       }
caa }:
// Intersecao de matroid sem pesos
// Dadas duas matroids M1 e M2 definidas sobre o mesmo
// conjunto I, retorna o maior subconjunto de I
// que eh independente tanto para M1 quanto para M2
// O(r^2*n)
// Matroid "pesada" deve ser a M2
132 template < typename Matroid1, typename Matroid2 >
801 vector < int > matroid_intersection(int n, Matroid1 M1, Matroid2 M2) {
f5b
        vector < bool > b(n);
        vector < int > I[2]:
a64
        bool converged = false;
a8b
```

```
0 c 1
        while (!converged) {
742
            I[0].clear(), I[1].clear();
99d
            for (int u = 0; u < n; u++) I[b[u]].push_back(u);</pre>
09d
            M1.build(I[1]), M2.build(I[1]);
289
            vector < bool > target(n), pushed(n);
26a
            queue < int > q:
5 c 5
            for (int u : I[0]) {
2b2
                 target[u] = M2.oracle(u);
                 if (M1.oracle(u)) pushed[u] = true, q.push(u);
c1b
0e6
3fe
            vector < int > p(n, -1);
07a
            converged = true;
402
            while (q.size()) {
be1
                int u = q.front(); q.pop();
5c6
                if (target[u]) {
101
                     converged = false;
                     for (int v = u; v != -1; v = p[v]) b[v] = !b[v];
c32
c2b
                     break;
a80
                }
e78
                for (int v : I[!b[u]]) if (!pushed[v]) {
34d
                     if ((b[u] and M1.oracle(u, v)) or (b[v] and
   M2.oracle(v, u)))
                         p[v] = u, pushed[v] = true, q.push(v);
bae
533
                }
1d9
            }
        }
5e7
b68
        return I[1];
381 }
// Intersecao de matroid com pesos
// Dadas duas matroids M1 e M2 e uma funcao de pesos w, todas
// um conjunto I retorna o maior subconjunto de I (desempatado pelo
   menor peso)
// que eh independente tanto para M1 quanto para M2
// A resposta eh construida incrementando o tamanho conjunto I de 1 em
// Se nao tiver custo negativo, nao precisa de SPFA
//
// O(r^3*n) com SPFA
// O(r^2*n*log(n)) com Dijkstra e potencial
42a template < typename T, typename Matroid1, typename Matroid2>
2b5 vector < int > weighted_matroid_intersection(int n, vector < T > w,
    Matroid1 M1, Matroid2 M2) {
        vector < bool > b(n), target(n), is_inside(n);
6c9
```

```
563
        vector < int > I[2], from(n);
        vector < pair < T, int >> d(n);
e35
169
        auto check_edge = [&](int u, int v) {
            return (b[u] and M1.oracle(u, v)) or (b[v] and
   M2.oracle(v, u));
253
        };
667
        while (true) {
742
            I[0].clear(), I[1].clear();
99d
            for (int u = 0; u < n; u++) I[b[u]].push_back(u);
            // I[1] contem o conjunto de tamanho I[1].size() de menor
                peso
            M1.build(I[1]), M2.build(I[1]);
687
            for (int u = 0: u < n: u++) {
                target[u] = false, is_inside[u] = false, from[u] = -1;
ea5
                d[u] = {numeric_limits <T>::max(), INF};
961
392
            }
8d3
            deque <T > q;
            sort(I[0].begin(), I[0].end(), [&](int i, int j){ return
476
   w[i] < w[i]; \});
            for (int u : I[0]) {
5 c 5
                target[u] = M2.oracle(u);
2b2
                if (M1.oracle(u)) {
5a7
                    if (is_inside[u]) continue;
4ef
                    d[u] = \{w[u], 0\};
7cc
                    if (!q.empty() and d[u] > d[q.front()])
   q.push_back(u);
655
                     else q.push_front(u);
4ae
                     is inside[u] = true:
                }
764
            }
add
402
            while (q.size()) {
97a
                int u = q.front(); q.pop_front();
                is_inside[u] = false;
6f3
57a
                for (int v : I[!b[u]]) if (check edge(u, v)) {
                     pair < T, int > nd(d[u].first + w[v], d[u].second +
9de
   1);
                    if (nd < d[v]) {</pre>
61b
                         from[v] = u, d[v] = nd;
6ac
bd7
                         if (is_inside[v]) continue;
                         if (q.size() and d[v] > d[q.front()])
   q.push_back(v);
                         else q.push_front(v);
275
                         is_inside[v] = true;
587
                    }
b3f
a3b
                }
            }
563
            pair <T, int > mini = pair(numeric_limits <T >:: max(), INF);
cc8
```

```
489
            int targ = -1;
259
            for (int u : I[0]) if (target[u] and d[u] < mini)</pre>
2b9
                mini = d[u], targ = u;
            if (targ != -1) for (int u = targ; u != -1; u = from[u])
e14
d89
                b[u] = !b[u], w[u] *= -1:
f97
            else break;
c7d
        }
b68
        return I[1];
8e7 }
```

#### $^{2}$ DP

#### 2.1 Divide and Conquer DP

```
// Particiona o arrav em k subarravs
// minimizando o somatorio das queries
// O(k n log n), assumindo quer query(1, r) eh O(1)
547 ll dp[MAX][2];
94b void solve(int k, int l, int r, int lk, int rk) {
de6
        if (1 > r) return;
109
        int m = (1+r)/2, p = -1;
        auto& ans = dp[m][k&1] = LINF;
d2b
        for (int i = max(m, lk): i <= rk: i++) {</pre>
6e2
            ll at = dp[i+1][\sim k\&1] + query(m, i);
07b
57d
            if (at < ans) ans = at, p = i;</pre>
8f5
1ee
        solve(k, l, m-1, lk, p), solve(k, m+1, r, p, rk);
d3e }
cf1 ll DC(int n, int k) {
321
        dp[n][0] = dp[n][1] = 0;
f27
        for (int i = 0; i < n; i++) dp[i][0] = LINF;</pre>
b76
        for (int i = 1: i \le k: i++) solve(i, 0, n-i, 0, n-i):
8e7
        return dp[0][k&1];
5e9 }
```

## 2.2 Longest Common Subsequence

// Computa a LCS entre dois arrays usando

```
// o algoritmo de Hirschberg para recuperar
// O(n*m), O(n+m) de memoria
eaf int lcs_s[MAX], lcs_t[MAX];
a6d int dp[2][MAX];
// dp[0][j] = max lcs(s[li...ri], t[lj, lj+j])
d12 void dp_top(int li, int ri, int lj, int rj) {
        memset(dp[0], 0, (rj-lj+1)*sizeof(dp[0][0]));
753
        for (int i = li; i <= ri; i++) {</pre>
9aa
            for (int j = rj; j >= lj; j--)
83b
                dp[0][i - 1i] = max(dp[0][i - 1i],
741
                (lcs_s[i] == lcs_t[j]) + (j > 1j ? dp[0][j-1 - 1j] :
   0));
            for (int j = lj+1; j <= rj; j++)</pre>
04c
939
                dp[0][i - 1i] = max(dp[0][i - 1i], dp[0][i-1 -1i]);
09f
        }
58f }
// dp[1][j] = max lcs(s[li...ri], t[lj+j, rj])
ca0 void dp_bottom(int li, int ri, int lj, int rj) {
0dd
        memset(dp[1], 0, (rj-lj+1)*sizeof(dp[1][0]));
3a2
        for (int i = ri; i >= li; i--) {
49c
            for (int j = lj; j <= rj; j++)</pre>
dbb
                dp[1][j - lj] = max(dp[1][j - lj],
                (lcs_s[i] == lcs_t[j]) + (j < rj ? dp[1][j+1 - lj] :
4da
   0)):
            for (int j = rj-1; j >= 1j; j--)
6ca
                dp[1][j - 1j] = max(dp[1][j - 1j], dp[1][j+1 - 1j]);
769
19b
        }
e8a }
93c void solve(vector<int>& ans. int li. int ri. int li. int ri) {
        if (li == ri){
2ad
49c
            for (int j = lj; j <= rj; j++)</pre>
f5b
                if (lcs_s[li] == lcs_t[j]){
a66
                    ans.push_back(lcs_t[i]);
c2b
                    break:
840
                }
505
            return;
126
534
        if (li == ri){
753
            for (int i = li; i <= ri; i++){</pre>
88f
                if (lcs_s[i] == lcs_t[li]){
                    ans.push_back(lcs_s[i]);
531
c2b
                    break:
```

```
68a
                }
a03
            }
505
            return;
76d
        }
        int mi = (li+ri)/2;
a57
        dp_top(li, mi, lj, rj), dp_bottom(mi+1, ri, lj, rj);
ade
d7a
        int i_{-} = 0, mx = -1;
        for (int j = lj-1; j <= rj; j++) {
aee
da8
            int val = 0;
2bb
            if (j >= lj) val += dp[0][j - lj];
b9e
            if (i < ri) val += dp[1][i+1 - li];</pre>
ba8
            if (val >= mx) mx = val, j_ = j;
14e
        }
6f1
        if (mx == -1) return;
        solve(ans, li, mi, lj, j_), solve(ans, mi+1, ri, j_+1, rj);
c2a
dd5 }
058 vector<int> lcs(const vector<int>& s, const vector<int>& t) {
        for (int i = 0; i < s.size(); i++) lcs_s[i] = s[i];
953
577
        for (int i = 0; i < t.size(); i++) lcs_t[i] = t[i];</pre>
dab
        vector < int > ans;
599
        solve(ans, 0, s.size()-1, 0, t.size()-1);
ba7
        return ans:
17c }
2.3 Mochila
// Resolve mochila, recuperando a resposta
//
// O(n * cap), O(n + cap) de memoria
add int v[MAX], w[MAX]; // valor e peso
582 int dp[2][MAX_CAP];
// DP usando os itens [1, r], com capacidade = cap
0d6 void get_dp(int x, int 1, int r, int cap) {
```

memset(dp[x], 0, (cap+1)\*sizeof(dp[x][0]));

| 5ab void solve(vector<int>& ans, int 1, int r, int cap) {

for (int i = 1;  $i \le r$ ; i++) for (int j = cap;  $j \ge 0$ ; j--)

if (i - w[i] >= 0) dp[x][j] = max(dp[x][j], v[i] + dp[x][j]

574

b95 }

- w[i]]);

```
893
        if (1 == r) {
            if (w[1] <= cap) ans.push_back(1);</pre>
9ff
505
            return;
        }
13a
        int m = (1+r)/2;
ee4
283
        get_dp(0, 1, m, cap), get_dp(1, m+1, r, cap);
056
        int left_cap = -1, opt = -INF;
c94
        for (int j = 0; j <= cap; j++)
2f2
            if (int at = dp[0][j] + dp[1][cap - j]; at > opt)
91d
                opt = at, left_cap = j;
da3
        solve(ans, 1, m, left_cap), solve(ans, m+1, r, cap - left_cap);
d75 }
0d7 vector<int> knapsack(int n, int cap) {
dab
        vector < int > ans;
1e0
        solve(ans, 0, n-1, cap);
ba7
        return ans;
e4d }
2.4 SOS DP
// \Omega(n 2^n)
// soma de sub-conjunto
e03 vector<ll> sos dp(vector<ll> f) {
       int N = __builtin_ctz(f.size());
        assert((1<<N) == f.size());
e59
       for (int i = 0; i < N; i++) for (int mask = 0; mask < (1<<N);
5a5
   mask++)
796
            if (mask>>i&1) f[mask] += f[mask^(1<<i)];</pre>
abe
        return f:
bec }
// soma de super-conjunto
e03 vector<ll> sos_dp(vector<ll> f) {
       int N = __builtin_ctz(f.size());
6c0
```

for (int i = 0; i < N; i++) for (int mask = 0; mask < (1<<N);

if  $(\sim mask >> i\&1)$  f[mask] += f[mask^(1<<ii)];

assert((1<<N) == f.size());

e59

5a5

a3c abe

dbd }

mask++)

return f:

#### 2.5 Subset sum

```
// Retorna max(x <= t tal que existe subset de w que soma x)
// O(n * max(w))
// O(max(w)) de memoria
efd int subset_sum(vector<int> w, int t) {
bb5
        int pref = 0, k = 0;
417
        while (k < w.size() and pref + w[k] <= t) pref += w[k++];</pre>
1e7
        if (k == w.size()) return pref;
        int W = *max_element(w.begin(), w.end());
444
44d
        vector \langle int \rangle last, dp(2*W, -1);
d7b
        dp[W - (t-pref)] = k;
        for (int i = k; i < w.size(); i++) {</pre>
54d
288
            last = dp:
15f
            for (int x = 0; x < W; x++) dp[x+w[i]] = max(dp[x+w[i]],
   last[x]);
17b
            for (int x = 2*W - 1; x > W; x--)
303
                 for (int j = max(0, last[x]); j < dp[x]; j++)</pre>
                     dp[x-w[j]] = max(dp[x-w[j]], j);
595
867
        }
2fb
        int ans = t;
1c1
        while (dp[W - (t-ans)] < 0) ans --;
ba7
        return ans;
d88 }
```

## 3 Matematica

#### 3.1 2-SAT

```
// solve() retorna um par, o first fala se eh possivel
// atribuir, o second fala se cada variavel eh verdadeira
//
// O(|V|+|E|) = O(#variaveis + #restricoes)

138 struct sat {
    e6c     int n, tot;
    789     vector<vector<int>> g;
    Oca     vector<int> vis, comp, id, ans;
4ce     stack<int> s;

141     sat() {}
    172     sat(int n_) : n(n_), tot(n), g(2*n) {}
```

```
f32
        int dfs(int i, int& t) {
cf0
            int lo = id[i] = t++;
            s.push(i), vis[i] = 2;
efc
            for (int j : g[i]) {
48e
                if (!vis[j]) lo = min(lo, dfs(j, t));
740
994
                else if (vis[j] == 2) lo = min(lo, id[j]);
            }
d64
3de
            if (lo == id[i]) while (1) {
                int u = s.top(); s.pop();
3c3
9c5
                vis[u] = 1, comp[u] = i;
91d
                if ((u>1) < n \text{ and } ans[u>1] == -1) ans[u>1] = \sim u\&1:
2ef
                if (u == i) break;
60d
            }
253
            return lo;
dec
        }
        void add_impl(int x, int y) { // x -> y = !x ou y
74a
            x = x >= 0 ? 2*x : -2*x-1;
26a
2b8
            y = y >= 0 ? 2*y : -2*y-1;
a1e
            g[x].push_back(y);
            g[y^1].push_back(x^1);
1e2
ef0
e85
        void add_cl(int x, int y) { // x ou y
0b5
            add_impl(\sim x, y);
254
487
        void add_xor(int x, int y) { // x xor y
0b7
            add_cl(x, y), add_cl(\simx, \simy);
9a1
978
        void add_eq(int x, int y) { // x = y
c86
            add_xor(\simx, y);
b91
b10
        void add_true(int x) { // x = T
18b
            add_impl(\sim x, x);
9e2
d14
        void at_most_one(vector<int> v) { // no max um verdadeiro
54d
            g.resize(2*(tot+v.size()));
f14
            for (int i = 0; i < v.size(); i++) {</pre>
8c9
                add_impl(tot+i, ~v[i]);
a8f
                if (i) {
b6a
                     add_impl(tot+i, tot+i-1);
                     add_impl(v[i], tot+i-1);
3d3
0f7
                }
            }
084
258
            tot += v.size();
        }
b00
```

```
a8e
         pair < bool, vector < int >> solve() {
27b
             ans = vector < int > (n, -1);
6bb
             int t = 0;
0de
             vis = comp = id = vector\langle int \rangle (2*tot, 0);
             for (int i = 0; i < 2*tot; i++) if (!vis[i]) dfs(i, t);</pre>
53c
             for (int i = 0; i < tot; i++)</pre>
f88
4c9
                  if (comp[2*i] == comp[2*i+1]) return {false, {}};
997
             return {true, ans};
7b3
ef6 };
```

## 3.2 Avaliacao de Interpolacao

```
// Dado 'n' pontos (i, y[i]), i \in [0, n),
// avalia o polinomio de grau n-1 que passa
// por esses pontos em 'x'
// Tudo modular, precisa do mint
//
// O(n)
ee8 mint evaluate_interpolation(int x, vector<mint> y) {
        int n = y.size();
184
        vector < mint > sulf(n+1, 1), fat(n, 1), ifat(n);
        for (int i = n-1; i \ge 0; i--) sulf[i] = sulf[i+1] * (x - i);
6fa
29b
        for (int i = 1; i < n; i++) fat[i] = fat[i-1] * i;</pre>
        ifat[n-1] = 1/fat[n-1];
0da
        for (int i = n-2: i >= 0: i--) ifat[i] = ifat[i+1] * (i + 1):
3db
ca1
        mint pref = 1, ans = 0;
5ea
        for (int i = 0; i < n; pref *= (x - i++)) {</pre>
42f
            mint num = pref * sulf[i+1];
b4e
            mint den = ifat[i] * ifat[n-1 - i];
0bd
            if ((n-1 - i)\%2) den *= -1:
03f
            ans += v[i] * num * den;
ce6
        }
ba7
        return ans;
4fe }
```

## 3.3 Berlekamp-Massey

```
// guess_kth(s, k) chuta o k-esimo (0-based) termo
// de uma recorrencia linear que gera s
// Para uma rec. lin. de ordem x, se passar 2x termos
// vai gerar a certa
// Usar aritmetica modular
// O(n^2 log k), em que n = |s|
b7c template < typename T> T evaluate (vector < T> c, vector < T> s, ll k) {
        int n = c.size();
9ee
        assert(c.size() <= s.size());</pre>
d09
        auto mul = [\&] (const vector T &a. const vector T &b) {
            vector<T> ret(a.size() + b.size() - 1):
564
d75
            for (int i = 0; i < a.size(); i++) for (int j = 0; j <
   b.size(); j++)
                ret[i+j] += a[i] * b[j];
cff
            for (int i = ret.size()-1; i \ge n; i--) for (int j = n-1;
83d
   i >= 0; i--)
                ret[i-j-1] += ret[i] * c[j];
112
            ret.resize(min<int>(ret.size(), n));
            return ret;
edf
3b9
        }:
        vector < T > a = n == 1 ? vector < T > ({c[0]}) : vector < T > ({0, 1}),
   x = \{1\}:
95f
        while (k) {
7f1
            if (k\&1) x = mul(x, a):
            a = mul(a, a), k >>= 1;
b28
8ea
       x.resize(n);
dd6
        T ret = 0:
ce8
e72
        for (int i = 0: i < n: i++) ret += x[i] * s[i]:
        return ret;
edf
7e2 }
192 template < typename T > vector < T > berlekamp_massey(vector < T > s) {
ce8
        int n = s.size(), l = 0, m = 1;
222
        vector < T > b(n), c(n);
        T ld = b[0] = c[0] = 1;
46e
        for (int i = 0; i < n; i++, m++) {</pre>
620
793
            T d = s[i]:
            for (int j = 1; j <= 1; j++) d += c[j] * s[i-j];</pre>
ab6
5f0
            if (d == 0) continue;
8b4
            vector <T> temp = c;
            T coef = d / 1d:
369
```

```
ba6
            for (int j = m; j < n; j++) c[j] -= coef * b[j-m];
88f
            if (2 * 1 \le i) 1 = i + 1 - 1, b = temp, 1d = d, m = 0;
76a
90c
        c.resize(1 + 1);
844
        c.erase(c.begin());
        for (T\& x : c) x = -x;
0dc
807
        return c:
4d9 }
2cf template < typename T> T guess_kth(const vector < T>& s, ll k) {
cc3
        auto c = berlekamp_massev(s);
96a
        return evaluate(c, s, k);
697 }
```

#### 3.4 Binomial Distribution

## 3.5 Convolucao de GCD / LCM

```
// O(n log(n))

// multiple_transform(a)[i] = \sum_d a[d * i]
bbe template < typename T > void multiple_transform(vector < T > & v, bool
    inv = false) {
64a     vector < int > I(v.size()-1);
847     iota(I.begin(), I.end(), 1);
674     if (inv) reverse(I.begin(), I.end());
dad     for (int i : I) for (int j = 2; i*j < v.size(); j++)
a8a     v[i] += (inv ? -1 : 1) * v[i*j];</pre>
```

```
338 }
// \gcd_{convolution(a, b)[k]} = \sum_{gcd(i, j)} = k} a_i * b_j
fe2 template < typename T > vector < T > gcd_convolution(vector < T > a,
   vector<T> b) {
bdf
        multiple_transform(a), multiple_transform(b);
799
        for (int i = 0; i < a.size(); i++) a[i] *= b[i];</pre>
        multiple_transform(a, true);
dea
3f5
        return a;
984 }
// divisor_transform(a)[i] = \sum_{d|i} a[i/d]
be7 template < typename T > void divisor_transform (vector < T > & v, bool inv
   = false) {
64a
        vector < int > I(v.size()-1);
847
        iota(I.begin(), I.end(), 1);
        if (!inv) reverse(I.begin(), I.end());
5ea
dad
        for (int i : I) for (int j = 2; i*j < v.size(); j++)</pre>
            v[i*j] += (inv ? -1 : 1) * v[i];
14f
aa7 }
// lcm_convolution(a, b)[k] = \sum_{lcm(i, j)} = k} a_i * b_j
b1b template < typename T > vector < T > lcm_convolution (vector < T > a,
   vector <T> b) {
3af
        divisor_transform(a), divisor_transform(b);
799
        for (int i = 0; i < a.size(); i++) a[i] *= b[i];</pre>
d8f
        divisor_transform(a, true);
3f5
        return a:
f5a }
3.6 Coprime Basis
// Dado um conjunto de elementos A constroi uma base B
// de fatores coprimos tal que todo elemento A[i]
// pode ser fatorado como A[i] = \prod B[j]^p_ij
// Sendo n o numero de inserts, a complexidade esperada fica
// O(n*(n*loglog(MAX) + log(MAX)^2))
// No pior caso, podemos trocar n*loglog(MAX) por
// 8n. se MAX <= 1e6
```

// 10n. se MAX <= 1e9

// 16n, se MAX <= 1e18

// 26n, se MAX <= 1e36

ebc template <typename T> struct coprime\_basis {

```
a00
        vector <T> basis;
        coprime_basis() {}
60e
055
        coprime_basis(vector<T> v) { for (T i : v) insert(i); }
845
        void insert(T z) {
сЗс
             int n = basis.size():
efe
             basis.push_back(z);
43c
             for (int i = n; i < basis.size(); i++) {</pre>
                 for (int j = (i != n) ? i+1 : 0; j < basis.size();</pre>
21 c
    j++) {
                     if (i == j) continue;
4ce
024
                     T \&x = basis[i]:
c91
                     if (x == 1) {
fac
                         j = INF;
5e2
                         continue;
                     }
6e0
544
                     T & y = basis[j];
3c9
                     T g = gcd(x, y);
e10
                     if (g == 1) continue;
                     y /= g, x /= g;
15b
8c6
                     basis.push_back(g);
                 }
069
422
fe8
             basis.erase(remove(basis.begin(), basis.end(), 1),
    basis.end()):
1a5
        }
4ba
        vector<int> factor(T x) {
21d
             vector < int > fat(basis.size());
6fd
             for (int i = 0; i < basis.size(); i++) {</pre>
25 c
                 while (x \% basis[i] == 0) x /= basis[i], fat[i]++;
8de
             }
6a7
             return fat:
b5d
671 };
3.7 Crivo de Eratosthenes
// "O" crivo
//
// Encontra maior divisor primo
// Um numero eh primo sse divi[x] == x
// fact fatora um numero <= lim
// A fatoracao sai ordenada
```

```
// crivo - O(n log(log(n)))
// fact - O(log(n))
f12 int divi[MAX];
fb9 void crivo(int lim) {
f53
        for (int i = 1; i <= lim; i++) divi[i] = 1;</pre>
d46
        for (int i = 2; i <= lim; i++) if (divi[i] == 1)</pre>
018
            for (int j = i; j <= lim; j += i) divi[j] = i;</pre>
349 }
470 void fact(vector<int>& v. int n) {
ac8
        if (n != divi[n]) fact(v, n/divi[n]);
ab4
        v.push_back(divi[n]);
1db }
// Crivo linear
//
// Mesma coisa que o de cima, mas tambem
// calcula a lista de primos
//
// O(n)
f12 int divi[MAX];
fd3 vector<int> primes;
fb9 void crivo(int lim) {
        divi[1] = 1:
f70
       for (int i = 2; i <= lim; i++) {</pre>
3eb
            if (divi[i] == 0) divi[i] = i, primes.push_back(i);
3ba
            for (int j : primes) {
522
                if (j > divi[i] or i*j > lim) break;
00b
                divi[i*j] = j;
            }
491
       }
85a
519 }
// Crivo de divisores
//
// Encontra numero de divisores
// ou soma dos divisores
// O(n log(n))
f12 int divi[MAX];
```

```
fb9 void crivo(int lim) {
        for (int i = 1; i <= lim; i++) divi[i] = 1;</pre>
        for (int i = 2; i <= lim; i++)</pre>
424
594
            for (int j = i; j <= lim; j += i) {</pre>
                // para numero de divisores
9e0
                divi[i]++:
                // para soma dos divisores
278
                divi[j] += i;
            }
c58
fc1 }
// Crivo de totiente
// Encontra o valor da funcao
// totiente de Euler
// O(n log(log(n)))
5f4 int tot[MAX];
fb9 void crivo(int lim) {
        for (int i = 1; i <= lim; i++) {</pre>
bc9
            tot[i] += i;
feb
            for (int j = 2*i; j <= lim; j += i)
837
                tot[i] -= tot[i];
        }
678
212 }
// Crivo de funcao de mobius
//
// O(n log(log(n)))
4e1 char meb[MAX];
fb9 void crivo(int lim) {
649
        for (int i = 2; i <= lim; i++) meb[i] = 2;</pre>
        meb[1] = 1;
ace
842
        for (int i = 2; i <= lim; i++) if (meb[i] == 2)</pre>
868
            for (int j = i; j <= lim; j += i) if (meb[j]) {</pre>
686
                if (meb[j] == 2) meb[j] = 1;
ae1
                 meb[j] *= j/i\%i ? -1 : 0;
97f
            }
9bc }
// Crivo linear de funcao multiplicativa
//
```

```
// Computa f(i) para todo 1 <= i <= n, sendo f
// uma funcao multiplicativa (se gcd(a,b) = 1,
// entao f(a*b) = f(a)*f(b)
// f_prime tem que computar f de um primo, e
// add_prime tem que computar f(p^(k+1)) dado f(p^k) e p
// Se quiser computar f(p^k) dado p \in k, usar os comentarios
//
// O(n)
fd3 vector < int > primes;
623 int f[MAX], pot[MAX];
//int expo[MAX];
5c4 void sieve(int lim) {
        // Funcoes para soma dos divisores:
fc9
        auto f_prime = [](int p) { return p+1; };
        auto add_prime = [](int fpak, int p) { return fpak*p+1; };
31c
        //auto f_pak = [](int p, int k) {};
02d
        f[1] = 1:
f70
        for (int i = 2; i <= lim; i++) {</pre>
            if (!pot[i]) {
e6b
e74
                primes.push_back(i);
                f[i] = f_prime(i), pot[i] = i;
f05
                //\expo[i] = 1;
b71
            }
3b9
            for (int p : primes) {
b9f
                if (i*p > lim) break;
569
                if (i\%p == 0) {
                    f[i*p] = f[i / pot[i]] * add_prime(f[pot[i]], p);
b97
                    // se for descomentar, tirar a linha de cima tambem
                    //f[i*p] = f[i / pot[i]] * f_pak(p, expo[i]+1);
                    //\expo[i*p] = \expo[i]+1;
                    pot[i*p] = pot[i] * p;
51f
c2b
                    break:
643
                } else {
9ef
                    f[i*p] = f[i] * f[p];
                    pot[i*p] = p;
638
                    //\expo[i*p] = 1;
6f7
                }
f31
            }
1bb
        }
350 }
```

```
3.8 Deteccao de ciclo - Tortoise and Hare
```

```
// Linear no tanto que tem que andar pra ciclar,
// O(1) de memoria
// Retorna um par com o tanto que tem que andar
// do fO ate o inicio do ciclo e o tam do ciclo
58d pair<11, 11> find_cycle() {
273
        11 \text{ tort} = f(f0):
b2b
        ll hare = f(f(f0));
b1b
        11 t = 0;
683
        while (tort != hare) {
b4d
            tort = f(tort);
4b2
            hare = f(f(hare));
c82
            t++:
93d
        }
0e8
        11 st = 0;
909
        tort = f0;
683
        while (tort != hare) {
b4d
            tort = f(tort):
1a2
            hare = f(hare);
397
            st++;
c91
        }
73d
        11 len = 1:
3cd
        hare = f(tort);
683
        while (tort != hare) {
1a2
            hare = f(hare):
040
            len++;
f1a
        }
ebd
        return {st, len};
899 }
3.9 Division Trick
// copiei do github do tfg50
```

## 3.10 Equacao Diofantina Linear

```
// Encontra o numero de solucoes de a*x + b*y = c,
// em que x \in [lx, rx] e y \in [ly, ry]
// Usar o comentario para recuperar as solucoes
// (note que o b ao final eh b/gcd(a, b))
// Cuidado com overflow! Tem que caber o quadrado dos valores
// O(log(min(a, b)))
c5e template < typename T > tuple < 11, T, T > ext_gcd(11 a, 11 b) {
        if (!a) return {b, 0, 1};
        auto [g, x, y] = ext_gcd < T > (b%a, a);
c4b
c59
        return \{g, y - b/a*x, x\};
8a8 }
// numero de solucoes de a*[lx, rx] + b*[ly, ry] = c
14c template < typename T = 11> // usar __int128 se for ate 1e18
2a4 ll diophantine(ll a, ll b, ll c, ll lx, ll rx, ll ly, ll ry) {
        if (lx > rx or ly > ry) return 0;
c80
        if (a == 0 \text{ and } b == 0) \text{ return } c ? 0 : (rx-lx+1)*(ry-ly+1);
a98
        auto [g, x, y] = ext_gcd < T > (abs(a), abs(b));
8ce
9c3
        if (c % g != 0) return 0;
249
        if (a == 0) return (rx-lx+1)*(ly <= c/b and c/b <= ry);
        if (b == 0) return (ry-ly+1)*(lx <= c/a and c/a <= rx);
4ce
        x *= a/abs(a) * c/g, y *= b/abs(b) * c/g, a /= g, b /= g;
fb1
b20
        auto shift = [\&](T qt) \{ x += qt*b, y -= qt*a; \};
efa
        auto test = [&](T& k, ll mi, ll ma, ll coef, int t) {
866
            shift((mi - k)*t / coef);
79d
            if (k < mi) shift(coef > 0 ? t : -t);
74d
            if (k > ma) return pair T, T > (rx+2, rx+1);
41f
            T x1 = x;
633
            shift((ma - k)*t / coef);
c5b
            if (k > ma) shift(coef > 0 ? -t : t);
4a9
            return pair<T, T>(x1, x);
8e1
        };
639
        auto [11, r1] = test(x, 1x, rx, b, 1);
38e
        auto [12, r2] = test(y, ly, ry, a, -1);
c43
        if (12 > r2) swap(12, r2);
50a
        T l = max(11, 12), r = min(r1, r2);
339
        if (1 > r) return 0;
        ll k = (r-1) / abs(b) + 1;
42f
839
        return k; // solucoes: x = 1 + [0, k)*|b|
98e }
```

#### 3.11 Euclides estendido

```
// Acha x e y tal que ax + by = mdc(a, b) (nao eh unico)
// Assume a, b >= 0
//
// O(log(min(a, b)))

2be tuple<ll, ll, ll> ext_gcd(ll a, ll b) {
3bd     if (!a) return {b, 0, 1};
550     auto [g, x, y] = ext_gcd(b%a, a);
c59     return {g, y - b/a*x, x};
354 }
```

## 3.12 Exponenciacao rapida

```
// (x^y mod m) em O(log(y))
03c ll pow(ll x, ll y, ll m) \{ // \text{ iterativo} \}
c85
        ll ret = 1:
1b8
        while (y) {
895
            if (y & 1) ret = (ret * x) % m;
23b
            y >>= 1;
cc5
             x = (x * x) % m;
        }
020
edf
        return ret;
12b }
03c ll pow(ll x, ll y, ll m) \{ // \text{ recursivo } \}
        if (!y) return 1;
426
        11 ans = pow(x*x\%m, y/2, m);
88d
        return v%2 ? x*ans%m : ans;
7d4 }
```

#### 3.13 Fast Walsh Hadamard Transform

```
// FWHT<'!'>(f) eh SOS DP
// FWHT<'&'>(f) eh soma de superset DP
// Se chamar com ^, usar tamanho potencia de 2!!
//
// O(n log(n))

382 template<char op, class T> vector<T> FWHT(vector<T> f, bool inv = false) {
```

```
b75
        int n = f.size();
                                                                           ee2
                                                                                   else r = r^(ord/n);
d78
        for (int k = 0; (n-1) >> k; k++) for (int i = 0; i < n; i++) if
                                                                           be4
                                                                                   roots[0] = 1:
   (i>>k&1) {
                                                                           078
                                                                           63f }
29e
            int j = i^{(1 << k)};
627
            if (op == '\^') f[i] += f[i], f[i] = f[i] - 2*f[i];
            if (op == ', ') f[i] += (inv ? -1 : 1) * f[j];
a38
93c
            if (op == '&') f[j] += (inv ? -1 : 1) * f[i];
                                                                               vector<int>& rev) {
1bb
        }
                                                                           bc7
578
        if (op == ', and inv) for (auto& i : f) i /= n;
                                                                               a[rev[i]]);
        return f;
                                                                                   int 1. r. m:
abe
                                                                           12b
50e }
                                                                           cb4
                                                                                   vector < T > roots(N);
                                                                           192
                                                                           0f4
                                                                                       get roots(f. n. roots):
                                                                           5dc
3.14 FFT
                                                                           432
                                                                                           while (m < n/2) {
                                                                           a88
// Chamar convolution com vector < complex < double >> para FFT
                                                                           297
// Precisa do mint para NTT
                                                                           254
                                                                                               a[r] = a[1] - t:
//
                                                                           b8f
                                                                                               a[1] = a[1] + t;
// O(n log(n))
                                                                           2c9
                                                                                               1++, r++, m++;
                                                                           d89
                                                                                          }
// Para FFT
                                                                           1fd
                                                                                       }
488 void get_roots(bool f, int n, vector<complex<double>>& roots) {
                                                                                   }
                                                                           185
f26
        const static double PI = acosl(-1);
                                                                           235
                                                                                   if (f) {
        for (int i = 0; i < n/2; i++) {</pre>
71a
                                                                           1c5
                                                                                       auto invN = T(1) / T(N);
b1e
            double alpha = i*((2*PI)/n);
                                                                           557
            if (f) alpha = -alpha;
1a1
                                                                           256
                                                                                   }
069
            roots[i] = {cos(alpha), sin(alpha)};
                                                                           1b1 }
804
       }
de5 }
                                                                              vector<T>& b) {
// Para NTT
                                                                           87a
9f7 template <int p>
                                                                                   int N = 1.size()+r.size()-1;
                                                                           e0a
97b void get_roots(bool f, int n, vector<mod_int<p>>& roots) {
                                                                           f03
                                                                                   int n = 1, \log n = 0;
1e6
        mod_int  r;
                                                                           0a4
de9
        int ord:
                                                                           808
                                                                                   vector<int> rev(n);
        if (p == 998244353) {
57a
                                                                           603
                                                                                   for (int i = 0; i < n; i++) {</pre>
9b6
            r = 102292;
                                                                           434
                                                                                       rev[i] = 0;
81b
            ord = (1 << 23);
                                                                           f44
        } else if (p == 754974721) {
121
                                                                           4ff
            r = 739831874:
43a
                                                                           256
                                                                                   }
            ord = (1 << 24);
f0a
                                                                                   assert(N <= n);
                                                                           143
d48
        } else if (p == 167772161) {
                                                                           fa4
                                                                                   l.resize(n):
a2a
            r = 243:
                                                                           7e4
                                                                                   r.resize(n);
033
            ord = (1 << 25);
                                                                           56e
                                                                                   fft(1, false, n, rev);
        } else assert(false);
5a4
                                                                           fcf
                                                                                   fft(r, false, n, rev);
                                                                           917
547
        if (f) r = r^(p - 1 - ord/n);
```

```
for (int i = 1; i < n/2; i++) roots[i] = roots[i-1]*r;</pre>
8a2 template < typename T > void fft(vector < T > & a, bool f, int N,
        for (int i = 0; i < N; i++) if (i < rev[i]) swap(a[i],</pre>
        for (int n = 2; n <= N; n *= 2) {</pre>
             for (int pos = 0; pos < N; pos += n) {</pre>
                 1 = pos + 0, r = pos + n/2, m = 0;
                     auto t = roots[m] * a[r];
             for (int i = 0; i < N; i++) a[i] = a[i] * invN;</pre>
bf5 template < typename T > vector < T > convolution (vector < T > & a,
        vector<T> l(a.begin(), a.end()), r(b.begin(), b.end());
        while (n <= N) n *= 2, log_n++;</pre>
            for (int j = 0; j < log_n; j++) if (i>>j&1)
                 rev[i] = 1 << (log_n-1-i);
        for (int i = 0; i < n; i++) l[i] *= r[i];</pre>
```

```
88b
        fft(l, true, n, rev);
5e1
        l.resize(N):
792
        return 1;
bd6 }
// NTT
6c8 template < int p, typename T > vector < mod_int < p >> ntt (vector < T > & a,
   vector<T>& b) {
        vector < mod_int < p >> A(a.begin(), a.end()), B(b.begin(),
d52
d29
        return convolution(A, B);
3bf }
// Convolucao de inteiro
// Precisa do CRT
// Tabela de valores:
// [0,1]
          - <int, 1>
// [-1e5, 1e5] - <11, 2>
// [-1e9, 1e9] - <__int128, 3>
b3c template < typename T, int mods >
eec vector<T> int_convolution(vector<int>& a, vector<int>& b) {
        static const int M1 = 998244353, M2 = 754974721, M3 =
fe8
   167772161:
bf5
        auto c1 = ntt < M1 > (a, b);
221
        auto c2 = (mods \ge 2 ? ntt \le M2 \ge (a. b) : vector \le mod int \le M2 \ge ()):
        auto c3 = (mods >= 3 ? ntt<M3>(a, b) : vector<mod_int<M3>>());
f9b
2da
        vector <T> ans:
5c5
        for (int i = 0; i < c1.size(); i++) {</pre>
c09
            crt < T > at(c1[i].v, M1);
            if (mods \ge 2) at = at * crt<T>(c2[i].v. M2):
316
            if (mods >= 3) at = at * crt<T>(c3[i].v, M3);
987
b2b
            ans.push_back(at.a);
26d
            if (at.a > at.m/2) ans.back() -= at.m;
b9f
ba7
        return ans;
5e8 }
3.15 Gauss
// Resolve sistema linear
// Retornar um par com o numero de solucoes
// e alguma solucao, caso exista
```

```
// O(n^2 * m)
67a template < typename T>
728 pair<int, vector<T>> gauss(vector<vector<T>> a, vector<T> b) {
        const double eps = 1e-6;
f92
        int n = a.size(), m = a[0].size();
2f0
        for (int i = 0; i < n; i++) a[i].push_back(b[i]);</pre>
        vector<int> where(m. -1);
3cb
237
        for (int col = 0, row = 0; col < m and row < n; col++) {
f05
             int sel = row:
b95
             for (int i=row: i<n: ++i)</pre>
e55
                 if (abs(a[i][col]) > abs(a[sel][col])) sel = i;
2c4
             if (abs(a[sel][col]) < eps) continue;</pre>
             for (int i = col; i <= m; i++)</pre>
1ae
dd2
                 swap(a[sel][i], a[row][i]);
             where[col] = row:
2c3
             for (int i = 0; i < n; i++) if (i != row) {</pre>
0 c 0
                 T c = a[i][col] / a[row][col]:
96c
d5c
                 for (int j = col; j <= m; j++)</pre>
                     a[i][j] -= a[row][j] * c;
c8f
490
            }
b70
             row++;
3d8
        }
b1d
        vector <T> ans(m. 0):
        for (int i = 0; i < m; i++) if (where[i] != -1)</pre>
e1a
12a
             ans[i] = a[where[i]][m] / a[where[i]][i];
        for (int i = 0; i < n; i++) {</pre>
603
501
            T sum = 0;
a75
             for (int j = 0; j < m; j++)</pre>
5a9
                 sum += ans[i] * a[i][i]:
             if (abs(sum - a[i][m]) > eps)
b1f
6cd
                 return pair(0, vector<T>());
ec9
        }
12e
        for (int i = 0; i < m; i++) if (where[i] == -1)</pre>
018
             return pair(INF, ans);
        return pair(1, ans);
280
292 }
3.16 Gauss - 72
```

// D eh dimensao do espaco vetorial

```
// add(v) - adiciona o vetor v na base (retorna se ele jah pertencia
   ao span da base)
// coord(v) - retorna as coordenadas (c) de v na base atual (basis^T.c
// recover(v) - retorna as coordenadas de v nos vetores na ordem em
   que foram inseridos
// coord(v).first e recover(v).first - se v pertence ao span
// Complexidade:
// add, coord, recover: O(D^2 / 64)
2a3 template <int D> struct Gauss_z2 {
3c1
        bitset <D> basis[D]. keep[D];
b16
        int rk, in;
482
        vector < int > id;
37f
        Gauss_z2 () : rk(0), in(-1), id(D, -1) {};
        bool add(bitset < D > v) {
04e
42c
            in++:
fb0
            bitset <D> k;
            for (int i = D - 1; i \ge 0; i--) if (v[i]) {
659
189
                if (basis[i][i]) v ^= basis[i], k ^= keep[i];
                else {
4e6
                    k[i] = true, id[i] = in, keep[i] = k;
                    basis[i] = v, rk++;
6ce
8a6
                    return true;
b34
                }
            }
09c
d1f
            return false;
58b
0f6
        pair < bool, bitset < D >> coord(bitset < D > v) {
944
            bitset <D> c:
659
            for (int i = D - 1; i \ge 0; i--) if (v[i]) {
                if (basis[i][i]) v ^= basis[i], c[i] = true;
a39
8af
                else return {false, bitset <D>()};
            }
a08
5db
            return {true, c};
a08
330
        pair < bool, vector < int >> recover(bitset < D > v) {
22e
            auto [span, bc] = coord(v);
            if (not span) return {false, {}};
af8
f79
            bitset <D> aux:
            for (int i = D - 1; i >= 0; i--) if (bc[i]) aux ^= keep[i];
5a0
            vector < int > oc;
ea9
            for (int i = D - 1; i >= 0; i--) if (aux[i])
   oc.push_back(id[i]);
```

```
001
            return {true, oc};
b75
d0a };
3.17 Integracao Numerica
// Metodo de Simpson 3/8
// Integra f no intervalo [a, b], erro cresce proporcional a (b - a)^5
676 const int N = 3*100; // multiplo de 3
287 ld integrate(ld a, ld b, function < ld(ld) > f) {
b4d
        ld s = 0, h = (b - a)/N;
        for (int i = 1; i < N; i++) s += f(a + i*h)*(i%3 ? 3 : 2);
067
        return (f(a) + s + f(b))*3*h/8:
0da
c7e }
3.18 Inverso Modular
// Computa o inverso de a modulo b
// Se b eh primo, basta fazer
// a^{(b-2)}
f0a ll inv(ll a, ll b) {
        return a > 1 ? b - inv(b\%a, a)*b/a : 1;
ae1
cf9 }
// computa o inverso modular de 1..MAX-1 modulo um primo
a88 ll inv[MAX]:
0f2 inv[1] = 1;
Ofa for (int i = 2; i < MAX; i++) inv[i] = MOD - MOD/i*inv[MOD%i]%MOD;
3.19 Karatsuba
// Os pragmas podem ajudar
// Para n \sim 2e5. roda em < 1 s
// O(n^1.58)
//#pragma GCC optimize("Ofast")
//#pragma GCC target ("avx,avx2")
```

77a template < typename T > void kar(T\* a, T\* b, int n, T\* r, T\* tmp) {

 $if (n \le 64)$ {

```
510
            for (int i = 0; i < n; i++) for (int j = 0; j < n; j++)
212
                 r[i+j] += a[i] * b[j];
505
            return;
        }
bb8
194
        int mid = n/2:
2d7
        T *atmp = tmp, *btmp = tmp+mid, *E = tmp+n;
4f1
        memset(E, 0, sizeof(E[0])*n):
        for (int i = 0; i < mid; i++) {</pre>
c65
c72
            atmp[i] = a[i] + a[i+mid];
4b9
            btmp[i] = b[i] + b[i+mid];
a3f
38a
        kar(atmp, btmp, mid, E, tmp+2*n);
b1e
        kar(a. b. mid. r. tmp+2*n):
        kar(a+mid, b+mid, mid, r+n, tmp+2*n);
229
c65
        for (int i = 0; i < mid; i++) {</pre>
735
            T \text{ temp} = r[i+mid];
de7
            r[i+mid] += E[i] - r[i] - r[i+2*mid];
            r[i+2*mid] += E[i+mid] - temp - r[i+3*mid];
f1e
f72
28f }
e38 template < typename T > vector < T > karatsuba (vector < T > a, vector < T > b)
   {
ba3
        int n = max(a.size(), b.size());
        while (n&(n-1)) n++;
a84
ca9
       a.resize(n). b.resize(n):
        vector \langle T \rangle ret(2*n), tmp(4*n);
ae0
644
        kar(&a[0], &b[0], n, &ret[0], &tmp[0]);
        return ret:
edf
f87 }
```

#### 3.20 Logaritmo Discreto

```
// Resolve logaritmo discreto com o algoritmo baby step giant step
// Encontra o menor x tal que a^x = b (mod m)
// Se nao tem. retorna -1
// O(sqrt(m) * log(sqrt(m))
d41
da8 int dlog(int b, int a, int m) {
9f8
        if (a == 0) return b ? -1 : 1: // caso nao definido
a6e
        a \%= m, b \%= m;
a 10
       int k = 1, shift = 0;
31e
        while (1) {
6e3
           int g = gcd(a, m);
```

```
d47
            if (g == 1) break;
d41
9bc
            if (b == k) return shift;
642
            if (b % g) return -1;
            b \neq g, m \neq g, shift++;
c36
            k = (11) k * a / g % m;
9ab
515
        }
d41
af7
        int sq = sqrt(m)+1, giant = 1;
        for (int i = 0; i < sq; i++) giant = (11) giant * a % m;</pre>
975
d41
0b5
        vector < pair < int , int >> baby;
33f
        for (int i = 0, cur = b; i \le sq; i++) {
496
            baby.emplace_back(cur, i);
16c
            cur = (11) cur * a % m;
622
eb4
        sort(baby.begin(), baby.end());
d41
9c9
        for (int j = 1, cur = k; j <= sq; j++) {</pre>
            cur = (11) cur * giant % m;
ace
78b
            auto it = lower_bound(baby.begin(), baby.end(), pair(cur,
   INF));
d26
            if (it != baby.begin() and (--it)->first == cur)
                return sq * j - it->second + shift;
ac3
b9d
        }
d41
daa
        return -1;
739 }
```

## 3.21 Miller-Rabin

```
// Testa se n eh primo, n \leq 3 * 10^18
// O(log(n)), considerando multiplicacao
// e exponenciacao constantes
d8b ll mul(ll a, ll b, ll m) {
        ll ret = a*b - ll((long double)1/m*a*b+0.5)*m;
074
        return ret < 0 ? ret+m : ret:</pre>
2f3 }
03c ll pow(ll x, ll y, ll m) {
13a
        if (!v) return 1;
dbc
        ll ans = pow(mul(x, x, m), y/2, m);
7fa
        return y%2 ? mul(x, ans, m) : ans;
539 }
```

```
1a2 bool prime(ll n) {
       if (n < 2) return 0;
1aa
       if (n <= 3) return 1;
       if (n % 2 == 0) return 0:
9de
        ll r = \_builtin\_ctzll(n - 1), d = n >> r;
f6a
        // com esses primos, o teste funciona garantido para n <= 2^64
        // funciona para n <= 3*10^24 com os primos ate 41
771
        for (int a : {2, 325, 9375, 28178, 450775, 9780504,
   795265022}) {
            11 x = pow(a, d, n);
da0
709
            if (x == 1 \text{ or } x == n - 1 \text{ or a } \% n == 0) continue:
4a2
            for (int j = 0; j < r - 1; j++) {
10f
                x = mul(x, x, n);
df0
                if (x == n - 1) break;
1ff
            if (x != n - 1) return 0;
e1b
e74
6a5
        return 1;
5ba }
```

## 3.22 Pollard's Rho Alg

```
// Usa o algoritmo de deteccao de ciclo de Flovd
// com uma otimizacao na qual o gcd eh acumulado
// A fatoracao nao sai necessariamente ordenada
// O algoritmo rho encontra um fator de n,
// e funciona muito bem quando n possui um fator pequeno
// Complexidades (considerando mul constante):
// rho - esperado O(n^{(1/4)}) no pior caso
// fact - esperado menos que O(n^{(1/4)} \log(n)) no pior caso
d8b ll mul(ll a, ll b, ll m) {
e7a
        11 \text{ ret} = a*b - 11((long double)1/m*a*b+0.5)*m;
074
        return ret < 0 ? ret+m : ret;</pre>
2f3 }
03c ll pow(ll x, ll y, ll m) {
       if (!y) return 1;
dbc
        ll ans = pow(mul(x, x, m), v/2, m);
7fa
        return y%2 ? mul(x, ans, m) : ans;
539 }
```

```
1a2 bool prime(ll n) {
        if (n < 2) return 0;
237
        if (n <= 3) return 1;
        if (n % 2 == 0) return 0;
f6a
        ll r = \_builtin\_ctzll(n - 1), d = n >> r;
771
        for (int a: {2, 325, 9375, 28178, 450775, 9780504,
   795265022}) {
da0
            11 x = pow(a, d, n);
            if (x == 1 \text{ or } x == n - 1 \text{ or a } \% n == 0) continue;
709
4a2
            for (int j = 0; j < r - 1; j++) {
10f
                x = mul(x, x, n):
df0
                if (x == n - 1) break;
1ff
e1b
            if (x != n - 1) return 0;
e74
6a5
        return 1;
5ba }
9cf ll rho(ll n) {
        if (n == 1 or prime(n)) return n;
0f9
f7c
        auto f = [n](11 x) \{ return mul(x, x, n) + 1; \};
8a5
        11 x = 0, y = 0, t = 30, prd = 2, x0 = 1, q;
533
        while (t % 40 != 0 or gcd(prd, n) == 1) {
8a0
            if (x==y) x = ++x0, y = f(x);
e13
            q = mul(prd, abs(x-y), n);
            if (q != 0) prd = q;
21f
450
            x = f(x), y = f(f(y)), t++;
379
002
        return gcd(prd, n);
523 }
5b7 vector<ll> fact(ll n) {
1b9
        if (n == 1) return {};
0ec
        if (prime(n)) return {n};
0ed
        11 d = rho(n);
1de
        vector < ll > l = fact(d), r = fact(n / d);
3af
        1.insert(1.end(), r.begin(), r.end());
792
        return 1:
902 }
```

#### 3.23 Produto de dois long long mod m

// 0(1)

```
d8b ll mul(ll a, ll b, ll m) { // a*b % m
        11 \text{ ret} = a*b - 11((long double)1/m*a*b+0.5)*m;
e7a
074
        return ret < 0 ? ret+m : ret;</pre>
2f3 }
3.24 Simplex
// Maximiza c^T x s.t. Ax <= b, x >= 0
// 0(2^n), porem executa em 0(n^3) no caso medio
395 const double eps = 1e-7;
493 namespace Simplex {
        vector < vector < double >> T:
69 c
14e
        int n, m;
43e
       vector<int> X. Y:
c51
        void pivot(int x, int y) {
8e6
             swap(X[y], Y[x-1]);
            for (int i = 0; i <= m; i++) if (i != y) T[x][i] /=
d03
   T[x][y];
33c
            T[x][y] = 1/T[x][y];
            for (int i = 0; i <= n; i++) if (i != x and abs(T[i][y]) >
38b
   eps) {
                 for (int j = 0; j <= m; j++) if (j != y) T[i][j] -=
774
   T[i][y] * T[x][j];
                 T[i][y] = -T[i][y] * T[x][y];
348
a7d
e05
        }
        // Retorna o par (valor maximo, vetor solucao)
6f8
        pair < double , vector < double >> simplex(
                 vector < vector < double >> A, vector < double >> b,
e9d
   vector < double > c) {
            n = b.size(), m = c.size();
5bb
            T = vector(n + 1, vector < double > (m + 1));
002
            X = vector<int>(m):
2d9
            Y = vector < int > (n):
0 c 2
            for (int i = 0; i < m; i++) X[i] = i;</pre>
115
            for (int i = 0: i < n: i++) Y[i] = i+m:</pre>
51f
```

for (int i = 0; i < m; i++) T[0][i] = -c[i];

for (int j = 0; j < m; j++) T[i+1][j] = A[i][j];

for (int i = 0: i < n: i++) {</pre>

T[i+1][m] = b[i];

5b5

603

ba6

eca

```
07c
667
            while (true) {
714
                int x = -1, y = -1;
2db
                double mn = -eps;
                for (int i = 1; i <= n; i++) if (T[i][m] < mn) mn =</pre>
c29
   T[i][m], x = i;
af2
                if (x < 0) break:
                for (int i = 0; i < m; i++) if (T[x][i] < -eps) { y = }
   i; break; }
                if (y < 0) return \{-1e18, \{\}\}; // sem solucao para Ax
4a6
   <= b
7fb
                pivot(x, y);
472
667
            while (true) {
714
                int x = -1, y = -1;
2db
                double mn = -eps;
                for (int i = 0; i < m; i++) if (T[0][i] < mn) mn =</pre>
562
   T[0][i], v = i;
                if (y < 0) break;
9ъ0
034
                mn = 1e200:
                for (int i = 1; i \le n; i++) if (T[i][y] > eps and
5af
   T[i][m] / T[i][y] < mn
48f
                    mn = T[i][m] / T[i][y], x = i;
53b
                if (x < 0) return {1e18, {}}; // c^T x eh ilimitado
7fb
                pivot(x, y);
81e
            }
290
            vector < double > r(m);
32f
            for(int i = 0; i < n; i++) if (Y[i] < m) r[Y[i]] =
   T[i+1][m]:
e59
            return {T[0][m], r};
        }
7a4
a64 }
3.25 Teorema Chines do Resto
// Combina equacoes modulares lineares: x = a (mod m)
// O m final eh o lcm dos m's, e a resposta eh unica mod o lcm
// Os m nao precisam ser coprimos
// Se nao tiver solucao, o 'a' vai ser -1
153 template < typename T > tuple < T, T, T > ext_gcd(T a, T b) {
```

3bd

550

c59

if (!a) return {b, 0, 1};

return  $\{g, y - b/a*x, x\};$ 

auto  $[g, x, y] = ext_gcd(b\%a, a);$ 

```
537 }
bfe template < typename T = 11 > struct crt {
627
        Ta, m;
        crt() : a(0), m(1) {}
5f3
7eb
        crt(T a_, T m_) : a(a_), m(m_) {}
        crt operator * (crt C) {
911
238
            auto [g, x, y] = ext_gcd(m, C.m);
            if ((a - C.a) % g) a = -1;
dc0
4f9
           if (a == -1 or C.a == -1) return crt(-1, 0);
d09
           T lcm = m/g*C.m;
eb2
           T ans = a + (x*(C.a-a)/g \% (C.m/g))*m;
d8d
            return crt((ans % lcm + lcm) % lcm, lcm);
1f2
       }
0d9 };
```

#### 3.26 Totiente

```
// O(sqrt(n))
a7e int tot(int n){
0f6
        int ret = n;
       for (int i = 2; i*i <= n; i++) if (n % i == 0) {
505
b0c
            while (n \% i == 0) n /= i;
            ret -= ret / i;
125
34a
       if (n > 1) ret -= ret / n;
af4
edf
        return ret;
fae }
```

## 4 Grafos

### 4.1 AGM Direcionada

```
// Fala o menor custo para selecionar arestas tal que
// o vertice 'r' alcance todos
// Se nao tem como, retorna LINF
//
// O(m log(n))
```

```
3c9 struct node {
f31
        pair<ll, int> val;
4e4
        ll lazy;
b19
        node *1, *r;
f93
        node() {}
c53
        node(pair < int, int > v) : val(v), lazy(0), l(NULL), r(NULL) {}
a9c
        void prop() {
768
            val.first += lazy;
b87
            if (1) 1->lazy += lazy;
d3b
            if (r) r->lazy += lazy;
c60
            lazv = 0:
05b
        }
296 };
de5 void merge(node*& a, node* b) {
c11
        if (!a) swap(a, b);
802
        if (!b) return;
626
        a->prop(), b->prop();
d04
        if (a->val > b->val) swap(a, b);
4b0
        merge(rand()%2 ? a->1 : a->r, b);
b82 }
d01 pair < ll, int > pop(node * & R) {
        R->prop();
e8f
22e
        auto ret = R->val;
af0
        node* tmp = R;
3f3
        merge(R->1, R->r);
6c9
        R = R - > 1:
3 e 4
        if (R) R->lazy -= ret.first;
7c3
        delete tmp;
edf
        return ret;
c4e }
6f6 void apaga(node * R) { if (R) apaga(R->1), apaga(R->r), delete R; }
f13 ll dmst(int n, int r, vector<pair<int, int>, int>>& ar) {
        vector < int > p(n); iota(p.begin(), p.end(), 0);
        function < int(int) > find = [&](int k) { return
   p[k] == k?k:p[k] = find(p[k]); };
        vector < node *> h(n);
        for (auto e : ar) merge(h[e.first.second], new node({e.second,
   e.first.first}));
fd1
        vector < int > pai(n, -1), path(n);
        pai[r] = r;
66e
04b
        11 \text{ ans} = 0;
603
        for (int i = 0; i < n; i++) { // vai conectando todo mundo
2a3
            int u = i. at = 0:
```

```
cae
            while (pai[u] == -1) {
daa
                if (!h[u]) { // nao tem
947
                    for (auto i : h) apaga(i);
                    return LINF;
77 c
dd1
167
                path[at++] = u, pai[u] = i;
55e
                auto [mi, v] = pop(h[u]);
                ans += mi;
                if (pai[u = find(v)] == i) { // ciclo
5e2
86f
                    while (find(v = path[--at]) != u)
621
                        merge(h[u], h[v]), h[v] = NULL, p[find(v)] = u;
57a
                    pai[u] = -1:
8b0
                }
           }
ce8
5df
       }
947
        for (auto i : h) apaga(i);
ba7
        return ans;
e02 }
```

#### 4.2 Articulation Points

```
// Computa os pontos de articulação (vertices criticos) de um grafo
// art[i] armazena o numero de novas componentes criadas ao deletar
// se art[i] >= 1, entao vertice i eh ponto de articulação
//
// O(n+m)
1a8 int n:
789 vector < vector < int >> g;
4ce stack < int > s;
b66 vector < int > id, art;
3e1 int dfs_art(int i, int& t, int p = -1) {
cf0
       int lo = id[i] = t++;
18e
        s.push(i);
       for (int j : g[i]) if (j != p) {
cac
            if (id[j] == -1) {
9a3
206
                int val = dfs_art(j, t, i);
               lo = min(lo, val);
0c3
588
                if (val >= id[i]) {
66a
                    art[i]++;
bd9
                    while (s.top() != j) s.pop();
```

```
2eb
                     s.pop();
1f3
                 // if (val > id[i]) aresta i-j eh ponte
238
328
             else lo = min(lo, id[j]);
762
3bd
        if (p == -1 and art[i]) art[i]--;
        return lo;
253
8e1 }
d79 void compute_art_points() {
        id = vector < int > (n, -1);
597
a62
        art = vector < int > (n. 0):
6bb
        int t = 0:
        for (int i = 0; i < n; i++) if (id[i] == -1)</pre>
d41
             dfs_art(i, t, -1);
625
379 }
```

#### 4.3 Bellman-Ford

```
// Calcula a menor distancia
// entre a e todos os vertices e
// detecta ciclo negativo
// Retorna 1 se ha ciclo negativo
// Nao precisa representar o grafo,
// soh armazenar as arestas
//
// O(nm)
14e int n, m;
248 int d[MAX]:
e93 vector <pair <int, int>> ar; // vetor de arestas
9e2 vector < int > w;
                                // peso das arestas
6be bool bellman_ford(int a) {
8ec
        for (int i = 0; i < n; i++) d[i] = INF;</pre>
8a8
        d[a] = 0;
        for (int i = 0: i <= n: i++)
4e3
891
            for (int j = 0; j < m; j++) {
6e4
                if (d[ar[j].second] > d[ar[j].first] + w[j]) {
705
                    if (i == n) return 1;
e93
                     d[ar[j].second] = d[ar[j].first] + w[j];
84b
                }
a82
            }
```

```
bb3 return 0;
6eb }
```

### 4.4 Block-Cut Tree

```
// Cria a block-cut tree, uma arvore com os blocos
// e os pontos de articulação
// Blocos sao componentes 2-vertice-conexos maximais
// Uma 2-coloracao da arvore eh tal que uma cor sao
// os blocos, e a outra cor sao os pontos de art.
// Funciona para grafo nao conexo
// art[i] responde o numero de novas componentes conexas
// criadas apos a remocao de i do grafo g
// Se art[i] >= 1, i eh ponto de articulação
//
// Para todo i <= blocks.size()</pre>
// blocks[i] eh uma componente 2-vertce-conexa maximal
// edgblocks[i] sao as arestas do bloco i
// tree[i] eh um vertice da arvore que corresponde ao bloco i
// pos[i] responde a qual vertice da arvore vertice i pertence
// Arvore tem no maximo 2n vertices
// O(n+m)
d10 struct block_cut_tree {
d8e
        vector < vector < int >> g, blocks, tree;
43b
        vector < vector < pair < int , int >>> edgblocks;
4ce
        stack<int> s;
6c0
        stack<pair<int, int>> s2;
        vector<int> id, art, pos;
2bb
763
        block_cut_tree(vector<vector<int>> g_) : g(g_) {
af1
            int n = g.size();
37a
            id.resize(n, -1), art.resize(n), pos.resize(n);
6f2
            build():
        }
6bd
df6
        int dfs(int i, int& t, int p = -1) {
cf0
            int lo = id[i] = t++;
18e
            s.push(i);
827
            if (p != -1) s2.emplace(i, p);
```

```
53f
            for (int j : g[i]) if (j != p and id[j] != -1)
   s2.emplace(i, j);
cac
            for (int j : g[i]) if (j != p) {
                if (id[j] == -1) {
9a3
121
                    int val = dfs(j, t, i);
0 c 3
                    lo = min(lo, val);
588
                    if (val >= id[i]) {
                         art[i]++:
66a
483
                         blocks.emplace_back(1, i);
110
                         while (blocks.back().back() != j)
138
                             blocks.back().push_back(s.top()), s.pop();
128
                         edgblocks.emplace_back(1, s2.top()), s2.pop();
47e
                         while (edgblocks.back().back() != pair(j, i))
bce
                             edgblocks.back().push_back(s2.top()),
   s2.pop();
870
                    // if (val > id[i]) aresta i-j eh ponte
85c
328
                else lo = min(lo, id[j]);
            }
344
3bd
            if (p == -1 and art[i]) art[i]--;
253
            return lo:
726
        }
        void build() {
0a8
6bb
            int t = 0;
            for (int i = 0; i < g.size(); i++) if (id[i] == -1) dfs(i,</pre>
abf
   t, -1);
56c
            tree.resize(blocks.size()):
f7d
            for (int i = 0; i < g.size(); i++) if (art[i])</pre>
965
                pos[i] = tree.size(), tree.emplace_back();
973
            for (int i = 0; i < blocks.size(); i++) for (int j :</pre>
   blocks[i]) {
403
                if (!art[i]) pos[i] = i;
                else tree[i].push_back(pos[j]),
101
   tree[pos[j]].push_back(i);
3df
        }
c03
056 };
```

#### 4.5 Blossom

```
// Matching maximo em grafo geral
//
// O(n^3)
// Se for bipartido, nao precisa da funcao
// 'contract', e roda em O(nm)
042 vector < int > g[MAX];
128 int match[MAX]; // match[i] = com quem i esta matchzado ou -1
1f1 int n, pai[MAX], base[MAX], vis[MAX];
26a queue < int > q;
107 void contract(int u, int v, bool first = 1) {
        static vector < bool > bloss;
165
fbe
        static int 1:
418
        if (first) {
a47
            bloss = vector < bool > (n, 0);
042
            vector < bool > teve(n, 0):
ddf
            int k = u; l = v;
31e
            while (1) {
297
                teve[k = base[k]] = 1;
                if (match[k] == -1) break;
116
                k = pai[match[k]];
dfa
68b
            }
d31
            while (!teve[l = base[l]]) l = pai[match[l]];
5d6
        while (base[u] != 1) {
2e9
            bloss[base[u]] = bloss[base[match[u]]] = 1;
e29
8fa
            pai[u] = v;
0b0
            v = match[u];
a51
            u = pai[match[u]];
58e
71c
        if (!first) return;
95e
        contract(v, u, 0);
        for (int i = 0; i < n; i++) if (bloss[base[i]]) {</pre>
6ee
594
            base[i] = 1;
ca7
            if (!vis[i]) q.push(i);
29a
            vis[i] = 1;
857
        }
e35 }
f10 int getpath(int s) {
88f
        for (int i = 0; i < n; i++) base[i] = i, pai[i] = -1, vis[i] =</pre>
   0:
ded
        vis[s] = 1; q = queue < int > (); q.push(s);
402
        while (q.size()) {
```

```
be1
            int u = q.front(); q.pop();
            for (int i : g[u]) {
bdc
7a2
                 if (base[i] == base[u] or match[u] == i) continue;
e35
                 if (i == s or (match[i] != -1 and pai[match[i]] != -1))
4f2
                     contract(u, i):
                 else if (pai[i] == -1) {
e2e
545
                     pai[i] = u;
f6a
                     if (match[i] == -1) return i;
818
                     i = match[i];
29d
                     vis[i] = 1; q.push(i);
90e
                }
0b5
            }
634
        }
daa
        return -1;
a16 }
83f int blossom() {
1a4
        int ans = 0:
315
        memset(match, -1, sizeof(match));
        for (int i = 0; i < n; i++) if (match[i] == -1)</pre>
2e3
f76
            for (int j : g[i]) if (match[j] == -1) {
1bc
                 match[i] = i;
f1d
                 match[j] = i;
Odf
                 ans++;
c2b
                 break;
723
            }
da8
        for (int i = 0; i < n; i++) if (match[i] == -1) {</pre>
7e3
            int j = getpath(i);
5f2
            if (j == -1) continue;
Odf
            ans++;
3a0
            while (j != -1) {
ef0
                int p = pai[j], pp = match[p];
348
                match[p] = j;
fe9
                match[j] = p;
55d
                 j = pp;
797
            }
f70
        }
ba7
        return ans;
fcd }
4.6 Centro de arvore
// Retorna o diametro e o(s) centro(s) da arvore
```

```
// Retorna o diametro e o(s) centro(s) da arvore
// Uma arvore tem sempre um ou dois centros e estes estao no meio do
    diametro
//
```

```
// O(n)
042 vector < int > g[MAX];
df1 int d[MAX], par[MAX];
544 pair <int, vector <int>> center() {
a95
        int f. df:
        function < void(int) > dfs = [&] (int v) {
36d
d47
            if (d[v] > df) f = v, df = d[v];
            for (int u : g[v]) if (u != par[v])
e68
1a5
                d[u] = d[v] + 1, par[u] = v, dfs(u);
90d
        }:
1b0
        f = df = par[0] = -1, d[0] = 0;
41e
        dfs(0);
c2d
        int root = f;
0f6
        f = df = par[root] = -1, d[root] = 0;
14e
        dfs(root):
761
        vector < int > c;
87e
        while (f != -1) {
            if (d[f] == df/2 \text{ or } d[f] == (df+1)/2) \text{ c.push_back}(f);
999
19c
            f = par[f];
3bf
        }
00f
        return {df, c};
9c7 }
4.7 Centroid
// Computa os 2 centroids da arvore
//
// O(n)
97a int n, subsize[MAX];
042 vector < int > g[MAX];
98f void dfs(int k, int p=-1) {
bd2
        subsize[k] = 1:
        for (int i : g[k]) if (i != p) {
6e5
801
            dfs(i, k):
            subsize[k] += subsize[i]:
1b2
5a5 }
```

2e8 int centroid(int k, int p=-1, int size=-1) {

```
e73
        if (size == -1) size = subsize[k];
8df
        for (int i : g[k]) if (i != p) if (subsize[i] > size/2)
            return centroid(i, k, size);
bab
839
        return k;
b6a }
f20 pair<int, int> centroids(int k=0) {
        dfs(k);
909
        int i = centroid(k), i2 = i;
        for (int j : g[i]) if (2*subsize[j] == subsize[k]) i2 = j;
668
0cb
        return {i, i2};
cf4 }
4.8 Centroid decomposition
// decomp(0, k) computa numero de caminhos com 'k' arestas
// Mudar depois do comentario
//
// O(n log(n))
042 vector <int> g[MAX];
ba8 int sz[MAX], rem[MAX];
747 void dfs(vector<int>& path, int i, int l=-1, int d=0) {
547
        path.push_back(d);
75f
        for (int j : g[i]) if (j != 1 and !rem[j]) dfs(path, j, i,
   d+1);
3e9 }
071 int dfs_sz(int i, int l=-1) {
02c
        sz[i] = 1:
        for (int j : g[i]) if (j != l and !rem[j]) sz[i] += dfs_sz(j,
e5c
 i);
191
        return sz[i];
86b }
85a int centroid(int i, int 1, int size) {
        for (int j : g[i]) if (j != 1 and !rem[j] and sz[j] > size / 2)
735
            return centroid(j, i, size);
d9a
        return i;
96e }
d79 ll decomp(int i, int k) {
106
        int c = centroid(i, i, dfs_sz(i));
a67
       rem[c] = 1;
```

```
// gasta O(n) aqui - dfs sem ir pros caras removidos
04b
        11 \text{ ans} = 0:
        vector < int > cnt(sz[i]);
020
878
        cnt[0] = 1;
        for (int j : g[c]) if (!rem[j]) {
0a8
5b4
             vector < int > path;
baf
             dfs(path, j);
             for (int d : path) if (0 \le k-d-1 \text{ and } k-d-1 \le sz[i])
1a1
285
                 ans += cnt[k-d-1];
             for (int d : path) cnt[d+1]++;
e8b
fa2
        }
1c1
        for (int j : g[c]) if (!rem[j]) ans += decomp(j, k);
3f1
        rem[c] = 0:
ba7
        return ans;
193 }
```

#### 4.9 Centroid Tree

```
// Constroi a centroid tree
// p[i] eh o pai de i na centroid-tree
// dist[i][k] = distancia na arvore original entre i
// e o k-esimo ancestral na arvore da centroid
// O(n log(n)) de tempo e memoria
845 vector < int > g[MAX], dist[MAX];
c1e int sz[MAX], rem[MAX], p[MAX];
071 int dfs_sz(int i, int l=-1) {
02c
        sz[i] = 1:
e5c
        for (int j : g[i]) if (j != l and !rem[j]) sz[i] += dfs_sz(j,
   i);
191
        return sz[i];
86b }
85a int centroid(int i, int 1, int size) {
        for (int j : g[i]) if (j != l and !rem[j] and sz[j] > size / 2)
994
735
            return centroid(j, i, size);
d9a
        return i;
96e }
324 void dfs_dist(int i, int 1, int d=0) {
541
        dist[i].push_back(d);
5a1
        for (int j : g[i]) if (j != l and !rem[j])
82a
            dfs_dist(j, i, d+1);
```

```
645 }
27e void decomp(int i, int l = -1) {
        int c = centroid(i, i, dfs_sz(i));
        rem[c] = 1, p[c] = 1:
1b9
534
        dfs_dist(c, c);
a2a
        for (int j : g[c]) if (!rem[j]) decomp(j, c);
ebd }
76c void build(int n) {
235
        for (int i = 0; i < n; i++) rem[i] = 0, dist[i].clear();
867
96b
        for (int i = 0: i < n: i++) reverse(dist[i].begin().
   dist[i].end()):
a78 }
```

# 4.10 Dijkstra

```
// encontra menor distancia de x
// para todos os vertices
// se ao final do algoritmo d[i] = LINF,
// entao x nao alcanca i
// O(m log(n))
eff ll d[MAX]:
c0d vector<pair<int, int>> g[MAX]; // {vizinho, peso}
1a8 int n;
abc void dijkstra(int v) {
22c
        for (int i = 0; i < n; i++) d[i] = LINF;</pre>
a7f
        d[v] = 0:
88c
        priority_queue < pair < ll, int >> pq;
b32
        pq.emplace(0, v);
265
         while (pq.size()) {
             auto [ndist, u] = pq.top(); pq.pop();
a25
953
             if (-ndist > d[u]) continue;
cda
             for (auto [idx, w] : g[u]) if (d[idx] > d[u] + w) {
331
                 d[idx] = d[u] + w:
a84
                 pg.emplace(-d[idx], idx);
c56
             }
e5c
        }
fec }
```

#### 4.11 Dinitz

```
// O(min(m * max flow, n^2 m))
// Grafo com capacidades 1: O(\min(m \text{ sqrt}(m), m * n^2(2/3)))
// Todo vertice tem grau de entrada ou saida 1: O(m sqrt(n))
472 struct dinitz {
61f
        const bool scaling = false; // com scaling -> 0(nm
   log(MAXCAP)),
                                     // com constante alta
206
        int lim;
670
        struct edge {
358
            int to, cap, rev, flow;
7f9
            bool res;
            edge(int to_, int cap_, int rev_, bool res_)
d36
a94
                : to(to_), cap(cap_), rev(rev_), flow(0), res(res_) {}
f70
        };
002
        vector < vector < edge >> g;
216
        vector < int > lev, beg;
a71
        11 F;
190
        dinitz(int n) : g(n), F(0) {}
087
        void add(int a, int b, int c) {
            g[a].emplace_back(b, c, g[b].size(), false);
bae
4c6
            g[b].emplace_back(a, 0, g[a].size()-1, true);
5c2
        }
123
        bool bfs(int s. int t) {
            lev = vector<int>(g.size(), -1); lev[s] = 0;
90f
64c
            beg = vector<int>(g.size(), 0);
8b2
            queue < int > q; q.push(s);
402
            while (q.size()) {
be1
                int u = q.front(); q.pop();
bd9
                for (auto& i : g[u]) {
dbc
                    if (lev[i.to] != -1 or (i.flow == i.cap)) continue;
b4f
                    if (scaling and i.cap - i.flow < lim) continue;</pre>
                    lev[i.to] = lev[u] + 1;
185
8ca
                    q.push(i.to);
                }
f97
e87
            return lev[t] != -1:
0de
742
dfb
        int dfs(int v. int s. int f = INF) {
            if (!f or v == s) return f:
50b
88f
            for (int& i = beg[v]; i < g[v].size(); i++) {</pre>
027
                auto& e = g[v][i];
206
                if (lev[e.to] != lev[v] + 1) continue;
ee0
                int foi = dfs(e.to, s, min(f, e.cap - e.flow));
```

```
749
                if (!foi) continue;
3c5
                e.flow += foi, g[e.to][e.rev].flow -= foi;
45c
                return foi;
            }
618
bb3
            return 0:
4b1
ff6
        11 max flow(int s. int t) {
a86
            for (lim = scaling ? (1<<30) : 1; lim; lim /= 2)</pre>
9d1
                while (bfs(s, t)) while (int ff = dfs(s, t)) F += ff;
4ff
            return F:
8b9
        }
86f }:
// Recupera as arestas do corte s-t
dbd vector<pair<int, int>> get_cut(dinitz& g, int s, int t) {
f07
        g.max_flow(s, t);
68 c
        vector < pair < int , int >> cut;
1b0
        vector<int> vis(g.g.size(), 0), st = {s};
321
        vis[s] = 1;
        while (st.size()) {
3 c 6
b17
            int u = st.back(); st.pop_back();
322
            for (auto e : g.g[u]) if (!vis[e.to] and e.flow < e.cap)</pre>
c17
                vis[e.to] = 1, st.push_back(e.to);
d14
        }
481
        for (int i = 0; i < g.g.size(); i++) for (auto e : g.g[i])</pre>
9d2
            if (vis[i] and !vis[e.to] and !e.res) cut.emplace_back(i,
   e.to):
d1b
        return cut:
1e8 }
4.12 Dominator Tree
// Codigo do Kawakami. Se vira pra usar ai
// build - O(m log(n))
// dominates - 0(1)
```

```
1a8 int n;
bbf namespace d_tree {
042
        vector < int > g[MAX];
        // The dominator tree
b39
        vector<int> tree[MAX]:
5af
        int dfs_l[MAX], dfs_r[MAX];
```

```
// Auxiliary data
        vector < int > rg[MAX], bucket[MAX];
a2e
3ef
        int idom[MAX], sdom[MAX], prv[MAX], pre[MAX];
44b
        int ancestor[MAX], label[MAX];
        vector<int> preorder;
563
76a
        void dfs(int v) {
6a1
            static int t = 0;
db6
            pre[v] = ++t;
767
            sdom[v] = label[v] = v;
a3d
            preorder.push_back(v);
d08
            for (int nxt: g[v]) {
56c
                if (sdom[nxt] == -1) {
                    prv[nxt] = v;
eed
900
                    dfs(nxt);
f48
2b5
                rg[nxt].push_back(v);
            }
5a1
        }
d6a
        int eval(int v) {
62e
c93
            if (ancestor[v] == -1) return v;
a75
            if (ancestor[ancestor[v]] == -1) return label[v];
            int u = eval(ancestor[v]);
f33
b49
            if (pre[sdom[u]] < pre[sdom[label[v]]]) label[v] = u;</pre>
66e
            ancestor[v] = ancestor[u]:
c24
            return label[v]:
0b9
        }
4b2
        void dfs2(int v) {
6a1
            static int t = 0;
330
            dfs_1[v] = t++;
5e0
            for (int nxt: tree[v]) dfs2(nxt);
8e2
            dfs_r[v] = t++;
cfa
        }
c2c
        void build(int s) {
            for (int i = 0; i < n; i++) {</pre>
603
                sdom[i] = pre[i] = ancestor[i] = -1;
e6f
2e1
                rg[i].clear();
50a
                tree[i].clear();
666
                bucket[i].clear();
3ba
772
            preorder.clear();
            dfs(s);
c6c
12b
            if (preorder.size() == 1) return;
            for (int i = int(preorder.size()) - 1; i >= 1; i--) {
3c7
6c6
                int w = preorder[i];
                for (int v: rg[w]) {
a52
                    int u = eval(v);
5c1
```

```
a17
                    if (pre[sdom[u]] < pre[sdom[w]]) sdom[w] = sdom[u];</pre>
018
                }
680
                bucket[sdom[w]].push_back(w);
                ancestor[w] = prv[w];
ea7
b99
                for (int v: bucket[prv[w]]) {
                    int u = eval(v):
5 c 1
                    idom[v] = (u == v) ? sdom[v] : u:
977
aff
2cc
                bucket[prv[w]].clear();
            }
0a3
d0c
            for (int i = 1; i < preorder.size(); i++) {</pre>
6c6
                int w = preorder[i];
14b
                if (idom[w] != sdom[w]) idom[w] = idom[idom[w]];
                tree[idom[w]].push_back(w);
32f
c58
8ac
            idom[s] = sdom[s] = -1;
1b6
            dfs2(s);
d09
        }
        // Whether every path from s to v passes through u
        bool dominates(int u, int v) {
490
c75
            if (pre[v] == -1) return 1; // vacuously true
2ea
            return dfs_l[u] <= dfs_l[v] && dfs_r[v] <= dfs_r[u];</pre>
332
        }
ce9 }:
4.13 Euler Path / Euler Cycle
// Para declarar: 'euler < true > E(n); ' se quiser
// direcionado e com 'n' vertices
// As funcoes retornam um par com um booleano
// indicando se possui o cycle/path que voce pediu,
// e um vector de {vertice, id da aresta para chegar no vertice}
// Se for get_path, na primeira posicao o id vai ser -1
// get_path(src) tenta achar um caminho ou ciclo euleriano
// comecando no vertice 'src'.
// Se achar um ciclo, o primeiro e ultimo vertice serao 'src'.
// Se for um P3, um possiveo retorno seria [0, 1, 2, 0]
// get_cycle() acha um ciclo euleriano se o grafo for euleriano.
// Se for um P3, um possivel retorno seria [0, 1, 2]
// (vertie inicial nao repete)
```

63f template <bool directed=false > struct euler {

//

1a8

// O(n+m)

int n;

```
vector < vector < pair < int , int >>> g;
4c0
d63
        vector < int > used;
30f
        euler(int n_) : n(n_), g(n) {}
50f
        void add(int a. int b) {
            int at = used.size();
4cd
c51
            used.push_back(0);
74e
            g[a].emplace_back(b, at);
fab
            if (!directed) g[b].emplace_back(a, at);
411
d41 #warning chamar para o src certo!
        pair < bool, vector < pair < int, int >>> get_path(int src) {
eed
baf
            if (!used.size()) return {true, {}};
b25
            vector < int > beg(n, 0);
4ec
            for (int& i : used) i = 0;
            // {{vertice, anterior}, label}
            vector<pair<pair<int, int>, int>> ret, st = {{{src, -1}},
363
   -1}}:
            while (st.size()) {
3c6
                int at = st.back().first.first;
8ff
002
                int& it = beg[at];
                while (it < g[at].size() and used[g[at][it].second])</pre>
8a1
   it++:
                if (it == g[at].size()) {
8e4
                    if (ret.size() and ret.back().first.second != at)
944
b82
                         return {false, {}}:
                    ret.push_back(st.back()), st.pop_back();
420
2c0
                } else {
                    st.push_back({{g[at][it].first, at},
   g[at][it].second});
                    used[g[at][it].second] = 1;
eb8
396
                }
            }
b3a
a19
            if (ret.size() != used.size()+1) return {false, {}};
            vector<pair<int, int>> ans;
f77
fdf
            for (auto i : ret) ans.emplace_back(i.first.first,
   i.second):
459
            reverse(ans.begin(), ans.end());
997
            return {true, ans};
844
9b6
        pair < bool, vector < pair < int, int >>> get_cycle() {
            if (!used.size()) return {true, {}};
baf
ad1
            int src = 0:
            while (!g[src].size()) src++;
34b
687
            auto ans = get_path(src);
            if (!ans.first or ans.second[0].first !=
33 c
   ans.second.back().first)
```

```
b82          return {false, {}};
350          ans.second[0].second = ans.second.back().second;
8b8          ans.second.pop_back();
ba7          return ans;
48f     }
711 };
```

## 4.14 Euler Tour Tree

```
// Mantem uma floresta enraizada dinamicamente
// e permite queries/updates em sub-arvore
//
// Chamar ETT E(n, v), passando n = numero de vertices
// e v = vector com os valores de cada vertice (se for vazio,
// constroi tudo com 0
//
// link(v, u) cria uma aresta de v pra u, de forma que u se torna
// o pai de v (eh preciso que v seja raiz anteriormente)
// cut(v) corta a resta de v para o pai
// query(v) retorna a soma dos valores da sub-arvore de v
// update(v, val) soma val em todos os vertices da sub-arvore de v
// update_v(v, val) muda o valor do vertice v para val
// is_in_subtree(v, u) responde se o vertice u esta na sub-arvore de v
// Tudo O(log(n)) com alta probabilidade
878 mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
9f9 template < typename T > struct ETT {
        // treap
3c9
        struct node {
ed1
            node *1, *r, *p;
fa4
            int pr, sz;
875
            T val, sub, lazy;
53e
            int id:
ffd
            bool f; // se eh o 'first'
            int qt_f; // numero de firsts na subarvore
5ef
            node(int id_, T v, bool f_ = 0) : l(NULL), r(NULL),
7a8
   p(NULL), pr(rng()),
62b
                sz(1), val(v), sub(v), lazy(), id(id_), f(f_),
   qt_f(f_) {}
a9c
            void prop() {
d09
                if (lazy != T()) {
021
                    if (f) val += lazy;
971
                    sub += lazy*sz;
```

```
b87
                     if (1) 1->lazy += lazy;
d3b
                     if (r) r->lazy += lazy;
30c
                }
bfd
                lazv = T();
0fc
            }
01e
            void update() {
8da
                 sz = 1, sub = val, gt f = f:
                if (1) 1->prop(), sz += 1->sz, sub += 1->sub, qt_f +=
   1->qt_f;
                if (r) r \rightarrow prop(), sz += r \rightarrow sz, sub += r \rightarrow sub, qt_f +=
117
   r->qt_f;
            }
ccb
bff
        }:
bb7
        node* root;
73c
        int size(node* x) { return x ? x->sz : 0; }
bcf
        void join(node* 1, node* r, node*& i) { // assume que 1 < r</pre>
986
            if (!1 or !r) return void(i = 1 ? 1 : r);
161
            1->prop(), r->prop();
            if (1->pr > r->pr) join(1->r, r, 1->r), 1->r->p = i = 1;
ff5
982
            else join(1, r->1, r->1), r->1->p = i = r;
bda
            i->update();
84d
        }
a20
        void split(node* i, node*& 1, node*& r, int v, int key = 0) {
26a
            if (!i) return void(r = 1 = NULL);
c89
            i->prop();
d9e
            if (kev + size(i->1) < v) {</pre>
                 split(i->r, i->r, r, v, key+size(i->l)+1), l = i;
448
a21
                if (r) r->p = NULL;
                if (i->r) i->r->p = i;
6e8
396
            } else {
98d
                 split(i->1, 1, i->1, v, key), r = i;
                if (1) 1 \rightarrow p = NULL:
5a3
                if (i->1) i->1->p = i;
899
            }
18b
bda
            i->update();
134
        }
ac7
        int get_idx(node* i) {
6cf
            int ret = size(i->1);
482
            for (; i->p; i = i->p) {
fbf
                node* pai = i->p;
                if (i != pai->1) ret += size(pai->1) + 1;
8a6
            }
e22
edf
            return ret;
479
048
        node* get_min(node* i) {
```

```
433
            if (!i) return NULL;
f8e
            return i->1 ? get_min(i->1) : i;
0de
        }
f03
        node* get_max(node* i) {
433
            if (!i) return NULL:
424
            return i->r ? get_max(i->r) : i;
e92
        }
        // fim da treap
        vector < node *> first, last;
4fb
f82
        ETT(int n, vector<T> v = {}) : root(NULL), first(n), last(n) {
с5е
            if (!v.size()) v = vector < T > (n):
            for (int i = 0; i < n; i++) {</pre>
603
                first[i] = last[i] = new node(i, v[i], 1);
a00
469
                join(root, first[i], root);
8ac
            }
ec3
83f
        ETT(const ETT& t) { throw logic_error("Nao copiar a ETT!"); }
c09
609
            vector < node *> q = {root};
402
            while (q.size()) {
                node* x = q.back(); q.pop_back();
e5d
ee9
                if (!x) continue;
1c7
                q.push_back(x->1), q.push_back(x->r);
bf0
                delete x:
653
            }
        }
672
153
        pair<int, int> get_range(int i) {
670
            return {get_idx(first[i]), get_idx(last[i])};
ada
        void link(int v, int u) { // 'v' tem que ser raiz
7af
            auto [lv. rv] = get range(v):
890
            int ru = get_idx(last[u]);
f13
4b4
            node* V;
df9
            node *L, *M, *R;
117
            split(root, M, R, rv+1), split(M, L, M, lv);
f1e
            V = M;
a28
            join(L, R, root);
            split(root, L, R, ru+1);
e66
367
            join(L, V, L);
            join(L, last[u] = new node(u, T() /* elemento neutro */),
7e8
   L);
a28
            join(L, R, root);
```

```
8d9
4e6
        void cut(int v) {
892
            auto [1, r] = get_range(v);
df9
            node *L. *M. *R:
            split(root, M, R, r+1), split(M, L, M, 1);
dca
de6
            node *LL = get max(L), *RR = get min(R);
710
            if (LL and RR and LL->id == RR->id) { // remove duplicata
e8b
                 if (last[RR->id] == RR) last[RR->id] = LL;
992
                 node *A. *B:
6b3
                 split(R, A, B, 1);
10c
                 delete A:
9d5
                 R = B:
7c0
a28
            join(L, R, root);
a0d
            join(root, M, root);
6ff
808
        T query(int v) {
892
            auto [1, r] = get_range(v);
df9
            node *L. *M. *R:
dca
            split(root, M, R, r+1), split(M, L, M, 1);
d43
            T ans = M->sub;
            join(L, M, M), join(M, R, root);
69d
ba7
            return ans;
ede
93b
        void update(int v, T val) { // soma val em todo mundo da
   subarvore
892
            auto [1, r] = get_range(v);
df9
            node *L, *M, *R;
dca
            split(root, M, R, r+1), split(M, L, M, 1);
409
            M->lazy += val;
69d
            join(L, M, M), join(M, R, root);
61c
129
        void update_v(int v, T val) { // muda o valor de v pra val
            int 1 = get_idx(first[v]);
ac1
df9
            node *L, *M, *R;
d0c
            split(root, M, R, l+1), split(M, L, M, 1);
25e
            M \rightarrow val = M \rightarrow sub = val;
69d
            join(L, M, M), join(M, R, root);
630
934
        bool is_in_subtree(int v, int u) { // se u ta na subtree de v
890
            auto [lv, rv] = get_range(v);
6ec
            auto [lu, ru] = get_range(u);
            return lv <= lu and ru <= rv;</pre>
732
a21
        }
355
        void print(node* i) {
```

## 4.15 Floyd-Warshall

```
// encontra o menor caminho entre todo
// par de vertices e detecta ciclo negativo
// returna 1 sse ha ciclo negativo
// d[i][i] deve ser 0
// para i != j, d[i][j] deve ser w se ha uma aresta
// (i, j) de peso w, INF caso contrario
//
// O(n^3)
1a8 int n;
ae5 int d[MAX][MAX];
73c bool floyd_warshall() {
        for (int k = 0: k < n: k++)
830
        for (int i = 0; i < n; i++)</pre>
f90
        for (int j = 0; j < n; j++)
            d[i][j] = min(d[i][j], d[i][k] + d[k][j]);
0ab
830
        for (int i = 0; i < n; i++)
            if (d[i][i] < 0) return 1;</pre>
753
        return 0;
bb3
192 }
```

# 4.16 Functional Graph

```
// rt[i] fala o ID da raiz associada ao vertice i
// d[i] fala a profundidade (0 sse ta no ciclo)
// pos[i] fala a posicao de i no array que eh a concat. dos ciclos
// build(f, val) recebe a funcao f e o custo de ir de
// i para f[i] (por default, val = f)
// f_k(i, k) fala onde i vai parar se seguir k arestas
// path(i, k) fala o custo (soma) seguir k arestas a partir de i
// Se quiser outra operacao, da pra alterar facil o codigo
```

```
// Codigo um pouco louco, tenho que admitir
// build - O(n)
// f_k - O(log(min(n, k)))
// path - O(log(min(n, k)))
6ef namespace func_graph {
        int n;
1a8
ce2
        int f[MAX], vis[MAX], d[MAX];
        int p[MAX], pp[MAX], rt[MAX], pos[MAX];
ebd
        int sz[MAX], comp;
6a9
        vector < vector < int >> ciclo;
405
        11 val[MAX], jmp[MAX], seg[2*MAX];
97c
        11 op(11 a, 11 b) { return a+b; }; // mudar a operacao aqui
27b
        void dfs(int i, int t = 2) {
9c9
            vis[i] = t;
f09
            if (vis[f[i]] >= 2) \{ // comeca ciclo - f[i] eh o rep.
                d[i] = 0, rt[i] = comp;
e0a
                sz[comp] = t - vis[f[i]] + 1;
74c
                p[i] = pp[i] = i, jmp[i] = val[i];
97b
                ciclo.emplace_back();
15c
bfb
                ciclo.back().push_back(i);
a22
            } else {
                if (!vis[f[i]]) dfs(f[i], t+1);
c16
8c0
                rt[i] = rt[f[i]];
                if (sz[comp]+1) { // to no ciclo
195
d0f
                    d[i] = 0:
97b
                    p[i] = pp[i] = i, jmp[i] = val[i];
bfb
                    ciclo.back().push_back(i);
                } else { // nao to no ciclo
c20
00d
                    d[i] = d[f[i]]+1, p[i] = f[i];
                    pp[i] = 2*d[pp[f[i]]] == d[pp[pp[f[i]]]]+d[f[i]] ?
   pp[pp[f[i]]] : f[i];
                    jmp[i] = pp[i] == f[i] ? val[i] : op(val[i],
114
   op(jmp[f[i]], jmp[pp[f[i]]]));
db8
                }
003
e4a
            if (f[ciclo[rt[i]][0]] == i) comp++; // fim do ciclo
29a
            vis[i] = 1;
0ba
        void build(vector<int> f_, vector<int> val_ = {}) {
1da
            n = f_size(), comp = 0;
bcb
            if (!val_.size()) val_ = f_;
527
830
            for (int i = 0; i < n; i++)</pre>
                f[i] = f_[i], val[i] = val_[i], vis[i] = 0, sz[i] = -1;
998
```

```
e74
            ciclo.clear();
158
            for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);</pre>
6bb
            for (auto& c : ciclo) {
daa
336
                reverse(c.begin(), c.end());
                for (int j : c) {
ea5
85b
                    pos[j] = t;
948
                     seg[n+t] = val[j];
c82
                    t++;
                }
25e
cbc
            for (int i = n-1; i; i--) seg[i] = op(seg[2*i],
   seg[2*i+1]):
90b
        }
        int f_k(int i, ll k) {
283
1b1
            while (d[i] and k) {
77b
                int big = d[i] - d[pp[i]];
                if (big <= k) k -= big, i = pp[i];</pre>
ded
                else k--, i = p[i];
584
09c
            }
77e
            if (!k) return i;
            return ciclo[rt[i]][(pos[i] - pos[ciclo[rt[i]][0]] + k) %
a 19
   sz[rt[i]];
f34
        }
047
        ll path(int i, ll k) {
3cf
            auto query = [&](int 1, int r) {
3e4
                11 a = 0:
47a
                for (1 += n, r += n; 1 <= r; ++1/=2, --r/=2) {
27 e
                    if (1\%2 == 1) q = op(q, seg[1]);
                    if (r\%2 == 0) q = op(q, seg[r]);
1f2
598
                }
bef
                return q;
6e1
            }:
b73
            11 \text{ ret} = 0;
1b1
            while (d[i] and k) {
77b
                int big = d[i] - d[pp[i]];
327
                if (big <= k) k -= big, ret = op(ret, jmp[i]), i =</pre>
   pp[i];
f9e
                 else k--, ret = op(ret, val[i]), i = p[i];
            }
7e3
еЗс
            if (!k) return ret;
a9e
            int first = pos[ciclo[rt[i]][0]], last =
   pos[ciclo[rt[i]].back()];
            // k/sz[rt[i]] voltas completas
            if (k/sz[rt[i]]) ret = op(ret, k/sz[rt[i]] * query(first,
430
```

```
last));
9af
            k %= sz[rt[i]];
e3c
            if (!k) return ret:
            int l = pos[i], r = first + (pos[i] - first + k - 1) %
8ea
   sz[rt[i]];
982
            if (1 <= r) return op(ret, query(1, r));</pre>
            return op(ret, op(query(1, last), query(first, r)));
687
380
        }
51f }
4.17 HLD - aresta
// SegTree de soma
// query / update de soma das arestas
// Complexidades:
// build - O(n)
// query_path - O(log^2 (n))
// update_path - O(log^2(n))
// query_subtree - O(log(n))
// update_subtree - O(log(n))
// namespace seg { ... }
826 namespace hld {
        vector < pair < int , int > > g[MAX];
c0d
e65
        int pos[MAX], sz[MAX];
7c0
        int sobe[MAX], pai[MAX];
096
        int h[MAX], v[MAX], t;
0ce
        void build_hld(int k, int p = -1, int f = 1) {
            v[pos[k] = t++] = sobe[k]; sz[k] = 1;
180
            for (auto& i : g[k]) if (i.first != p) {
418
dd2
                auto [u, w] = i;
                sobe[u] = w; pai[u] = k;
a76
0 c 1
                h[u] = (i == g[k][0] ? h[k] : u);
                build_hld(u, k, f); sz[k] += sz[u];
da7
                if (sz[u] > sz[g[k][0].first] or g[k][0].first == p)
865
9a3
                    swap(i, g[k][0]);
804
667
            if (p*f == -1) build_hld(h[k] = k, -1, t = 0);
4dd
1f8
        void build(int root = 0) {
a34
            t = 0;
```

```
295
            build_hld(root);
c83
            seg::build(t, v);
ea2
        }
3fc
        ll query_path(int a, int b) {
2d5
            if (a == b) return 0:
            if (pos[a] < pos[b]) swap(a, b);</pre>
aa1
29b
            if (h[a] == h[b]) return seg::query(pos[b]+1, pos[a]);
fca
            return seg::query(pos[h[a]], pos[a]) +
   query_path(pai[h[a]], b);
87f
920
        void update_path(int a, int b, int x) {
d54
            if (a == b) return;
            if (pos[a] < pos[b]) swap(a, b);</pre>
aa1
            if (h[a] == h[b]) return (void)seg::update(pos[b]+1,
881
   pos[a], x);
701
            seg::update(pos[h[a]], pos[a], x); update_path(pai[h[a]],
   b, x);
dbf
d0a
        11 query_subtree(int a) {
b9f
            if (sz[a] == 1) return 0;
2f6
            return seg::query(pos[a]+1, pos[a]+sz[a]-1);
77f
        }
        void update_subtree(int a, int x) {
acc
a5a
            if (sz[a] == 1) return;
9cd
            seg::update(pos[a]+1, pos[a]+sz[a]-1, x);
a46
        }
7be
        int lca(int a, int b) {
aa1
            if (pos[a] < pos[b]) swap(a, b);
            return h[a] == h[b] ? b : lca(pai[h[a]], b);
ca5
219
        }
599 }
4.18 HLD - vertice
// SegTree de soma
// query / update de soma dos vertices
//
// Complexidades:
// build - O(n)
// \text{ query_path - O(log^2 (n))}
// update_path - O(log^2 (n))
// query_subtree - O(log(n))
// update_subtree - O(log(n))
```

```
// namespace seg { ... }
826 namespace hld {
        vector < int > g[MAX];
042
        int pos[MAX], sz[MAX];
e65
        int peso[MAX], pai[MAX];
bd4
096
        int h[MAX], v[MAX], t;
        void build_hld(int k, int p = -1, int f = 1) {
Осе
            v[pos[k] = t++] = peso[k]; sz[k] = 1;
b18
b94
            for (auto& i : g[k]) if (i != p) {
78d
                pai[i] = k;
26e
                h[i] = (i == g[k][0] ? h[k] : i);
                build_hld(i, k, f); sz[k] += sz[i];
193
                if (sz[i] > sz[g[k][0]] or g[k][0] == p) swap(i,
cd1
   g[k][0]);
d94
            if (p*f == -1) build_hld(h[k] = k, -1, t = 0);
667
a83
1f8
        void build(int root = 0) {
            t = 0;
a34
            build_hld(root);
295
c83
            seg::build(t, v);
ea2
3fc
        11 query_path(int a, int b) {
aa1
            if (pos[a] < pos[b]) swap(a, b);</pre>
4bf
            if (h[a] == h[b]) return seg::query(pos[b], pos[a]);
            return seg::query(pos[h[a]], pos[a]) +
fca
   query_path(pai[h[a]], b);
c17
920
        void update_path(int a, int b, int x) {
            if (pos[a] < pos[b]) swap(a, b);</pre>
aa1
198
            if (h[a] == h[b]) return (void) seg::update(pos[b], pos[a],
   x);
701
            seg::update(pos[h[a]], pos[a], x); update_path(pai[h[a]],
   b, x);
421
        11 query_subtree(int a) {
d0a
            return seg::query(pos[a], pos[a]+sz[a]-1);
b3e
ba2
        }
        void update_subtree(int a, int x) {
acc
a22
            seg::update(pos[a], pos[a]+sz[a]-1, x);
480
7be
        int lca(int a, int b) {
```

```
if (pos[a] < pos[b]) swap(a, b);</pre>
aa1
            return h[a] == h[b] ? b : lca(pai[h[a]], b);
ca5
219
       }
de3 }
4.19 HLD sem Update
// query de min do caminho
//
// Complexidades:
// build - O(n)
// query_path - O(log(n))
826 namespace hld {
c0d
        vector<pair<int, int> > g[MAX];
        int pos[MAX], sz[MAX];
e65
        int sobe[MAX], pai[MAX];
7c0
096
        int h[MAX], v[MAX], t;
ea2
        int men[MAX], seg[2*MAX];
        void build_hld(int k, int p = -1, int f = 1) {
0ce
180
            v[pos[k] = t++] = sobe[k]; sz[k] = 1;
418
            for (auto& i : g[k]) if (i.first != p) {
1f5
                sobe[i.first] = i.second; pai[i.first] = k;
                h[i.first] = (i == g[k][0] ? h[k] : i.first);
6fa
                men[i.first] = (i == g[k][0] ? min(men[k], i.second) :
87b
   i.second);
                build_hld(i.first, k, f); sz[k] += sz[i.first];
4b2
bc3
                if (sz[i.first] > sz[g[k][0].first] or g[k][0].first
   == p)
9a3
                    swap(i, g[k][0]);
ea4
667
            if (p*f == -1) build_hld(h[k] = k, -1, t = 0);
8ec
1f8
        void build(int root = 0) {
a34
            t = 0:
295
            build_hld(root);
            for (int i = 0; i < t; i++) seg[i+t] = v[i];</pre>
3ae
            for (int i = t-1; i; i--) seg[i] = min(seg[2*i],
   seg[2*i+1]);
```

ea5

f04

490

aa1

}

int query\_path(int a, int b) {

if (a == b) return INF:

if (pos[a] < pos[b]) swap(a, b);

### 4.20 Isomorfismo de arvores

```
// thash() retorna o hash da arvore (usando centroids como vertices
   especiais).
// Duas arvores sao isomorfas sse seu hash eh o mesmo
//
// O(|V|.log(|V|))
91f map < vector < int >, int > mphash;
df6 struct tree {
1a8
        int n:
        vector < vector < int >> g;
789
347
        vector < int > sz, cs;
        tree(int n_{-}): n(n_{-}), g(n_{-}), sz(n_{-}) {}
1b5
76b
        void dfs_centroid(int v, int p) {
588
            sz[v] = 1;
fa7
            bool cent = true:
            for (int u : g[v]) if (u != p) {
18e
365
                 dfs_centroid(u, v), sz[v] += sz[u];
e90
                if(sz[u] > n/2) cent = false;
ece
            }
1f6
            if (cent and n - sz[v] \le n/2) cs.push_back(v);
368
        int fhash(int v, int p) {
784
544
            vector < int > h;
            for (int u : g[v]) if (u != p) h.push_back(fhash(u, v));
332
1c9
            sort(h.begin(), h.end());
            if (!mphash.count(h)) mphash[h] = mphash.size();
3ac
bbc
            return mphash[h];
748
        }
38f
        11 thash() {
23a
            cs.clear():
3a5
            dfs_centroid(0, -1);
16d
            if (cs.size() == 1) return fhash(cs[0], -1);
```

## 4.21 Kosaraju

```
// O(n + m)
1a8 int n;
042 vector <int> g[MAX];
58d vector < int > gi[MAX]; // grafo invertido
c5a int vis[MAX];
ee6 stack<int> S:
a52 int comp[MAX]; // componente conexo de cada vertice
1ca void dfs(int k) {
        vis[k] = 1;
59a
54f
        for (int i = 0; i < (int) g[k].size(); i++)</pre>
            if (!vis[g[k][i]]) dfs(g[k][i]);
845
58f
        S.push(k);
89c }
436 void scc(int k, int c) {
        vis[k] = 1;
59a
52c
        comp[k] = c;
        for (int i = 0; i < (int) gi[k].size(); i++)</pre>
ff0
            if (!vis[gi[k][i]]) scc(gi[k][i], c);
bf6
088 }
db8 void kosaraju() {
        for (int i = 0; i < n; i++) vis[i] = 0;</pre>
158
        for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);</pre>
991
        for (int i = 0; i < n; i++) vis[i] = 0;</pre>
d32
        while (S.size()) {
70b
            int u = S.top();
7de
            S.pop();
f43
            if (!vis[u]) scc(u, u);
        }
207
e21 }
```

## 4.22 Kruskal

```
// Gera e retorna uma AGM e seu custo total a partir do vetor de
   arestas (edg)
// do grafo
//
// O(m log(m) + m a(m))
// 864875
1b9 vector<tuple<int, int, int>> edg; // {peso,[x,y]}
// DSU em O(a(n))
4a6 void dsu_build();
d78 int find(int a);
369 void unite(int a, int b);
c67 pair<ll, vector<tuple<int, int, int>>> kruskal(int n) {
8d2
        dsu_build(n);
        sort(edg.begin(), edg.end());
e31
854
        11 cost = 0;
979
        vector<tuple<int, int, int>> mst;
fea
       for (auto [w,x,y] : edg) if (find(x) != find(y)) {
9de
            mst.emplace_back(w, x, y);
45f
            cost += w:
05a
            unite(x,y);
ca2
5df
        return {cost, mst};
b6a }
4.23 Kuhn
// Computa matching maximo em grafo bipartido
// 'n' e 'm' sao quantos vertices tem em cada particao
// chamar add(i, j) para add aresta entre o cara i
// da particao A, e o cara j da particao B
// (entao i < n, j < m)
// Para recuperar o matching, basta olhar 'ma' e 'mb'
// 'recover' recupera o min vertex cover como um par de
// {caras da particao A, caras da particao B}
//
// O(|V| * |E|)
// Na pratica, parece rodar tao rapido quanto o Dinic
878 mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
```

6c6 struct kuhn {

```
14e
        int n, m;
789
        vector < vector < int >> g;
d3f
        vector < int > vis, ma, mb;
40e
        kuhn(int n_, int m_) : n(n_), m(m_), g(n),
            vis(n+m), ma(n, -1), mb(m, -1) {}
8af
ba6
        void add(int a, int b) { g[a].push_back(b); }
        bool dfs(int i) {
caf
29a
            vis[i] = 1;
29Ъ
            for (int j : g[i]) if (!vis[n+j]) {
8 c 9
                 vis[n+i] = 1:
2cf
                 if (mb[j] == -1 or dfs(mb[j])) {
bfe
                     ma[i] = j, mb[j] = i;
8a6
                     return true;
b17
                }
82a
            }
d1f
            return false;
4ef
        }
bf7
        int matching() {
1ae
            int ret = 0, aum = 1;
5a8
            for (auto& i : g) shuffle(i.begin(), i.end(), rng);
392
            while (aum) {
618
                for (int j = 0; j < m; j++) vis[n+j] = 0;
c5d
                aum = 0:
                for (int i = 0; i < n; i++)</pre>
830
0.1f
                     if (ma[i] == -1 and dfs(i)) ret++, aum = 1;
085
            }
edf
            return ret;
        }
2ee
b0d };
ebf pair<vector<int>, vector<int>> recover(kuhn& K) {
e80
        K.matching();
50c
        int n = K.n, m = K.m;
        for (int i = 0; i < n+m; i++) K.vis[i] = 0;</pre>
9d0
bde
        for (int i = 0; i < n; i++) if (K.ma[i] == -1) K.dfs(i);
8ad
        vector<int> ca. cb:
576
        for (int i = 0; i < n; i++) if (!K.vis[i]) ca.push_back(i);</pre>
        for (int i = 0; i < m; i++) if (K.vis[n+i]) cb.push_back(i);</pre>
f24
aad
        return {ca, cb};
55f }
```

# 4.24 LCA com binary lifting

```
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b, lca(a, b) = a
// MAX2 = ceil(log(MAX))
//
// Complexidades:
// build - O(n log(n))
// lca - O(log(n))
677 vector < vector < int > > g(MAX);
41c int n, p;
e75 int pai[MAX2][MAX];
999 int in[MAX], out[MAX];
1ca void dfs(int k) {
fdf
        in[k] = p++;
54f
        for (int i = 0; i < (int) g[k].size(); i++)</pre>
            if (in[g[k][i]] == -1) {
9b7
ba6
                pai[0][g[k][i]] = k;
c38
                dfs(g[k][i]);
            }
e2d
26f
        out[k] = p++;
691 }
c11 void build(int raiz) {
        for (int i = 0; i < n; i++) pai[0][i] = i;</pre>
c63
        p = 0, memset(in, -1, sizeof in);
        dfs(raiz):
ecb
        // pd dos pais
511
        for (int k = 1; k < MAX2; k++) for (int i = 0; i < n; i++)
d38
            pai[k][i] = pai[k - 1][pai[k - 1][i]];
530 }
00f bool anc(int a. int b) { // se a eh ancestral de b
        return in[a] <= in[b] and out[a] >= out[b];
bfe
2d6 }
7be int lca(int a, int b) {
86d
       if (anc(a, b)) return a;
e52
        if (anc(b, a)) return b;
        // sobe a
        for (int k = MAX2 - 1; k >= 0; k--)
f70
acf
            if (!anc(pai[k][a], b)) a = pai[k][a];
847
        return pai[0][a];
5c4 }
```

```
// Alternativamente:
// 'binary lifting' gastando O(n) de memoria
// Da pra add folhas e fazer queries online
// 3 vezes o tempo do binary lifting normal
// build - O(n)
// kth, lca, dist - O(log(n))
9c6 int d[MAX], p[MAX], pp[MAX];
d40 void set_root(int i) { p[i] = pp[i] = i, d[i] = 0; }
e9d void add_leaf(int i, int u) {
e0b
        p[i] = u, d[i] = d[u]+1;
b15
        pp[i] = 2*d[pp[u]] == d[pp[pp[u]]]+d[u] ? pp[pp[u]] : u;
33f }
c37 int kth(int i, int k) {
        int dd = max(0, d[i]-k):
        while (d[i] > dd) i = d[pp[i]] >= dd ? pp[i] : p[i];
935
d9a
        return i:
f3c }
7be int lca(int a, int b) {
a69
        if (d[a] < d[b]) swap(a, b);</pre>
6cd
        while (d[a] > d[b]) a = d[pp[a]] >= d[b] ? pp[a] : p[a];
984
        while (a != b) {
932
            if (pp[a] != pp[b]) a = pp[a], b = pp[b];
e7c
            else a = p[a], b = p[b];
        }
4ea
3f5
        return a;
21d }
4fe int dist(int a, int b) { return d[a]+d[b]-2*d[lca(a,b)]; }
042 vector < int > g[MAX];
3ab void build(int i, int pai=-1) {
5cf
        if (pai == -1) set_root(i);
15f
        for (int j : g[i]) if (j != pai) {
d31
            add_leaf(j, i);
b21
            build(j, i);
        }
43b
74a }
```

#### 4.25 LCA com HLD

```
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b. lca(a, b) = a
// Para buildar pasta chamar build(root)
// anc(a, b) responde se 'a' eh ancestral de 'b'
//
// Complexidades:
// build - O(n)
// lca - O(log(n))
// anc - 0(1)
042 vector < int > g[MAX];
713 int pos[MAX], h[MAX], sz[MAX];
ff1 int pai[MAX], t;
8bf void build(int k, int p = -1, int f = 1) {
        pos[k] = t++; sz[k] = 1;
bce
        for (int& i : g[k]) if (i != p) {
e26
78d
            pai[i] = k:
26e
            h[i] = (i == g[k][0] ? h[k] : i);
            build(i, k, f); sz[k] += sz[i];
cb8
cd1
           if (sz[i] > sz[g[k][0]] or g[k][0] == p) swap(i, g[k][0]);
917
3da
        if (p*f == -1) t = 0, h[k] = k, build(k, -1, 0);
1b9 }
7be int lca(int a, int b) {
        if (pos[a] < pos[b]) swap(a, b);</pre>
ca5
        return h[a] == h[b] ? b : lca(pai[h[a]], b);
219 }
00f bool anc(int a, int b) {
        return pos[a] <= pos[b] and pos[b] <= pos[a]+sz[a]-1;</pre>
db5
272 }
4.26 LCA com RMQ
// Assume que um vertice eh ancestral dele mesmo, ou seja,
// se a eh ancestral de b. lca(a, b) = a
// dist(a, b) retorna a distancia entre a e b
//
// Complexidades:
// build - O(n)
```

```
// lca - O(1)
// dist - 0(1)
1a5 template < typename T > struct rmg {
517
        vector <T> v:
fcc
        int n; static const int b = 30;
        vector < int > mask, t;
70e
18e
        int op(int x, int y) { return v[x] < v[y] ? x : y; }
        int msb(int x) { return __builtin_clz(1)-__builtin_clz(x); }
ee1
6ad
        rmq() {}
43c
        rmq(const \ vector < T > \& \ v_) : v(v_), n(v.size()), mask(n), t(n) {
2e5
            for (int i = 0, at = 0; i < n; mask[i++] = at |= 1) {
                at = (at << 1) &((1 << b) -1):
a61
76a
                while (at and op(i, i-msb(at&-at)) == i) at ^= at&-at;
53c
243
            for (int i = 0; i < n/b; i++) t[i] =
   b*i+b-1-msb(mask[b*i+b-1]);
            for (int j = 1; (1<<j) <= n/b; j++) for (int i = 0;
39d
   i+(1<< i) <= n/b: i++)
                t[n/b*j+i] = op(t[n/b*(j-1)+i],
   t[n/b*(j-1)+i+(1<<(j-1))]);
2d3
        int small(int r, int sz = b) { return
c92
   r-msb(mask[r]&((1<<sz)-1));}
b7a
        T query(int 1, int r) {
            if (r-l+1 <= b) return small(r, r-l+1);</pre>
27b
7bf
            int ans = op(small(1+b-1), small(r));
e80
            int x = 1/b+1, y = r/b-1;
e25
            if (x <= y) {
                int j = msb(y-x+1);
a4e
002
                ans = op(ans, op(t[n/b*j+x], t[n/b*j+y-(1<<j)+1]));
4b6
ba7
            return ans:
6bf
        }
021 };
065 namespace lca {
042
        vector < int > g[MAX];
8ec
        int v[2*MAX], pos[MAX], dep[2*MAX];
8bd
        int t:
2de
        rmq<int> RMQ;
4cf
        void dfs(int i, int d = 0, int p = -1) {
c97
            v[t] = i, pos[i] = t, dep[t++] = d;
            for (int j : g[i]) if (j != p) {
cac
                dfs(i, d+1, i):
8ec
```

```
cf2
                v[t] = i, dep[t++] = d;
            }
843
d6a
        }
789
        void build(int n, int root) {
            t = 0:
a34
            dfs(root):
14e
3f4
            RMQ = rmq < int > (vector < int > (dep. dep+2*n-1)):
657
        }
7be
        int lca(int a, int b) {
            a = pos[a], b = pos[b];
ab7
9c0
            return v[RMQ.query(min(a, b), max(a, b))];
5db
b5d
        int dist(int a. int b) {
670
            return dep[pos[a]] + dep[pos[b]] - 2*dep[pos[lca(a, b)]];
5b7
        }
645 }
```

## 4.27 Line Tree

```
// Reduz min-query em arvore para RMQ
// Se o grafo nao for uma arvore, as queries
// sao sobre a arvore geradora maxima
// Queries de minimo
//
// build - O(n log(n))
// query - O(log(n))
1a8 int n;
3ae namespace linetree {
f37
        int id[MAX], seg[2*MAX], pos[MAX];
43f
        vector < int > v[MAX], val[MAX];
430
        vector < pair < int , pair < int , int > > ar;
        void add(int a, int b, int p) { ar.push_back({p, {a, b}}); }
dc6
0a8
        void build() {
b09
            sort(ar.rbegin(), ar.rend());
            for (int i = 0; i < n; i++) id[i] = i, v[i] = \{i\},
   val[i].clear():
            for (auto i : ar) {
8bb
c91
                int a = id[i.second.first], b = id[i.second.second];
                if (a == b) continue;
c58
                if (v[a].size() < v[b].size()) swap(a, b);</pre>
fb8
                for (auto j : v[b]) id[j] = a, v[a].push_back(j);
482
                val[a].push_back(i.first);
78b
                for (auto j : val[b]) val[a].push_back(j);
```

```
e39
                v[b].clear(), val[b].clear();
012
            }
8e8
            vector < int > vv;
            for (int i = 0; i < n; i++) for (int j = 0; j < 1
2ce
   v[i].size(); j++) {
                pos[v[i][j]] = vv.size();
e52
941
                if (j + 1 < v[i].size()) vv.push_back(val[i][j]);</pre>
                else vv.push_back(0);
1cb
475
            for (int i = n; i < 2*n; i++) seg[i] = vv[i-n];</pre>
bb4
69e
            for (int i = n-1; i; i--) seg[i] = min(seg[2*i],
   seg[2*i+1]);
9fe
        }
4ea
        int query(int a, int b) {
596
            if (id[a] != id[b]) return 0; // nao estao conectados
ab7
            a = pos[a], b = pos[b];
d11
            if (a > b) swap(a, b);
199
            b--:
38a
            int ans = INF;
            for (a += n, b += n; a \le b; ++a/=2, --b/=2) ans =
513
   min({ans, seg[a], seg[b]});
            return ans:
ba7
        }
952
00f };
4.28 Link-cut Tree
// Todas as operacoes sao O(log(n)) amortizado
1ef namespace lct {
3c9
        struct node {
19f
            int p, ch[2];
```

```
// Link-cut tree padrao
             node() \{ p = ch[0] = ch[1] = -1; \}
062
f43
        };
5f3
        node t[MAX];
971
        bool is_root(int x) {
657
             return t[x].p == -1 or (t[t[x].p].ch[0] != x and
    t[t[x].p].ch[1] != x);
cf1
ed6
        void rotate(int x) {
497
            int p = t[x].p, pp = t[p].p;
fc4
             if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
```

```
251
            bool d = t[p].ch[0] == x;
            t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
461
a76
            if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
8fa
            t[x].p = pp, t[p].p = x;
49b
        }
07с
        void splay(int x) {
18c
            while (!is root(x)) {
497
                int p = t[x].p, pp = t[p].p;
0 c 5
                if (!is_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch[0]
   == x) ? x : p);
                rotate(x);
64f
            }
4fa
f16
        int access(int v) {
0eb
            int last = -1:
01a
            for (int w = v; w+1; last = w, splay(v), w = t[v].p)
024
                splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
3d3
            return last;
0a4
e89
        int find_root(int v) {
5e3
            access(v):
3de
            while (t[v].ch[0]+1) v = t[v].ch[0];
f05
            return splay(v), v;
ee7
        }
142
        void link(int v, int w) { // v deve ser raiz
5e3
            access(v):
10d
            t[v].p = w;
c56
4e6
        void cut(int v) { // remove aresta de v pro pai
5e3
            access(v);
264
            t[v].ch[0] = t[t[v].ch[0]].p = -1;
5f5
        }
        int lca(int v, int w) {
bbb
948
            return access(v). access(w):
b6d
e4e }
4.29 Link-cut Tree - aresta
// Valores nas arestas
// rootify(v) torna v a raiz de sua arvore
// query(v, w) retorna a soma do caminho v--w
// update(v, w, x) soma x nas arestas do caminho v--w
//
// Todas as operacoes sao O(log(n)) amortizado
```

```
1ef namespace lct {
        struct node {
3c9
19f
            int p, ch[2];
810
            ll val, sub;
            bool rev:
aa6
04a
            int sz, ar;
4e4
            ll lazv:
f93
            node() {}
7a8
            node(int v, int ar_) :
            p(-1), val(v), sub(v), rev(0), sz(ar_{-}), ar(ar_{-}), lazy(0) {
546
b07
                ch[0] = ch[1] = -1;
53b
6e0
        }:
c53
        node t[2*MAX]; // MAXN + MAXQ
99e
        map<pair<int, int>, int> aresta;
e4d
        int sz;
95a
        void prop(int x) {
dc1
            if (t[x].lazy) {
25 e
                if (t[x].ar) t[x].val += t[x].lazy;
2ab
                t[x].sub += t[x].lazy*t[x].sz;
                if (t[x].ch[0]+1) t[t[x].ch[0]].lazy += t[x].lazy;
edc
942
                if (t[x].ch[1]+1) t[t[x].ch[1]].lazy += t[x].lazy;
1ba
            }
aa2
            if (t[x].rev) {
f95
                swap(t[x].ch[0], t[x].ch[1]);
379
                if (t[x].ch[0]+1) t[t[x].ch[0]].rev ^= 1;
c3d
                if (t[x].ch[1]+1) t[t[x].ch[1]].rev ^= 1;
50e
230
            t[x].lazy = 0, t[x].rev = 0;
f9d
       }
564
        void update(int x) {
            t[x].sz = t[x].ar, t[x].sub = t[x].val;
1a3
            for (int i = 0; i < 2; i++) if (t[x].ch[i]+1) {
8ca
621
                prop(t[x].ch[i]);
c4f
                t[x].sz += t[t[x].ch[i]].sz;
269
                t[x].sub += t[t[x].ch[i]].sub;
400
            }
       }
28b
971
        bool is root(int x) {
            return t[x].p == -1 or (t[t[x].p].ch[0] != x and
   t[t[x].p].ch[1] != x);
       }
cf1
ed6
        void rotate(int x) {
497
            int p = t[x].p, pp = t[p].p;
fc4
            if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
```

```
251
            bool d = t[p].ch[0] == x;
            t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
461
a76
            if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
            t[x].p = pp, t[p].p = x;
8fa
444
            update(p), update(x);
f31
238
        int splay(int x) {
18c
            while (!is_root(x)) {
497
                int p = t[x].p, pp = t[p].p;
77b
                if (!is_root(p)) prop(pp);
be5
                prop(p), prop(x);
                if (!is\_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch[0]
   == x) ? x : p):
64f
                rotate(x);
            }
72c
aab
            return prop(x), x;
08f
       }
       int access(int v) {
f16
0eb
            int last = -1;
            for (int w = v; w+1; update(last = w), splay(v), w =
d9f
   t[v].p)
024
                splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
3d3
            return last;
294
9f1
        void make_tree(int v, int w=0, int ar=0) { t[v] = node(w, ar);
  }
e89
        int find_root(int v) {
13f
            access(v), prop(v);
            while (t[v].ch[0]+1) v = t[v].ch[0], prop(v);
9f0
637
            return splay(v);
16a
82f
        bool conn(int v, int w) {
2cf
            access(v), access(w);
b9b
            return v == w ? true : t[v].p != -1;
ec0
277
        void rootify(int v) {
5e3
            access(v):
a02
            t[v].rev ^= 1;
a05
971
        11 query(int v, int w) {
b54
            rootify(w), access(v);
249
            return t[v].sub;
652
        void update(int v, int w, int x) {
3fa
b54
            rootify(w), access(v);
            t[v].lazy += x;
12c
74 f
        }
```

```
204
        void link_(int v, int w) {
821
            rootify(w);
389
            t[w].p = v;
523
        }
6b8
        void link(int v, int w, int x) { // v--w com peso x
379
            int id = MAX + sz++;
110
            aresta[make_pair(v, w)] = id;
a88
            make_tree(id, x, 1);
c88
            link_(v, id), link_(id, w);
58 c
e63
        void cut_(int v, int w) {
b54
            rootify(w), access(v);
264
            t[v].ch[0] = t[t[v].ch[0]].p = -1:
7cd
        void cut(int v, int w) {
031
b0f
            int id = aresta[make_pair(v, w)];
a4a
            cut_(v, id), cut_(id, w);
840
bbb
        int lca(int v, int w) {
5e3
            access(v):
a8b
            return access(w);
524
        }
9ce }
```

### 4.30 Link-cut Tree - vertice

```
// Valores nos vertices
// make_tree(v, w) cria uma nova arvore com um
// vertice soh com valor 'w'
// rootify(v) torna v a raiz de sua arvore
// query(v, w) retorna a soma do caminho v--w
// update(v, w, x) soma x nos vertices do caminho v--w
// Todas as operacoes sao O(\log(n)) amortizado
1ef namespace lct {
3c9
        struct node {
19f
            int p, ch[2];
810
            ll val, sub;
            bool rev;
aa6
e4d
            int sz:
4 e 4
            ll lazy;
f93
            node() {}
aa0
            node(int v) : p(-1), val(v), sub(v), rev(0), sz(1),
   lazy(0) {
b07
                ch[0] = ch[1] = -1;
```

```
c4e
2b7
        };
5f3
        node t[MAX];
        void prop(int x) {
95a
dc1
            if (t[x].lazy) {
9f7
                t[x].val += t[x].lazy, t[x].sub += t[x].lazy*t[x].sz;
edc
                if (t[x].ch[0]+1) t[t[x].ch[0]].lazy += t[x].lazy;
                if (t[x].ch[1]+1) t[t[x].ch[1]].lazy += t[x].lazy;
942
e26
            }
aa2
            if (t[x].rev) {
f95
                swap(t[x].ch[0], t[x].ch[1]);
                if (t[x].ch[0]+1) t[t[x].ch[0]].rev ^= 1;
379
                if (t[x].ch[1]+1) t[t[x].ch[1]].rev ^= 1;
c3d
50e
230
            t[x].lazy = 0, t[x].rev = 0;
        }
c62
564
        void update(int x) {
            t[x].sz = 1, t[x].sub = t[x].val:
ec2
            for (int i = 0; i < 2; i++) if (t[x].ch[i]+1) {
8ca
621
                prop(t[x].ch[i]);
                t[x].sz += t[t[x].ch[i]].sz;
c4f
269
                t[x].sub += t[t[x].ch[i]].sub;
400
            }
da7
        }
971
        bool is_root(int x) {
657
            return t[x].p == -1 or (t[t[x].p].ch[0] != x and
   t[t[x].p].ch[1] != x);
cf1
       }
        void rotate(int x) {
ed6
497
            int p = t[x].p, pp = t[p].p;
            if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
fc4
            bool d = t[p].ch[0] == x:
251
            t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
461
            if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
a76
8fa
            t[x].p = pp, t[p].p = x;
444
            update(p), update(x);
f31
238
        int splay(int x) {
            while (!is root(x)) {
18c
497
                int p = t[x].p, pp = t[p].p;
77b
                if (!is_root(p)) prop(pp);
be5
                prop(p), prop(x);
                if (!is_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch[0]
0 c 5
   == x) ? x : p);
                rotate(x):
64f
```

```
72c
aab
            return prop(x), x;
08f
        }
f16
        int access(int v) {
            int last = -1:
0eb
d9f
            for (int w = v; w+1; update(last = w), splay(v), w =
   t[v].p)
                splay(w), t[w].ch[1] = (last == -1 ? -1 : v);
024
3d3
            return last;
294
        }
f17
        void make_tree(int v, int w) { t[v] = node(w); }
e89
        int find root(int v) {
13f
            access(v). prop(v):
            while (t[v].ch[0]+1) v = t[v].ch[0], prop(v);
9f0
637
            return splay(v);
        }
16a
f94
        bool connected(int v, int w) {
            access(v), access(w);
2cf
b9b
            return v == w ? true : t[v].p != -1;
ec6
277
        void rootify(int v) {
5e3
            access(v):
            t[v].rev ^= 1;
a02
a05
        }
971
        11 query(int v, int w) {
b54
            rootifv(w). access(v):
            return t[v].sub;
249
652
        }
        void update(int v, int w, int x) {
3fa
b54
            rootify(w), access(v);
            t[v].lazy += x;
12c
74f
        }
142
        void link(int v, int w) {
821
            rootifv(w):
            t[w].p = v;
389
8a8
        }
031
        void cut(int v, int w) {
b54
            rootify(w), access(v);
264
            t[v].ch[0] = t[t[v].ch[0]].p = -1;
d9a
        }
        int lca(int v, int w) {
bbb
5e3
            access(v);
a8b
            return access(w);
524
        }
f9f }
```

### 4.31 Max flow com lower bound

```
// add(a, b, l, r):
// adiciona aresta de a pra b, onde precisa passar f de fluxo, l <= f
// add(a, b, c):
// adiciona aresta de a pra b com capacidade c
// Mesma complexidade do Dinitz
cd5 struct lb_max_flow : dinitz {
5ce
        vector < int > d:
d8c
        lb_max_flow(int n) : dinitz(n + 2), d(n, 0) {}
b12
        void add(int a, int b, int 1, int r) {
c97
            d[a] -= 1:
f1b
            d[b] += 1;
4c0
            dinitz::add(a, b, r - 1);
ed4
087
        void add(int a, int b, int c) {
            dinitz::add(a, b, c);
0f3
039
7a1
        bool has_circulation() {
50c
            int n = d.size();
854
            11 cost = 0;
603
            for (int i = 0; i < n; i++) {</pre>
c69
                if (d[i] > 0) {
                    cost += d[i]:
f56
57a
                    dinitz::add(n, i, d[i]);
c72
                } else if (d[i] < 0) {</pre>
b76
                    dinitz::add(i, n+1, -d[i]);
b42
                }
            }
676
067
            return (dinitz::max_flow(n, n+1) == cost);
110
7bd
        bool has_flow(int src, int snk) {
387
            dinitz::add(snk, src, INF);
e40
            return has_circulation();
cc1
        11 max_flow(int src, int snk) {
            if (!has_flow(src, snk)) return -1;
ee8
            dinitz::F = 0:
4ad
fe5
            return dinitz::max_flow(src, snk);
619
        }
e8f };
```

## 4.32 MinCostMaxFlow

```
// min_cost_flow(s, t, f) computa o par (fluxo, custo)
// com max(fluxo) <= f que tenha min(custo)</pre>
// min_cost_flow(s, t) -> Fluxo maximo de custo minimo de s pra t
// Se for um dag, da pra substituir o SPFA por uma DP pra nao
// pagar O(nm) no comeco
// Se nao tiver aresta com custo negativo, nao precisa do SPFA
// O(nm + f * m log n)
123 template < typename T > struct mcmf {
670
        struct edge {
b75
            int to, rev, flow, cap; // para, id da reversa, fluxo,
    capacidade
7f9
            bool res; // se eh reversa
            T cost; // custo da unidade de fluxo
635
            edge(): to(0), rev(0), flow(0), cap(0), cost(0),
   res(false) {}
1d7
            edge(int to_, int rev_, int flow_, int cap_, T cost_, bool
   res )
                : to(to_), rev(rev_), flow(flow_), cap(cap_),
   res(res_), cost(cost_) {}
723
        };
002
        vector<vector<edge>> g;
168
        vector<int> par_idx, par;
f1e
        T inf;
a03
        vector <T> dist;
        mcmf(int n) : g(n), par_idx(n), par(n),
    inf(numeric_limits <T>::max()/3) {}
91c
        void add(int u, int v, int w, T cost) { // de u pra v com cap
   w e custo cost
            edge a = edge(v, g[v].size(), 0, w, cost, false);
2fc
            edge b = edge(u, g[u].size(), 0, 0, -cost, true);
234
b24
            g[u].push_back(a);
c12
            g[v].push_back(b);
0ed
8bc
        vector<T> spfa(int s) { // nao precisa se nao tiver custo
   negativo
871
            deque < int > q;
3d1
            vector < bool > is_inside(g.size(), 0);
577
            dist = vector <T>(g.size(), inf);
```

```
a93
            dist[s] = 0;
a30
            q.push_back(s);
            is_inside[s] = true;
ecb
            while (!q.empty()) {
14d
b1e
                 int v = q.front();
                 q.pop_front();
ced
                is_inside[v] = false;
48d
76e
                for (int i = 0; i < g[v].size(); i++) {</pre>
9d4
                     auto [to, rev, flow, cap, res, cost] = g[v][i];
e61
                     if (flow < cap and dist[v] + cost < dist[to]) {</pre>
                         dist[to] = dist[v] + cost;
943
                         if (is_inside[to]) continue;
ed6
020
                         if (!q.empty() and dist[to] > dist[q.front()])
   q.push_back(to);
                         else q.push_front(to);
b33
                         is inside[to] = true:
b52
2d1
                     }
8cd
f2c
            }
8d7
            return dist;
96c
2a2
        bool dijkstra(int s, int t, vector <T>& pot) {
489
            priority_queue <pair <T, int>, vector <pair <T, int>>,
   greater<>> q;
            dist = vector <T>(g.size(), inf);
577
a93
            dist[s] = 0;
115
            q.emplace(0, s);
402
            while (q.size()) {
91b
                 auto [d, v] = q.top();
833
                 q.pop();
                if (dist[v] < d) continue;</pre>
68b
76e
                for (int i = 0; i < g[v].size(); i++) {</pre>
                     auto [to, rev, flow, cap, res, cost] = g[v][i];
9d4
                     cost += pot[v] - pot[to];
e8c
                     if (flow < cap and dist[v] + cost < dist[to]) {</pre>
e61
943
                         dist[to] = dist[v] + cost;
                         q.emplace(dist[to], to);
441
88ъ
                         par_idx[to] = i, par[to] = v;
873
                     }
                }
de3
9d4
1d4
            return dist[t] < inf;</pre>
        }
c68
```

```
3d2
        pair < int , T> min_cost_flow(int s, int t, int flow = INF) {
3dd
            vector <T> pot(g.size(), 0);
            pot = spfa(s); // mudar algoritmo de caminho minimo aqui
9e4
d22
            int f = 0:
ce8
            T ret = 0:
4a0
            while (f < flow and dijkstra(s, t, pot)) {</pre>
bda
                 for (int i = 0; i < g.size(); i++)</pre>
                     if (dist[i] < inf) pot[i] += dist[i];</pre>
d2a
71b
                 int mn flow = flow - f. u = t:
045
                 while (u != s){
90f
                     mn_flow = min(mn_flow,
07d
                         g[par[u]][par_idx[u]].cap -
   g[par[u]][par_idx[u]].flow);
3d1
                     u = par[u];
                }
935
1f2
                ret += pot[t] * mn_flow;
476
                 u = t:
                 while (u != s) {
045
e09
                     g[par[u]][par_idx[u]].flow += mn_flow;
498
                     g[u][g[par[u]][par_idx[u]].rev].flow -= mn_flow;
3d1
                     u = par[u];
bcc
                }
04d
                 f += mn_flow;
36d
            }
15b
            return make_pair(f, ret);
        }
cc3
        // Opcional: retorna as arestas originais por onde passa flow
        vector<pair<int,int>> recover() {
182
24a
            vector < pair < int , int >> used;
2a4
            for (int i = 0; i < g.size(); i++) for (edge e : g[i])</pre>
                 if(e.flow == e.cap && !e.res) used.push_back({i,
587
   e.to}):
f6b
            return used;
390
697 };
```

### 4.33 Prufer code

```
// Traduz de lista de arestas para prufer code
// e vice-versa
// Os vertices tem label de O a n-1
// Todo arrav com n-2 posicoes e valores de
// O a n-1 sao prufer codes validos
//
// O(n)
47d vector < int > to_prufer (vector < pair < int , int >> tree) {
        int n = tree.size()+1;
1fa
2cf
        vector < int > d(n, 0);
4aa
        vector < vector < int >> g(n);
        for (auto [a, b] : tree) d[a]++, d[b]++,
f87
f60
            g[a].push_back(b), g[b].push_back(a);
        vector < int > pai(n, -1);
c5a
        queue < int > q; q.push(n-1);
260
402
        while (q.size()) {
be1
            int u = q.front(); q.pop();
34c
            for (int v : g[u]) if (v != pai[u])
9c9
                pai[v] = u, q.push(v);
70d
        }
399
        int idx, x;
897
        idx = x = find(d.begin(), d.end(), 1) - d.begin();
4b8
        vector<int> ret;
b28
       for (int i = 0: i < n-2: i++) {
            int y = pai[x];
d4b
e81
            ret.push_back(y);
            if (--d[y] == 1 \text{ and } y < idx) x = y;
666
367
            else idx = x = find(d.begin()+idx+1, d.end(), 1) -
   d.begin();
5f9
edf
        return ret;
d3b }
4d8 vector < pair < int , int >> from prufer (vector < int > p) {
455
        int n = p.size()+2;
        vector < int > d(n, 1);
126
        for (int i : p) d[i]++;
650
85b
        p.push_back(n-1);
399
        int idx, x:
897
        idx = x = find(d.begin(), d.end(), 1) - d.begin();
1df
        vector < pair < int , int >> ret;
b06
        for (int y : p) {
dab
            ret.push_back({x, y});
666
            if (--d[y] == 1 \text{ and } y < idx) x = y;
```

```
367
            else idx = x = find(d.begin()+idx+1, d.end(), 1) -
   d.begin();
c3b
        }
edf
        return ret;
765 }
4.34 Sack (DSU em arvores)
// Responde queries de todas as sub-arvores
// offline
//
// O(n log(n))
6bf int sz[MAX], cor[MAX], cnt[MAX];
042 vector <int> g[MAX];
6df void build(int k, int d=0) {
e8f
        sz[k] = 1;
        for (auto& i : g[k]) {
01a
30f
            build(i, d+1): sz[k] += sz[i]:
925
            if (sz[i] > sz[g[k][0]]) swap(i, g[k][0]);
011
        }
189 }
74f void compute(int k, int x, bool dont=1) {
de9
        cnt[cor[k]] += x;
828
        for (int i = dont; i < g[k].size(); i++)</pre>
b5c
            compute(g[k][i], x, 0);
896 }
dc4 void solve(int k, bool keep=0) {
32a
        for (int i = int(g[k].size())-1; i >= 0; i--)
b4c
            solve(g[k][i], !i);
4a0
        compute(k, 1);
        // agora cnt[i] tem quantas vezes a cor
        // i aparece na sub-arvore do k
830
        if (!keep) compute(k, -1, 0);
8bc }
```

# 4.35 Stable Marriage

```
// Emparelha todos os elementos de A com elementos de B
// de forma que nao exista um par x \in A, y \in B
// e x nao pareado com y tal que x prefira parear com y
// e y prefira parear com x.
// a[i] contem os elementos de B ordenados por preferencia de i
// b[j] contem os elementos de A ordenados por preferencia de j
// |A| <= |B|
                                                                            cf0
// Retorna um vetor v de tamanho |A| onde v[i] guarda o match de i.
                                                                            18e
//
                                                                            0c2
// O(|A| * |B|)
                                                                            48e
380 vector<int> stable_marriage(vector<vector<int>> &a,
                                                                            740
   vector < vector < int >> &b) {
                                                                            994
        int n = a.size(), m = b.size();
652
                                                                            d64
        assert(a[0].size() == m and b[0].size() == n and n <= m);
83e
        vector < int > match(m, -1), it(n, 0);
017
        vector inv_b(m, vector<int>(n));
e6f
                                                                            3de
        for (int i = 0; i < m; i++) for (int j = 0; j < n; j++)
a34
                                                                            3c3
9f2
            inv_b[i][b[i][j]] = j;
                                                                            9c5
                                                                            2ef
26a
        queue < int > q;
                                                                            266
5af
        for (int i = 0; i < n; i++) q.push(i);</pre>
402
        while (q.size()) {
                                                                            253
379
            int i = q.front(); q.pop();
                                                                            38a }
            int j = a[i][it[i]];
4b8
57c
            if (match[j] == -1) match[j] = i;
02d
            else if (inv_b[j][i] < inv_b[j][match[j]]) {</pre>
                                                                            991
5d1
                q.emplace(match[j]);
e7d
                it[match[j]]++;
                                                                            3be
f1d
                match[j] = i;
                                                                            ea1 }
bc4
            } else q.emplace(i), it[i]++;
258
        vector < int > ret(n):
825
d72
        for (int i = 0; i < m; i++) if (match[i] != -1) ret[match[i]]</pre>
   = i:
edf
        return ret;
Off }
4.36 Tarjan para SCC
```

```
// O(n + m)
```

```
042 vector < int > g[MAX];
4ce stack<int> s;
a42 int vis[MAX], comp[MAX];
3fd int id[MAX];
// se quiser comprimir ciclo ou achar ponte em grafo nao direcionado,
// colocar um if na dfs para nao voltar pro pai da DFS tree
f32 int dfs(int i, int& t) {
        int lo = id[i] = t++;
        s.push(i);
        vis[i] = 2;
        for (int j : g[i]) {
            if (!vis[j]) lo = min(lo, dfs(j, t));
            else if (vis[j] == 2) lo = min(lo, id[j]);
        }
        // aresta de i pro pai eh uma ponte (no caso nao direcionado)
        if (lo == id[i]) while (1) {
            int u = s.top(); s.pop();
            vis[u] = 1, comp[u] = i;
            if (u == i) break;
        }
        return lo;
f93 void tarjan(int n) {
        int t = 0:
        for (int i = 0; i < n; i++) vis[i] = 0;</pre>
        for (int i = 0; i < n; i++) if (!vis[i]) dfs(i, t);</pre>
4.37 Topological Sort
// Retorna uma ordenacaoo topologica de g
// Se g nao for DAG retorna um vetor vazio
// O(n + m)
042 vector <int> g[MAX];
b6a vector<int> topo_sort(int n) {
46e
        vector < int > ret(n,-1), vis(n,0);
```

```
f51
        int pos = n-1, dag = 1;
        function < void(int) > dfs = [&](int v) {
36d
            vis[v] = 1;
cca
440
            for (auto u : g[v]) {
                if (vis[u] == 1) dag = 0;
152
532
                else if (!vis[u]) dfs(u);
e37
            }
d44
            ret[pos--] = v, vis[v] = 2;
57e
        };
158
        for (int i = 0; i < n; i++) if (!vis[i]) dfs(i);</pre>
d8f
        if (!dag) ret.clear();
edf
        return ret;
d6b }
4.38 Vertex cover
// Encontra o tamanho do vertex cover minimo
// Da pra alterar facil pra achar os vertices
// Parece rodar com < 2 s pra N = 90
//
// \Omega(n * 1.38^n)
76a namespace cover {
5a4
        const int MAX = 96:
        vector < int > g[MAX];
042
823
        bitset < MAX > bs[MAX];
1a8
       int n:
697
        void add(int i, int j) {
bd0
            if (i == j) return;
78c
            n = max({n, i+1, j+1});
200
            bs[i][j] = bs[j][i] = 1;
203
        }
        int rec(bitset < MAX > m) {
6c0
1a4
            int ans = 0:
            for (int x = 0; x < n; x++) if (m[x]) {
25b
002
                bitset < MAX > comp;
4bf
                function < void(int) > dfs = [&](int i) {
                     comp[i] = 1, m[i] = 0;
b96
0c3
                    for (int j : g[i]) if (m[j]) dfs(j);
815
                };
963
                dfs(x);
```

```
d34
                int ma, deg = -1, cvc = 1;
417
                for (int i = 0; i < n; i++) if (comp[i]) {</pre>
d0b
                     int d = (bs[i]&comp).count();
18a
                     if (d \le 1) cyc = 0;
c1f
                     if (d > deg) deg = d, ma = i;
d8e
269
                if (deg <= 2) { // caminho ou ciclo</pre>
340
                     ans += (comp.count() + cyc) / 2;
5e2
                     continue;
702
                }
3f9
                comp[ma] = 0;
                // ou ta no cover. ou nao ta no cover
1dd
                ans += min(1 + rec(comp), deg + rec(comp & <math>\sim bs[ma]));
            }
6e6
ba7
            return ans;
2ec
        }
        int solve() {
f5c
3c5
            bitset < MAX > m;
603
            for (int i = 0; i < n; i++) {</pre>
939
                m[i] = 1;
f90
                for (int j = 0; j < n; j++)
741
                     if (bs[i][j]) g[i].push_back(j);
13e
            }
4f9
            return rec(m);
708
        }
9c5 }
4.39 Virtual Tree
// Comprime uma arvore dado um conjunto S de vertices, de forma que
// o conjunto de vertices da arvore comprimida contenha S e seja
// minimal e fechado sobre a operação de LCA
// Se |S| = k, a arvore comprimida tem menos que 2k vertices
// As arestas de virt possuem a distancia do vertice ate o vizinho
// Retorna a raiz da virtual tree
// lca::pos deve ser a ordem de visitacao no dfs
// voce pode usar o LCAcomHLD, por exemplo
// O(k log(k))
b36 vector<pair<int, int>> virt[MAX];
d41 #warning lembrar de buildar o LCA antes
c14 int build_virt(vector<int> v) {
```

```
b46
        auto cmp = [&](int i, int j) { return lca::pos[i] <</pre>
   lca::pos[i]; };
        sort(v.begin(), v.end(), cmp);
074
        for (int i = v.size()-1; i; i--) v.push_back(lca::lca(v[i],
   v[i-1])):
        sort(v.begin(), v.end(), cmp);
074
d76
        v.erase(unique(v.begin(), v.end()), v.end());
        for (int i = 0; i < v.size(); i++) virt[v[i]].clear();</pre>
37 c
197
        for (int i = 1; i < v.size(); i++) virt[lca::lca(v[i-1],</pre>
   v[i])].clear():
        for (int i = 1; i < v.size(); i++) {</pre>
ad7
51b
            int parent = lca::lca(v[i-1], v[i]);
290
            int d = lca::dist(parent. v[i]):
d41 #warning soh to colocando aresta descendo
4d0
            virt[parent].emplace_back(v[i], d);
fe5
        }
832
        return v[0];
142 }
```

# 5 Problemas

## 5.1 Angle Range Intersection

```
// Computa intersecao de angulos
// Os angulos (arcos) precisam ter comprimeiro < pi
// (caso contrario a intersecao eh estranha)
//
// Tudo 0(1)
32a struct angle_range {
75e
        static constexpr ld ALL = 1e9, NIL = -1e9;
395
c77
        angle_range() : 1(ALL), r(ALL) {}
        angle_range(ld l_, ld r_) : l(l_), r(r_) { fix(l), fix(r); }
894
        void fix(ld& theta) {
4ee
            if (theta == ALL or theta == NIL) return;
da7
323
            if (theta > 2*pi) theta -= 2*pi;
868
            if (theta < 0) theta += 2*pi;</pre>
625
2ee
        bool empty() { return 1 == NIL; }
931
        bool contains(ld q) {
40f
            fix(q);
4d7
            if (1 == ALL) return true;
```

```
fec
            if (1 == NIL) return false;
6a6
            if (1 < r) return 1 < q and q < r;
075
            return q > 1 or q < r;</pre>
800
        }
        friend angle_range operator &(angle_range p, angle_range q) {
9c7
743
            if (p.1 == ALL or q.1 == NIL) return q;
20f
            if (q.l == ALL or p.l == NIL) return p;
7d5
            if (p.1 > p.r \text{ and } q.1 > q.r) \text{ return } \{\max(p.1, q.1),
   min(p.r, q.r)};
            if (q.1 > q.r) swap(p.1, q.1), swap(p.r, q.r);
aa6
8b8
            if (p.1 > p.r) {
249
                 if (q.r > p.l) return {max(q.l, p.l) , q.r};
6f7
                 else if (q.l < p.r) return \{q.l, min(q.r, p.r)\};
270
                 return {NIL, NIL};
            }
337
5a8
            if (max(p.1, q.1) > min(p.r, q.r)) return {NIL, NIL};
bcb
            return {max(p.1, q.1), min(p.r, q.r)};
142
5e1 };
```

## 5.2 Area da Uniao de Retangulos

```
// O(n log(n))
// 5d8d2f
aa4 namespace seg {
        pair < int , 11 > seg[4*MAX];
b1b
        ll lazy[4*MAX], *v;
1a8
        int n;
e01
        pair<int, ll> merge(pair<int, ll> l, pair<int, ll> r){
719
            if (1.second == r.second) return {1.first+r.first,
   1.second}:
53b
            else if (l.second < r.second) return l;</pre>
            else return r;
aa0
        }
d82
        pair<int, ll> build(int p=1, int l=0, int r=n-1) {
6fc
3c7
            lazv[p] = 0:
            if (1 == r) return seg[p] = {1, v[1]};
bf8
ee4
            int m = (1+r)/2:
            return seg[p] = merge(build(2*p, 1, m), build(2*p+1, m+1,
432
   r));
f94
d9e
        void build(int n2, 11* v2) {
680
            n = n2, v = v2;
```

```
6f2
            build();
f8a
ceb
        void prop(int p, int l, int r) {
208
            seg[p].second += lazy[p];
2c9
            if (1 != r) lazy[2*p] += lazy[p], lazy[2*p+1] += lazy[p];
3c7
            lazv[p] = 0;
bf2
        }
693
        pair < int, 11 > query (int a, int b, int p=1, int 1=0, int r=n-1)
  {
            prop(p, 1, r);
6b9
527
            if (a <= 1 and r <= b) return seg[p];</pre>
9b7
            if (b < 1 or r < a) return {0, LINF};</pre>
ee4
            int m = (1+r)/2:
            return merge(query(a, b, 2*p, 1, m), query(a, b, 2*p+1,
   m+1, r));
786
       }
07 c
        pair < int, 11 > update(int a, int b, int x, int p=1, int 1=0,
   int r=n-1) {
            prop(p, 1, r);
6b9
            if (a \le 1 \text{ and } r \le b) {
9a3
                lazy[p] += x;
b94
6b9
                prop(p, 1, r);
                return seg[p];
534
821
            }
e9f
            if (b < l or r < a) return seg[p];</pre>
ee4
            int m = (1+r)/2:
086
            return seg[p] = merge(update(a, b, x, 2*p, 1, m),
579
                    update(a, b, x, 2*p+1, m+1, r));
c65
        }
043 };
eb5 ll seg_vec[MAX];
8be 11 area_sq(vector<pair<int, int>, pair<int, int>>> &sq){
        vector<pair<int, int>, pair<int, int>>> up;
28 c
60a
        for (auto it : sq){
619
            int x1, y1, x2, y2;
ae0
            tie(x1, v1) = it.first;
68e
            tie(x2, y2) = it.second;
80f
            up.push_back(\{\{x1+1, 1\}, \{y1, y2\}\}\});
            up.push_back({{x2+1, -1}, {y1, y2}});
aee
6c3
092
        sort(up.begin(), up.end());
        memset(seg_vec, 0, sizeof seg_vec);
049
6fe
        11 H_MAX = MAX;
156
        seg::build(H_MAX-1, seg_vec);
7ba
        auto it = up.begin();
```

```
04b
        11 \text{ ans} = 0;
f14
        while (it != up.end()){
07f
            11 L = (*it).first.first;
718
            while (it != up.end() && (*it).first.first == L){
127
                int x, inc, y1, y2;
                tie(x, inc) = it->first;
d35
d3d
                tie(y1, y2) = it -> second;
5d1
                seg::update(y1+1, y2, inc);
40d
                it++;
9b1
            }
852
            if (it == up.end()) break;
d8a
            11 R = (*it).first.first;
f59
            11 W = R-L:
efd
            auto jt = seg::query(0, H_MAX-1);
91a
            11 H = H_MAX - 1;
e8a
            if (jt.second == 0) H -= jt.first;
8df
            ans += W*H:
5c8
        }
ba7
        return ans;
385 }
```

# 5.3 Area Maxima de Histograma

```
// Assume que todas as barras tem largura 1,
// e altura dada no vetor v
//
// O(n)
15e ll area(vector<int> v) {
b73
        ll ret = 0:
4ce
        stack<int> s;
        // valores iniciais pra dar tudo certo
447
        v.insert(v.begin(), -1);
d56
        v.insert(v.end(), -1);
1f8
        s.push(0);
        for(int i = 0; i < (int) v.size(); i++) {</pre>
0be
78e
            while (v[s.top()] > v[i]) {
265
                11 h = v[s.top()]; s.pop();
de1
                ret = max(ret, h * (i - s.top() - 1));
            }
40a
18e
            s.push(i);
020
        }
edf
        return ret;
```

e43 }

## 5.4 Binomial modular

```
// Computa C(n, k) mod m em O(m + log(m) log(n))
// = O(rapido)
97c ll divi[MAX];
398 ll expo(ll a, ll b, ll m) {
        if (!b) return 1;
1 c 1
        ll ans = expo(a*a\%m, b/2, m);
399
751
        if (b\%2) ans *= a;
2e9
        return ans%m;
754 }
f0a ll inv(ll a, ll b){
bca
        return 1 < a ? b - inv(b\%a,a)*b/a : 1:
041 }
153 template < typename T> tuple < T, T, T> ext_gcd(T a, T b) {
        if (!a) return {b, 0, 1};
3bd
550
        auto [g, x, y] = ext_gcd(b\%a, a);
        return \{g, y - b/a*x, x\};
c59
537 }
bfe template < typename T = 11 > struct crt {
627
        Ta, m;
        crt(): a(0), m(1) {}
5f3
7eb
        crt(T a_, T m_) : a(a_), m(m_) {}
911
        crt operator * (crt C) {
238
            auto [g, x, y] = ext_gcd(m, C.m);
dc0
            if ((a - C.a) \% g) a = -1;
            if (a == -1 or C.a == -1) return crt(-1, 0);
4f9
d09
            T lcm = m/g*C.m;
            T ans = a + (x*(C.a-a)/g \% (C.m/g))*m;
eb2
d8d
            return crt((ans % lcm + lcm) % lcm, lcm);
1f2
        }
0d9 };
6f2 pair<11, 11> divide_show(11 n, int p, int k, int pak) {
        if (n == 0) return {0, 1};
4f7
d02
        11 blocos = n/pak, falta = n%pak;
2ce
        ll periodo = divi[pak], resto = divi[falta];
616
        11 r = expo(periodo, blocos, pak)*resto%pak;
```

```
auto rec = divide_show(n/p, p, k, pak);
445
        ll v = n/p + rec.first;
a51
        r = r*rec.second % pak;
bb9
90f
        return {y, r};
533 }
6e6 ll solve_pak(ll n, ll x, int p, int k, int pak) {
        divi[0] = 1:
f2b
        for (int i = 1; i <= pak; i++) {</pre>
901
            divi[i] = divi[i-1];
840
            if (i%p) divi[i] = divi[i] * i % pak;
51a
        }
4ac
        auto dn = divide_show(n, p, k, pak), dx = divide_show(x, p, k,
   pak),
162
             dnx = divide_show(n-x, p, k, pak);
768
        11 y = dn.first-dx.first-dnx.first, r =
             (dn.second*inv(dx.second, pak)%pak)*inv(dnx.second,
b64
   pak)%pak;
035
        return expo(p, y, pak) * r % pak;
d78 }
9dd ll solve(ll n, ll x, int mod) {
490
        vector < pair < int , int >> f;
c3b
        int mod2 = mod;
7b4
        for (int i = 2: i*i \le mod 2: i++) if (mod 2\%i = = 0) {
aff
75b
            while (mod2\%i==0) mod2 /= i, c++;
2a1
            f.push_back({i, c});
fe7
Off
        if (mod2 > 1) f.push_back({mod2, 1});
e96
        crt ans(0, 1):
a13
        for (int i = 0; i < f.size(); i++) {</pre>
702
            int pak = 1;
7 e 4
            for (int j = 0; j < f[i].second; j++) pak *= f[i].first;</pre>
            ans = ans * crt(solve_pak(n, x, f[i].first, f[i].second,
304
   pak), pak);
7fd
        }
5fb
        return ans.a;
689 }
5.5 Closest pair of points
```

// O(nlogn)

```
915 pair <pt, pt > closest_pair_of_points(vector <pt > v) {
3d2
        int n = v.size();
fca
        sort(v.begin(), v.end());
        for (int i = 1; i < n; i++) if (v[i] == v[i-1]) return</pre>
31c
   {v[i-1], v[i]};
c20
        auto cmp_y = [&](const pt &1, const pt &r) {
            if (1.v != r.v) return 1.v < r.v;</pre>
b53
920
            return 1.x < r.x;</pre>
55a
        }:
62e
        set < pt, decltype(cmp_v) > s(cmp_v);
3d9
        int 1 = 0, r = -1:
6a2
        11 d2 min = numeric limits < ll >:: max();
4d5
        pt pl, pr;
bd1
        const int magic = 5;
a55
        while (r+1 < n) {
7f1
            auto it = s.insert(v[++r]).first;
c92
            int cnt = magic/2;
            while (cnt-- and it != s.begin()) it--;
773
a01
            cnt = 0:
d68
            while (cnt++ < magic and it != s.end()) {</pre>
f19
                 if (!((*it) == v[r])) {
                     11 d2 = dist2(*it, v[r]);
67e
74e
                     if (d2_min > d2) {
229
                         d2 \min = d2:
841
                         pl = *it:
4f2
                         pr = v[r];
                     }
7d9
                 }
10a
40d
                 it++;
801
            while (1 < r \text{ and } sq(v[1].x-v[r].x) > d2_min)
eb0
   s.erase(v[1++]):
de6
        return {pl, pr};
c74
f90 }
```

## 5.6 Coloracao de Grafo de Intervalo

```
// Colore os intervalos com o numero minimo
// de cores de tal forma que dois intervalos
// que se interceptam tem cores diferentes
// As cores vao de 1 ate n
//
// O(n log(n))
```

```
615 vector<int> coloring(vector<pair<int, int>>& v) {
3d2
        int n = v.size();
c08
        vector<pair<int, pair<int, int>>> ev;
        for (int i = 0; i < n; i++) {</pre>
603
            ev.push_back({v[i].first, {1, i}});
150
            ev.push_back({v[i].second, {0, i}});
cda
6a4
        }
        sort(ev.begin(), ev.end());
49e
360
        vector < int > ans(n), avl(n);
        for (int i = 0; i < n; i++) avl.push_back(n-i);</pre>
265
4bf
        for (auto i : ev) {
cbe
            if (i.second.first == 1) {
021
                ans[i.second.second] = avl.back():
                avl.pop_back();
a00
e98
            } else avl.push_back(ans[i.second.second]);
3a6
ba7
        return ans;
83a }
```

## 5.7 Conectividade Dinamica DC

```
// Offline com Divide and Conquer e
// DSU com rollback
// O(n log^2(n))
8f2 typedef pair <int, int> T;
1cd namespace data {
553
        int n, ans;
573
        int p[MAX], sz[MAX];
        stack<int> S;
ee6
e5b
        void build(int n2) {
1e3
            n = n2:
8a6
            for (int i = 0; i < n; i++) p[i] = i, sz[i] = 1;</pre>
0b2
             ans = n:
        }
cba
1 b 1
        int find(int k) {
006
             while (p[k] != k) k = p[k];
839
            return k;
c1e
        }
072
        void add(T x) {
700
            int a = x.first, b = x.second;
605
            a = find(a), b = find(b);
843
            if (a == b) return S.push(-1);
e7d
            ans --;
```

```
3c6
            if (sz[a] > sz[b]) swap(a, b);
4c2
            S.push(a);
582
            sz[b] += sz[a];
84b
            p[a] = b;
        }
e1a
5eb
        int query() {
ba7
            return ans:
35c
        }
        void rollback() {
5cf
465
            int u = S.top(); S.pop();
61c
            if (u == -1) return;
270
            sz[p[u]] -= sz[u];
546
            p[u] = u:
0df
            ans++;
       }
456
568 };
357 int ponta[MAX]; // outra ponta do intervalo ou -1 se for query
4f0 int ans[MAX], n, q;
487 T qu[MAX];
47b void solve(int l = 0, int r = q-1) {
        if (1 >= r) {
0b1
8c0
            ans[1] = data::query(); // agora a estrutura ta certa
505
            return:
f77
962
        int m = (1+r)/2, qnt = 1;
fc7
        for (int i = m+1; i <= r; i++) if (ponta[i]+1 and ponta[i] < 1)</pre>
37d
            data::add(qu[i]), qnt++;
221
        solve(1, m);
593
        while (--qnt) data::rollback();
a2c
        for (int i = 1; i <= m; i++) if (ponta[i]+1 and ponta[i] > r)
37d
            data::add(qu[i]), qnt++;
37b
        solve(m+1, r):
        while (qnt--) data::rollback();
281
0d4 }
    Conectividade Dinamica LCT
// Offline com link-cut trees
// O(n log(n))
1ef namespace lct {
3 c 9
        struct node {
19f
            int p, ch[2];
a2a
            int val, sub;
```

```
aa6
            bool rev;
f93
            node() {}
54e
            node(int v) : p(-1), val(v), sub(v), rev(0) { ch[0] =}
   ch[1] = -1; }
       };
cac
c53
        node t[2*MAX]: // MAXN + MAXQ
99e
        map<pair<int, int>, int> aresta;
e4d
        int sz;
95a
        void prop(int x) {
aa2
            if (t[x].rev) {
f95
                swap(t[x].ch[0], t[x].ch[1]);
379
                if (t[x].ch[0]+1) t[t[x].ch[0]].rev ^= 1;
c3d
                if (t[x].ch[1]+1) t[t[x].ch[1]].rev ^= 1;
50e
693
            t[x].rev = 0;
750
        }
        void update(int x) {
564
            t[x].sub = t[x].val:
e8d
            for (int i = 0; i < 2; i++) if (t[x].ch[i]+1) {
8ca
621
                prop(t[x].ch[i]);
78d
                t[x].sub = min(t[x].sub, t[t[x].ch[i]].sub);
3e4
            }
9bf
        }
971
        bool is root(int x) {
657
            return t[x].p == -1 or (t[t[x].p].ch[0] != x and
   t[t[x].p].ch[1] != x);
       }
cf1
ed6
        void rotate(int x) {
497
            int p = t[x].p, pp = t[p].p;
fc4
            if (!is_root(p)) t[pp].ch[t[pp].ch[1] == p] = x;
251
            bool d = t[p].ch[0] == x;
            t[p].ch[!d] = t[x].ch[d], t[x].ch[d] = p;
461
a76
            if (t[p].ch[!d]+1) t[t[p].ch[!d]].p = p;
8fa
            t[x].p = pp, t[p].p = x;
            update(p), update(x);
444
f31
       }
238
        int splay(int x) {
18c
            while (!is_root(x)) {
497
                int p = t[x].p, pp = t[p].p;
77b
                if (!is_root(p)) prop(pp);
be5
                prop(p), prop(x);
                if (!is\_root(p)) rotate((t[pp].ch[0] == p)^(t[p].ch[0]
0c5
   == x) ? x : p);
64f
                rotate(x):
72c
            }
```

```
aab
            return prop(x), x;
08f
f16
       int access(int v) {
0eb
            int last = -1:
            for (int w = v; w+1; update(last = w), splay(v), w =
d9f
   t[v].p)
024
                splav(w), t[w], ch[1] = (last == -1 ? -1 : v):
3d3
            return last;
294
        void make_tree(int v, int w=INF) { t[v] = node(w); }
952
82f
        bool conn(int v, int w) {
2cf
            access(v), access(w);
b9b
            return v == w ? true : t[v].p != -1;
ec0
277
        void rootify(int v) {
5e3
            access(v):
a02
            t[v].rev ^= 1;
a05
        int query(int v, int w) {
a1d
b54
            rootify(w), access(v);
            return t[v].sub;
249
c28
        void link_(int v, int w) {
204
821
            rootify(w);
389
            t[w].p = v;
523
        void link(int v, int w, int x) { // v--w com peso x
6b8
379
            int id = MAX + sz++;
110
            aresta[make_pair(v, w)] = id;
ab6
            make_tree(id, x);
            link_(v, id), link_(id, w);
c88
984
        void cut_(int v, int w) {
e63
            rootifv(w). access(v):
b54
            t[v].ch[0] = t[t[v].ch[0]].p = -1;
264
7cd
031
        void cut(int v, int w) {
b0f
            int id = aresta[make_pair(v, w)];
a4a
            cut_(v, id), cut_(id, w);
840
        }
0d7 }
893 void dyn_conn() {
c5f
        int n, q; cin >> n >> q;
d6e
        vector < int > p(2*q, -1); // outra ponta do intervalo
       for (int i = 0; i < n; i++) lct::make_tree(i);</pre>
b4f
fbf
        vector < pair < int , int >> qu(q);
```

```
139
        map<pair<int, int>, int> m;
abf
        for (int i = 0; i < q; i++) {</pre>
3c2
            char c; cin >> c;
ef6
            if (c == '?') continue;
602
            int a. b: cin >> a >> b: a--. b--:
d11
            if (a > b) swap(a, b);
8a1
            qu[i] = \{a, b\}:
            if (c == '+') {
8d7
94b
                p[i] = i+q, p[i+q] = i;
                m[make_pair(a, b)] = i;
906
8a0
            } else {
412
                int j = m[make_pair(a, b)];
ac2
                p[i] = j, p[j] = i;
0da
        }
9e5
447
        int ans = n;
        for (int i = 0; i < q; i++) {</pre>
abf
87d
            if (p[i] == -1) {
                cout << ans << endl; // numero de comp conexos</pre>
886
5e2
                continue:
            }
b35
69d
            int a = qu[i].first, b = qu[i].second;
            if (p[i] > i) { // +
c4d
ac5
                if (lct::conn(a, b)) {
18f
                    int mi = lct::query(a, b);
993
                    if (p[i] < mi) {</pre>
dd3
                         p[p[i]] = p[i];
5e2
                         continue;
474
                    }
6f7
                    lct::cut(qu[p[mi]].first, qu[p[mi]].second), ans++;
6ea
                    p[mi] = mi;
9a9
d1d
                lct::link(a, b, p[i]), ans--;
            } else if (p[i] != i) lct::cut(a, b), ans++: // -
9d0
c03
56a }
     Conj. Indep. Maximo com Peso em Grafo de Intervalo
```

```
// Retorna os indices ordenados dos intervalos selecionados
// Se tiver empate, retorna o que minimiza o comprimento total
//
// O(n log(n))
31e vector<int> ind_set(vector<tuple<int, int, int>>& v) {
b27
        vector<tuple<int, int, int>> w;
```

```
f 14
                       for (int i = 0; i < v.size(); i++) {</pre>
e85
                                   w.push_back(tuple(get<0>(v[i]), 0, i));
6f0
                                   w.push_back(tuple(get<1>(v[i]), 1, i));
17f
                       }
                       sort(w.begin(), w.end());
d1d
844
                       vector < int > nxt(v.size());
c22
                       vector < pair < 11, int >> dp(v.size());
0eb
                       int last = -1;
723
                       for (auto [fim, t, i] : w) {
25a
                                   if (t == 0) {
4ca
                                               nxt[i] = last;
5e2
                                               continue:
5fd
                                   }
78b
                                   dp[i] = \{0, 0\};
                                   if (last != -1) dp[i] = max(dp[i], dp[last]);
cb8
                                   pair < 11, int > pega = \{get < 2 > (v[i]), -(get < 1 > (v[i]) - (get < 1 > (v[i]) - (get < 1 > (v[i])) - (get < 1 > (v[i]) - (get < 1 > (v[i])) - (get < 1
911
          get < 0 > (v[i]) + 1);
                                   if (nxt[i] != -1) pega.first += dp[nxt[i]].first,
5d3
          pega.second += dp[nxt[i]].second;
b08
                                   if (pega > dp[i]) dp[i] = pega;
                                   else nxt[i] = last;
7cb
381
                                   last = i;
b7c
                       }
977
                       pair < 11, int > ans = \{0, 0\};
919
                       int idx = -1:
                       for (int i = 0; i < v.size(); i++) if (dp[i] > ans) ans =
         dp[i], idx = i;
                       vector < int > ret:
4b8
                       while (idx != -1) {
fdd
d69
                                   if (get < 2 > (v[idx]) > 0 and
                                               (nxt[idx] == -1 or get<1>(v[nxt[idx]]) <</pre>
a05
          get <0>(v[idx]))) ret.push_back(idx);
e4f
                                   idx = nxt[idx]:
042
                       sort(ret.begin(), ret.end());
0ea
                       return ret;
edf
c4d }
5.10 Convex Hull Dinamico
```

```
// insert - O(log n) amortizado
// is_inside - O(log n)

Ob9 struct upper {
af8    set<pt> se;
```

```
80ъ
        set <pt>::iterator it;
        int is_under(pt p) { // 1 -> inside ; 2 -> border
25 c
            it = se.lower_bound(p);
fe0
633
            if (it == se.end()) return 0;
            if (it == se.begin()) return p == *it ? 2 : 0;
a94
ca0
            if (ccw(p, *it, *prev(it))) return 1;
            return ccw(p, *prev(it), *it) ? 0 : 2;
402
dba
        void insert(pt p) {
eaa
712
            if (is_under(p)) return;
a86
            if (it != se.end()) while (next(it) != se.end() and
   !ccw(*next(it), *it, p))
316
                it = se.erase(it);
            if (it != se.begin()) while (--it != se.begin() and
be3
   !ccw(p, *it, *prev(it)))
                it = se.erase(it):
316
0c8
            se.insert(p);
        }
5da
750 };
06f struct dyn_hull {
d93
        upper U, L;
333
        int is_inside(pt p) {
632
            int u = U.is_under(p), l = L.is_under({-p.x, -p.y});
            if (!u or !1) return 0;
4 c c
fc0
            return max(u, 1);
478
        }
        void insert(pt p) {
eaa
86c
            U.insert(p);
925
            L.insert({-p.x, -p.y});
64b
        }
285
       int size() {
7c2
            int ans = U.se.size() + L.se.size();
1c9
            return ans <= 2 ? ans/2 : ans-2;</pre>
ad5
65e };
```

# 5.11 Distancia maxima entre dois pontos

```
// max_dist2(v) - 0(n log(n))
// max_dist_manhattan - 0(n)
```

```
// Quadrado da Distancia Euclidiana (precisa copiar convex_hull, ccw e
                                                                            a58
                                                                            69a
859 ll max_dist2(vector<pt> v) {
                                                                            d1f
        v = convex_hull(v);
                                                                            4f2
        if (v.size() <= 2) return dist2(v[0], v[1%v.size()]);</pre>
a14
                                                                            460
04b
                                                                            efe
323
       int n = v.size(), j = 0;
                                                                            d6f
        for (int i = 0; i < n; i++) {</pre>
603
            while (!ccw(v[(i+1)%n]-v[i], pt(0, 0), v[(j+1)%n]-v[j])) j
057
   = (j+1)%n;
            ans = \max(\{ans, dist2(v[i], v[j]), dist2(v[(i+1)%n],
                                                                            080
   v[i])});
        }
1f6
ba7
        return ans;
bda }
// Distancia de Manhattan
c51 template < typename T> T max_dist_manhattan(vector < pair < T, T>> v) {
8eb
        T min_sum, max_sum, min_dif, max_dif;
        min_sum = max_sum = v[0].first + v[0].second;
4f5
        min_dif = max_dif = v[0].first - v[0].second;
271
c25
        for (auto [x, v] : v) {
            min_sum = min(min_sum, x+y);
1cb
683
            max_sum = max(max_sum, x+y);
782
            min_dif = min(min_dif, x-y);
                                                                            def
af7
            max_dif = max(max_dif, x-y);
                                                                            3a1
e3a
9f0
        return max(max sum - min sum. max dif - min dif):
4e9 }
5.12 Distinct Range Query
                                                                            686
// build - O(n (log n + log(sigma)))
                                                                            0a8
// query - O(log(sigma))
                                                                            3e1
                                                                            78a
789 namespace perseg { };
                                                                            edf
53d int qt[MAX];
                                                                            d03
                                                                            5cb
edc void build(vector<int>& v) {
                                                                            af6
3d2
        int n = v.size():
                                                                            d3f
                                                                            7 c.9
        perseg::build(n);
663
        map < int , int > last;
```

05e

603

817

int at = 0:

for (int i = 0; i < n; i++) {</pre>

if (last.count(v[i])) {

# 5.13 Distinct Range Query com Update

```
// build - O(n log(n))
// \text{ query - } O(\log^2(n))
// update - O(log^2(n))
774 #include <ext/pb_ds/assoc_container.hpp>
30f #include <ext/pb_ds/tree_policy.hpp>
0d7 using namespace __gnu_pbds;
4fc template <class T>
        using ord_set = tree<T, null_type, less<T>, rb_tree_tag,
        tree_order_statistics_node_update>;
042 int v[MAX], n, nxt[MAX], prv[MAX];
f60 map<int, set<int> > ocor;
e04 namespace bit {
        ord_set < pair < int , int >> bit [MAX];
        void build() {
             for (int i = 1; i <= n; i++) bit[i].insert({nxt[i-1],</pre>
   i-1});
             for (int i = 1; i <= n; i++) {</pre>
                 int j = i + (i\&-i);
                 if (j <= n) for (auto x : bit[i]) bit[j].insert(x);</pre>
             }
        }
        int pref(int p, int x) {
             int ret = 0:
bbf
            for (; p; p -= p\&-p) ret += bit[p].order_of_key({x, -INF});
edf
             return ret:
0e1
        }
d50
        int query(int 1, int r, int x) {
```

```
e55
            return pref(r+1, x) - pref(1, x);
9b4
ff3
        void update(int p, int x) {
f17
            int p2 = p;
5ed
            for (p++; p <= n; p += p&-p) {</pre>
                 bit[p].erase({nxt[p2], p2});
ca8
f6b
                 bit [p]. insert (\{x, p2\});
            }
3df
151
        }
c63 }
0a8 void build() {
383
        for (int i = 0: i < n: i++) nxt[i] = INF:</pre>
        for (int i = 0; i < n; i++) prv[i] = -INF;</pre>
7b3
d07
        vector < pair < int , int >> t;
348
        for (int i = 0; i < n; i++) t.push_back({v[i], i});</pre>
        sort(t.begin(), t.end());
3fd
603
        for (int i = 0; i < n; i++) {</pre>
            if (i and t[i].first == t[i-1].first)
b40
565
                 prv[t[i].second] = t[i-1].second;
            if (i+1 < n and t[i].first == t[i+1].first)</pre>
a8b
12f
                 nxt[t[i].second] = t[i+1].second;
        }
48d
a23
        for (int i = 0; i < n; i++) ocor[v[i]].insert(i);</pre>
1d7
        bit::build();
d44 }
aae void muda(int p, int x) {
        bit::update(p, x);
f92
c3d
        nxt[p] = x;
97c }
4ea int query(int a, int b) {
        return b-a+1 - bit::query(a, b, b+1);
a0a
511 }
ff3 void update(int p, int x) { // mudar valor na pos. p para x
c0b
        if (prv[p] > -INF) muda(prv[p], nxt[p]);
        if (nxt[p] < INF) prv[nxt[p]] = prv[p];</pre>
4ae
5bf
        ocor[v[p]].erase(p);
4b4
        if (!ocor[x].size()) {
19d
            muda(p, INF);
8d4
            prv[p] = -INF;
f6c
        } else if (*ocor[x].rbegin() < p) {</pre>
```

```
5b5
            int i = *ocor[x].rbegin();
f64
            prv[p] = i;
19d
            muda(p, INF);
5f2
            muda(i, p);
f36
       } else {
d46
            int i = *ocor[x].lower_bound(p);
33f
            if (prv[i] > -INF) {
f17
                muda(prv[i], p);
8f9
                prv[p] = prv[i];
            } else prv[p] = -INF;
bc4
523
            prv[i] = p;
597
            muda(p, i);
2ac
        v[p] = x; ocor[x].insert(p);
c96
38e }
```

#### 5.14 Dominator Points

```
// Se um ponto A tem ambas as coordenadas >= B, dizemos
// que A domina B
// is_dominated(p) fala se existe algum ponto no conjunto
// que domina p
// insert(p) insere p no conjunto
// (se p for dominado por alguem, nao vai inserir)
// o multiset 'quina' guarda informacao sobre os pontos
// nao dominados por um elemento do conjunto que nao dominam
// outro ponto nao dominado por um elemento do conjunto
// No caso, armazena os valores de x+y esses pontos
//
// Complexidades:
// is_dominated - O(log(n))
// insert - O(log(n)) amortizado
// query - 0(1)
e2a struct dominator_points {
        set<pair<int, int>> se;
baf
        multiset < int > quina;
4dd
        bool is_dominated(pair<int, int> p) {
a85
            auto it = se.lower_bound(p);
80f
            if (it == se.end()) return 0:
633
            return it->second >= p.second;
ab4
28f
99b
        void mid(pair<int, int> a, pair<int, int> b, bool rem) {
29a
            pair < int , int > m = {a.first+1, b.second+1};
b19
            int val = m.first + m.second;
```

```
638
            if (!rem) quina.insert(val);
731
            else quina.erase(quina.find(val));
241
7 c 4
        bool insert(pair<int, int> p) {
fb4
            if (is_dominated(p)) return 0;
80f
            auto it = se.lower_bound(p);
ca9
            if (it != se.begin() and it != se.end())
d4a
                mid(*prev(it), *it, 1);
1fa
            while (it != se.begin()) {
049
                it--:
23c
                if (it->second > p.second) break;
b86
                if (it != se.begin()) mid(*prev(it), *it, 1);
316
                it = se.erase(it):
acd
            }
433
            it = se.insert(p).first;
            if (it != se.begin()) mid(*prev(it), *it, 0);
69e
96d
            if (next(it) != se.end()) mid(*it, *next(it), 0);
6a5
            return 1;
       }
688
5eb
        int query() {
956
            if (!quina.size()) return INF;
add
            return *quina.begin();
       }
b8b
09f };
```

## 5.15 DP de Dominação 3D

```
// Computa para todo ponto i,
// dp[i] = 1 + max_{j dominado por i} dp[j]
// em que ser dominado eh ter as 3 coordenadas menores
// Da pra adaptar facil para outras dps
//
// O(n log^2 n), O(n) de memoria
c53 void lis2d(vector<vector<tuple<int, int, int>>>& v, vector<int>&
   dp, int 1, int r) {
       if (1 == r) {
893
            for (int i = 0; i < v[1].size(); i++) {</pre>
56f
                int ii = get <2>(v[1][i]);
8b5
                dp[ii] = max(dp[ii], 1);
1ce
4b0
            }
505
            return;
3e4
ee4
        int m = (1+r)/2;
62b
        lis2d(v, dp, 1, m);
```

```
325
        vector<tuple<int, int, int>> vv[2];
d44
        vector < int > Z;
871
        for (int i = 1; i <= r; i++) for (auto it : v[i]) {</pre>
2ef
            vv[i > m].push_back(it);
042
            Z.push_back(get<1>(it));
0d1
e9f
        sort(vv[0].begin(), vv[0].end());
9b5
        sort(vv[1].begin(), vv[1].end());
0d1
        sort(Z.begin(), Z.end());
573
        auto get_z = [&](int z) { return lower_bound(Z.begin(),
   Z.end(), z) - Z.begin(); };
        vector < int > bit(Z.size());
c51
181
        int i = 0;
e9a
        for (auto [y, z, id] : vv[1]) {
            while (i < vv[0].size() and get<0>(vv[0][i]) < y) {</pre>
6bd
397
                auto [v2, z2, id2] = vv[0][i++];
ea0
                for (int p = get_z(z2)+1; p <= Z.size(); p += p\&-p)
300
                     bit[p-1] = max(bit[p-1], dp[id2]);
            }
82c
d3b
            int q = 0;
fd9
            for (int p = get_z(z); p; p -= p\&-p) q = max(q, bit[p-1]);
614
            dp[id] = max(dp[id], q + 1);
acc
c25
        lis2d(v, dp, m+1, r);
4d6 }
4de vector < int > solve (vector < tuple < int , int , int >> v) {
        int n = v.size();
        vector<tuple<int, int, int, int>> vv;
cd4
        for (int i = 0; i < n; i++) {</pre>
603
9be
            auto [x, y, z] = v[i];
            vv.emplace_back(x, y, z, i);
5bb
64c
        }
        sort(vv.begin(), vv.end());
bd3
e11
        vector < vector < tuple < int , int , int >>> V;
603
        for (int i = 0; i < n; i++) {</pre>
a5b
            int j = i;
808
            V.emplace_back();
            while (j < n and get <0>(vv[j]) == get <0>(vv[i])) {
c01
                auto [x, y, z, id] = vv[j++];
ba6
                V.back().emplace_back(y, z, id);
cbb
            }
8bd
452
            i = j-1;
ac4
        }
388
        vector < int > dp(n);
```

```
839 lis2d(V, dp, 0, V.size()-1);
898 return dp;
b0a }
```

## 5.16 Gray Code

```
// Gera uma permutacao de 0 a 2^n-1, de forma que
// duas posicoes adjacentes diferem em exatamente 1 bit
//
// O(2^n)

df6 vector<int> gray_code(int n) {
73f vector<int> ret(1<<n);
for (int i = 0; i < (1<<n); i++) ret[i] = i^(i>>1);
edf return ret;
840 }
```

# 5.17 Half-plane intersection

```
// Cada half-plane eh identificado por uma reta e a regiao ccw a ela
// O(n log n)
f4f vector <pt> hp_intersection(vector <line> &v) {
9bc
                              deque < pt > dq = \{\{INF, INF\}, \{-INF, INF\}, \{-INF, -INF\}, \{INF, \}, \{-INF, -INF\}, \{INF, -INF\}, 
             -INF}};
d41 #warning considerar trocar por compare_angle
                              sort(v.begin(), v.end(), [&](line r, line s) { return
             angle(r.q-r.p) < angle(s.q-s.p); });
                              for(int i = 0; i < v.size() and dq.size() > 1; i++) {
5e9
c69
                                              pt p1 = dq.front(), p2 = dq.back();
                                              while (dq.size() and !ccw(v[i].p, v[i].q, dq.back()))
6c6
47b
                                                             p1 = dq.back(), dq.pop_back();
                                              while (dg.size() and !ccw(v[i].p, v[i].q, dq.front()))
0a2
                                                             p2 = dq.front(), dq.pop_front();
7cf
4d9
                                              if (!dq.size()) break;
606
                                              if (p1 == dg.front() and p2 == dg.back()) continue;
c9b
                                              dq.push_back(inter(v[i], line(dq.back(), p1)));
65 c
                                              dq.push_front(inter(v[i], line(dq.front(), p2)));
```

```
// O(n log n)
f18 void down(vector<int>& v, int n, int i) {
e1f
        while ((i = 2*i+1) < n) {
583
            if (i+1 < n and v[i] < v[i+1]) i++;</pre>
b27
            if (v[i] < v[(i-1)/2]) break;
322
            swap(v[i], v[(i-1)/2]);
        }
170
724 }
eb6 void heap_sort(vector<int>& v) {
        int n = v.size();
3d2
61d
        for (int i = n/2-1; i \ge 0; i--) down(v, n, i);
917
        for (int i = n-1; i > 0; i--)
37f
            swap(v[0], v[i]), down(v, i, 0);
b33 }
```

# 5.19 Hungaro

```
// Resolve o problema de assignment (matriz n x n)
// Colocar os valores da matriz em 'a' (pode < 0)</pre>
// assignment() retorna um par com o valor do
// assignment minimo, e a coluna escolhida por cada linha
//
// O(n^3)
a6a template < typename T > struct hungarian {
1a8
        int n;
        vector < vector < T >> a;
a08
f36
         vector<T> u, v;
        vector < int > p, way;
5ff
f1e
        T inf:
        hungarian(int n_{-}): n(n_{-}), u(n+1), v(n+1), p(n+1), way(n+1) {
c3f
b2f
             a = vector < vector < T >> (n, vector < T > (n));
1f3
             inf = numeric_limits <T>::max();
78f
        }
```

```
d67
        pair <T, vector <int >> assignment() {
            for (int i = 1; i <= n; i++) {</pre>
78a
8c9
                 p[0] = i;
                 int i0 = 0;
625
                 vector <T> minv(n+1, inf);
ce7
                 vector < int > used(n+1, 0);
241
016
                 do {
472
                     used[i0] = true;
d24
                     int i0 = p[j0], j1 = -1;
                     T delta = inf:
7 e 5
9ac
                     for (int j = 1; j <= n; j++) if (!used[j]) {</pre>
7bf
                         T cur = a[i0-1][j-1] - u[i0] - v[j];
9f2
                         if (cur < minv[j]) minv[j] = cur, way[j] = j0;</pre>
821
                         if (minv[j] < delta) delta = minv[j], j1 = j;</pre>
                     }
4d1
f63
                     for (int j = 0; j \le n; j++)
2c5
                         if (used[i]) u[p[i]] += delta, v[i] -= delta;
                         else minv[j] -= delta;
6ec
6d4
                     i0 = i1;
f4f
                 } while (p[j0] != 0);
016
                 do {
                     int j1 = wav[j0];
4c5
0d7
                     p[j0] = p[j1];
6d4
                     j0 = j1;
886
                 } while (j0);
38d
306
            vector < int > ans(n);
6db
            for (int j = 1; j \le n; j++) ans [p[j]-1] = j-1;
            return make_pair(-v[0], ans);
da3
979
        }
64c };
5.20 Inversion Count
// Computa o numero de inversoes para transformar
// l em r (se nao tem como, retorna -1)
// O(n log(n))
37b template < typename T > 1l inv_count(vector < T > 1, vector < T > r = {}) {
bb6
        if (!r.size()) {
796
            r = 1:
```

sort(r.begin(), r.end());

int n = l.size();

vector < int > v(n), bit(n);

1bc dfb

874

8c0

```
4e9
        vector < pair < T, int >> w;
61c
        for (int i = 0; i < n; i++) w.push_back({r[i], i+1});</pre>
d1d
        sort(w.begin(), w.end());
603
        for (int i = 0; i < n; i++) {</pre>
            auto it = lower_bound(w.begin(), w.end(), make_pair(l[i],
bf3
   0));
1bf
            if (it == w.end() or it->first != 1[i]) return -1: // nao
962
            v[i] = it->second;
            it->second = -1:
6.0
964
        }
04b
        11 \text{ ans} = 0:
45b
        for (int i = n-1; i >= 0; i--) {
2d9
            for (int j = v[i]-1; j; j -= j\&-j) ans += bit[j];
3a1
            for (int j = v[i]; j < n; j += j\&-j) bit[j]++;
ebe
        }
ba7
        return ans;
eef }
5.21 LIS - recupera
// Calcula e retorna uma LIS
// O(n.log(n))
121 template < typename T> vector < T> lis(vector < T>& v) {
1fa
        int n = v.size(), m = -1;
f0c
        vector <T> d(n+1. INF):
        vector < int > l(n);
aec
007
        d[0] = -INF:
603
        for (int i = 0; i < n; i++) {</pre>
            // Para non-decreasing use upper_bound()
            int t = lower_bound(d.begin(), d.end(), v[i]) - d.begin();
4fd
            d[t] = v[i], l[i] = t, m = max(m, t);
3ad
89 c
        }
4ff
        int p = n;
5a9
        vector <T> ret;
cdf
        while (p--) if (l[p] == m) {
883
            ret.push_back(v[p]);
76b
f83
        }
969
        reverse(ret.begin(),ret.end());
```

```
edf
        return ret;
474 }
5.22 LIS - tamanho
// Calcula o tamanho da LIS
// O(n log(n))
84b template < typename T > int lis(vector < T > &v){
        vector <T> ans:
2da
5e0
        for (T t : v){
            // Para non-decreasing use upper_bound()
            auto it = lower_bound(ans.begin(), ans.end(), t);
fe6
            if (it == ans.end()) ans.push_back(t);
d7f
b94
            else *it = t:
1f5
1eb
        return ans.size();
402 }
     Minimum Enclosing Circle
// O(n) com alta probabilidade
```

```
22c const double EPS = 1e-12;
878 mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
b2a struct pt {
662
        double x, y;
be7
        pt(double x_{=} = 0, double y_{=} = 0) : x(x_{=}), y(y_{=}) {}
        pt operator + (const pt& p) const { return pt(x+p.x, y+p.y); }
b23
        pt operator - (const pt& p) const { return pt(x-p.x, y-p.y); }
        pt operator * (double c) const { return pt(x*c, y*c); }
254
701
        pt operator / (double c) const { return pt(x/c, y/c); }
54d };
2f9 double dot(pt p, pt q) { return p.x*q.x+p.y*q.y; }
dd5 double cross(pt p, pt q) { return p.x*q.y-p.y*q.x; }
e7c double dist(pt p, pt q) { return sqrt(dot(p-q, p-q)); }
3f4 pt center(pt p, pt q, pt r) {
       pt a = p-r, b = q-r;
5d9
e84
        pt c = pt(dot(a, p+r)/2, dot(b, q+r)/2);
```

```
return pt(cross(c, pt(a.y, b.y)), cross(pt(a.x, b.x), c)) /
   cross(a, b);
fc8 }
aa8 struct circle {
        pt cen;
c12
        double r:
898
        circle(pt cen_, double r_) : cen(cen_), r(r_) {}
83 c
        circle(pt a, pt b, pt c) {
            cen = center(a, b, c);
13d
1f1
            r = dist(cen, a);
bc1
cd5
        bool inside(pt p) { return dist(p, cen) < r+EPS; }</pre>
2a6 };
806 circle minCirc(vector<pt> v) {
f21
        shuffle(v.begin(), v.end(), rng);
        circle ret = circle(pt(0, 0), 0);
ae0
618
        for (int i = 0; i < v.size(); i++) if (!ret.inside(v[i])) {</pre>
            ret = circle(v[i]. 0):
16a
f11
            for (int j = 0; j < i; j++) if (!ret.inside(v[j])) {</pre>
                 ret = circle((v[i]+v[j])/2, dist(v[i], v[j])/2);
881
                for (int k = 0; k < j; k++) if (!ret.inside(v[k]))
b8c
43f
                     ret = circle(v[i], v[j], v[k]);
5f8
            }
6a1
        }
edf
        return ret;
eba }
5.24 Minkowski Sum
// Computa A+B = \{a+b : a \setminus A, b \setminus B\}, em que
// A e B sao poligonos convexos
// A+B eh um poligono convexo com no max |A|+|B| pontos
```

```
//
// O(|A|+|B|)
539 vector <pt > minkowski (vector <pt > p, vector <pt > q) {
051
        auto fix = [](vector<pt>& P) {
515
             rotate(P.begin(), min_element(P.begin(), P.end()),
   P.end()):
018
            P.push_back(P[0]), P.push_back(P[1]);
f24
        };
889
        fix(p), fix(q);
8af
        vector<pt> ret;
692
        int i = 0, j = 0;
```

```
2ee
        while (i < p.size()-2 or j < q.size()-2) {</pre>
            ret.push_back(p[i] + q[j]);
898
            auto c = ((p[i+1] - p[i]) ^ (q[j+1] - q[j]));
732
           if (c >= 0) i = min<int>(i+1, p.size()-2);
ebc
           if (c <= 0) j = min<int>(j+1, q.size()-2);
81e
9ff
edf
        return ret:
d7c }
c3e ld dist_convex(vector<pt> p, vector<pt> q) {
dc2
        for (pt& i : p) i = i * -1;
44c
        auto s = minkowski(p, q);
95d
       if (inpol(s, pt(0, 0))) return 0;
6a5
       return 1:
921
       ld ans = DINF;
073
       for (int i = 0; i < s.size(); i++) ans = min(ans,</pre>
f04
                disttoseg(pt(0, 0), line(s[(i+1)%s.size()], s[i])));
ba7
        return ans;
2f5 }
5.25 MO
// Para ter o bound abaixo, escolher
// SQ = n / sqrt(q)
//
// O(n * sqrt(q))
0d2 const int MAX = 1e5+10;
6ff const int SQ = sqrt(MAX);
b69 int v[MAX];
b65 int ans, freq[MAX];
9da inline void insert(int p) {
       int o = v[p];
ae0
591
       freq[o]++;
992
        ans += (freq[o] == 1);
21d }
a25 inline void erase(int p) {
ae0
       int o = v[p]:
        ans -= (freq[o] == 1);
ba2
       freq[o]--;
dc7 }
e51 inline ll hilbert(int x, int y) {
```

```
71e
        static int N = 1 << (_builtin_clz(0) - _builtin_clz(MAX));</pre>
100
        int rx, ry, s;
b72
        11 d = 0;
43b
        for (s = N/2; s > 0; s /= 2) {
c95
            rx = (x \& s) > 0, ry = (y \& s) > 0;
            d += s * 11(s) * ((3 * rx) ^ ry);
еЗе
d2e
            if (rv == 0) {
5aa
                if (rx == 1) x = N-1 - x, y = N-1 - y;
9dd
                swap(x, y);
            }
e2d
888
        }
be2
        return d;
7fa }
bac #define HILBERT true
617 vector<int> MO(vector<pair<int, int>> &q) {
c3b
        ans = 0;
        int m = q.size();
c23
3f8
        vector < int > ord(m);
        iota(ord.begin(), ord.end(), 0);
be8
6a6 #if HILBERT
8c4
        vector<ll> h(m);
        for (int i = 0; i < m; i++) h[i] = hilbert(q[i].first,</pre>
74c
   q[i].second);
075
        sort(ord.begin(), ord.end(), [&](int 1, int r) { return h[1] <</pre>
   h[r]: }):
8c1 #else
d01
        sort(ord.begin(), ord.end(), [&](int 1, int r) {
            if (q[1].first / SQ != q[r].first / SQ) return q[1].first
   < q[r].first;
            if ((q[1].first / SQ) % 2) return q[1].second >
   q[r].second;
            return q[1].second < q[r].second;</pre>
a66
bec
        }):
f2e #endif
435
        vector < int > ret(m);
3d9
        int 1 = 0, r = -1;
8ъ0
        for (int i : ord) {
6c6
            int ql, qr;
4f5
            tie(ql, qr) = q[i];
026
            while (r < qr) insert(++r);</pre>
232
            while (1 > q1) insert(--1);
75e
            while (1 < q1) erase(1++);</pre>
fe8
            while (r > qr) erase(r--);
381
            ret[i] = ans:
116
        }
```

```
edf
        return ret;
fb7 }
5.26 MO - DSU
// Dado uma lista de arestas de um grafo, responde
// para cada query(1, r), quantos componentes conexos
// o grafo tem se soh considerar as arestas 1, 1+1, ..., r
// Da pra adaptar pra usar MO com qualquer estrutura rollbackavel
// O(m sqrt(q) log(n))
8d3 struct dsu {
553
        int n, ans;
2e3
        vector < int > p, sz;
        stack<int> S;
ee6
4b8
        dsu(int n_{-}) : n(n_{-}), ans(n), p(n), sz(n) {
            for (int i = 0; i < n; i++) p[i] = i, sz[i] = 1;</pre>
8a6
aae
        }
1 b 1
        int find(int k) {
006
            while (p[k] != k) k = p[k];
839
            return k:
c1e
        }
553
        void add(pair<int, int> x) {
700
            int a = x.first, b = x.second;
605
            a = find(a), b = find(b);
843
            if (a == b) return S.push(-1);
e7d
            ans - -:
3c6
            if (sz[a] > sz[b]) swap(a, b);
4c2
            S.push(a);
582
            sz[b] += sz[a];
84b
            p[a] = b;
720
        }
        int query() { return ans; }
35c
        void rollback() {
5cf
465
            int u = S.top(); S.pop();
            if (u == -1) return;
61c
270
            sz[p[u]] -= sz[u]:
            p[u] = u;
546
Odf
            ans++:
456
        }
9c1 };
1a8 int n;
e93 vector<pair<int, int>> ar;
```

```
// 9d242b
617 vector <int > MO(vector <pair <int, int >> &q) {
        int SQ = sqrt(q.size()) + 1;
        int m = q.size();
c23
3f8
        vector<int> ord(m);
be8
        iota(ord.begin(), ord.end(), 0);
        sort(ord.begin(), ord.end(), [&](int 1, int r) {
d01
9c9
                 if (q[1].first / SQ != q[r].first / SQ) return
   q[1].first < q[r].first;</pre>
a66
                 return q[1].second < q[r].second;</pre>
b90
                 }):
435
        vector < int > ret(m);
3bd
        dsu small(n);
        for (int i = 0; i < m; i++) {</pre>
dd5
5ec
             auto [1, r] = q[ord[i]];
             if (1 / SQ == r / SQ) {
acc
                 for (int k = 1; k <= r; k++) small.add(ar[k]);</pre>
00c
                 ret[ord[i]] = small.query();
b99
                 for (int k = 1; k <= r; k++) small.rollback();</pre>
64 e
259
            }
        }
6b0
dd5
        for (int i = 0; i < m; i++) {</pre>
176
             dsu D(n):
ae9
             int fim = q[ord[i]].first/SQ*SQ + SQ - 1;
e25
             int last r = fim:
ebc
             int j = i-1;
00c
             while (j+1 < m and g[ord[j+1]].first / SQ ==</pre>
    q[ord[i]].first / SQ) {
a0e
                 auto [1, r] = q[ord[++j]];
f58
                 if (1 / SQ == r / SQ) continue:
59b
                 while (last_r < r) D.add(ar[++last_r]);</pre>
2cf
                 for (int k = 1; k <= fim; k++) D.add(ar[k]);</pre>
9b2
                 ret[ord[j]] = D.query();
572
                 for (int k = 1; k \le fim; k++) D.rollback();
9c8
             }
bdf
             i = j;
        }
e99
edf
        return ret;
9d2 }
```

#### 5.27 MO em Arvores

```
// Problema que resolve: https://www.spoj.com/problems/COT2/
// Complexidade sendo c = O(update) e SQ = sqrt(n):
// O((n + a) * sart(n) * c)
1bc const int MAX = 40010, SQ = 400;
042 vector < int > g[MAX];
c54 namespace LCA { ... }
249 int in[MAX], out[MAX], vtx[2 * MAX];
81b bool on [MAX];
4c3 int dif, freq[MAX];
9e2 vector < int > w;
d9a void dfs(int v, int p, int &t) {
659
        vtx[t] = v, in[v] = t++;
        for (int u : g[v]) if (u != p) {
c53
            dfs(u, v, t);
e0f
217
        vtx[t] = v, out[v] = t++;
42b }
e5f void update(int p) { // faca alteracoes aqui
bbc
        int v = vtx[p];
0ec
        if (not on[v]) { // insere vtx v
31c
            dif += (freq[w[v]] == 0);
b20
            freq[w[v]]++;
cf7
4e6
        else { // retira o vertice v
0a9
            dif -= (freq[w[v]] == 1);
fd3
            freq[w[v]]--;
2c8
73e
        on[v] = not on[v];
ea9 }
a3a vector < tuple < int, int, int >> build_queries (const vector < pair < int,
   int>>& q) {
        LCA::build(0):
ea6
f77
        vector<tuple<int, int, int>> ret;
aa9
        for (auto [1, r] : q){
d24
            if (in[r] < in[l]) swap(l, r);</pre>
6f9
            int p = LCA::lca(l, r);
```

```
826
            int init = (p == 1) ? in[1] : out[1];
07a
            ret.emplace_back(init, in[r], in[p]);
b0e
        }
edf
        return ret;
8e6 }
f31 vector<int> mo_tree(const vector<pair<int, int>>& vq){
        int t = 0:
dab
        dfs(0, -1, t);
        auto q = build_queries(vq);
af1
f48
        vector<int> ord(q.size());
be8
        iota(ord.begin(), ord.end(), 0);
d01
        sort(ord.begin(), ord.end(), [&] (int 1, int r) {
d8d
             int bl = get<0>(q[1]) / SQ, br = <math>get<0>(q[r]) / SQ;
596
            if (bl != br) return bl < br;</pre>
            else if (bl % 2 == 1) return get<1>(q[1]) < get<1>(q[r]);
158
f1d
            else return get<1>(q[1]) > get<1>(q[r]);
0a8
        });
80e
        memset(freq, 0, sizeof freq);
bf6
        dif = 0:
ff2
        vector < int > ret(q.size());
3d9
        int 1 = 0, r = -1:
        for (int i : ord) {
8b0
3c7
            auto [ql, qr, qp] = q[i];
af7
            while (r < qr) update(++r);</pre>
d6b
            while (1 > q1) update(--1);
            while (1 < q1) update(1++);</pre>
951
6a1
            while (r > qr) update(r--);
3d8
            if (qp < 1 \text{ or } qp > r)  { // se LCA estah entre as pontas
74b
                 update(qp);
                 ret[i] = dif;
2e1
74b
                 update(qp);
e83
0fe
            else ret[i] = dif;
0fd
        }
edf
        return ret;
48d }
     Palindromic Factorization
```

// Precisa da eertree

```
// Computa o numero de formas de particionar cada
// prefixo da string em strings palindromicas
// O(n log n), considerando alfabeto O(1)
070 struct eertree { ... };
0e7 ll factorization(string s) {
b19
        int n = s.size(), sz = 2;
580
        eertree PT(n):
147
        vector \langle int \rangle diff (n+2), slink (n+2), sans (n+2), dp (n+1);
0ec
        dp[0] = 1;
78a
        for (int i = 1: i <= n: i++) {
c58
            PT.add(s[i-1]):
a7c
            if (PT.size()+2 > sz) {
6c4
                diff[sz] = PT.len[sz] - PT.len[PT.link[sz]];
241
                if (diff[sz] == diff[PT.link[sz]])
                    slink[sz] = slink[PT.link[sz]];
d6f
f53
                else slink[sz] = PT.link[sz];
eb9
                sz++:
f6a
            }
            for (int v = PT.last; PT.len[v] > 0; v = slink[v]) {
911
                sans[v] = dp[i - (PT.len[slink[v]] + diff[v])];
297
85d
                if (diff[v] == diff[PT.link[v]])
f20
                     sans[v] = (sans[v] + sans[PT.link[v]]) % MOD;
071
                dp[i] = (dp[i] + sans[v]) % MOD;
            }
e5e
fc0
5f0
        return dp[n];
3a7 }
```

# 5.29 Parsing de Expressao

```
// Operacoes associativas a esquerda por default
// Para mudar isso, colocar em r_assoc
// Operacoes com maior prioridade sao feitas primeiro

cc1 bool blank(char c) {
f34     return c == ' ';
ec3 }

8e4 bool is_unary(char c) {
f9c     return c == '+' or c == '-';
b6b }

76d bool is_op(char c) {
```

```
010
        if (is_unary(c)) return true;
        return c == '*' or c == '/' or c == '+' or c == '-';
31c
4e4 }
fa3 bool r_assoc(char op) {
        // operator unario - deve ser assoc. a direita
cf0
        return op < 0;</pre>
c5c }
79d int priority(char op) {
        // operator unario - deve ter precedencia maior
103
        if (op < 0) return INF;</pre>
        if (op == '*' or op == '/') return 2;
727
        if (op == '+' or op == '-') return 1;
439
        return -1;
daa
966 }
c15 void process_op(stack<int>& st, stack<int>& op) {
88c
        char o = op.top(); op.pop();
91c
        if (o < 0) {
4e6
            o *= -1:
            int 1 = st.top(); st.pop();
1e2
Off
            if (o == '+') st.push(1);
7e9
            if (o == '-') st.push(-1);
320
        } else {
            int r = st.top(); st.pop();
14c
1e2
            int 1 = st.top(); st.pop();
1 e 4
            if (o == '*') st.push(1 * r);
f55
            if (o == '/') st.push(1 / r);
605
            if (o == '+') st.push(1 + r);
c40
            if (o == '-') st.push(1 - r);
        }
0aa
2b2 }
439 int eval(string& s) {
        stack<int> st, op;
212
d0c
        bool un = true;
1cf
        for (int i = 0; i < s.size(); i++) {</pre>
68d
            if (blank(s[i])) continue;
139
            if (s[i] == '(') {
367
                op.push('(');
99d
                un = true;
b88
            } else if (s[i] == ')') {
709
                while (op.top() != '(') process_op(st, op);
75e
                op.pop();
```

```
ce2
                un = false:
003
            } else if (is_op(s[i])) {
4d0
                char o = s[i];
                if (un and is_unary(o)) o *= -1;
37 c
                while (op.size() and (
ae3
                             (!r_assoc(o) and priority(op.top()) >=
   priority(o)) or
                             (r_assoc(o) and priority(op.top()) >
c41
   priority(o))))
                    process_op(st, op);
                op.push(o);
c00
99d
                un = true:
196
            } else {
da8
                int val = 0:
c2b
                while (i < s.size() and isalnum(s[i]))</pre>
                    val = val * 10 + s[i++] - '0':
8a3
169
                i--;
25d
                st.push(val);
ce2
                un = false;
            }
442
b19
        }
7f6
        while (op.size()) process_op(st, op);
123
        return st.top();
05c }
```

# 5.30 RMQ com Divide and Conquer

```
// Responde todas as queries em
// O(n log(n))
f74 typedef pair <pair <int, int>, int> iii;
7c6 #define f first
Oab #define s second
87d int n, q, v[MAX];
e3f iii qu[MAX];
aeb int ans[MAX], pref[MAX], sulf[MAX];
0e3 void solve(int l=0, int r=n-1, int ql=0, int qr=q-1) {
8a3
        if (1 > r or q1 > qr) return;
        int m = (1+r)/2:
        int qL = partition(qu+ql, qu+qr+1, [=](iii x){return x.f.s <</pre>
1b1
   m;}) - qu;
        int qR = partition(qu+qL, qu+qr+1, [=](iii x){return x.f.f
   <=m;}) - qu;
```

```
pref[m] = sulf[m] = v[m];
3cd
9f9
        for (int i = m-1; i >= 1; i--) pref[i] = min(v[i], pref[i+1]);
        for (int i = m+1; i <= r; i++) sulf[i] = min(v[i], sulf[i-1]);</pre>
ea8
        for (int i = qL; i < qR; i++)</pre>
b2a
f3a
            ans[qu[i].s] = min(pref[qu[i].f.f], sulf[qu[i].f.s]);
364
        solve(l, m-1, ql, qL-1), solve(m+1, r, qR, qr);
13e }
5.31 Segment Intersection
// Verifica, dado n segmentos, se existe algum par de segmentos
// que se intersecta
//
// O(n log n)
6e0 bool operator < (const line& a, const line& b) { // comparador pro
   sweepline
        if (a.p == b.p) return ccw(a.p, a.q, b.q);
191
        if (!eq(a.p.x, a.q.x)) and (eq(b.p.x, b.q.x)) or a.p.x+eps <
231
   b.p.x))
780
            return ccw(a.p, a.q, b.p);
dc0
        return ccw(a.p, b.q, b.p);
e36 }
8e2 bool has intersection(vector<line> v) {
        auto intersects = [&](pair<line, int> a, pair<line, int> b) {
a08
            return interseg(a.first, b.first);
3e6
        };
e1b
        vector<pair<pt, pair<int, int>>> w;
        for (int i = 0; i < v.size(); i++) {</pre>
f14
876
            if (v[i].q < v[i].p) swap(v[i].p, v[i].q);</pre>
            w.push_back({v[i].p, {0, i}});
e1d
034
            w.push_back({v[i].q, {1, i}});
220
        }
d1d
        sort(w.begin(), w.end());
7f2
        set < pair < line, int >> se;
e58
        for (auto i : w) {
            line at = v[i.second.second];
bfd
292
            if (i.second.first == 0) {
145
                auto nxt = se.lower_bound({at, i.second.second});
d1e
                if (nxt != se.end() and intersects(*nxt, {at,
   i.second.second})) return 1;
```

```
257
                if (nxt != se.begin() and intersects(*(--nxt), {at,
   i.second.second})) return 1;
                se.insert({at, i.second.second});
78f
08b
884
                auto nxt = se.upper_bound({at, i.second.second}), cur
   = nxt, prev = --cur;
b64
                if (nxt != se.end() and prev != se.begin()
                    and intersects(*nxt, *(--prev))) return 1;
                se.erase(cur);
cca
            }
e27
a00
       }
bb3
        return 0;
196 }
```

## 5.32 Sequencia de de Brujin

```
// Se passar sem o terceiro parametro, gera um vetor com valores
// em [0, k) de tamanho k^n de forma que todos os subarrays ciclicos
// de tamanho n ocorrem exatamente uma vez
// Se passar com um limite lim, gera o menor vetor com valores
// em [0, k) que possui lim subarrays de tamanho n distintos
// (assume que lim <= k^n)</pre>
//
// Linear no tamanho da resposta
860 vector<int> de_brujin(int n, int k, int lim = INF) {
b55
        if (k == 1) return vector<int>(lim == INF ? 1 : n, 0);
5f6
        vector<int> 1 = {0}, ret; // 1 eh lyndon word
667
        while (true) {
c86
            if (1.size() == 0) {
1b9
                if (lim == INF) break;
daf
                1.push_back(0);
            }
            if (n % 1.size() == 0) for (int i : 1) {
686
728
                ret.push_back(i);
c99
                if (ret.size() == n+lim-1) return ret;
            }
56e
            int p = 1.size();
630
            while (1.size() < n) 1.push_back(1[1.size()%p]);</pre>
905
e7f
            while (1.size() and 1.back() == k-1) 1.pop_back();
            if (1.size()) 1.back()++;
88a
2ef
edf
        return ret:
197 }
```

#### 5.33 Shortest Addition Chain

```
// Computa o menor numero de adicoes para construir
// cada valor. comecando com 1 (e podendo salvar variaveis)
// Retorna um par com a dp e o pai na arvore
// A arvore eh tao que o taminho da raiz (1) ate x
// contem os valores que devem ser criados para gerar x
// A profundidade de x na arvore eh dp[x]
// DP funciona para ateh 300, mas a arvore soh funciona
// para ateh 148
// recuperacao certa soh ateh 148 (erra para 149, 233, 298)
3de pair < vector < int >, vector < int >> addition_chain() {
        int MAX = 301;
875
        vector < int > dp(MAX), p(MAX);
        for (int n = 2; n < MAX; n++) {
1ab
            pair < int , int > val = {INF , -1};
7c0
            for (int i = 1; i < n; i++) for (int j = i; j; j = p[j])
212
94a
                if (j == n-i) val = min(val, pair(dp[i]+1, i));
eb3
            tie(dp[n], p[n]) = val;
efe
            if (n == 9) p[n] = 8;
            if (n == 149 \text{ or } n == 233) dp[n]--;
ba1
bcd
717
        return {dp, p};
84f }
```

# 5.34 Simple Polygon

```
// Verifica se um poligono com n pontos eh simples
//
// O(n log n)
6e0 bool operator < (const line& a, const line& b) { // comparador pro
   sweepline
191
        if (a.p == b.p) return ccw(a.p, a.q, b.q);
        if (!eq(a.p.x, a.q.x)) and (eq(b.p.x, b.q.x)) or a.p.x+eps <
231
   b.p.x))
780
            return ccw(a.p, a.q, b.p);
dc0
        return ccw(a.p, b.q, b.p);
e36 }
6f3 bool simple(vector<pt> v) {
576
        auto intersects = [&](pair<line, int> a, pair<line, int> b) {
e72
            if ((a.second+1)%v.size() == b.second or
80e
                (b.second+1)%v.size() == a.second) return false;
```

```
a08
            return interseg(a.first, b.first);
1c5
41a
        vector<line> seg;
        vector < pair < pt , pair < int , int >>> w;
e1b
f14
        for (int i = 0; i < v.size(); i++) {</pre>
            pt at = v[i], nxt = v[(i+1)%v.size()];
0a8
828
            if (nxt < at) swap(at, nxt);</pre>
937
            seg.push_back(line(at, nxt));
f7e
            w.push_back({at, {0, i}});
            w.push_back({nxt, {1, i}});
69 c
            // casos degenerados estranhos
            if (isinseg(v[(i+2)%v.size()], line(at, nxt))) return 0;
88d
            if (isinseg(v[(i+v.size()-1)%v.size()], line(at, nxt)))
   return 0;
cba
d1d
        sort(w.begin(), w.end());
7f2
        set < pair < line, int >> se;
        for (auto i : w) {
e58
ff8
            line at = seg[i.second.second];
292
            if (i.second.first == 0) {
                auto nxt = se.lower_bound({at, i.second.second});
145
                if (nxt != se.end() and intersects(*nxt, {at,
7 c 4
   i.second.second})) return 0;
                if (nxt != se.begin() and intersects(*(--nxt), {at,
b34
   i.second.second})) return 0;
78f
                se.insert({at, i.second.second});
537
            } else {
884
                auto nxt = se.upper_bound({at, i.second.second}), cur
   = nxt, prev = --cur;
b64
                if (nxt != se.end() and prev != se.begin()
                     and intersects(*nxt, *(--prev))) return 0;
403
                se.erase(cur);
cca
            }
7be
d17
        }
6a5
        return 1;
af3 }
5.35 Sweep Direction
// Passa por todas as ordenacoes dos pontos definitas por "direcoes"
// Assume que nao existem pontos coincidentes
```

```
// Passa por todas as ordenacoes dos pontos definitas por "direcoes"
// Assume que nao existem pontos coincidentes
//
// O(n^2 log n)

4b8 void sweep_direction(vector<pt> v) {
3d2    int n = v.size();
```

```
163
        sort(v.begin(), v.end(), [](pt a, pt b) {
3a5
            if (a.x != b.x) return a.x < b.x;</pre>
572
            return a.v > b.v;
79a
        });
        vector < int > at(n);
b89
        iota(at.begin(), at.end(), 0);
516
b79
        vector<pair<int, int>> swapp;
25e
        for (int i = 0; i < n; i++) for (int j = i+1; j < n; j++)
95f
            swapp.push_back({i, j}), swapp.push_back({j, i});
269
        sort(swapp.begin(), swapp.end(), [&](auto a, auto b) {
134
            pt A = rotate90(v[a.first] - v[a.second]);
247
            pt B = rotate90(v[b.first] - v[b.second]):
            if (quad(A) == quad(B) and !sarea2(pt(0, 0), A, B)) return
615
   a < b;
            return compare_angle(A, B);
224
5e7
        });
        for (auto par : swapp) {
4e6
e24
            assert(abs(at[par.first] - at[par.second]) == 1);
a 96
            int 1 = min(at[par.first], at[par.second]),
                r = n-1 - max(at[par.first], at[par.second]);
0.43
            // l e r sao quantos caras tem de cada lado do par de
                pontos
            // (cada par eh visitado duas vezes)
9cf
            swap(v[at[par.first]], v[at[par.second]]);
1c0
            swap(at[par.first], at[par.second]);
        }
241
6bb }
```

# 5.36 Triangulação de Delaunay

```
// Computa a triangulacao de Delaunay, o dual
// do diagrama de Voronoi (a menos de casos degenerados)
// Retorna um grafo indexado pelos indices dos pontos, e as arestas
// sao as arestas da triangulacao
// As arestas partindo de um vertice ja vem ordenadas por angulo,
// ou seja, se o vertice v nao esta no convex hull, (v, v_i, v_{i+1})
// eh um triangulo da triangulacao, em que v_i eh o i-esimo vizinho
// Usa o alg d&c, precisa representar MAX_COOR^4, por isso __int128
// pra aguentar valores ateh 1e9
//
// Propriedades:
// 1 - O grafo tem no max 3n-6 arestas
// 2 - Para todo triangulo, a circunf. que passa pelos 3 pontos
// nao contem estritamente nenhum ponto
// 3 - A MST euclidiana eh subgrafo desse grafo
```

```
// 4 - Cada ponto eh vizinho do ponto mais proximo dele
// O(n log n)
2ad typedef struct QuadEdge* Q;
ba5 struct QuadEdge {
53e
        int id:
114
        pt o;
        Q rot, nxt;
41e
3e5
       bool used;
3fc
        QuadEdge(int id_ = -1, pt o_ = pt(INF, INF)) :
            id(id_), o(o_), rot(nullptr), nxt(nullptr), used(false) {}
4ba
        Q rev() const { return rot->rot; }
00f
        Q next() const { return nxt; }
сЗс
        Q prev() const { return rot->next()->rot; }
188
0d4
        pt dest() const { return rev()->o; }
828 };
91b Q edge(pt from, pt to, int id_from, int id_to) {
        Q e1 = new QuadEdge(id_from, from);
c6e
61b
        Q e2 = new QuadEdge(id_to, to);
8f6
        Q e3 = new QuadEdge;
5ca
        Q e4 = new QuadEdge;
e69
        tie(e1->rot, e2->rot, e3->rot, e4->rot) = \{e3, e4, e2, e1\}:
        tie(e1->nxt, e2->nxt, e3->nxt, e4->nxt) = \{e1, e2, e4, e3\};
f22
1ad
        return e1:
c70 }
d8d void splice(Q a, Q b) {
a6f
        swap(a->nxt->rot->nxt, b->nxt->rot->nxt);
        swap(a->nxt, b->nxt);
da4
a58 }
167 void del_edge(Q& e, Q ne) { // delete e and assign e <- ne
        splice(e, e->prev());
cc0
        splice(e->rev(), e->rev()->prev());
7ea
        delete e->rev()->rot. delete e->rev():
524
        delete e->rot; delete e;
6b2
        e = ne:
18b }
d08 Q conn(Q a, Q b) {
cc5
        Q = edge(a->dest(), b->o, a->rev()->id, b->id);
        splice(e, a->rev()->prev());
f2b
d37
        splice(e->rev(), b):
```

```
6bf
        return e;
f78 }
d64 bool in_c(pt a, pt b, pt c, pt p) { // p ta na circunf. (a, b, c) ?
        _{-}int128 p2 = p*p, A = a*a - p2, B = b*b - p2, C = c*c - p2;
        return sarea2(p, a, b) * C + sarea2(p, b, c) * A + sarea2(p,
   c, a) * B > 0;
b54 }
540 pair < Q, Q > build_tr(vector < pt > & p, int 1, int r) {
09d
        if (r-1+1 <= 3) {
2eb
            Q = edge(p[1], p[1+1], 1, 1+1), b = edge(p[1+1], p[r],
   1+1. r):
912
            if (r-1+1 == 2) return \{a, a->rev()\};
0ec
            splice(a->rev(), b);
сЗс
            11 ar = sarea2(p[1], p[1+1], p[r]);
1af
            Q c = ar ? conn(b, a) : 0;
            if (ar >= 0) return {a, b->rev()};
021
9db
            return {c->rev(), c};
        }
bce
        int m = (1+r)/2:
ee4
328
        auto [la, ra] = build_tr(p, l, m);
        auto [lb, rb] = build_tr(p, m+1, r);
b93
667
        while (true) {
b99
            if (ccw(lb->o, ra->o, ra->dest())) ra = ra->rev()->prev();
458
            else if (ccw(lb->o, ra->o, lb->dest())) lb =
   1b->rev()->next():
f97
            else break:
24a
ca5
        Q b = conn(lb->rev(), ra);
        auto valid = [&](Q e) { return ccw(e->dest(), b->dest(),
713
   b->o); };
        if (ra->o == la->o) la = b->rev():
ee1
        if (1b->o == rb->o) rb = b:
63f
        while (true) {
667
            Q L = b - > rev() - > next();
71e
            if (valid(L)) while (in_c(b->dest(), b->o, L->dest(),
d11
   L->next()->dest()))
1c0
                del_edge(L, L->next());
c76
            Q R = b - > prev();
            if (valid(R)) while (in_c(b->dest(), b->o, R->dest(),
2b0
   R->prev()->dest()))
541
                del_edge(R, R->prev());
            if (!valid(L) and !valid(R)) break;
аЗа
            if (!valid(L) or (valid(R) and in_c(L->dest(), L->o, R->o,
ccd
   R->dest())))
36c
                b = conn(R, b \rightarrow rev());
```

```
666
            else b = conn(b -> rev(), L -> rev());
94d
a2b
        return {la, rb};
689 }
b58 vector < vector < int >> delaunay (vector < pt > v) {
3d2
        int n = v.size();
        auto tmp = v;
397
135
        vector < int > idx(n);
295
        iota(idx.begin(), idx.end(), 0);
        sort(idx.begin(), idx.end(), [&](int 1, int r) { return v[1] <</pre>
fe9
   v[r]; });
       for (int i = 0: i < n: i++) v[i] = tmp[idx[i]]:
5d8
780
        assert(unique(v.begin(), v.end()) == v.end());
4aa
        vector < vector < int >> g(n);
        bool col = true;
4ec
        for (int i = 2; i < n; i++) if (sarea2(v[i], v[i-1], v[i-2]))
   col = false:
        if (col) {
bf5
            for (int i = 1; i < n; i++)</pre>
aa4
                g[idx[i-1]].push_back(idx[i]),
   g[idx[i]].push_back(idx[i-1]);
96b
            return g;
0ae
d36
        Q e = build_tr(v, 0, n-1).first;
113
        vector < Q > edg = {e};
        for (int i = 0; i < edg.size(); e = edg[i++]) {</pre>
5d1
3ed
            for (Q at = e: !at->used: at = at->next()) {
                at->used = true;
60d
cf8
                g[idx[at->id]].push_back(idx[at->rev()->id]);
                edg.push_back(at->rev());
15d
9f2
            }
d19
        }
96b
        return g;
b43 }
      Triangulos em Grafos
// get_triangles(i) encontra todos os triangulos ijk no grafo
// Custo nas arestas
```

```
// get_triangles(i) encontra todos os triangulos ijk no grafo
// Custo nas arestas
// retorna {custo do triangulo, {j, k}}
//
// O(m sqrt(m) log(n)) se chamar para todos os vertices
cod vector<pair<int, int>> g[MAX]; // {para, peso}
```

```
d41 #warning o 'g' deve estar ordenado
9a5 vector<pair<int, pair<int, int>>> get_triangles(int i) {
        vector<pair<int, pair<int, int>>> tri;
771
        for (pair<int, int> j : g[i]) {
b23
            int a = i, b = j.first;
2b3
            if (g[a].size() > g[b].size()) swap(a, b);
6dd
eb0
            for (pair < int, int > c : g[a]) if (c.first != b and c.first
   > j.first) {
525
                auto it = lower_bound(g[b].begin(), g[b].end(),
   make_pair(c.first, -INF));
f55
                if (it == g[b].end() or it->first != c.first) continue;
                tri.push_back({j.second+c.second+it->second, {a == i ?
0aa
   b : a. c.first}}):
b5e
            }
7e1
f5e
        return tri;
036 }
```

# 6 Estruturas

### 6.1 BIT

```
// BIT de soma 1-based, v 0-based
// Para mudar o valor da posicao p para x,
// faca: poe(x - query(p, p), p)
// l_bound(x) retorna o menor p tal que
// query(1, p+1) > x (0 based!)
//
// Complexidades:
// build - O(n)
// poe - O(\log(n))
// query - O(log(n))
// l_bound - O(log(n))
1a8 int n;
7f4 int bit[MAX];
b69 int v[MAX];
0a8 void build() {
b91
        bit[0] = 0:
        for (int i = 1; i <= n; i++) bit[i] = v[i - 1];
33c
78a
        for (int i = 1; i <= n; i++) {</pre>
edf
            int j = i + (i \& -i);
```

```
if (j <= n) bit[j] += bit[i];</pre>
b8a
b39
5dc }
// soma x na posicao p
235 void poe(int x, int p) {
        for (; p <= n; p += p & -p) bit[p] += x;
ec3 }
// soma [1, p]
Obf int pref(int p) {
       int ret = 0;
805
        for (; p; p -= p & -p) ret += bit[p];
edf
        return ret:
601 }
// soma [a, b]
4ea int query(int a, int b) {
        return pref(b) - pref(a - 1);
70c
3dd }
e4a int l_bound(ll x) {
1ba
        int p = 0;
676
        for (int i = MAX2; i+1; i--) if (p + (1 << i) <= n
729
            and bit [p + (1 << i)] <= x) x -= bit <math>[p += (1 << i)];
74e
        return p;
c89 }
6.2 BIT 2D
// BIT de soma, update incrementa posicao
// Tem que construir com um vetor com todos os pontos
// que vc quer um dia atualizar (os pontos q vc vai chamar update)
//
// Complexidades:
// construir - O(n log(n))
// update e query - O(log^2(n))
a6b template < class T = int > struct bit2d {
acf
        vector <T> X;
a84
        vector < vector < T >> Y . t:
709
        int ub(vector<T>& v, T x) {
dde
            return upper_bound(v.begin(), v.end(), x) - v.begin();
9cc
```

bit2d(vector<pair<T, T>> v) {

5cb

```
2e1
            for (auto [x, y] : v) X.push_back(x);
            sort(X.begin(), X.end());
fd4
            X.erase(unique(X.begin(), X.end()), X.end());
1ee
d56
            t.resize(X.size() + 1):
d12
            Y.resize(t.size());
3d0
            sort(v.begin(), v.end(), [](auto a, auto b) {
                return a.second < b.second; });</pre>
e8f
            for (auto [x, y] : v) for (int i = ub(X, x); i < t.size();</pre>
961
   i += i\&-i)
b75
                if (!Y[i].size() or Y[i].back() != y)
   Y[i].push_back(y);
7c7
            for (int i = 0; i < t.size(); i++) t[i].resize(Y[i].size()</pre>
   + 1);
8cc
        void update(T x, T y, T v) {
e78
            for (int i = ub(X, x); i < t.size(); i += i\&-i)
2a9
                for (int j = ub(Y[i], y); j < t[i].size(); j += j&-j)</pre>
cd2
   t[i][j] += v;
        }
533
        T query(T x, T y) {
5d2
966
            T ans = 0:
c54
            for (int i = ub(X, x); i; i -= i\&-i)
                for (int j = ub(Y[i], y); j; j -= j\&-j) ans += t[i][j];
4fb
ba7
            return ans:
62d
46d
        T query(T x1, T y1, T x2, T y2) {
            return query (x2, y2) - query (x2, y1-1) - query (x1-1, y2)
   v2) + querv(x1-1, v1-1);
232
6a7 }:
6.3 BIT com update em range
// Operacoes 0-based
// query(1, r) retorna a soma de v[1..r]
// update(l, r, x) soma x em v[l..r]
//
// Complexidades:
// build - O(n)
// query - O(log(n))
// update - O(log(n))
```

```
e04 namespace bit {
        11 bit[2][MAX+2];
3ba
1a8
        int n;
        void build(int n2, int* v) {
61c
1e3
            n = n2;
535
            for (int i = 1; i <= n; i++)
edd
                bit [1] [min(n+1, i+(i\&-i))] += bit [1][i] += v[i-1];
db0
637
        ll get(int x, int i) {
b73
           11 \text{ ret} = 0;
360
            for (; i; i -= i&-i) ret += bit[x][i];
edf
            return ret:
99c
        }
20 c
        void add(int x, int i, ll val) {
503
            for (; i <= n; i += i&-i) bit[x][i] += val;</pre>
bf6
        }
162
        11 get2(int p) {
c7c
            return get(0, p) * p + get(1, p);
153
02a
        11 query(int 1, int r) {
ff5
            return get2(r+1) - get2(1);
633
089
        void update(int 1, int r, ll x) {
e5f
            add(0, 1+1, x), add(0, r+2, -x);
f58
            add(1, 1+1, -x*1), add(1, r+2, x*(r+1));
e5f
        }
f91 }:
```

## 6.4 BIT-Sort Tree

```
// Tipo uma MergeSort Tree usando Bit
// Apesar da complexidade ser pior, fica melhor na pratica.
//
// query(l, r, k) retorna o numero de elementos menores que k
// no intervalo [l, r]
//
// Usa O(n log(n)) de memoria
//
// Complexidades:
// construir - O(n log^2(n))
// query - O(log^2(n))

6fa template < typename T > struct ms_bit {
1a8    int n;
b2f    vector < vector < T >> bit;
```

```
ms\_bit(vector < T > \& v) : n(v.size()), bit(n+1) {
899
830
            for (int i = 0; i < n; i++)</pre>
d51
                 for (int j = i+1; j \le n; j += j\&-j)
                     bit[j].push_back(v[i]);
dad
            for (int i = 1; i <= n; i++)</pre>
535
eec
                 sort(bit[i].begin(), bit[i].end());
b4d
        }
257
        int p_query(int i, T k) {
7c9
            int ret = 0;
be8
            for (i++: i: i -= i&-i)
1bd
                ret += lower_bound(bit[i].begin(), bit[i].end(), k) -
    bit[i].begin();
edf
            return ret;
6f9
        }
690
        int query(int 1, int r, T k) {
83d
            return p_query(r, k) - p_query(l-1, k);
bcc
        }
8d0 }:
6.5 Convex Hull Trick Dinamico
// para double, use LINF = 1/.0, div(a, b) = a/b
// update(x) atualiza o ponto de intersecao da reta x
// overlap(x) verifica se a reta x sobrepoe a proxima
// add(a, b) adiciona reta da forma ax + b
// query(x) computa maximo de ax + b para entre as retas
// O(log(n)) amortizado por insercao
// O(log(n)) por query
72c struct Line {
073
        mutable 11 a, b, p;
8e3
        bool operator < (const Line& o) const { return a < o.a; }</pre>
abf
        bool operator<(ll x) const { return p < x; }</pre>
469 };
326 struct dynamic_hull : multiset <Line, less <>> {
33a
        11 div(11 a. 11 b) {
            return a / b - ((a ^ b) < 0 and a % b);
a 20
a8a
        }
bbb
        void update(iterator x) {
```

if (next(x) == end()) x->p = LINF;

b2a

```
772
             else if (x->a == next(x)->a) x->p = x->b >= next(x)->b ?
   LINF : -LINF;
             else x \rightarrow p = div(next(x) \rightarrow b - x \rightarrow b, x \rightarrow a - next(x) \rightarrow a);
424
        }
0 c 4
        bool overlap(iterator x) {
71c
f18
             update(x):
             if (next(x) == end()) return 0;
cfa
a4a
             if (x->a == next(x)->a) return x->b >= next(x)->b;
             return x - p >= next(x) - p;
d40
901
        }
176
        void add(ll a. ll b) {
             auto x = insert({a, b, 0});
1c7
4ab
             while (overlap(x)) erase(next(x)), update(x);
            if (x != begin() and !overlap(prev(x))) x = prev(x),
dbc
   update(x);
             while (x != begin() and overlap(prev(x)))
0fc
                 x = prev(x), erase(next(x)), update(x);
4d2
        }
48f
        11 query(11 x) {
4ad
229
             assert(!empty());
7d1
             auto 1 = *lower_bound(x);
d41 #warning cuidado com overflow!
            return 1.a * x + 1.b:
        }
3f5
8f2 }:
6.6 Convex Hull Trick Estatico
```

```
// adds tem que serem feitos em ordem de slope
// queries tem que ser feitas em ordem de x
//
// linear
4b5 struct CHT {
942
        int it;
ac1
       vector < ll > a, b;
       CHT():it(0){}
45e
0bb
       ll eval(int i. ll x){
93d
            return a[i]*x + b[i];
b2a
63a
        bool useless(){
           int sz = a.size();
a20
           int r = sz-1, m = sz-2, 1 = sz-3;
35f
```

```
d41 #warning cuidado com overflow!
 d71
             return (b[1] - b[r])*(a[m] - a[1]) <
413
                 (b[1] - b[m])*(a[r] - a[1]);
 a0c
        }
 bf4
         void add(ll A, ll B){
 7f5
             a.push_back(A); b.push_back(B);
 565
             while (!a.empty()){
 233
                 if ((a.size() < 3) || !useless()) break;</pre>
 ecb
                 a.erase(a.end() - 2);
                 b.erase(b.end() - 2);
 568
b21
             }
 165
         }
 81b
         ll get(ll x){
             it = min(it, int(a.size()) - 1);
 d27
 46a
             while (it+1 < a.size()){</pre>
3c4
                 if (eval(it+1, x) > eval(it, x)) it++;
 f97
                 else break;
 fe9
 420
             return eval(it, x);
 88a
        }
 303 };
6.7 DSU
// Une dois conjuntos e acha a qual conjunto um elemento pertence por
    seu id
// find e unite: O(a(n)) \sim = O(1) amortizado
8d3 struct dsu {
         vector<int> id, sz;
 b33
         dsu(int n) : id(n), sz(n, 1) { iota(id.begin(), id.end(), 0); }
         int find(int a) { return a == id[a] ? a : id[a] = find(id[a]);
 0cf
  }
 440
         void unite(int a, int b) {
 605
             a = find(a), b = find(b);
 d54
             if (a == b) return;
956
             if (sz[a] < sz[b]) swap(a, b);
             sz[a] += sz[b], id[b] = a;
ea7
8e1 };
// DSU de bipartido
```

```
// Une dois vertices e acha a qual componente um vertice pertence
// Informa se a componente de um vertice e bipartida
// find e unite: O(log(n))
8d3 struct dsu {
        vector < int > id, sz, bip, c;
5b4
        dsu(int n) : id(n), sz(n, 1), bip(n, 1), c(n) {
db8
            iota(id.begin(), id.end(), 0);
f25
       }
ef0
        int find(int a) { return a == id[a] ? a : find(id[a]); }
        int color(int a) { return a == id[a] ? c[a] : c[a] ^
f30
   color(id[a]); }
440
        void unite(int a, int b) {
            bool change = color(a) == color(b);
263
605
            a = find(a), b = find(b);
a89
            if (a == b) {
                if (change) bip[a] = 0;
4ed
505
                return;
32d
           }
956
            if (sz[a] < sz[b]) swap(a, b);
efe
            if (change) c[b] = 1;
            sz[a] += sz[b], id[b] = a, bip[a] &= bip[b];
2cd
       }
22b
118 };
// DSU Persistente
//
// Persistencia parcial, ou seja, tem que ir
// incrementando o 't' no une
//
// find e unite: O(log(n))
8d3 struct dsu {
33c
        vector < int > id, sz, ti;
        dsu(int n) : id(n), sz(n, 1), ti(n, -INF) {
733
db8
            iota(id.begin(), id.end(), 0);
       }
aad
5e6
       int find(int a, int t) {
```

```
6ba
            if (id[a] == a or ti[a] > t) return a;
ea5
            return find(id[a], t);
6cb
       }
        void unite(int a, int b, int t) {
fa0
            a = find(a, t), b = find(b, t);
84f
d54
            if (a == b) return:
956
            if (sz[a] < sz[b]) swap(a, b);</pre>
35d
            sz[a] += sz[b], id[b] = a, ti[b] = t;
513
        }
6c6 };
// DSU com rollback
//
// checkpoint(): salva o estado atual de todas as variaveis
// rollback(): retorna para o valor das variaveis para
// o ultimo checkpoint
//
// Sempre que uma variavel muda de valor, adiciona na stack
// find e unite: O(log(n))
// checkpoint: 0(1)
// rollback: O(m) em que m e o numero de vezes que alguma
// variavel mudou de valor desde o ultimo checkpoint
8d3 struct dsu {
825
        vector < int > id, sz;
27 c
        stack<stack<pair<int&, int>>> st;
98d
        dsu(int n) : id(n), sz(n, 1) {
1cc
            iota(id.begin(), id.end(), 0), st.emplace();
8cd
        }
bdf
        void save(int &x) { st.top().emplace(x, x); }
30d
        void checkpoint() { st.emplace(); }
5cf
        void rollback() {
ba9
            while(st.top().size()) {
6bf
                auto [end, val] = st.top().top(); st.top().pop();
149
                end = val:
f9a
25a
            st.pop();
        }
3c6
        int find(int a) { return a == id[a] ? a : find(id[a]); }
ef0
```

```
440     void unite(int a, int b) {
605          a = find(a), b = find(b);
d54          if (a == b) return;
956          if (sz[a] < sz[b]) swap(a, b);
803          save(sz[a]), save(id[b]);
6d0          sz[a] += sz[b], id[b] = a;
1b9     }
c6e };</pre>
```

### 6.8 Li-Chao Tree

```
// Adiciona retas (ax+b), e computa o minimo entre as retas
// em um dado 'x'
// Cuidado com overflow!
// Se tiver overflow, tenta comprimir o 'x' ou usar
// convex hull trick
//
// O(log(MA-MI)), O(n) de memoria
5b0 template < 11 MI = 11(-1e9), 11 MA = 11(1e9) > struct lichao {
        struct line {
b3a
12d
            ll a, b;
cef
            array<int, 2> ch;
fdf
            line(ll a_- = 0, ll b_- = LINF):
423
                a(a_{-}), b(b_{-}), ch(\{-1, -1\})  {}
888
            11 operator ()(11 x) { return a*x + b; }
d1d
        };
17b
        vector<line> ln;
df8
        int ch(int p, int d) {
e85
            if (ln[p].ch[d] == -1) {
9af
                ln[p].ch[d] = ln.size();
cdc
                ln.emplace_back();
bc1
            }
ef2
            return ln[p].ch[d];
86a
021
        lichao() { ln.emplace_back(); }
        void add(line s, ll l=MI, ll r=MA, int p=0) {
c33
            11 m = (1+r)/2;
3e3
911
            bool L = s(1) < ln[p](1):
            bool M = s(m) < ln[p](m);
d37
03b
            bool R = s(r) < ln[p](r);
825
            if (M) swap(ln[p], s), swap(ln[p].ch, s.ch);
cac
            if (s.b == LINF) return;
f6d
            if (L != M) add(s, l, m-1, ch(p, 0));
```

```
898
            else if (R != M) add(s, m+1, r, ch(p, 1));
76e
        }
092
        11 query(int x, 11 1=MI, 11 r=MA, int p=0) {
11b
            11 m = (1+r)/2, ret = ln[p](x);
9db
            if (ret == LINF) return ret;
529
            if (x < m) return min(ret, query(x, 1, m-1, ch(p, 0)));
81a
            return min(ret, query(x, m+1, r, ch(p, 1)));
fba
        }
59b };
6.9 Li-Chao Tree - Lazv
// Sendo N = MA-MI:
// insert(\{a, b\}) minimiza tudo com ax+b - 0(\log N)
// insert(\{a, b\}, 1, r) minimiza com ax+b no range [1, r] - 0(\log^2 N)
// shift({a. b}) soma ax+b em tudo - O(1)
// shift({a, b}, l, r) soma ax+b no range [l, r] - O(log^2 N)
// query(x) retorna o valor da posicao x - O(\log N)
// No inicio eh tudo LINF, se inserir {0, 0} fica tudo 0
// O(n log N) de memoria ; O(n) de memoria se nao usar as operacoes de
    range
41c template < int MI = int(-1e9), int MA = int(1e9) > struct lichao {
b3a
        struct line {
12d
            ll a. b:
            11 la, lb; // lazy
158
cef
            array < int, 2 > ch;
fdf
            line(ll a_- = 0, ll b_- = LINF):
b09
                a(a_{-}), b(b_{-}), la(0), lb(0), ch(\{-1, -1\})  {}
888
            11 operator ()(11 x) { return a*x + b; }
92e
        }:
17b
        vector<line> ln;
df8
        int ch(int p, int d) {
e85
            if (ln[p].ch[d] == -1) {
9af
                ln[p].ch[d] = ln.size();
cdc
                ln.emplace_back();
bc1
ef2
            return ln[p].ch[d];
86a
        }
021
        lichao() { ln.emplace_back(); }
ceb
        void prop(int p, int 1, int r) {
ff8
            if (ln[p].la == 0 and ln[p].lb == 0) return;
```

```
1d3
            ln[p].a += ln[p].la, ln[p].b += ln[p].lb;
579
            if (1 != r) {
b9e
                int pl = ch(p, 0), pr = ch(p, 1);
0d7
                 ln[pl].la += ln[p].la, ln[pl].lb += ln[p].lb;
fa8
                 ln[pr].la += ln[p].la, ln[pr].lb += ln[p].lb;
77f
01e
            ln[p].la = ln[p].lb = 0:
89b
        }
        11 query(int x, int p=0, int l=MI, int r=MA) {
c06
6b9
            prop(p, 1, r);
6f3
            ll ret = ln[p](x);
33b
            if (\ln \lceil p \rceil, ch \lceil 0 \rceil == -1 \text{ and } \ln \lceil p \rceil, ch \lceil 1 \rceil == -1) return ret:
            int m = 1 + (r-1)/2;
90d
da9
            if (x \le m) return min(ret, query(x, ch(p, 0), 1, m));
c55
            return min(ret, query(x, ch(p, 1), m+1, r));
953
        }
        void push(line s, int p, int l, int r) {
5df
6b9
            prop(p, 1, r);
90d
            int m = 1 + (r-1)/2;
911
            bool L = s(1) < ln[p](1);
            bool M = s(m) < ln[p](m);
d37
03b
            bool R = s(r) < ln[p](r);
c3f
            if (M) swap(ln[p].a, s.a), swap(ln[p].b, s.b);
cac
            if (s.b == LINF) return;
c49
            if (L != M) push(s, ch(p, 0), 1, m);
29e
            else if (R != M) push(s, ch(p, 1), m+1, r);
ceb
        void insert(line s, int a=MI, int b=MA, int p=0, int l=MI, int
   r=MA) {
6b9
            prop(p, 1, r);
2d3
            if (a \le 1 \text{ and } r \le b) \text{ return push}(s, p, 1, r);
            if (b < 1 or r < a) return;</pre>
1dd
            int m = 1 + (r-1)/2;
90d
f1e
            insert(s, a, b, ch(p, 0), 1, m);
952
            insert(s, a, b, ch(p, 1), m+1, r);
375
        }
97a
        void shift(line s, int a=MI, int b=MA, int p=0, int l=MI, int
   r=MA) {
6b9
            prop(p, 1, r);
            int m = 1 + (r-1)/2:
90d
            if (a <= 1 and r <= b) {</pre>
9a3
ada
                 ln[p].la += s.a, ln[p].lb += s.b;
505
                 return:
570
            }
```

```
1dd
            if (b < 1 or r < a) return;
fdd
            if (ln[p].b != LINF) {
751
                push(ln[p], ch(p, 0), 1, m);
ade
                push(ln[p], ch(p, 1), m+1, r);
c2f
                ln[p].a = 0, ln[p].b = LINF;
199
a04
            shift(s, a, b, ch(p, 0), 1, m);
e7d
            shift(s, a, b, ch(p, 1), m+1, r);
d43
       }
285 };
```

# 6.10 MergeSort Tree

```
// Se for construida sobre um array:
//
        count(i, j, a, b) retorna quantos
//
        elementos de v[i..j] pertencem a [a, b]
//
        report(i, j, a, b) retorna os indices dos
//
        elementos de v[i..j] que pertencem a [a, b]
        retorna o vetor ordenado
// Se for construida sobre pontos (x, y):
        count(x1, x2, y1, y2) retorna quantos pontos
//
//
        pertencem ao retangulo (x1, y1), (x2, y2)
//
        report(x1, x2, y1, y2) retorna os indices dos pontos que
//
        pertencem ao retangulo (x1, y1), (x2, y2)
//
        retorna os pontos ordenados lexicograficamente
//
        (assume x1 \le x2, v1 \le v2)
//
// kth(y1, y2, k) retorna o indice do ponto com k-esimo menor
// x dentre os pontos que possuem y em [y1, y2] (0 based)
// Se quiser usar para achar k-esimo valor em range, construir
// com ms_tree t(v, true), e chamar kth(1, r, k)
//
// Usa O(n log(n)) de memoria
// Complexidades:
// construir - O(n log(n))
// count - O(log(n))
// report - O(\log(n) + k) para k indices retornados
// kth - O(log(n))
c6c template <typename T = int> struct ms_tree {
        vector<tuple<T, T, int>> v;
6f7
1a8
5ee
        vector < vector < tuple < T, T, int >>> t; // {y, idx, left}
6ae
        vector <T> vv;
```

```
78c
        ms_tree(vector<pair<T, T>>& vv) : n(vv.size()), t(4*n), vv(n) {
             for (int i = 0; i < n; i++) v.push_back({vv[i].first,</pre>
e80
   vv[i].second, i});
            sort(v.begin(), v.end());
fca
            build(1, 0, n-1);
224
            for (int i = 0; i < n; i++) vy[i] = get<0>(t[1][i+1]);
01a
45e
        }
        ms_tree(vector<T>& vv, bool inv = false) { // inv: inverte
   indice e valor
            vector < pair < T, T >> v2;
8e8
            for (int i = 0; i < vv.size(); i++)</pre>
e1e
                 inv ? v2.push_back({vv[i], i}) : v2.push_back({i,
   vv[i]}):
             *this = ms_tree(v2);
cca
f23
        }
        void build(int p, int l, int r) {
2c6
             t[p].push_back({get<0>(v[1]), get<0>(v[r]), 0}); //
   \{\min x. \max x. 0\}
            if (1 == r) return t[p].push_back({get<1>(v[1]),
5c8
   get <2>(v[1]), 0});
            int m = (1+r)/2;
ee4
            build(2*p, 1, m), build(2*p+1, m+1, r);
bd9
32d
            int L = 0, R = 0;
a 0.3
             while (t[p].size() \le r-l+1) {
68e
                 int left = get <2>(t[p].back());
                 if (L > m-1 \text{ or } (R+m+1 \le r \text{ and } t[2*p+1][1+R] \le r
   t[2*p][1+L])) {
8cf
                     t[p].push_back(t[2*p+1][1 + R++]);
                     get < 2 > (t[p].back()) = left;
da0
                     continue;
5e2
ce0
249
                 t[p].push_back(t[2*p][1 + L++]);
                 get < 2 > (t[p].back()) = left + 1;
339
            }
208
2eb
        }
        int get_1(T y) { return lower_bound(vy.begin(), vy.end(), y) -
   vy.begin(); }
        int get_r(T y) { return upper_bound(vy.begin(), vy.end(), y) -
   vy.begin(); }
        int count(T x1, T x2, T y1, T y2) {
f62
            function<int(int, int, int)> dfs = [&](int p, int 1, int
902
   r) {
                 if (1 == r \text{ or } x2 < get<0>(t[p][0]) \text{ or } get<1>(t[p][0])
7 c 6
   < x1) return 0:
```

```
2bb
                 if (x1 \le get<0>(t[p][0]) and get<1>(t[p][0]) \le x2)
   return r-1:
                 int nl = get<2>(t[p][1]), nr = get<2>(t[p][r]);
784
                 return dfs(2*p, nl, nr) + dfs(2*p+1, l-nl, r-nr);
eb6
122
            }:
7cb
            return dfs(1, get_l(y1), get_r(y2));
f65
        }
002
        vector<int> report(T x1, T x2, T y1, T y2) {
4b8
            vector < int > ret;
85 e
            function < void(int, int, int) > dfs = [&](int p, int 1, int
   r) {
                 if (1 == r \text{ or } x2 < get<0>(t[p][0]) \text{ or } get<1>(t[p][0])
882
   < x1) return:
8da
                 if (x1 \le get<0>(t[p][0]) and get<1>(t[p][0]) \le x2) {
e00
                     for (int i = 1; i < r; i++)</pre>
   ret.push_back(get<1>(t[p][i+1]));
505
                     return;
067
784
                 int nl = get < 2 > (t[p][1]), nr = get < 2 > (t[p][r]);
                 dfs(2*p, nl, nr), dfs(2*p+1, l-nl, r-nr);
194
12b
            };
8ad
            dfs(1, get_l(y1), get_r(y2));
edf
            return ret;
668
        }
985
        int kth(T y1, T y2, int k) {
902
            function < int(int, int, int) > dfs = [&](int p, int 1, int
   r) {
150
                 if (k >= r-1) {
941
                     k -= r-1:
                     return -1;
daa
                 }
b8d
8da
                 if (r-l == 1) return get<1>(t[p][1+1]);
784
                 int n1 = get<2>(t[p][1]), nr = get<2>(t[p][r]);
072
                 int left = dfs(2*p. nl. nr);
                 if (left != -1) return left;
3b6
04d
                 return dfs(2*p+1, l-nl, r-nr);
a1b
            };
7cb
             return dfs(1, get_l(v1), get_r(v2));
635
1ce };
```

# 6.11 Min queue - deque

```
// Tudo O(1) amortizado
| 1dc template < class T> struct minqueue {
```

```
2d8
        deque<pair<T, int>> q;
3fc
        void push(T x) {
            int ct = 1;
56e
953
            while (q.size() and x < q.front().first)</pre>
                ct += q.front().second, q.pop_front();
75f
987
            q.emplace_front(x, ct);
        }
e8d
42d
        void pop() {
aa2
            if (q.back().second > 1) q.back().second--;
c51
            else q.pop_back();
5fd
ea6
        T min() { return q.back().first; }
c13 };
6.12 Min queue - stack
// Tudo O(1) amortizado
```

```
557 template < class T > struct minstack {
81f
        stack<pair<T, T>> s;
3fc
        void push(T x) {
12b
            if (!s.size()) s.push({x, x});
9d9
            else s.emplace(x, std::min(s.top().second, x));
f8d
4f0
       T top() { return s.top().first; }
94a
        T pop() {
            T ans = s.top().first;
1f2
2eb
            s.pop();
ba7
            return ans;
013
614
        int size() { return s.size(); }
        T min() { return s.top().second; }
13b
4c0 };
1dc template < class T> struct minqueue {
cdc
        minstack <T> s1, s2;
        void push(T x) { s1.push(x); }
c96
        void move() {
            if (s2.size()) return;
d92
            while (s1.size()) {
7ae
                T x = s1.pop();
489
                s2.push(x);
656
            }
```

```
ef1
787
        T front() { return move(), s2.top(); }
        T pop() { return move(), s2.pop(); }
23a
7f3
        int size() { return s1.size()+s2.size(); }
19c
        T min() {
            if (!s1.size()) return s2.min();
cd6
58e
            else if (!s2.size()) return s1.min();
31d
            return std::min(s1.min(), s2.min());
9c7
        }
6d3 };
6.13 Order Statistic Set
// Funciona do C++11 pra cima
774 #include <ext/pb_ds/assoc_container.hpp>
30f #include <ext/pb_ds/tree_policy.hpp>
0d7 using namespace __gnu_pbds;
4fc template <class T>
def
        using ord_set = tree<T, null_type, less<T>, rb_tree_tag,
3a1
        tree_order_statistics_node_update>;
// para declarar:
// ord set <int> s:
// coisas do set normal funcionam:
// for (auto i : s) cout << i << endl:
// cout << s.size() << endl;</pre>
// k-esimo maior elemento O(log|s|):
// k=0: menor elemento
// cout << *s.find_by_order(k) << endl;</pre>
// quantos sao menores do que k O(log|s|):
// cout << s.order_of_key(k) << endl;</pre>
// Para fazer um multiset, tem que
// usar ord_set<pair<int, int>> com o
// segundo parametro sendo algo para diferenciar
// os ementos iguais.
// s.order_of_kev({k, -INF}) vai retornar o
// numero de elementos < k
6.14 Priority Queue DS
// Mantem updates aplicados em uma estrutura de dados
```

// que permita rollback e nao seja amortizada.

#### 91

```
// Cada update possui uma prioridade,
// sendo possivel remover o update com maior prioridade.
// Os updates devem ser comutativos, ou seja, o estado
// da estrutura deve ser o mesmo independente da ordem
// que eles sejam aplicados.
//
// Complexidades:
// update - O(log(n) + T(n))
// query - T(n)
// pop - O(log(n) * T(n)) amortizado
//
// onde T(n) eh a complexidade do update
// assumes all priorities are distinct
945 template < typename DS, typename UPD > struct priority_queue_ds {
df4
a7e
        vector < tuple < UPD, int, int >> upd; // {u, p, idx_in_pos}
866
        set < pair < int , int >> st;
        vector < int > pos;
927
cf0
        priority_queue_ds(int n) : D(n) {}
6af
        void update(UPD u, int p) {
9ab
            D.update(u);
d07
            st.emplace(p, pos.size());
6ca
            upd.emplace_back(u, p, pos.size());
e3d
            pos.push_back(upd.size() - 1);
6af
        }
427
        int query(int a) {
            return D.find(a);
aa3
2d3
        }
42d
        void pop() {
25f
            int k = 1, min_p; // k = number of pops we will do
43e
            vector < tuple < UPD, int, int >> small, big;
639
            auto it = st.end();
231
            for (int qt = 0; qt++ < (k+1)/2;) {
049
                min_p = it->first;
3ab
                int i = pos[it->second];
80f
                if (qt > 1) big.push_back(upd[i]);
e82
84b
                k = max<int>(k, upd.size() - i);
            }
b9a
b3d
            for (int i = 0; i < k; i++) {</pre>
a62
                D.rollback():
```

```
8b6
                auto [u, p, idx] = upd.rbegin()[i];
864
                if (p < min_p) small.emplace_back(u, p, idx);</pre>
588
            }
            st.erase(prev(st.end()));
23e
623
            upd.erase(upd.end() - k, upd.end());
            small.insert(small.end(), big.rbegin(), big.rend());
a25
06f
            for (auto [u, p, idx] : small) {
9ab
                D.update(u);
c8e
                upd.emplace_back(u, p, idx);
a7d
                pos[idx] = upd.size() - 1;
ec7
            }
bd1
        }
54a };
6.15 Range color
// update(l, r, c) colore o range [l, r] com a cor c,
// e retorna os ranges que foram coloridos {1, r, cor}
// query(i) returna a cor da posicao i
//
// Complexidades (para q operacoes):
// update - O(log(q)) amortizado
// query - O(log(q))
df6 template < typename T > struct color {
f0c
        set < tuple < int , int , T >> se;
071
        vector<tuple<int, int, T>> update(int 1, int r, T val) {
9 c 4
            auto it = se.upper_bound({r, INF, val});
753
            if (it != se.begin() and get<1>(*prev(it)) > r) {
e91
                auto [L, R, V] = *--it;
3f0
                se.erase(it);
bfd
                se.emplace(L, r, V), se.emplace(r+1, R, V);
683
d9e
            it = se.lower_bound({1, -INF, val});
516
            if (it != se.begin() and get<1>(*prev(it)) >= 1) {
                auto [L, R, V] = *--it:
e91
3f0
                se.erase(it);
75a
                se.emplace(L, 1-1, V), it = se.emplace(1, R, V).first;
b65
            }
d7b
            vector<tuple<int, int, T>> ret;
```

for (; it != se.end() and get<0>(\*it) <= r; it =</pre>

ret.push\_back(\*it);

7 a 1

8c0

se.erase(it))

```
b4a
            se.emplace(1, r, val);
                                                                            62a
edf
            return ret;
                                                                            093
                                                                                    T query(int 1, int r) { return v[index_query(1, r)]; }
b6c
                                                                            bab };
ff9
        T query(int i) {
c31
            auto it = se.upper_bound({i, INF, T()});
            if (it == se.begin() or get<1>(*--it) < i) return -1; //</pre>
8e7
                                                                            6.17 SegTreap
   nao tem
53d
            return get <2>(*it);
                                                                            // Muda uma posicao do plano, e faz query de operacao
daf
        }
                                                                            // associativa e comutativa em retangulo
9e9 };
                                                                            // Mudar ZERO e op
                                                                            // Esparso nas duas coordenadas, inicialmente eh tudo ZERO
6.16 RMQ < O(n), O(1) > - min queue
                                                                            // Para query com distancia de manhattan <= d, faca
                                                                            // nx = x+y, ny = x-y
                                                                            // Update em (nx, ny), query em ((nx-d, ny-d), (nx+d, ny+d))
// O(n) pra buildar, query O(1)
                                                                            //
// Se tiver varios minimos, retorna
// o de menor indice
                                                                            // Valores no X tem que ser de O ateh NX
                                                                            // Para q operacoes, usa O(q log(NX)) de memoria, e as
1a5 template < typename T > struct rmq {
                                                                            // operacoes custa O(log(q) log(NX))
517
        vector <T> v;
fcc
        int n: static const int b = 30:
                                                                            55b const int ZERO = INF:
70e
        vector < int > mask, t;
                                                                            560 const int op(int 1, int r) { return min(1, r); }
183
        int op(int x, int y) { return v[x] \le v[y] ? x : y; }
                                                                            878 mt19937 rng((int)
        int msb(int x) { return __builtin_clz(1)-__builtin_clz(x); }
                                                                                chrono::steady_clock::now().time_since_epoch().count());
ee1
        int small(int r, int sz = b) { return
   r-msb(mask[r]&((1<<sz)-1)); }
                                                                            aa1 template < typename T > struct treap {
                                                                                    struct node {
6ad
        rmq() {}
                                                                            3c9
43c
        rmq(const \ vector < T > \& \ v_) : v(v_), n(v.size()), mask(n), t(n) 
                                                                            b19
                                                                                        node *1, *r;
            for (int i = 0, at = 0; i < n; mask[i++] = at |= 1) {</pre>
2e5
                                                                            ee1
                                                                                        int p;
                at = (at << 1) &((1 << b) -1);
                                                                            850
a61
                                                                                        pair < 11, 11 > idx; // {y, x}
                while (at and op(i-msb(at&-at), i) == i) at ^= at&-at;
                                                                            36d
                                                                                        T val, mi;
c00
c2f
            }
                                                                            bc2
                                                                                         node(11 x, 11 y, T val_) : 1(NULL), r(NULL), p(rng()),
ea4
            for (int i = 0; i < n/b; i++) t[i] = small(b*i+b-1);
                                                                            1 b 5
                                                                                             idx(pair(y, x)), val(val_), mi(val) {}
            for (int j = 1; (1<<j) <= n/b; j++) for (int i = 0;
                                                                            01e
                                                                                        void update() {
   i+(1<< j) <= n/b; i++)
                                                                                             mi = val:
                                                                            d6e
                t[n/b*j+i] = op(t[n/b*(j-1)+i],
                                                                            182
                                                                                             if (1) mi = op(mi, 1->mi);
   t[n/b*(j-1)+i+(1<<(j-1))]);
                                                                            b68
                                                                                            if (r) mi = op(mi, r->mi);
41a
        }
                                                                            282
                                                                                        }
e34
        int index_query(int 1, int r) {
                                                                            6e1
                                                                                    };
            if (r-l+1 <= b) return small(r, r-l+1);</pre>
27b
e80
            int x = 1/b+1, y = r/b-1;
                                                                            bb7
                                                                                    node* root:
fd3
            if (x > y) return op(small(1+b-1), small(r));
            int j = msb(y-x+1);
                                                                            84b
a4e
                                                                                    treap() { root = NULL; }
            int ans = op(small(1+b-1), op(t[n/b*j+x],
                                                                            cec
                                                                                    \simtreap() {
   t[n/b*j+y-(1<<j)+1]));
                                                                            609
                                                                                        vector < node *> q = {root};
be6
            return op(ans, small(r));
                                                                            402
                                                                                        while (q.size()) {
```

```
e5d
                node* x = q.back(); q.pop_back();
ee9
                if (!x) continue;
1c7
                q.push_back(x->1), q.push_back(x->r);
bf0
                delete x;
653
            }
50e
        treap(treap&& t) : treap() { swap(root, t.root); }
225
bcf
        void join(node* 1, node* r, node*& i) { // assume que 1 < r</pre>
986
            if (!1 or !r) return void(i = 1 ? 1 : r);
80e
            if (1->p > r->p) join(1->r, r, 1->r), i = 1;
fa0
            else join(1, r->1, r->1), i = r;
bda
            i->update():
671
        }
        void split(node* i, node*& 1, node*& r, pair<11, 11> idx) {
c82
26a
            if (!i) return void(r = 1 = NULL);
13c
            if (i->idx < idx) split(i->r, i->r, r, idx), l = i;
d26
            else split(i \rightarrow 1, l, i \rightarrow 1, idx), r = i;
            i->update();
bda
143
        }
d3b
        void update(ll x, ll y, T v) {
df9
            node *L, *M, *R;
            split(root, M, R, pair(y, x+1)), split(M, L, M, pair(y,
8b2
   x));
1e4
            if (M) M->val = M->mi = v;
9e5
            else M = new node(x, y, v);
69d
            join(L, M, M), join(M, R, root);
58e
91b
        T query(ll ly, ll ry) {
df9
            node *L, *M, *R;
            split(root, M, R, pair(ry, LINF)), split(M, L, M, pair(ly,
1c0
   0));
            T \text{ ret} = M ? M -> mi : ZERO:
0f7
            join(L, M, M), join(M, R, root);
69d
            return ret;
edf
        }
1ae
bdf };
46a template < typename T > struct segtreap {
c4f
        vector < treap < T >> seg;
6e7
        vector < int > ch[2];
e4e
        ll NX;
        segtreap(ll NX_) : seg(1), NX(NX_) { ch[0].push_back(-1).}
253
   ch[1].push_back(-1); }
a71
        int get_ch(int i, int d){
```

```
if (ch[d][i] == -1) {
e51
2d6
                ch[d][i] = seg.size();
23e
                seg.emplace_back();
842
                ch[0].push_back(-1), ch[1].push_back(-1);
3e1
968
            return ch[d][i];
        }
bb6
10c
        T query(11 lx, 11 rx, 11 ly, 11 ry, int p, 11 l, 11 r) {
            if (rx < 1 or r < 1x) return ZERO;</pre>
003
fOf
            if (lx <= l and r <= rx) return seg[p].query(ly, ry);</pre>
e6a
            11 m = 1 + (r-1)/2:
            return op(query(lx, rx, ly, ry, get_ch(p, 0), l, m),
354
                    query(lx, rx, ly, ry, get_ch(p, 1), m+1, r));
060
a5e
        }
f48
        T query(ll lx, ll rx, ll ly, ll ry) { return query(lx, rx, ly,
   ry, 0, 0, NX); }
249
        void update(ll x, ll y, T val, int p, ll l, ll r) {
73c
            if (1 == r) return seg[p].update(x, y, val);
e6a
            11 m = 1 + (r-1)/2:
            if (x <= m) update(x, y, val, get_ch(p, 0), 1, m);</pre>
cc5
5a2
            else update(x, y, val, get_ch(p, 1), m+1, r);
980
            seg[p].update(x, y, val);
cc2
517
        void update(11 x, 11 y, T val) { update(x, y, val, 0, 0, NX); }
40a }:
```

# 6.18 SegTree

```
d22
        ll build(int p=1, int l=0, int r=n-1) {
3c7
            lazy[p] = 0;
6cd
            if (1 == r) return seg[p] = v[1];
            int m = (1+r)/2;
ee4
193
            return seg[p] = build(2*p, 1, m) + build(2*p+1, m+1, r);
c71
0d8
        void build(int n2. int* v2) {
680
            n = n2, v = v2;
6f2
            build();
acb
ceb
        void prop(int p, int l, int r) {
cdf
            seg[p] += lazy[p]*(r-l+1);
2c9
            if (1 != r) lazy[2*p] += lazy[p], lazy[2*p+1] += lazy[p];
            lazv[p] = 0;
3c7
c10
2c3
        11 query(int a, int b, int p=1, int l=0, int r=n-1) {
6b9
            prop(p, 1, r);
527
            if (a <= 1 and r <= b) return seg[p];</pre>
            if (b < 1 or r < a) return 0;</pre>
786
            int m = (1+r)/2:
ee4
            return query(a, b, 2*p, 1, m) + query(a, b, 2*p+1, m+1, r);
b1f
4c5
        ll update(int a, int b, int x, int p=1, int l=0, int r=n-1) {
cfb
6b9
            prop(p, 1, r);
9a3
            if (a <= 1 and r <= b) {</pre>
b94
                lazy[p] += x;
6b9
                prop(p, 1, r);
534
               return seg[p];
821
            }
e9f
            if (b < 1 or r < a) return seg[p];</pre>
            int m = (1+r)/2:
ee4
fdb
            return seg[p] = update(a, b, x, 2*p, 1, m) +
                update(a, b, x, 2*p+1, m+1, r);
7fd
       }
75c
0af };
// Se tiver uma seg de max, da pra descobrir em O(log(n))
// o primeiro e ultimo elemento >= val numa range:
// primeira posicao >= val em [a, b] (ou -1 se nao tem)
119 int get_left(int a, int b, int val, int p=1, int l=0, int r=n-1) {
6b9
        prop(p, 1, r);
       if (b < l or r < a or seg[p] < val) return -1;</pre>
f38
205
       if (r == 1) return 1;
       int m = (1+r)/2:
ee4
753
       int x = get_left(a, b, val, 2*p, 1, m);
```

```
50e
        if (x != -1) return x;
сЗс
        return get_left(a, b, val, 2*p+1, m+1, r);
68c }
// ultima posicao >= val em [a, b] (ou -1 se nao tem)
992 int get_right(int a, int b, int val, int p=1, int l=0, int r=n-1) {
6b9
        prop(p, 1, r):
f38
        if (b < l or r < a or seg[p] < val) return -1;</pre>
205
        if (r == 1) return 1;
        int m = (1+r)/2:
ee4
1b1
        int x = get_right(a, b, val, 2*p+1, m+1, r);
50e
        if (x != -1) return x;
6a7
        return get_right(a, b, val, 2*p, 1, m);
1b7 }
// Se tiver uma seg de soma sobre um array nao negativo v, da pra
// descobrir em O(\log(n)) o maior j tal que v[i]+v[i+1]+...+v[j-1] <
   val
6a9 int lower_bound(int i, ll& val, int p, int l, int r) {
6b9
        prop(p, 1, r);
        if (r < i) return n;</pre>
6e8
        if (i <= l and seg[p] < val) {</pre>
b5d
bff
            val -= seg[p];
041
            return n;
634
Зсе
        if (1 == r) return 1:
ee4
        int m = (1+r)/2:
        int x = lower_bound(i, val, 2*p, 1, m);
514
ee0
        if (x != n) return x;
8b9
        return lower_bound(i, val, 2*p+1, m+1, r);
2b8 }
6.19 SegTree 2D Iterativa
```

```
// Consultas O-based
// Um valor inicial em (x, y) deve ser colocado em seg[x+n][y+n]
// Query: soma do retangulo ((x1, y1), (x2, y2))
// Update: muda o valor da posicao (x, y) para val
// Nao pergunte como que essa coisa funciona
// Para query com distancia de manhattan <= d, faca
// nx = x+v, nv = x-v
// Update em (nx, ny), query em ((nx-d, ny-d), (nx+d, ny+d))
//
// Se for de min/max, pode tirar os if's da 'query', e fazer
// sempre as 4 operacoes. Fica mais rapido
```

```
// Complexidades:
                                                                               // build - O(n)
// Complexidades:
// build - O(n^2)
                                                                               // query - 0(log(n))
// query - O(log^2(n))
                                                                               // update - O(log^2 (n)) amortizado
// update - 0(log^2(n))
                                                                               // (se nao usar updatesum, fica log(n) amortizado)
731 int seg[2*MAX][2*MAX], n;
                                                                               7c6 #define f first
                                                                               Oab #define s second
0a8 void build() {
        for (int x = 2*n; x; x--) for (int y = 2*n; y; y--) {
                                                                               f39 namespace beats {
             if (x < n) seg[x][y] = seg[2*x][y] + seg[2*x+1][y];
                                                                               3c9
                                                                                        struct node {
c81
fe9
             if (y < n) seg[x][y] = seg[x][2*y] + seg[x][2*y+1];
                                                                               526
                                                                                            int tam:
d51
                                                                               125
                                                                                            ll sum, lazy; // lazy pra soma
499 }
                                                                               4f3
                                                                                            ll mi1, mi2, mi; // mi = #mi1
                                                                                c61
                                                                                            ll ma1, ma2, ma; // ma = #ma1
251 int query(int x1, int y1, int x2, int y2) {
827
        int ret = 0, y3 = y1 + n, y4 = y2 + n;
                                                                               426
                                                                                            node(11 x = 0) {
        for (x1 += n, x2 += n; x1 <= x2; ++x1 /= 2, --x2 /= 2)
83e
                                                                               ba6
                                                                                                 sum = mi1 = ma1 = x;
0f2
             for (y1 = y3, y2 = y4; y1 \le y2; ++y1 /= 2, --y2 /= 2) {
                                                                               b29
                                                                                                 mi2 = LINF, ma2 = -LINF;
554
                                                                                                 mi = ma = tam = 1:
                 if (x1\%2 == 1 \text{ and } y1\%2 == 1) \text{ ret } += \text{seg}[x1][y1];
                                                                               62c
6b0
                 if (x1\%2 == 1 \text{ and } y2\%2 == 0) \text{ ret } += \text{seg}[x1][y2];
                                                                               c60
                                                                                                 lazy = 0;
c01
                 if (x2\%2 == 0 \text{ and } y1\%2 == 1) \text{ ret } += \text{seg}[x2][y1];
                                                                               b00
                                                                               770
5d4
                 if (x2\%2 == 0 \text{ and } y2\%2 == 0) \text{ ret } += \text{seg}[x2][y2];
                                                                                            node(const node& 1, const node& r) {
2d0
             }
                                                                               a95
                                                                                                 sum = 1.sum + r.sum, tam = 1.tam + r.tam;
                                                                                                 lazv = 0:
                                                                               c60
edf
                                                                               797
                                                                                                 if (1.mi1 > r.mi1) {
        return ret;
ff1 }
                                                                               230
                                                                                                     mi1 = r.mi1, mi = r.mi;
                                                                               ea2
                                                                                                     mi2 = min(1.mi1, r.mi2);
                                                                                                 } else if (1.mi1 < r.mi1) {</pre>
767 void update(int x, int y, int val) {
                                                                               f1e
66a
        int v2 = v += n;
                                                                               e34
                                                                                                     mi1 = 1.mi1, mi = 1.mi;
        for (x += n; x; x /= 2, y = y2) {
                                                                                                     mi2 = min(r.mi1, 1.mi2);
192
                                                                               4b3
970
             if (x >= n) seg[x][y] = val;
                                                                               ef2
                                                                                                 } else {
             else seg[x][y] = seg[2*x][y] + seg[2*x+1][y];
                                                                               a39
                                                                                                     mi1 = 1.mi1, mi = 1.mi+r.mi;
ba9
                                                                               83d
                                                                                                     mi2 = min(1.mi2. r.mi2):
             while (y /= 2) seg[x][y] = seg[x][2*y] + seg[x][2*y+1];
3b1
                                                                               a92
        }
d8d
                                                                               cd0
                                                                                                 if (1.ma1 < r.ma1) {</pre>
62e }
                                                                               6a0
                                                                                                     ma1 = r.ma1, ma = r.ma;
                                                                               96d
                                                                                                     ma2 = max(1.ma1, r.ma2);
                                                                               3c0
                                                                                                 } else if (1.ma1 > r.ma1) {
                                                                               ae0
                                                                                                     ma1 = l.ma1, ma = l.ma;
      SegTree Beats
6.20
                                                                               2ca
                                                                                                     ma2 = max(r.ma1, 1.ma2);
                                                                               da8
                                                                                                 } else {
// \text{ query(a, b)} - \{\{\min(v[a..b]), \max(v[a..b])\}, \sup(v[a..b])\}
                                                                               db2
                                                                                                     ma1 = 1.ma1, ma = 1.ma+r.ma;
// updatemin(a, b, x) faz com que v[i] <- min(v[i], x),</pre>
                                                                               c05
                                                                                                     ma2 = max(1.ma2, r.ma2);
// para i em [a, b]
                                                                               11c
                                                                                                 }
// updatemax faz o mesmo com max, e updatesum soma x
                                                                               1ba
                                                                                            }
// em todo mundo do intervalo [a, b]
                                                                               4b4
                                                                                            void setmin(ll x) {
//
```

```
55e
                 if (x \ge ma1) return:
463
                 sum += (x - ma1)*ma:
be5
                 if (mi1 == ma1) mi1 = x;
                 if (mi2 == ma1) mi2 = x;
0a0
b81
                 ma1 = x:
            }
0 c 3
6cb
            void setmax(ll x) {
                 if (x <= mi1) return:</pre>
e25
7e8
                 sum += (x - mi1)*mi;
                 if (ma1 == mi1) ma1 = x:
Obb
c32
                 if (ma2 == mi1) ma2 = x;
1ff
                 mi1 = x:
a86
            }
4cf
            void setsum(ll x) {
fe8
                 mi1 += x, mi2 += x, ma1 += x, ma2 += x;
620
                 sum += x*tam;
c46
                 lazv += x;
            }
b53
47f
        };
62b
        node seg[4*MAX];
052
        int n, *v;
93b
        node build(int p=1, int l=0, int r=n-1) {
d84
            if (1 == r) return seg[p] = {v[1]};
ee4
            int m = (1+r)/2:
3d6
            return seg[p] = {build(2*p, 1, m), build(2*p+1, m+1, r)};
444
860
        void build(int n2, int* v2) {
680
            n = n2, v = v2;
6f2
            build():
acb
        void prop(int p, int 1, int r) {
ceb
8ce
            if (1 == r) return;
            for (int k = 0; k < 2; k++) {
abd
d07
                 if (seg[p].lazy) seg[2*p+k].setsum(seg[p].lazy);
843
                 seg[2*p+k].setmin(seg[p].ma1);
f79
                 seg[2*p+k].setmax(seg[p].mi1);
585
431
            seg[p].lazv = 0;
7ee
        pair < pair < 11 , 11 > , 11 > query (int a, int b, int p=1, int 1=0,
055
   int r=n-1) {
            if (b < l or r < a) return {{LINF, -LINF}, 0};</pre>
e07
            if (a \le 1 \text{ and } r \le b) \text{ return } \{\{seg[p].mi1, seg[p].ma1\},
9be
   seg[p].sum};
            prop(p, 1, r);
6b9
```

```
ee4
             int m = (1+r)/2:
e6f
             auto L = query(a, b, 2*p, 1, m), R = query(a, b, 2*p+1,
   m+1, r);
96d
             return {{min(L.f.f, R.f.f), max(L.f.s, R.f.s)}, L.s+R.s};
e9d
2c8
         node updatemin(int a, int b, ll x, int p=1, int l=0, int
   r=n-1) {
744
             if (b < l or r < a or seg[p].ma1 <= x) return seg[p];</pre>
309
             if (a \le 1 \text{ and } r \le b \text{ and } seg[p].ma2 < x) {
                 seg[p].setmin(x);
ccd
534
                 return seg[p];
bbf
             }
6b9
             prop(p, 1, r);
ee4
             int m = (1+r)/2:
96a
             return seg[p] = \{updatemin(a, b, x, 2*p, 1, m),
                              updatemin(a, b, x, 2*p+1, m+1, r)};
4db
        }
aad
044
         node updatemax(int a, int b, ll x, int p=1, int l=0, int
   r=n-1) {
b59
             if (b < l or r < a or seg[p].mi1 >= x) return seg[p];
             if (a \le 1 \text{ and } r \le b \text{ and } seg[p].mi2 > x) {
a9e
                 seg[p].setmax(x);
e8a
534
                 return seg[p];
e9b
             }
6b9
             prop(p, 1, r);
ee4
             int m = (1+r)/2:
ee3
             return seg[p] = \{updatemax(a, b, x, 2*p, 1, m),
98b
                              updatemax(a, b, x, 2*p+1, m+1, r)};
323
        }
         node updatesum(int a, int b, ll x, int p=1, int l=0, int
aee
   r=n-1) {
e9f
             if (b < 1 or r < a) return seg[p];</pre>
             if (a \le 1 \text{ and } r \le b)
9a3
8f4
                 seg[p].setsum(x):
534
                 return seg[p];
             }
596
6b9
             prop(p, 1, r);
             int m = (1+r)/2;
ee4
7b6
             return seg[p] = \{updatesum(a, b, x, 2*p, 1, m),
                              updatesum(a, b, x, 2*p+1, m+1, r)};
483
        }
111
0d2 };
```

# 6.21 SegTree Colorida

// Cada posicao tem um valor e uma cor

```
// O construtor receve um vector de {valor, cor}
// e o numero de cores (as cores devem estar em [0, c-1])
// query(c, a, b) retorna a soma dos valores
// de todo mundo em [a, b] que tem cor c
// update(c, a, b, x) soma x em todo mundo em
// [a, b] que tem cor c
// paint(c1, c2, a, b) faz com que todo mundo
// em [a, b] que tem cor c1 passe a ter cor c2
// Complexidades:
// construir - O(n log(n)) espaco e tempo
// query - O(log(n))
// update - O(log(n))
// paint - O(log(n)) amortizado
04f struct seg_color {
3c9
        struct node {
b19
            node *1, *r;
0f9
            int cnt;
9ca
            ll val, lazy;
            node() : 1(NULL), r(NULL), cnt(0), val(0), lazy(0) {}
277
            void update() {
01e
                cnt = 0, val = 0;
d0a
                for (auto i : {1, r}) if (i) {
bc4
c89
                    i->prop();
281
                    cnt += i->cnt, val += i->val;
                }
68d
554
            }
a9c
            void prop() {
2dd
                if (!lazy) return;
3f7
                val += lazy*(ll)cnt;
b64
                for (auto i : {1, r}) if (i) i->lazy += lazy;
c60
                lazv = 0:
e24
            }
514
        };
1a8
        int n;
9b0
        vector < node *> seg;
6e0
        seg_color(vector<pair<int, int>>& v, int c) : n(v.size()),
   seg(c, NULL) {
830
            for (int i = 0; i < n; i++)</pre>
                seg[v[i].second] = insert(seg[v[i].second], i,
9b7
   v[i].first, 0, n-1);
94a
3c7
        \simseg_color() {
dde
            queue < node *> q;
```

```
3a6
            for (auto i : seg) q.push(i);
402
            while (q.size()) {
20b
                auto i = q.front(); q.pop();
dab
                if (!i) continue;
7c7
                q.push(i->1), q.push(i->r);
5се
                delete i;
c60
            }
        }
139
        node* insert(node* at, int idx, int val, int l, int r) {
40b
1a4
            if (!at) at = new node();
232
            if (l == r) return at->cnt = 1, at->val = val, at;
ee4
            int m = (1+r)/2:
137
            if (idx <= m) at->1 = insert(at->1, idx, val, 1, m);
3e6
            else at->r = insert(at->r, idx, val, m+1, r);
            return at->update(), at;
cff
d6e
        }
870
        11 query(node* at, int a, int b, int l, int r) {
61b
            if (!at or b < 1 or r < a) return 0;
d9f
            at->prop();
ch2
            if (a <= l and r <= b) return at->val;
            int m = (1+r)/2;
ee4
4c4
            return query(at->1, a, b, 1, m) + query(at->r, a, b, m+1,
   r);
8c3
        11 query(int c, int a, int b) { return query(seg[c], a, b, 0,
e54
   n-1); }
91c
        void update(node* at, int a, int b, int x, int l, int r) {
            if (!at or b < l or r < a) return;</pre>
fba
d9f
            at->prop();
9a3
            if (a \le 1 \text{ and } r \le b) {
e9a
                at -> lazv += x;
cb2
                return void(at->prop());
051
            }
ee4
            int m = (1+r)/2;
0ъ0
            update(at->1, a, b, x, 1, m), update(at->r, a, b, x, m+1,
   r):
7b4
            at->update();
9fd
a40
        void update(int c, int a, int b, int x) { update(seg[c], a, b,
   x, 0, n-1); }
70c
        void paint(node*& from, node*& to, int a, int b, int l, int r)
   {
10f
            if (to == from or !from or b < l or r < a) return;</pre>
e85
            from ->prop();
889
            if (to) to->prop();
            if (a \le 1 \text{ and } r \le b)
9a3
```

```
24d
                if (!to) {
38f
                    to = from;
140
                    from = NULL;
505
                    return;
e5f
                int m = (1+r)/2;
ee4
1cb
                paint(from->1, to->1, a, b, 1, m), paint(from->r,
   to->r, a, b, m+1, r);
                to->update();
72d
270
                delete from:
140
                from = NULL;
505
                return;
a0e
            }
            if (!to) to = new node();
019
ee4
            int m = (1+r)/2;
            paint(from->1, to->1, a, b, 1, m), paint(from->r, to->r,
1cb
   a, b, m+1, r);
45a
            from ->update(), to ->update();
4aa
        void paint(int c1, int c2, int a, int b) { paint(seg[c1],
   seg[c2], a, b, 0, n-1); }
293 };
6.22 SegTree Esparsa - Lazy
// Query: soma do range [a. b]
// Update: flipa os valores de [a, b]
// O MAX tem q ser Q log N para Q updates
// Complexidades:
// build - O(1)
// query - 0(log(n))
// update - O(log(n))
aa4 namespace seg {
6de
        int seg[MAX], lazy[MAX], R[MAX], L[MAX], ptr;
        int get_l(int i){
e9a
3db
            if (L[i] == 0) L[i] = ptr++;
            return L[i];
a96
b6e
943
        int get_r(int i){
71b
            if (R[i] == 0) R[i] = ptr++;
283
            return R[i];
43a
        }
```

e71

void build() { ptr = 2; }

```
void prop(int p, int 1, int r) {
ceb
b77
             if (!lazv[p]) return;
76c
             seg[p] = r-l+1 - seg[p];
             if (1 != r) lazy[get_l(p)]^=lazy[p],
213
    lazy[get_r(p)]^=lazy[p];
3c7
             lazy[p] = 0;
        }
20b
158
        int query(int a, int b, int p=1, int l=0, int r=N-1) {
6b9
             prop(p, 1, r);
786
             if (b < 1 \text{ or } r < a) \text{ return } 0;
             if (a <= 1 and r <= b) return seg[p];</pre>
527
ee4
             int m = (1+r)/2;
818
             return query(a, b, get_l(p), l, m)+query(a, b, get_r(p),
   m+1, r);
0d9
       }
51f
        int update(int a, int b, int p=1, int l=0, int r=N-1) {
6b9
             prop(p, 1, r);
e9f
             if (b < 1 or r < a) return seg[p];</pre>
             if (a \le 1 \text{ and } r \le b) \{
9a3
ab6
                 lazv[p] ^= 1;
6b9
                 prop(p, 1, r);
534
                 return seg[p];
8e4
             }
ee4
             int m = (1+r)/2:
             return seg[p] = update(a, b, get_l(p), l, m)+update(a, b,
    get_r(p), m+1, r);
       }
1dc
dc3 };
6.23 SegTree Esparsa - O(q) memoria
// Query: min do range [a, b]
// Update: troca o valor de uma posicao
// Usa O(q) de memoria para q updates
//
// Complexidades:
// query - O(log(n))
// update - O(log(n))
13d template < typename T> struct seg {
3c9
        struct node {
d53
             node* ch[2];
```

```
970
            char d;
ca0
            T v;
c4e
            T mi;
            node(int d_, T v_, T val) : d(d_), v(v_) {
d4e
e71
                ch[0] = ch[1] = NULL:
d6e
                mi = val;
065
            node(node* x) : d(x->d), v(x->v), mi(x->mi) {
b32
c99
                ch[0] = x -> ch[0], ch[1] = x -> ch[1];
cb7
            }
01e
            void update() {
909
                mi = numeric_limits <T>::max();
                for (int i = 0; i < 2; i++) if (ch[i])</pre>
151
                     mi = min(mi, ch[i]->mi);
b5a
fe3
            }
530
        };
bb7
        node* root:
9c5
        char n;
ba7
        seg() : root(NULL), n(0) {}
512
        \simseg() {
4c0
            std::vector<node*> q = {root};
402
            while (q.size()) {
e5d
                node* x = q.back(); q.pop_back();
ee9
                if (!x) continue:
                q.push_back(x->ch[0]), q.push_back(x->ch[1]);
73f
bf0
                delete x;
            }
d3e
d8c
        }
1a6
        char msb(T v, char l, char r) { // msb in range (l, r]
            for (char i = r; i > 1; i--) if (v>>i&1) return i;
8e4
daa
            return -1:
688
        }
430
        void cut(node* at, T v, char i) {
677
            char d = msb(v ^a at -> v, at -> d, i);
23b
            if (d == -1) return; // no need to split
            node* nxt = new node(at):
ebf
d43
            at -> ch[v>>d&1] = NULL;
34f
            at -> ch[!(v>>d&1)] = nxt;
150
            at -> d = d:
0b3
        }
6e5
        node* update(node* at, T idx, T val, char i) {
```

```
c8c
             if (!at) return new node(-1, idx, val);
d67
             cut(at, idx, i);
1a2
             if (at -> d == -1) { // leaf }
792
                 at->mi = val;
ce6
                 return at;
a6f
b29
             bool dir = idx>>at->d&1;
c8f
             at - ch[dir] = update(at - ch[dir], idx, val, at - ch-1);
7b4
             at->update();
             return at:
ce6
76d
        }
85 c
         void update(T idx, T val) {
8f4
             while (idx >> n) n++:
61e
             root = update(root, idx, val, n-1);
        }
79d
9d8
         T query(node* at, T a, T b, T l, T r, char i) {
df0
             if (!at or b < l or r < a) return numeric_limits<T>::max();
fd3
             if (a <= l and r <= b) return at->mi;
             T m = 1 + (r-1)/2:
841
c85
             if (at->d < i) {</pre>
c59
                 if ((at->v>>i&1) == 0) return query(at, a, b, 1, m,
   i-1):
                 else return query(at, a, b, m+1, r, i-1);
ca4
934
373
             return min(query(at->ch[0], a, b, 1, m, i-1),
    query(at->ch[1], a, b, m+1, r, i-1));
2db
         T query(T 1, T r) { return query(root, 1, r, 0, (T(1) \le n) - 1,
034
    n-1); }
d7f };
6.24 SegTree Iterativa
// Consultas 0-based
// Valores iniciais devem estar em (seg[n], ..., seg[2*n-1])
// Query: soma do range [a, b]
// Update: muda o valor da posicao p para x
//
// Complexidades:
// build - O(n)
// query - O(log(n))
// update - O(log(n))
6a4 int seg[2 * MAX];
la8 int n;
```

```
0a8 void build() {
d15
        for (int i = n - 1; i; i--) seg[i] = seg[2*i] + seg[2*i+1];
9a8 }
4ea int query(int a, int b) {
7c9
       int ret = 0:
728
       for (a += n, b += n; a <= b; ++a /= 2, --b /= 2) {
            if (a % 2 == 1) ret += seg[a];
4ea
            if (b \% 2 == 0) ret += seg[b];
244
ac0
edf
       return ret;
24a }
ff3 void update(int p, int x) {
        seg[p += n] = x;
37d
c8c
        while (p /= 2) seg[p] = seg[2*p] + seg[2*p+1];
02d }
```

## 6.25 SegTree Iterativa com Lazy Propagation

```
// Query: soma do range [a, b]
// Update: soma x em cada elemento do range [a, b]
// Para mudar, mudar as funcoes junta, poe e query
// LOG = ceil(log2(MAX))
// Complexidades:
// build - O(n)
// query - 0(log(n))
// update - O(log(n))
aa4 namespace seg {
6db
        11 seg[2*MAX], lazy[2*MAX];
1a8
        int n;
9b3
        ll junta(ll a, ll b) {
534
            return a+b;
e26
       }
        // soma x na posicao p de tamanho tam
1b4
        void poe(int p, ll x, int tam, bool prop=1) {
517
            seg[p] += x*tam;
6ae
            if (prop and p < n) lazy[p] += x;</pre>
8bc
        // atualiza todos os pais da folha p
```

```
b1e
        void sobe(int p) {
d5a
            for (int tam = 2; p /= 2; tam *= 2) {
4ca
                seg[p] = junta(seg[2*p], seg[2*p+1]);
388
                poe(p, lazy[p], tam, 0);
            }
acd
        }
b76
        // propaga o caminho da raiz ate a folha p
        void prop(int p) {
a0a
076
            int tam = 1 << (LOG-1);</pre>
0a8
            for (int s = LOG; s; s--, tam /= 2) {
4b1
                int i = p >> s;
27 c
                if (lazv[i]) {
860
                    poe(2*i, lazy[i], tam);
e38
                    poe(2*i+1, lazy[i], tam);
b97
                    lazy[i] = 0;
de8
                }
3ed
            }
e29
        }
61c
        void build(int n2, int* v) {
            n = n2:
1e3
95f
            for (int i = 0; i < n; i++) seg[n+i] = v[i];
            for (int i = n-1; i; i--) seg[i] = junta(seg[2*i],
c41
   seg[2*i+1]):
f4c
            for (int i = 0; i < 2*n; i++) lazy[i] = 0;</pre>
8bb
        }
4f3
        11 query(int a, int b) {
b73
            ll ret = 0;
b48
            for (prop(a+=n), prop(b+=n); a \le b; ++a/=2, --b/=2) {
a8e
                if (a%2 == 1) ret = junta(ret, seg[a]);
                if (b%2 == 0) ret = junta(ret, seg[b]);
c58
510
            }
edf
            return ret;
38b
        }
        void update(int a, int b, int x) {
a28
c2d
            int a2 = a += n, b2 = b += n, tam = 1;
Off
            for (; a <= b; ++a/=2, --b/=2, tam *= 2) {
32a
                if (a\%2 == 1) poe(a, x, tam);
                if (b\%2 == 0) poe(b, x, tam);
9da
9bc
0f7
            sobe(a2), sobe(b2);
adc
        }
6dc }:
```

## 6.26 SegTree PA

```
// Segtree de PA
// update_set(1, r, A, R) seta [1, r] para PA(A, R),
// update_add soma PA(A, R) em [1, r]
// query(l, r) retorna a soma de [l, r]
// PA(A, R) eh a PA: [A+R, A+2R, A+3R, ...]
// Complexidades:
// construir - O(n)
// update_set, update_add, query - O(log(n))
dc7 struct seg_pa {
350
        struct Data {
8f5
            ll sum:
662
            11 set_a, set_r, add_a, add_r;
9b7
            Data(): sum(0), set_a(LINF), set_r(0), add_a(0), add_r(0)
   {}
eb6
        };
        vector < Data > seg;
16a
1a8
        int n;
d45
        seg_pa(int n_) {
e95
            n = n_{-};
            seg = vector < Data > (4*n);
fc3
ce0
       }
ceb
        void prop(int p, int l, int r) {
d5a
            int tam = r-1+1:
c3f
            11 &sum = seg[p].sum, &set_a = seg[p].set_a, &set_r =
   seg[p].set_r,
a1b
                \&add_a = seg[p].add_a, \&add_r = seg[p].add_r;
c02
            if (set a != LINF) {
660
                set_a += add_a, set_r += add_r;
                sum = set_a*tam + set_r*tam*(tam+1)/2;
06e
579
                if (1 != r) {
                    int m = (1+r)/2;
ee4
                    seg[2*p].set_a = set_a;
886
358
                    seg[2*p].set_r = set_r;
                    seg[2*p].add_a = seg[2*p].add_r = 0;
ed6
f0c
                    seg[2*p+1].set_a = set_a + set_r * (m-l+1);
471
                    seg[2*p+1].set_r = set_r;
d48
                    seg[2*p+1].add_a = seg[2*p+1].add_r = 0;
```

```
a97
                }
823
                set_a = LINF, set_r = 0;
953
                 add_a = add_r = 0;
            } else if (add_a or add_r) {
da7
18b
                sum += add_a*tam + add_r*tam*(tam+1)/2;
579
                if (1 != r) {
ee4
                     int m = (1+r)/2:
ffO
                     seg[2*p].add_a += add_a;
                     seg[2*p].add_r += add_r;
ec0
06c
                     seg[2*p+1].add_a += add_a + add_r * (m-l+1);
a6d
                     seg[2*p+1].add r += add r:
8af
953
                add_a = add_r = 0;
ab7
            }
07f
        }
        int inter(pair<int, int> a, pair<int, int> b) {
0b7
            if (a.first > b.first) swap(a, b);
98c
eef
            return max(0, min(a.second, b.second) - b.first + 1);
628
        11 set(int a, int b, ll aa, ll rr, int p, int l, int r) {
be1
6b9
            prop(p, 1, r);
457
            if (b < 1 or r < a) return seg[p].sum;</pre>
9a3
            if (a \le 1 \text{ and } r \le b)
91c
                seg[p].set_a = aa;
774
                seg[p].set_r = rr;
6b9
                prop(p, 1, r);
254
                return seg[p].sum;
            }
8ee
ee4
            int m = (1+r)/2;
963
            int tam_1 = inter({1, m}, {a, b});
c34
            return seg[p].sum = set(a, b, aa, rr, 2*p, 1, m) +
                 set(a, b, aa + rr * tam_l, rr, 2*p+1, m+1, r);
365
8e2
        }
f55
        void update_set(int 1, int r, 11 aa, 11 rr) {
6f7
            set(1, r, aa, rr, 1, 0, n-1);
913
5f6
        11 add(int a, int b, ll aa, ll rr, int p, int l, int r) {
6b9
            prop(p, 1, r):
            if (b < l or r < a) return seg[p].sum;</pre>
457
9a3
            if (a \le 1 \text{ and } r \le b) \{
359
                 seg[p].add_a += aa;
1ee
                 seg[p].add_r += rr;
6b9
                prop(p, 1, r);
254
                return seg[p].sum;
```

```
d19
ee4
            int m = (1+r)/2:
963
            int tam_1 = inter({1, m}, {a, b});
            return seg[p].sum = add(a, b, aa, rr, 2*p, 1, m) +
586
695
                add(a, b, aa + rr * tam_l, rr, 2*p+1, m+1, r);
904
848
        void update_add(int 1, int r, 11 aa, 11 rr) {
            add(1, r, aa, rr, 1, 0, n-1);
afa
81e
f45
        ll query(int a, int b, int p, int l, int r) {
6b9
            prop(p, 1, r);
786
            if (b < 1 or r < a) return 0;</pre>
e9a
            if (a <= 1 and r <= b) return seg[p].sum:
ee4
            int m = (1+r)/2:
b1f
            return query(a, b, 2*p, 1, m) + query(a, b, 2*p+1, m+1, r);
f6e
bfc
        11 query(int 1, int r) { return query(1, r, 1, 0, n-1); }
bc4 }:
6.27 SegTree Persistente
```

```
// SegTree de soma, update de somar numa posicao
//
// query(a, b, t) retorna a query de [a, b] na versao t
// update(a, x, t) faz um update v[a]+=x a partir da
// versao de t. criando uma nova versao e retornando seu id
// Por default, faz o update a partir da ultima versao
//
// build - O(n)
// query - 0(log(n))
// update - 0(log(n))
54a const int MAX = 1e5+10, UPD = 1e5+10, LOG = 18;
6de const int MAXS = 2*MAX+UPD*LOG;
f6e namespace perseg {
bd6
       11 seg[MAXS];
f4e
       int rt[UPD], L[MAXS], R[MAXS], cnt, t;
052
       int n, *v;
3c4
        ll build(int p, int l, int r) {
            if (1 == r) return seg[p] = v[1];
6cd
855
           L[p] = cnt++, R[p] = cnt++;
ee4
            int m = (1+r)/2;
275
           return seg[p] = build(L[p], 1, m) + build(R[p], m+1, r);
39d
        }
```

```
8b0
        void build(int n2, int* v2) {
680
            n = n2, v = v2;
856
            rt[0] = cnt++;
            build(0, 0, n-1);
c50
a2e
        }
f45
        11 query(int a, int b, int p, int l, int r) {
786
            if (b < 1 \text{ or } r < a) \text{ return } 0:
527
            if (a <= 1 and r <= b) return seg[p];</pre>
            int m = (1+r)/2;
ee4
            return query(a, b, L[p], 1, m) + query(a, b, R[p], m+1, r);
1 ed
4d2
        }
182
        11 query(int a, int b, int tt) {
c13
            return querv(a, b, rt[tt], 0, n-1);
726
        }
bb3
        11 update(int a, int x, int lp, int p, int l, int r) {
            if (1 == r) return seg[p] = seg[lp]+x;
747
ee4
            int m = (1+r)/2;
            if (a <= m)
ab8
b48
                return seg[p] = update(a, x, L[lp], L[p]=cnt++, l, m)
    + seg[R[p]=R[lp]];
            return seg[p] = seg[L[p]=L[lp]] + update(a, x, R[lp],
   R[p]=cnt++, m+1, r);
788
        }
6f6
        int update(int a, int x, int tt=t) {
ab3
            update(a, x, rt[tt], rt[++t]=cnt++, 0, n-1);
e0d
            return t:
d63
        }
26f }:
6.28 SegTree Persistente com Lazy
// Nao propaga, meio estranho de mexer, mas da
// query(a, b, t) retorna a query de [a, b] na versao t
// update(a, b, x, t) faz um update v[a..b]+=x a partir da
// versao de t, criando uma nova versao e retornando seu id
// Por default, faz o update a partir da ultima versao
//
// build - O(n)
// query - O(log(n))
// update - O(log(n))
```

54a const int MAX = 1e5+10, UPD = 1e5+10, LOG = 18;

ab3 const int MAXS = 2\*MAX + 4\*UPD\*LOG:

f6e namespace perseg {

```
9eb
        int seg[MAXS];
f4e
        int rt[UPD], L[MAXS], R[MAXS], cnt, t;
052
        int n, *v;
        int build(int p, int 1, int r) {
adf
            if (1 == r) return seg[p] = v[1];
6cd
855
            L[p] = cnt++, R[p] = cnt++;
            int m = (1+r)/2;
ee4
01d
            return seg[p] = max(build(L[p], 1, m), build(R[p], m+1,
   r));
ffd
       }
0d8
        void build(int n2, int *v2) {
680
            n = n2, v = v2:
            rt[0] = cnt++;
856
            build(0, 0, n-1);
c50
       }
a2e
976
        int query(int a, int b, int p, int l, int r) {
27b
            if (b < 1 or r < a) return -INF:
793
            if (a <= 1 and r <= b) return lazy[p] + seg[p];</pre>
            int m = (1+r)/2;
ee4
7a2
            int ret = lazy[p] + max(query(a, b, L[p], 1, m), query(a,
   b, R[p], m+1, r));
            return ret;
edf
9a7
442
        int query(int a, int b, int tt) {
c13
            return query(a, b, rt[tt], 0, n-1);
        }
a05
bc1
        int update(int a, int b, int x, int lp, int p, int l, int r) {
            tie(seg[p], lazy[p], L[p], R[p]) = {seg[lp], lazy[lp],
   L[lp], R[lp]};
            if (b < l or r < a) return seg[p] + lazy[p];</pre>
847
            if (a <= 1 and r <= b) return seg[p] + (lazy[p] += x);
32a
ee4
            int m = (1+r)/2:
            seg[p] = max(update(a, b, x, L[lp], L[p] = cnt++, l, m),
24a
bdb
                         update(a, b, x, R[lp], R[p] = cnt++, m+1, r);
            lazv[p] = lazv[lp];
1ed
1b7
            return seg[p] + lazv[p];
877
cbf
        int update(int a, int b, int x, int tt=t) {
            assert(tt <= t):
aa8
            update(a, b, x, rt[tt], rt[++t]=cnt++, 0, n-1);
661
e0d
            return t:
       }
aad
f27 };
```

# 6.29 Sparse Table

```
// Resolve RMQ
// MAX2 = log(MAX)
//
// Complexidades:
// build - O(n log(n))
// query - 0(1)
cca namespace sparse {
710
        int m[MAX2][MAX], n;
61c
        void build(int n2, int* v) {
1e3
            n = n2:
78e
            for (int i = 0; i < n; i++) m[0][i] = v[i];</pre>
            for (int j = 1; (1<<j) <= n; j++) for (int i = 0; i+(1<<<math>j)
a1c
   <= n: i++)
5d5
                m[j][i] = min(m[j-1][i], m[j-1][i+(1<<(j-1))]);
cae
        }
4ea
        int query(int a, int b) {
            int j = __builtin_clz(1) - __builtin_clz(b-a+1);
ee5
dc3
            return min(m[j][a], m[j][b-(1<<j)+1]);</pre>
fba
        }
7aa }
```

# 6.30 Sparse Table Disjunta

```
// Resolve qualquer operacao associativa
// MAX2 = log(MAX)
//
// Complexidades:
// build - O(n log(n))
// query - 0(1)
cca namespace sparse {
9bf
        int m[MAX2][2*MAX], n, v[2*MAX];
5f7
        int op(int a, int b) { return min(a, b); }
0d8
        void build(int n2, int* v2) {
1e3
            n = n2;
df4
            for (int i = 0; i < n; i++) v[i] = v2[i];</pre>
            while (n&(n-1)) n++;
a84
3d2
            for (int j = 0; (1<<j) < n; j++) {
                 int len = 1<<j;</pre>
1c0
d9b
                for (int c = len; c < n; c += 2*len) {
332
                     m[j][c] = v[c], m[j][c-1] = v[c-1];
668
                     for (int i = c+1; i < c+len; i++) m[j][i] =</pre>
   op(m[j][i-1], v[i]);
```

```
432
                    for (int i = c-2; i >= c-len; i--) m[j][i] =
   op(v[i], m[j][i+1]);
eda
               }
            }
f4d
ce3
9e3
        int query(int 1, int r) {
f13
            if (1 == r) return v[1];
            int j = __builtin_clz(1) - __builtin_clz(1^r);
e6d
d67
            return op(m[j][1], m[j][r]);
a7b
       }
fd8 }
6.31 Splay Tree
```

```
// SEMPRE QUE DESCER NA ARVORE, DAR SPLAY NO
// NODE MAIS PROFUNDO VISITADO
// Todas as operacoes sao O(\log(n)) amortizado
// Se quiser colocar mais informação no node,
// mudar em 'update'
538 template < typename T > struct splaytree {
3c9
        struct node {
183
            node *ch[2], *p;
e4d
            int sz;
f48
            T val;
da0
            node(T v) {
                ch[0] = ch[1] = p = NULL;
696
a26
                sz = 1;
250
                val = v:
2d0
            }
01e
            void update() {
a26
                sz = 1;
c7c
                for (int i = 0; i < 2; i++) if (ch[i]) {</pre>
d5f
                     sz += ch[i]->sz;
486
                }
            }
f45
        };
aa3
bb7
        node* root;
fbc
        splaytree() { root = NULL; }
214
        splaytree(const splaytree& t) {
cbf
            throw logic_error("Nao copiar a splaytree!");
1 f 1
891
        \simsplaytree() {
609
            vector < node *> q = {root};
```

```
402
            while (q.size()) {
e5d
                 node* x = q.back(); q.pop_back();
ee9
                 if (!x) continue;
73f
                 q.push_back(x->ch[0]), q.push_back(x->ch[1]);
bf0
                 delete x;
d3e
            }
        }
837
94f
        void rotate(node* x) { // x vai ficar em cima
d9b
            node *p = x->p, *pp = p->p;
ecf
            if (pp) pp->ch[pp->ch[1] == p] = x;
286
            bool d = p - ch[0] = x;
d63
            p \rightarrow ch[!d] = x \rightarrow ch[d], x \rightarrow ch[d] = p:
bad
            if (p->ch[!d]) p->ch[!d]->p = p;
fc2
            x->p = pp, p->p = x;
            p->update(), x->update();
1ea
007
        }
3fa
        node* splay(node* x) {
a39
            if (!x) return x;
4ea
            root = x:
3cf
            while (x->p) {
d9b
                 node *p = x->p, *pp = p->p;
                 if (!pp) return rotate(x), x; // zig
359
еЗс
                 if ((pp->ch[0] == p)^(p->ch[0] == x))
a2b
                     rotate(x), rotate(x); // zigzag
4b2
                 else rotate(p), rotate(x); // zigzig
028
ea5
            return x:
379
        }
313
        node* insert(T v, bool lb=0) {
b64
            if (!root) return lb ? NULL : root = new node(v);
002
            node *x = root, *last = NULL;;
31e
            while (1) {
5d7
                bool d = x -> val < v:
0fd
                 if (!d) last = x;
c2e
                 if (x->val == v) break;
c16
                 if (x->ch[d]) x = x->ch[d];
4e6
                 else {
dea
                     if (lb) break:
                     x - ch[d] = new node(v);
055
99c
                     x - ch[d] - p = x;
30e
                     x = x - ch[d];
c2b
                     break;
                }
68a
1ab
            }
0b6
            splay(x);
61c
            return lb ? splay(last) : x;
```

```
622
        int size() { return root ? root->sz : 0; }
сОс
2ca
        int count(T v) { return insert(v, 1) and root->val == v; }
111
        node* lower_bound(T v) { return insert(v, 1); }
26b
        void erase(T v) {
446
            if (!count(v)) return;
bce
            node *x = root. *1 = x -> ch[0]:
268
            if (!1) {
8b1
                root = x -> ch[1];
32e
                if (root) root->p = NULL;
8f3
                return delete x;
a86
            }
5e7
            root = 1. 1 - > p = NULL:
902
            while (1->ch[1]) 1 = 1->ch[1];
bab
            splay(1);
            1 - ch[1] = x - ch[1];
f0e
7d9
            if (1->ch[1]) 1->ch[1]->p = 1;
bf0
            delete x;
62a
            1->update();
007
        }
24a
        int order_of_key(T v) {
62b
            if (!lower_bound(v)) return root ? root->sz : 0;
1cc
            return root -> ch [0] ? root -> ch [0] -> sz : 0;
b00
        }
db6
        node* find_by_order(int k) {
084
            if (k >= size()) return NULL;
52f
            node* x = root;
31e
            while (1) {
                if (x->ch[0] \text{ and } x->ch[0]->sz >= k+1) x = x->ch[0];
20f
4e6
                else {
                     if (x->ch[0]) k -= x->ch[0]->sz;
a1c
                    if (!k) return splay(x);
1dc
                    k--, x = x-> ch[1];
eb8
                }
            }
e01
0de
        }
        T min() {
19c
52f
            node* x = root;
6f6
            while (x->ch[0]) x = x->ch[0]; // max -> ch[1]
3e9
            return splay(x)->val;
        }
70e
4ff };
```

# 6.32 Splay Tree Implicita

```
// vector da NASA
```

```
// Um pouco mais rapido q a treap
// O construtor a partir do vector
// eh linear, todas as outras operacoes
// custam O(log(n)) amortizado
081 template < typename T > struct splay {
3c9
        struct node {
183
            node *ch[2], *p;
e4d
            int sz;
875
            T val, sub, lazy;
aa6
            bool rev;
da0
            node(T v) {
696
                 ch[0] = ch[1] = p = NULL;
a26
                 sz = 1:
1e4
                 sub = val = v;
c60
                lazy = 0;
b67
                 rev = false;
            }
48f
a9c
            void prop() {
0ec
                 if (lazy) {
924
                     val += lazy, sub += lazy*sz;
091
                     if (ch[0]) ch[0]->lazy += lazy;
                     if (ch[1]) ch[1]->lazy += lazy;
1a8
a98
                }
1bb
                if (rev) {
80a
                     swap(ch[0], ch[1]);
628
                     if (ch[0]) ch[0]->rev ^= 1;
adc
                     if (ch[1]) ch[1]->rev ^= 1;
30a
a32
                lazv = 0, rev = 0;
            }
6bf
01e
            void update() {
0 c 3
                 sz = 1, sub = val;
c7c
                 for (int i = 0; i < 2; i++) if (ch[i]) {
                     ch[i]->prop();
05f
d5f
                     sz += ch[i]->sz;
4a1
                     sub += ch[i] -> sub;
6 c 1
                }
e98
            }
b4a
        };
bb7
        node* root;
5d9
        splay() { root = NULL; }
9b1
        splay(node* x) {
4ea
            root = x;
32e
            if (root) root->p = NULL;
```

```
371
1b7
        splay(vector < T > v) { // O(n)}
950
            root = NULL;
806
            for (T i : v) {
2a0
                node* x = new node(i):
bd1
                x - ch[0] = root;
37a
                if (root) root->p = x;
                root = x;
4ea
a0a
                root ->update();
            }
17 c
c6b
a9e
        splay(const splay& t) {
e62
            throw logic_error("Nao copiar a splay!");
d4d
        }
5ab
        \simsplay() {
609
            vector < node *> q = {root};
402
            while (q.size()) {
                node* x = q.back(); q.pop_back();
e5d
ee9
                if (!x) continue;
73f
                q.push_back(x->ch[0]), q.push_back(x->ch[1]);
bf0
                delete x;
d3e
            }
        }
a1c
73c
        int size(node* x) { return x ? x->sz : 0; }
94f
        void rotate(node* x) { // x vai ficar em cima
            node *p = x - p, *pp = p - p;
d9b
ecf
            if (pp) pp -> ch[pp -> ch[1] == p] = x:
            bool d = p - ch[0] == x;
286
d63
            p -> ch[!d] = x -> ch[d], x -> ch[d] = p;
            if (p->ch[!d]) p->ch[!d]->p = p;
bad
fc2
            x->p = pp, p->p = x;
            p->update(), x->update();
1ea
007
        }
6a0
        node* splaya(node* x) {
a39
            if (!x) return x;
be6
            root = x, x->update();
3cf
            while (x->p) {
                node *p = x->p, *pp = p->p;
d9b
359
                if (!pp) return rotate(x), x; // zig
                if ((pp->ch[0] == p)^(p->ch[0] == x))
e3c
a2b
                    rotate(x), rotate(x); // zigzag
4b2
                else rotate(p), rotate(x); // zigzig
            }
028
ea5
            return x;
21a
a7f
        node* find(int v) {
```

```
a2e
             if (!root) return NULL;
52f
             node *x = root:
6cd
             int key = 0;
31e
             while (1) {
857
                 x->prop():
ba1
                 bool d = kev + size(x->ch[0]) < v;
877
                 if (\text{key} + \text{size}(x->\text{ch}[0]) != v \text{ and } x->\text{ch}[d]) {
15e
                      if (d) key += size(x->ch[0])+1;
30e
                      x = x - ch[d];
a30
                 } else break:
3c3
152
             return splaya(x);
f19
c0c
        int size() { return root ? root->sz : 0; }
        void join(splay<T>& 1) { // assume que l < *this</pre>
c26
690
             if (!size()) swap(root, l.root);
579
             if (!size() or !l.size()) return;
             node* x = 1.root:
bee
31e
             while (1) {
857
                 x->prop();
34d
                 if (!x->ch[1]) break;
bd8
                 x = x -> ch[1];
fa3
             }
147
             1.splaya(x), root->prop(), root->update();
42b
             x - ch[1] = root, x - ch[1] - p = x;
0aa
             root = 1.root, 1.root = NULL:
a0a
             root ->update();
7e6
        }
5ed
        node* split(int v) { // retorna os elementos < v</pre>
             if (v <= 0) return NULL;</pre>
398
             if (v >= size()) {
060
f87
                 node* ret = root;
950
                 root = NULL:
8c9
                 ret ->update();
                 return ret;
edf
             }
d0f
adc
             find(v);
a59
             node* 1 = root -> ch[0];
4df
             root -> ch [0] = NULL:
5a3
             if (1) 1->p = NULL;
             root ->update();
a0a
792
             return 1;
826
        T& operator [](int i) {
511
9d4
             find(i);
ae0
             return root -> val:
829
        }
```

```
231
        void push_back(T v) { // 0(1)
                                                                             658
                                                                                         node() : l(NULL), r(NULL), cnt(0) {}
a01
            node* r = new node(v);
                                                                             01e
                                                                                         void update() {
0de
            r \rightarrow ch[0] = root;
                                                                             a01
                                                                                             cnt = 0;
            if (root) root->p = r;
b11
                                                                             d8a
                                                                                             if (1) cnt += 1->cnt;
b13
            root = r, root->update();
                                                                             e49
                                                                                             if (r) cnt += r->cnt;
315
                                                                             74d
b7a
        T query(int 1, int r) {
                                                                             84f
                                                                                     };
95f
            splay <T> M(split(r+1));
5ff
            splay <T> L(M.split(1));
                                                                             bb7
                                                                                     node* root;
d1c
            T ans = M.root->sub;
                                                                             60
                                                                                     T N:
49c
            M.join(L), join(M);
ba7
            return ans;
                                                                            f34
                                                                                     sms() : root(NULL), N(0) {}
ca3
                                                                             83b
                                                                                     sms(T v) : sms() { while (v >= N) N = 2*N+1; }
41f
        void update(int 1, int r, T s) {
                                                                            5e1
                                                                                     sms(const sms& t) : root(NULL), N(t.N) {
95f
            splay <T > M(split(r+1));
                                                                             3af
                                                                                         for (SIZE_T i = 0; i < t.size(); i++) {</pre>
5ff
            splay <T> L(M.split(1));
                                                                             a0f
                                                                                             T at = t[i];
996
            M.root->lazy += s;
                                                                             e6d
                                                                                             SIZE_T qt = t.count(at);
49c
            M.join(L), join(M);
                                                                             a43
                                                                                             insert(at, qt);
9e9
                                                                            f42
                                                                                             i += qt-1;
                                                                                         }
8c1
        void reverse(int 1, int r) {
                                                                            1e9
                                                                                     }
95f
            splay <T > M(split(r+1));
                                                                             ea8
5ff
            splay <T> L(M.split(1));
                                                                             a96
                                                                                     sms(initializer_list<T> v) : sms() { for (T i : v) insert(i); }
                                                                            2dd
945
            M.root->rev ^= 1;
                                                                                     \simsms() {
49c
                                                                             609
                                                                                         vector < node *> q = {root};
            M.join(L), join(M);
c1a
                                                                             402
                                                                                         while (q.size()) {
2fb
        void erase(int 1, int r) {
                                                                             e5d
                                                                                             node* x = q.back(); q.pop_back();
95f
            splay <T> M(split(r+1));
                                                                             ee9
                                                                                             if (!x) continue;
5ff
            splay <T > L(M.split(1));
                                                                            1c7
                                                                                             q.push_back(x->1), q.push_back(x->r);
dcc
                                                                            bf0
            join(L);
                                                                                             delete x;
68e
        }
                                                                             653
                                                                                         }
                                                                                     }
a35 };
                                                                             f0d
                                                                                     friend void swap(sms& a, sms& b) {
                                                                            fdc
                                                                             49e
                                                                                         swap(a.root, b.root), swap(a.N, b.N);
6.33 Split-Merge Set
                                                                             984
                                                                             83e
                                                                                     sms& operator =(const sms& v) {
// Representa um conjunto de inteiros nao negativos
                                                                             768
                                                                                         sms tmp = v;
// Todas as operacoes custam O(log(N)),
                                                                             420
                                                                                         swap(tmp, *this);
// em que N = maior elemento do set,
                                                                             357
                                                                                         return *this;
// exceto o merge, que custa O(\log(N)) amortizado
                                                                             e9b
                                                                                     }
// Usa O(min(N, n log(N))) de memoria, sendo 'n' o
                                                                             d06
                                                                                     SIZE_T size() const { return root ? root->cnt : 0; }
// numero de elementos distintos no set
                                                                             17f
                                                                                     SIZE_T count(node* x) const { return x ? x->cnt : 0; }
                                                                             75a
                                                                                     void clear() {
2dc template < typename T, bool MULTI = false, typename SIZE_T = int > struct
                                                                             0a0
                                                                                         sms tmp;
   sms {
                                                                             4ac
                                                                                         swap(*this, tmp);
3c9
        struct node {
                                                                            fcb
                                                                                     }
b19
            node *1, *r;
                                                                             a06
                                                                                     void expand(T v) {
15f
            SIZE_T cnt;
```

```
for (; N < v; N = 2*N+1) if (root) {
bc3
                node* nroot = new node();
63 c
956
                nroot ->1 = root;
897
                root = nroot;
a0a
                root ->update();
dd9
9f0
        }
b14
        node* insert(node* at, T idx, SIZE_T qt, T 1, T r) {
            if (!at) at = new node();
1a4
893
            if (1 == r) {
435
                at->cnt += qt;
beb
                if (!MULTI) at->cnt = 1:
ce6
                return at;
            }
a53
841
            T m = 1 + (r-1)/2;
a02
            if (idx \le m) at->1 = insert(at->1, idx, qt, 1, m);
8d9
            else at->r = insert(at->r, idx, qt, m+1, r);
cff
            return at ->update(), at;
83b
        void insert(T v, SIZE_T qt=1) { // insere 'qt' ocorrencias de
cf7
  , ,, ,
882
            if (qt <= 0) return erase(v, -qt);</pre>
72b
            assert(v >= 0);
f52
            expand(v);
5e9
            root = insert(root, v, qt, 0, N);
f62
        }
f06
        node* erase(node* at, T idx, SIZE_T qt, T 1, T r) {
28 c
            if (!at) return at;
54b
            if (1 == r) at->cnt = at->cnt < qt ? 0 : at->cnt - qt;
4e6
            else {
                T m = 1 + (r-1)/2;
841
                if (idx \le m) at->1 = erase(at->1, idx, gt, 1, m):
281
                else at->r = erase(at->r, idx, qt, m+1, r);
ba1
7b4
                at->update();
d3d
            }
135
            if (!at->cnt) delete at, at = NULL;
ce6
            return at;
e1f
        void erase(T v, SIZE_T qt=1) { // remove 'qt' ocorrencias de
43d
   , <sub>v</sub> ,
9c3
            if (v < 0 \text{ or } v > N \text{ or } !qt) \text{ return};
            if (qt < 0) insert(v, -qt);</pre>
9dc
b1d
            root = erase(root, v, qt, 0, N);
b32
8d6
        void erase_all(T v) { // remove todos os 'v'
```

```
if (v < 0 \text{ or } v > N) return;
347
9f2
            root = erase(root, v, numeric_limits < SIZE_T >:: max(), 0, N);
569
        }
        SIZE_T count(node* at, T a, T b, T l, T r) const {
0fe
            if (!at or b < 1 or r < a) return 0;</pre>
61b
            if (a <= l and r <= b) return at->cnt:
0fe
            T m = 1 + (r-1)/2;
841
84a
            return count(at->1, a, b, 1, m) + count(at->r, a, b, m+1,
  r);
4e6
        }
0a9
        SIZE_T count(T v) const { return count(root, v, v, 0, N); }
ffc
        SIZE T order of kev(T v) { return count(root, 0, v-1, 0, N): }
        SIZE_T lower_bound(T v) { return order_of_key(v); }
df2
        const T operator [](SIZE_T i) const { // i-esimo menor elemento
e68
809
            assert(i >= 0 and i < size());
            node* at = root:
c43
4a5
            T 1 = 0, r = N;
            while (1 < r) {</pre>
40c
                T m = 1 + (r-1)/2;
841
5c2
                if (count(at->1) > i) at = at->1, r = m;
4e6
                else {
b4a
                    i -= count(at->1);
ded
                    at = at->r; 1 = m+1;
fa6
                }
41a
792
            return 1;
67f
        }
        node* merge(node* 1, node* r) {
78c
347
            if (!1 or !r) return 1 ? 1 : r;
504
            if (!1->1 \text{ and } !1->r) { // folha}
599
                if (MULTI) 1->cnt += r->cnt;
55d
                delete r;
792
                return 1;
92c
            }
f58
            1->1 = merge(1->1, r->1), 1->r = merge(1->r, r->r);
f4f
            1->update(), delete r;
792
            return 1;
06a
f59
        void merge(sms& s) { // mergeia dois sets
068
            if (N > s.N) swap(*this, s);
785
            expand(s.N);
938
            root = merge(root, s.root);
            s.root = NULL:
ee2
2f6
        }
```

```
dc6
        node* split(node*& x, SIZE_T k) {
7ca
            if (k <= 0 or !x) return NULL;</pre>
6d0
            node* ret = new node():
386
            if (!x->l \text{ and } !x->r) x->cnt -= k, ret->cnt += k;
4e6
85e
                if (k \le count(x->1)) ret->1 = split(x->1, k):
                else {
4e6
06f
                     ret->r = split(x->r, k - count(x->1));
                     swap(x->1, ret->1);
cfd
63b
                }
674
                ret ->update(), x ->update();
379
            }
d5b
            if (!x->cnt) delete x, x = NULL;
edf
            return ret;
f18
        }
02b
        void split(SIZE_T k, sms& s) { // pega os 'k' menores
e63
            s.clear():
            s.root = split(root, min(k, size()));
6e5
e3c
            s.N = N:
9a6
        }
        // pega os menores que 'k'
        void split_val(T k, sms& s) { split(order_of_key(k), s); }
131
2d2 };
```

# 6.34 SQRT Tree

```
// RMQ em O(log log n) com O(n log log n) pra buildar
// Funciona com qualquer operacao associativa
// Tao rapido quanto a sparse table, mas usa menos memoria
// (log log (1e9) < 5, entao a query eh praticamente O(1))
// build - O(n log log n)
// query - O(log log n)
97a namespace sqrtTree {
052
        int n, *v;
        int pref[4][MAX], sulf[4][MAX], getl[4][MAX], entre[4][MAX],
   sz[4]:
5f7
        int op(int a, int b) { return min(a, b); }
        inline int getblk(int p, int i) { return (i-getl[p][i])/sz[p];
c72
  }
2c6
        void build(int p, int l, int r) {
bc8
            if (1+1 >= r) return;
368
            for (int i = 1; i <= r; i++) getl[p][i] = 1;</pre>
```

```
f16
            for (int L = 1; L \le r; L += sz[p]) {
191
                int R = min(L+sz[p]-1, r);
89 c
                pref[p][L] = v[L], sulf[p][R] = v[R];
59f
                for (int i = L+1; i <= R; i++) pref[p][i] =</pre>
   op(pref[p][i-1], v[i]);
                for (int i = R-1; i >= L; i--) sulf[p][i] = op(v[i],
d9a
   sulf[p][i+1]):
221
                build(p+1, L, R);
c7b
695
            for (int i = 0; i <= sz[p]; i++) {</pre>
ca5
                int at = entre[p][l+i*sz[p]+i] = sulf[p][l+i*sz[p]];
759
                for (int j = i+1; j <= sz[p]; j++)</pre>
   entre[p][1+i*sz[p]+j] = at =
23a
                         op(at, sulf[p][1+j*sz[p]]);
c51
            }
        }
861
        void build(int n2, int* v2) {
8b0
            n = n2, v = v2:
680
            for (int p = 0; p < 4; p++) sz[p] = n2 = sqrt(n2);
44c
c50
            build(0, 0, n-1);
        }
940
9e3
        int query(int 1, int r) {
792
            if (1+1 >= r) return 1 == r ? v[1] : op(v[1], v[r]);
1ba
            int p = 0;
            while (getblk(p, 1) == getblk(p, r)) p++;
4ba
9e4
            int ans = sulf[p][1], a = getblk(p, 1)+1, b = getblk(p, 1)
   r)-1;
8bf
            if (a <= b) ans = op(ans, entre[p][getl[p][1]+a*sz[p]+b]);</pre>
dea
            return op(ans, pref[p][r]);
589
        }
8ff }
6.35 Treap
// Todas as operacoes custam
```

```
b19
             node *1, *r;
284
             int p, sz;
36d
            T val, mi;
4c7
             node(T v) : 1(NULL), r(NULL), p(rng()), sz(1), val(v),
   mi(v) {}
            void update() {
01e
a26
                 sz = 1:
d6e
                 mi = val:
bd7
                 if (1) sz += 1->sz, mi = min(mi, 1->mi);
                 if (r) sz += r \rightarrow sz, mi = min(mi, r \rightarrow mi);
a 54
660
            }
c1b
        };
bb7
        node* root;
84b
        treap() { root = NULL; }
2d8
        treap(const treap& t) {
465
             throw logic_error("Nao copiar a treap!");
1e9
cec
        \simtreap() {
609
             vector < node *> q = {root};
402
             while (q.size()) {
e5d
                 node* x = q.back(); q.pop_back();
ee9
                 if (!x) continue;
1c7
                 q.push_back(x->1), q.push_back(x->r);
bf0
                 delete x:
            }
653
50e
        }
73c
        int size(node* x) { return x ? x->sz : 0; }
        int size() { return size(root); }
b2b
bcf
        void join(node* 1, node* r, node*& i) { // assume que 1 < r</pre>
986
             if (!l or !r) return void(i = 1 ? 1 : r);
80e
             if (1->p > r->p) join(1->r, r, 1->r), i = 1:
            else join(1, r->1, r->1), i = r;
fa0
bda
            i->update();
671
        }
        void split(node* i, node*& 1, node*& r, T v) {
ece
26a
             if (!i) return void(r = 1 = NULL);
f05
             if (i\rightarrow val < v) split(i\rightarrow r, i\rightarrow r, r, v), l = i;
             else split(i \rightarrow 1, 1, i \rightarrow 1, v), r = i;
807
            i->update();
bda
2cd
        void split_leg(node* i, node*& 1, node*& r, T v) {
3fc
26a
            if (!i) return void(r = 1 = NULL);
181
             if (i->val <= v) split_leq(i->r, i->r, r, v), l = i;
58f
             else split_leq(i->1, 1, i->1, v), r = i;
```

```
bda
            i->update();
70f
        }
e13
        int count(node* i, T v) {
6b4
            if (!i) return 0;
352
            if (i->val == v) return 1;
0.68
            if (v < i->val) return count(i->1, v);
4d0
            return count(i->r, v);
5e6
        }
26d
        void index_split(node* i, node*& 1, node*& r, int v, int key =
   0) {
26a
            if (!i) return void(r = 1 = NULL);
c10
            if (key + size(i->1) < v) index_split(i->r, i->r, r, v,
   kev+size(i->1)+1). 1 = i:
            else index_split(i->1, l, i->1, v, key), r = i;
e5a
bda
            i->update();
       }
ccf
a1f
        int count(T v) {
e06
            return count(root, v);
980
        }
c27
        void insert(T v) {
980
            if (count(v)) return;
031
            node *L, *R;
d42
            split(root, L, R, v);
585
            node* at = new node(v);
59f
            join(L, at, L);
a28
            join(L, R, root);
37 c
        }
26b
        void erase(T v) {
df9
            node *L, *M, *R;
b6b
            split_leq(root, M, R, v), split(M, L, M, v);
            if (M) delete M:
f17
f38
            M = NULL;
a28
            join(L, R, root);
b92
       }
        void meld(treap& t) { // segmented merge
e77
4a6
            node *L = root, *R = t.root;
950
            root = NULL;
6b1
            while (L or R) {
fe2
                if (!L or (L and R and L->mi > R->mi)) std::swap(L, R);
5e1
                if (!R) join(root, L, root), L = NULL;
                else if (L->mi == R->mi) {
3c9
a76
                    node* LL;
439
                    split(L, LL, L, R->mi+1);
359
                    delete LL;
2a3
                } else {
a76
                    node* LL;
537
                    split(L, LL, L, R->mi);
```

# 6.36 Treap Implicita

```
// Todas as operacoes custam
// O(log(n)) com alta probabilidade
878 mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
aa1 template < typename T > struct treap {
        struct node {
3c9
b19
            node *1, *r;
284
            int p, sz;
875
            T val, sub, lazy;
aa6
            bool rev;
            node(T v) : l(NULL), r(NULL), p(rng()), sz(1), val(v),
   sub(v), lazy(0), rev(0) {}
            void prop() {
a9c
0ec
                 if (lazy) {
924
                     val += lazy, sub += lazy*sz;
b87
                     if (1) 1->lazy += lazy;
                     if (r) r->lazy += lazy;
d3b
                 }
cea
                 if (rev) {
1bb
e4f
                     swap(1, r);
dc8
                     if (1) 1->rev ^= 1;
f2f
                     if (r) r->rev ^= 1;
                 }
3e5
a32
                 lazy = 0, rev = 0;
            }
ca6
01e
            void update() {
0 c 3
                 sz = 1, sub = val;
a09
                 if (1) 1->prop(), sz += 1->sz, sub += 1->sub;
                 if (r) r \rightarrow prop(), sz += r \rightarrow sz, sub += r \rightarrow sub;
095
360
            }
d37
        };
bb7
        node* root;
84b
        treap() { root = NULL; }
```

```
2d8
        treap(const treap& t) {
465
             throw logic_error("Nao copiar a treap!");
1e9
        }
cec
        \simtreap() {
609
             vector < node *> q = {root};
402
             while (q.size()) {
e5d
                 node* x = q.back(); q.pop_back();
                 if (!x) continue;
ee9
1c7
                 q.push_back(x->1), q.push_back(x->r);
                 delete x;
bf0
653
            }
50e
        }
73c
        int size(node* x) { return x ? x->sz : 0; }
b2b
        int size() { return size(root); }
bcf
        void join(node* 1, node* r, node*& i) { // assume que 1 < r</pre>
986
             if (!1 or !r) return void(i = 1 ? 1 : r);
161
             1->prop(), r->prop();
             if (1->p > r->p) join(1->r, r, 1->r), i = 1;
80e
fa0
             else join(1, r->1, r->1), i = r;
bda
             i->update();
b57
        void split(node* i, node*& 1, node*& r, int v, int key = 0) {
a20
26a
             if (!i) return void(r = 1 = NULL);
c89
             i->prop();
5bd
             if (\text{key} + \text{size}(i\rightarrow 1) < v) split(i\rightarrow r, i\rightarrow r, r, v,
    key+size(i->1)+1), 1 = i;
219
             else split(i->1, l, i->1, v, key), r = i;
bda
             i->update();
d37
        }
231
        void push_back(T v) {
2e0
             node* i = new node(v);
7ab
             join(root, i, root);
46b
        }
        T query(int 1, int r) {
b7a
df9
             node *L, *M, *R;
dca
             split(root, M, R, r+1), split(M, L, M, 1);
d43
             T ans = M -> sub;
69d
             join(L, M, M), join(M, R, root);
ba7
             return ans;
1f7
        }
41f
        void update(int 1, int r, T s) {
df9
             node *L, *M, *R;
dca
             split(root, M, R, r+1), split(M, L, M, 1);
8f6
             M \rightarrow lazy += s;
69d
             join(L, M, M), join(M, R, root);
29f
        }
```

## 6.37 Treap Persistent Implicita

```
// Todas as operacoes custam
// O(log(n)) com alta probabilidade
6cf mt19937_64 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
3c9 struct node {
b19
        node *1. *r:
       ll sz, val, sub;
304
       node(11 v) : 1(NULL), r(NULL), sz(1), val(v), sub(v) {}
        node(node* x) : 1(x->1), r(x->r), sz(x->sz), val(x->val),
   sub(x->sub) {}
       void update() {
01e
0c3
           sz = 1, sub = val;
77e
          if (1) sz += 1->sz, sub += 1->sub;
d6e
           if (r) sz += r->sz, sub += r->sub:
124
            sub %= MOD;
       }
472
95f };
bc9 ll size(node* x) { return x ? x->sz : 0; }
761 void update(node* x) { if (x) x->update(); }
828 node* copy(node* x) { return x ? new node(x) : NULL; }
b02 node* join(node* 1, node* r) {
e1f
        if (!1 or !r) return 1 ? copy(1) : copy(r);
48b
        node* ret;
49f
        if (rng() % (size(l) + size(r)) < size(l)) {</pre>
7eb
            ret = copy(1);
cc1
            ret->r = join(ret->r, r);
784
        } else {
4c5
            ret = copy(r);
551
            ret ->1 = join(1, ret ->1);
7a0
74f
        return update(ret), ret;
2cc }
```

```
723 void split(node* x, node*& 1, node*& r, 11 \text{ v}, 11 \text{ key} = 0) {
        if (!x) return void(l = r = NULL);
421
        if (key + size(x->1) < v) {
b4b
72f
            1 = copy(x);
d70
            split(1->r, 1->r, r, v, key+size(1->1)+1);
710
        } else {
303
            r = copy(x);
417
            split(r->1, l, r->l, v, key);
3d1
da2
        update(1), update(r);
666 }
f9e vector < node *> treap;
139 void init(const vector<ll>& v) {
bbd
        treap = {NULL};
969
        for (auto i : v) treap[0] = join(treap[0], new node(i));
286 }
6.38 Wavelet Tree
// Usa O(sigma + n log(sigma)) de memoria,
// onde sigma = MAXN - MINN
// Depois do build, o v fica ordenado
// count(i, j, x, y) retorna o numero de elementos de
// v[i, j) que pertencem a [x, y]
// kth(i, j, k) retorna o elemento que estaria
// na poscicao k-1 de v[i, j), se ele fosse ordenado
// sum(i, j, x, y) retorna a soma dos elementos de
// v[i, j) que pertencem a [x, y]
// sumk(i, j, k) retorna a soma dos k-esimos menores
// elementos de v[i, j) (sum(i, j, 1) retorna o menor)
//
// Complexidades:
// build - O(n log(sigma))
// count - O(log(sigma))
// kth - O(log(sigma))
// sum - O(log(sigma))
// sumk - O(log(sigma))
597 int n, v[MAX];
578 vector <int > esq[4*(MAXN-MINN)], pref[4*(MAXN-MINN)];
f8d void build(int b = 0, int e = n, int p = 1, int l = MINN, int r =
   MAXN) {
```

```
int m = (1+r)/2; esq[p].push_back(0); pref[p].push_back(0);
58f
        for (int i = b; i < e; i++) {</pre>
f2f
6b9
            esq[p].push_back(esq[p].back()+(v[i]<=m));</pre>
26f
            pref[p].push_back(pref[p].back()+v[i]);
206
8ce
        if (1 == r) return;
3a7
        int m2 = stable partition(v+b, v+e, [=](int i){return i <=
347
        build(b, m2, 2*p, 1, m), build(m2, e, 2*p+1, m+1, r);
Ofb }
540 int count(int i, int j, int x, int y, int p = 1, int l = MINN, int
   r = MAXN)
2ad
        if (y < 1 \text{ or } r < x) \text{ return } 0;
4db
        if (x <= 1 and r <= y) return j-i;
        int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
ddc
        return count(ei, ej, x, y, 2*p, l, m)+count(i-ei, j-ej, x, y,
   2*p+1, m+1, r);
3cf }
f62 int kth(int i, int j, int k, int p=1, int l = MINN, int r = MAXN) {
        if (1 == r) return 1;
Зсе
        int m = (1+r)/2, ei = esq[p][i], ej = esq[p][j];
ddc
585
        if (k <= ej-ei) return kth(ei, ej, k, 2*p, 1, m);</pre>
28b
        return kth(i-ei, j-ej, k-(ej-ei), 2*p+1, m+1, r);
8b6 }
f2c int sum(int i, int j, int x, int y, int p = 1, int l = MINN, int r
   = MAXN)
        if (y < 1 \text{ or } r < x) return 0;
2ad
        if (x <= 1 and r <= y) return pref[p][j]-pref[p][i];</pre>
ddc
        int m = (1+r)/2, ei = esq[p][i], ei = esq[p][i];
        return sum(ei, ej, x, y, 2*p, 1, m) + sum(i-ei, j-ej, x, y,
   2*p+1. m+1. r):
b6d }
b84 int sumk(int i, int j, int k, int p = 1, int l = MINN, int r =
   MAXN) {
8a1
        if (1 == r) return 1*k:
ddc
        int m = (1+r)/2, ei = esq[p][i], ei = esq[p][i];
        if (k <= ej-ei) return sumk(ei, ej, k, 2*p, 1, m);</pre>
50c
        return pref[2*p][ej]-pref[2*p][ei]+sumk(i-ei, j-ej, k-(ej-ei),
4c9
   2*p+1, m+1, r);
940 }
```

# 7 Strings

#### 7.1 Aho-corasick

```
// guery retorna o somatorio do numero de matches de
// todas as stringuinhas na stringona
//
// insert - O(|s| log(SIGMA))
// build - O(N), onde N = somatorio dos tamanhos das strings
// query - O(|s|)
eal namespace aho {
807
        map < char , int > to[MAX];
        int link[MAX], idx, term[MAX], exit[MAX], sobe[MAX];
c87
        void insert(string& s) {
bfc
            int at = 0:
05e
b4f
            for (char c : s) {
b68
                auto it = to[at].find(c);
1c9
                if (it == to[at].end()) at = to[at][c] = ++idx;
361
                else at = it->second;
ff4
            term[at]++. sobe[at]++:
142
6eb
        }
d41 #warning nao esquece de chamar build() depois de inserir
0a8
        void build() {
            queue < int > q;
26a
537
            q.push(0);
            link[0] = exit[0] = -1;
dff
402
            while (q.size()) {
379
                int i = q.front(); q.pop();
3c4
                for (auto [c, j] : to[i]) {
5da
                     int 1 = link[i];
                     while (1 != -1 and !to[1].count(c)) 1 = link[1];
102
7a5
                     link[i] = 1 == -1 ? 0 : to[1][c];
3ab
                     exit[j] = term[link[j]] ? link[j] : exit[link[j]];
6f2
                     if (exit[j]+1) sobe[j] += sobe[exit[j]];
113
                     q.push(j);
f1d
                }
            }
367
768
        }
bc0
        int query(string& s) {
86d
            int at = 0, ans = 0;
b4f
            for (char c : s){
1ca
                while (at != -1 and !to[at].count(c)) at = link[at];
5b9
                at = at == -1 ? 0 : to[at][c];
```

```
2b1
                ans += sobe[at];
            }
b85
ba7
            return ans;
038
       }
a30 }
7.2 eertree
// Constroi a eertree, caractere a caractere
// Inicializar com a quantidade de caracteres maxima
// size() retorna a quantidade de substrings pal. distintas
// depois de chamar propagate(), cada substring palindromica
// ocorre qt[i] vezes. O propagate() retorna o numero de
// substrings pal. com repeticao
//
// O(n) amortizado, considerando alfabeto O(1)
8eb struct eertree {
7cc
       vector < vector < int >> t;
42e
       int n. last. sz:
745
        vector < int > s, len, link, qt;
d36
        eertree(int N) {
ec8
            t = vector(N+2, vector(26, int()));
            s = len = link = qt = vector < int > (N+2);
cee
cd1
            s[0] = -1:
            link[0] = 1, len[0] = 0, link[1] = 1, len[1] = -1;
288
688
            sz = 2, last = 0, n = 1;
2 b 8
       }
244
        void add(char c) {
692
            s[n++] = c -= 'a';
34f
            while (s[n-len[last]-2] != c) last = link[last];
289
            if (!t[last][c]) {
                int prev = link[last];
dab
                while (s[n-len[prev]-2] != c) prev = link[prev];
553
fb2
                link[sz] = t[prev][c];
3f5
                len[sz] = len[last]+2;
                t[last][c] = sz++;
1f8
            }
f8b
344
            qt[last = t[last][c]]++;
b1d
f17
        int size() { return sz-2; }
2af
        11 propagate() {
b73
            11 \text{ ret} = 0;
```

for (int i = n; i > 1; i--) {

ebb

```
qt[link[i]] += qt[i];
fd3
db5
                 ret += qt[i];
074
            }
edf
             return ret;
        }
ef6
a2e };
7.3 KMP
// matching(s, t) retorna os indices das ocorrencias
// de s em t
// autKMP constroi o automato do KMP
// Complexidades:
// pi - O(n)
// match - 0(n + m)
// construir o automato - O(|sigma|*n)
// n = |padrao| e m = |texto|
ea8 template < typename T > vector < int > pi(T s) {
019
        vector < int > p(s.size());
725
        for (int i = 1, j = 0; i < s.size(); i++) {</pre>
             while (j \text{ and } s[j] != s[i]) j = p[j-1];
a51
973
             if (s[j] == s[i]) j++;
f8c
             p[i] = j;
e0a
        }
74e
        return p;
f50 }
c10 template < typename T > vector < int > matching (T& s, T& t) {
658
         vector<int> p = pi(s), match;
a1b
         for (int i = 0, j = 0; i < t.size(); i++) {</pre>
6be
             while (j \text{ and } s[j] != t[i]) j = p[j-1];
c4d
             if (s[j] == t[i]) j++;
310
             if (j == s.size()) match.push_back(i-j+1), j = p[j-1];
028
        }
ed8
        return match;
c82 }
a2d struct KMPaut : vector<vector<int>> {
47 c
         KMPaut(){}
         KMPaut (string& s) : vector < vector < int >> (26,
    vector < int > (s.size()+1)) {
503
             vector < int > p = pi(s);
04b
             auto& aut = *this;
4fa
             aut[s[0]-'a'][0] = 1;
```

# 7.4 Manacher

```
// manacher recebe um vetor de T e retorna o vetor com tamanho dos
   palindromos
// ret[2*i] = tamanho do maior palindromo centrado em i
// \text{ ret}[2*i+1] = \text{tamanho maior palindromo centrado em i e i+1}
// Complexidades:
// manacher - O(n)
// palindrome - <0(n), 0(1)>
// pal_end - O(n)
28a template < typename T > vector < int > manacher (const T& s) {
18f
        int 1 = 0, r = -1, n = s.size();
        vector < int > d1(n), d2(n);
fc9
603
        for (int i = 0; i < n; i++) {</pre>
821
            int k = i > r ? 1 : min(d1[l+r-i], r-i);
61a
            while (i+k < n && i-k >= 0 && s[i+k] == s[i-k]) k++;
            d1[i] = k--:
61e
9f6
            if (i+k > r) l = i-k, r = i+k;
950
e03
        1 = 0, r = -1;
        for (int i = 0: i < n: i++) {
603
a64
            int k = i > r ? 0 : min(d2[1+r-i+1], r-i+1); k++;
2c6
            while (i+k \le n \&\& i-k \ge 0 \&\& s[i+k-1] == s[i-k]) k++:
            d2[i] = --k;
eaa
26d
            if (i+k-1 > r) l = i-k, r = i+k-1;
4fe
        vector<int> ret(2*n-1):
c41
        for (int i = 0; i < n; i++) ret[2*i] = 2*d1[i]-1;</pre>
e6b
        for (int i = 0; i < n-1; i++) ret[2*i+1] = 2*d2[i+1];
e1d
edf
        return ret;
ebb }
// verifica se a string s[i..j] eh palindromo
cac template < typename T > struct palindrome {
f97
        vector < int > man;
b2d
        palindrome(const T& s) : man(manacher(s)) {}
9d7
        bool query(int i, int j) {
```

```
bad
            return man[i+j] >= j-i+1;
1e7
60c };
// tamanho do maior palindromo que termina em cada posicao
7cb template < typename T > vector < int > pal_end(const T& s) {
e57
        vector < int > ret(s.size());
fde
        palindrome <T> p(s);
d51
        ret[0] = 1;
        for (int i = 1; i < s.size(); i++) {</pre>
88e
a32
            ret[i] = min(ret[i-1]+2, i+1);
6ea
            while (!p.query(i-ret[i]+1, i)) ret[i]--;
78e
edf
        return ret;
8bd }
7.5 Min/max suffix/cyclic shift
// Computa o indice do menor/maior sufixo/cyclic shift
// da string, lexicograficamente
//
// O(n)
016 template < typename T > int max_suffix(T s, bool mi = false) {
        s.push_back(*min_element(s.begin(), s.end())-1);
476
        int ans = 0;
1a4
        for (int i = 1; i < s.size(); i++) {</pre>
88e
eec
            int j = 0;
            while (ans+j < i and s[i+j] == s[ans+j]) j++;
708
7a2
            if (s[i+j] > s[ans+j]) {
b52
                 if (!mi or i != s.size()-2) ans = i;
e51
            } else if (j) i += j-1;
69c
        }
ba7
        return ans;
f2a }
a1a template < typename T> int min_suffix(T s) {
        for (auto& i : s) i *= -1;
76b
09d
        s.push_back(*max_element(s.begin(), s.end())+1);
925
        return max_suffix(s, true);
ec0 }
97c template < typename T > int max_cyclic_shift(T s) {
163
        int n = s.size();
1ad
        for (int i = 0; i < n; i++) s.push_back(s[i]);</pre>
20a
        return max_suffix(s);
```

```
d34 }
08a template < typename T> int min_cyclic_shift(T s) {
        for (auto& i : s) i *= -1;
76b
        return max_cyclic_shift(s);
7be
c7a }
    String Hashing
// Complexidades:
// construtor - O(|s|)
// operator() - 0(1)
878 mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
463 int uniform(int 1, int r) {
a7f
        uniform_int_distribution < int > uid(1, r);
f54
        return uid(rng);
d9e }
9e0 template <int MOD> struct str_hash { // 116fcb
c63
        static int P:
dcf
        vector<ll> h, p;
        str_hash(string s) : h(s.size()), p(s.size()) {
ea8
            p[0] = 1, h[0] = s[0];
7a2
ad7
            for (int i = 1; i < s.size(); i++)</pre>
                p[i] = p[i - 1]*P\%MOD, h[i] = (h[i - 1]*P + s[i])\%MOD;
84 c
1ef
        11 operator()(int 1, int r) { // retorna hash s[1...r]
af7
749
            ll hash = h[r] - (1 ? h[1 - 1] *p[r - 1 + 1]%MOD : 0);
            return hash < 0 ? hash + MOD : hash;</pre>
dfd
3ba
        }
116 }:
217 template < int MOD > int str_hash < MOD > :: P = uniform (256, MOD - 1); //
   1 > |sigma|
     String Hashing - modulo 2<sup>61</sup> - 1
// Quase duas vezes mais lento
// Complexidades:
```

// build - O(|s|)
// operator() - O(1)

#### 798 mt19937 64 rng(chrono::steady\_clock::now().time\_since\_epoch().count()); f89 ll uniform(ll l, ll r) { 969 uniform\_int\_distribution < ll > uid(1, r); f54 return uid(rng); cac } d7d struct str hash { static 11 P; c20 dcf vector<11> h, p; str\_hash(string s) : h(s.size()), p(s.size()) { ea8 7a2 p[0] = 1, h[0] = s[0];ad7 for (int i = 1: i < s.size(): i++)</pre> p[i] = mulmod(p[i - 1], P), h[i] = (mulmod(h[i - 1], P) + s[i])%MOD; 507 } 11 operator()(int 1, int r) { // retorna hash s[1...r] af7 538 ll hash = h[r] - (1 ? mulmod(h[1 - 1], p[r - 1 + 1]) : 0); dfd return hash < 0 ? hash + MOD : hash;</pre> 544 } 148 }: 6c5 ll str\_hash::P = uniform(256, MOD - 1); // 1 > |sigma| 7.8 Suffix Array - O(n log n) // kasai recebe o suffix array e calcula lcp[i], // o lcp entre s[sa[i],...,n-1] e s[sa[i+1],...,n-1] // // Complexidades: // suffix\_array - O(n log(n)) // kasai - O(n) 733 vector <int > suffix\_array(string s) { 117

9d0 const 11 MOD = (111<<61) - 1;

11 m = 11\*h2 + 12\*h1, h = h1\*h2;

const static 11 LOWER = (111<<30) - 1, GET31 = (111<<31) - 1;

 $11 \ 11 = a\&LOWER$ , h1 = a>>30, 12 = b&LOWER, h2 = b>>30;

ans = (ans&MOD) + (ans>>61), ans = (ans&MOD) + (ans>>61);

11 ans = 11\*12 + (h>>1) + ((h&1)<<60) + (m>>31) +

e38 ll mulmod(ll a, ll b) {

((m&GET31) << 30) + 1;

return ans - 1;

410

d54

784

1dd

c0f f98 }

```
b38
        s += "$":
        int n = s.size(), N = max(n, 260);
043
2f3
        vector < int > sa(n), ra(n);
        for(int i = 0; i < n; i++) sa[i] = i, ra[i] = s[i];
29b
        for (int k = 0; k < n; k ? k *= 2 : k++) {
0a2
            vector < int > nsa(sa), nra(n), cnt(N);
5ce
            for(int i = 0; i < n; i++) nsa[i] = (nsa[i]-k+n)%n,
   cnt[ra[i]]++:
            for(int i = 1; i < N; i++) cnt[i] += cnt[i-1];</pre>
4c4
            for(int i = n-1; i+1; i--) sa[--cnt[ra[nsa[i]]] = nsa[i];
            for(int i = 1, r = 0; i < n; i++) nra[sa[i]] = r +=</pre>
   ra[sa[i]] !=
                 ra[sa[i-1]] or ra[(sa[i]+k)%n] != ra[(sa[i-1]+k)%n];
26b
            ra = nra:
            if (ra[sa[n-1]] == n-1) break:
d5e
11e
057
        return vector < int > (sa.begin()+1, sa.end());
ff3 }
481 vector < int > kasai(string s, vector < int > sa) {
232
        int n = s.size(), k = 0;
408
        vector < int > ra(n), lcp(n);
676
        for (int i = 0; i < n; i++) ra[sa[i]] = i;</pre>
740
        for (int i = 0: i < n: i++, k -= !!k) {
            if (ra[i] == n-1) { k = 0; continue; }
199
1de
            int j = sa[ra[i]+1];
            while (i+k < n \text{ and } j+k < n \text{ and } s[i+k] == s[j+k]) k++;
891
d98
            lcp[ra[i]] = k;
a07
        }
5ed
        return lcp;
fbe }
7.9 Suffix Array - O(n)
// Rapidao
// Computa o suffix array em 'sa', o rank em 'rnk'
// e o lcp em 'lcp'
// query(i, j) retorna o LCP entre s[i..n-1] e s[j..n-1]
// Complexidades
// O(n) para construir
// query - 0(1)
```

```
1a5 template < typename T > struct rmq {
517
        vector <T> v;
fcc
        int n; static const int b = 30;
        vector < int > mask, t;
70e
183
        int op(int x, int y) { return v[x] \le v[y] ? x : y; }
        int msb(int x) { return __builtin_clz(1)-__builtin_clz(x); }
ee1
c92
        int small(int r, int sz = b) { return
   r-msb(mask[r]&((1<<sz)-1)); }
6ad
        rmq() {}
43c
        rmq(const\ vector < T > \&\ v_) : v(v_), n(v.size()), mask(n), t(n) {
2e5
            for (int i = 0, at = 0; i < n; \max\{i++\} = at |= 1) {
                at = (at <<1) &((1 << b) -1):
a61
c00
                while (at and op(i-msb(at&-at), i) == i) at ^= at&-at;
c2f
ea4
            for (int i = 0; i < n/b; i++) t[i] = small(b*i+b-1);
            for (int j = 1; (1<<j) <= n/b; j++) for (int i = 0;
39d
   i+(1<< j) <= n/b; i++)
                t[n/b*j+i] = op(t[n/b*(j-1)+i],
ba5
   t[n/b*(j-1)+i+(1<<(j-1))]);
41a
e34
        int index_query(int 1, int r) {
27b
            if (r-l+1 \le b) return small(r, r-l+1);
e80
            int x = 1/b+1, y = r/b-1;
fd3
            if (x > y) return op(small(l+b-1), small(r));
            int j = msb(y-x+1);
a4e
ea3
            int ans = op(small(l+b-1), op(t[n/b*j+x],
   t[n/b*j+y-(1<<j)+1]));
be6
            return op(ans, small(r));
62a
093
        T query(int 1, int r) { return v[index_query(1, r)]; }
bab }:
9d7 struct suffix_array {
ac0
        string s;
1a8
        int n;
5b4
        vector < int > sa, cnt, rnk, lcp;
2de
        rmq<int> RMQ;
d6e
        bool cmp(int a1, int b1, int a2, int b2, int a3=0, int b3=0) {
            return a1 != b1 ? a1 < b1 : (a2 != b2 ? a2 < b2 : a3 < b3);
91d
82d
4a4
        template < typename T > void radix(int* fr, int* to, T* r, int N,
   int k) {
c17
            cnt = vector < int > (k+1, 0);
            for (int i = 0; i < N; i++) cnt[r[fr[i]]]++;</pre>
bac
```

```
703
            for (int i = 1; i <= k; i++) cnt[i] += cnt[i-1];</pre>
                                                                             eb6
                                                                                          if (N\%3==1) for (int i = 0; i < N; i++) sa[i] = sa[i+1];
            for (int i = N-1; i+1; i--) to[--cnt[r[fr[i]]]] = fr[i];
                                                                                     }
000
                                                                             ee6
6f3
        }
d66
        void rec(vector<int>& v, int k) {
                                                                             938
                                                                                      suffix_array(const string& s_) : s(s_), n(s.size()), sa(n+3),
a76
            auto &tmp = rnk, &m0 = lcp;
                                                                                              cnt(n+1), rnk(n), lcp(n-1) {
                                                                             e62
            int N = v.size()-3, sz = (N+2)/3, sz2 = sz+N/3;
3a9
                                                                             9fe
                                                                                          vector < int > v(n+3);
7f8
            vector < int > R(sz2+3);
                                                                             f9b
                                                                                          for (int i = 0: i < n: i++) v[i] = i:
74f
            for (int i = 1, j = 0; j < sz2; i += i%3) R[j++] = i;
                                                                                          radix(&v[0], &rnk[0], &s[0], n, 256);
                                                                             eba
                                                                             e6d
                                                                                          int dif = 1;
            radix(&R[0], &tmp[0], &v[0]+2, sz2, k);
                                                                             830
                                                                                          for (int i = 0: i < n: i++)</pre>
b30
207
            radix(\&tmp[0], \&R[0], \&v[0]+1, sz2, k);
                                                                             419
                                                                                              v[rnk[i]] = dif += (i and s[rnk[i]] != s[rnk[i-1]]);
5f1
            radix(&R[0], &tmp[0], &v[0]+0, sz2, k);
                                                                             7cf
                                                                                          if (n \ge 2) rec(v, dif);
                                                                             fb9
                                                                                          sa.resize(n):
af5
            int dif = 0:
            int 10 = -1, 11 = -1, 12 = -1;
                                                                                          for (int i = 0; i < n; i++) rnk[sa[i]] = i;</pre>
ed9
                                                                             76f
d81
            for (int i = 0; i < sz2; i++) {</pre>
                                                                             892
                                                                                          for (int i = 0, k = 0; i < n; i++, k -= !!k) {
                if (v[tmp[i]] != 10 or v[tmp[i]+1] != 11 or
                                                                             668
                                                                                              if (rnk[i] == n-1) {
8de
                                                                                                  k = 0:
   v[tmp[i]+2] != 12)
                                                                             5a4
                     10 = v[tmp[i]], 11 = v[tmp[i]+1], 12 =
                                                                             5e2
b43
                                                                                                  continue;
                                                                                              }
   v[tmp[i]+2], dif++;
                                                                             9df
                if (tmp[i]%3 == 1) R[tmp[i]/3] = dif;
199
                                                                             39a
                                                                                              int j = sa[rnk[i]+1];
                 else R[tmp[i]/3+sz] = dif;
                                                                             891
                                                                                              while (i+k < n \text{ and } j+k < n \text{ and } s[i+k] == s[j+k]) k++;
1f5
            }
d18
                                                                             825
                                                                                              lcp[rnk[i]] = k;
                                                                             a3e
                                                                                          }
47 f
            if (dif < sz2) {
                                                                             9ff
                                                                                          RMQ = rmq < int > (lcp);
146
                rec(R. dif):
                                                                             9a8
                                                                                     }
                for (int i = 0; i < sz2; i++) R[sa[i]] = i+1;</pre>
746
105
            else\ for\ (int\ i = 0;\ i < sz2;\ i++)\ sa[R[i]-1] = i;
                                                                             588
                                                                                     int query(int i, int j) {
                                                                             d97
                                                                                          if (i == j) return n-i;
6f4
            for (int i = 0, j = 0; j < sz2; i++) if (sa[i] < sz)
                                                                             223
                                                                                          i = rnk[i], j = rnk[j];
   tmp[j++] = 3*sa[i];
                                                                                          return RMQ.query(min(i, j), max(i, j)-1);
                                                                             c3a
7ce
            radix(&tmp[0], &m0[0], &v[0], sz, k);
                                                                             940
74d
            for (int i = 0; i < sz2; i++)
                                                                             71c
                                                                                     pair<int, int> next(int L, int R, int i, char c) {
                                                                                          int l = L, r = R+1;
                 sa[i] = sa[i] < sz ? 3*sa[i]+1 : 3*(sa[i]-sz)+2:
с9е
                                                                             024
                                                                                          while (1 < r) {
                                                                             40c
332
            int at = sz2+sz-1, p = sz-1, p2 = sz2-1;
                                                                             ee4
                                                                                             int m = (1+r)/2:
1c9
            while (p \ge 0 \text{ and } p2 \ge 0) {
                                                                             e7e
                                                                                              if (i+sa[m] >= n or s[i+sa[m]] < c) l = m+1;
                if ((sa[p2]%3==1 and cmp(v[m0[p]], v[sa[p2]],
                                                                             ef3
                                                                                              else r = m;
   R[m0[p]/3],
                                                                             ebe
                     R[sa[p2]/3+sz])) or (sa[p2]%3==2 and cmp(v[m0[p]],
                                                                             575
                                                                                          if (1 == R+1 \text{ or } s[i+sa[1]] > c) \text{ return } \{-1, -1\};
Осе
   v[sa[p2]],
                                                                             eb7
                                                                                          L = 1:
                     v[m0[p]+1], v[sa[p2]+1], R[m0[p]/3+sz],
af6
   R[sa[p2]/3+1]))
                                                                             9e2
                                                                                          1 = L, r = R+1;
                     sa[at--] = sa[p2--];
                                                                                          while (1 < r) {
300
                                                                             40c
cb0
                 else sa[at--] = m0[p--];
                                                                             ee4
                                                                                             int m = (1+r)/2;
            }
                                                                                              if (i+sa[m] >= n or s[i+sa[m]] <= c) l = m+1;</pre>
214
                                                                             1a1
f2b
            while (p >= 0) sa[at--] = m0[p--];
                                                                                              else r = m:
                                                                             ef3
```

```
b5b
56a
            R = 1-1:
e13
            return {L, R};
71b
        }
        // quantas vezes 't' ocorre em 's' - O(|t| log n)
        int count_substr(string& t) {
66d
b2b
            int L = 0, R = n-1:
            for (int i = 0; i < t.size(); i++) {</pre>
c9d
de0
                tie(L, R) = next(L, R, i, t[i]);
                if (L == -1) return 0;
4fc
cff
            }
fbf
            return R-L+1;
aaa
        // exemplo de f que resolve o problema
           https://codeforces.com/edu/course/2/lesson/2/5/practice/contes/t/269656/problem/D
57e
        ll f(ll k) { return k*(k+1)/2: }
       11 dfs(int L, int R, int p) { // dfs na suffix tree chamado em
e68
   pre ordem
            int ext = L != R ? RMQ.query(L, R-1) : n - sa[L];
c54
            // Tem 'ext - p' substrings diferentes que ocorrem 'R-L+1'
               vezes
            // O LCP de todas elas eh 'ext'
f80
            ll ans = (ext-p)*f(R-L+1);
            // L eh terminal, e folha sse L == R
            if (sa[L]+ext == n) L++;
63c
            // se for um SA de varias strings separadas como s#t$u&,
               usar no lugar do if de cima
            // (separadores < 'a', diferentes e inclusive no final)
            // while (L <= R && (sa[L]+ext == n || s[sa[L]+ext] <
               'a')) {
            // L++:
            // }
add
            while (L <= R) {
                int idx = L != R ? RMQ.index_query(L, R-1) : -1;
5a8
                if (idx == -1 or lcp[idx] != ext) idx = R;
5ef
478
                ans += dfs(L, idx, ext);
28d
                L = idx+1;
            }
590
ba7
            return ans;
```

```
e21
       }
        // sum over substrings: computa, para toda substring t
           distinta de s.
        // \sum f(# ocorrencias de t em s) - 0 (n)
        11 sos() { return dfs(0, n-1, 0); }
ca8
6fa }:
```

# 7.10 Suffix Array Dinamico

```
// Mantem o suffix array, lcp e rank de uma string,
// premitindo push_front e pop_front
// O operador [i] return um par com sa[i] e lcp[i]
// lcp[i] tem o lcp entre sa[i] e sa[i-1] (lcp[0] = 0)
// Complexidades:
// Construir sobre uma string de tamanho n: O(n log n)
// push_front e pop_front: O(log n) amortizado
2fe struct dyn_sa {
3c9
        struct node {
1d4
            int sa, lcp;
ed1
            node *1, *r, *p;
f0d
            int sz, mi;
17b
            node(int sa_, int lcp_, node* p_) : sa(sa_), lcp(lcp_),
543
                1(NULL), r(NULL), p(p_), sz(1), mi(lcp) {}
            void update() {
01e
58f
                sz = 1, mi = lcp;
bd7
                if (1) sz += 1->sz, mi = min(mi, 1->mi);
a54
                if (r) sz += r->sz, mi = min(mi, r->mi);
27 c
574
        };
bb7
        node* root;
295
        vector<ll> tag; // tag of a suffix (reversed id)
ac0
        string s; // reversed
cf4
        dyn_sa() : root(NULL) {}
e45
        dyn_sa(string s_) : dyn_sa() {
            reverse(s_.begin(), s_.end());
ae4
519
            for (char c : s_) push_front(c);
2a7
        }
a86
        \sim dvn_sa() {
609
            vector < node *> q = {root};
402
            while (q.size()) {
e5d
                node* x = q.back(); q.pop_back();
```

```
ee9
                 if (!x) continue;
                 q.push_back(x->1), q.push_back(x->r);
1c7
bf0
                 delete x;
653
            }
        }
8c1
73c
        int size(node* x) { return x ? x->sz : 0; }
        int mirror(int i) { return s.size()-1 - i; }
08e
580
        bool cmp(int i, int j) {
a 29
            if (s[i] != s[j]) return s[i] < s[j];</pre>
5b4
            if (i == 0 \text{ or } j == 0) \text{ return } i < j;
988
            return tag[i-1] < tag[j-1];</pre>
9fd
919
        void fix_path(node* x) { while (x) x->update(), x = x->p; }
        void flatten(vector < node * > & v, node * x) {
245
8c8
            if (!x) return;
e96
            flatten(v, x->1);
2a2
            v.push_back(x);
            flatten(v, x->r);
42d
01f
        }
964
        void build(vector<node*>& v, node*& x, node* p, int L, int R,
   11 1, 11 r) {
            if (L > R) return void(x = NULL);
04c
331
            int M = (L+R)/2;
3e3
            11 m = (1+r)/2;
7e5
            x = v[M]:
63e
            x->p = p;
bb3
            tag[x->sa] = m:
            build(v, x->1, x, L, M-1, 1, m-1), build(v, x->r, x, M+1,
ae0
   R, m+1, r);
            x->update();
ca8
a3a
82f
        void fix(node*& x, node* p, ll l, ll r) {
            if (3*max(size(x->1), size(x->r)) \le 2*size(x)) return
7f0
   x->update();
            vector < node *> v;
3d1
            flatten(v, x);
0cc
ea9
            build(v, x, p, 0, v.size()-1, 1, r);
b86
b19
        node* next(node* x) {
728
            if (x->r) {
a91
                 x = x - > r;
347
                 while (x->1) x = x->1;
ea5
                 return x;
e7d
402
            while (x->p \text{ and } x->p->r == x) x = x->p;
137
            return x->p;
```

```
48b
        }
b68
        node* prev(node* x) {
e41
            if (x->1) {
a26
                x = x -> 1;
93c
                 while (x->r) x = x->r;
ea5
                 return x;
9be
            }
6a1
            while (x->p \text{ and } x->p->l == x) x = x->p;
137
            return x->p;
        }
73e
4f7
        int get_lcp(node* x, node* y) {
75a
            if (!x or !v) return 0: // change defaut value here
            if (s[x->sa] != s[y->sa]) return 0;
e51
            if (x->sa == 0 \text{ or } y->sa == 0) \text{ return } 1;
843
4d0
            return 1 + query(mirror(x->sa-1), mirror(y->sa-1));
8d6
        }
        void add_suf(node*& x, node* p, int id, ll l, ll r) {
ad6
91e
            if (!x) {
8e3
                x = new node(id, 0, p);
                 node *prv = prev(x), *nxt = next(x);
8e2
65d
                 int lcp_cur = get_lcp(prv, x), lcp_nxt = get_lcp(x,
   nxt);
                 if (nxt) nxt->lcp = lcp_nxt, fix_path(nxt);
ca3
71f
                 x \rightarrow lcp = lcp_cur;
7b4
                 tag[id] = (1+r)/2;
ca8
                 x->update();
505
                 return;
d0e
4a3
            if (cmp(id, x->sa)) add_suf(x->1, x, id, 1, tag[x->sa]-1);
            else add_suf(x->r, x, id, tag[x->sa]+1, r);
c3a
3db
            fix(x, p, 1, r);
c98
        }
ec2
        void push_front(char c) {
            s += c;
cc7
            tag.push_back(-1);
493
            add_suf(root, NULL, s.size() - 1, 0, 1e18);
05e
1f2
        }
7f3
        void rem_suf(node*& x, int id) {
            if (x->sa != id) {
6cf
864
                if (tag[id] < tag[x->sa]) return rem_suf(x->1, id);
e6f
                 return rem_suf(x->r, id);
            }
2ae
2cf
            node* nxt = next(x);
            if (nxt) nxt->lcp = min(nxt->lcp, x->lcp), fix_path(nxt);
09b
```

```
b20
             node *p = x \rightarrow p, *tmp = x;
             if (!x->1 \text{ or } !x->r) {
f3f
2fd
                 x = x - > 1 ? x - > 1 : x - > r;
753
                 if (x) x->p = p;
696
             } else {
7f7
                 for (tmp = x->1, p = x; tmp->r; tmp = tmp->r) p = tmp;
f2a
                 x->sa = tmp->sa, x->lcp = tmp->lcp:
482
                 if (tmp -> 1) tmp -> 1 -> p = p;
14c
                 if (p->1 == tmp) p->1 = tmp->1;
                 else p - > r = tmp - > 1;
a 94
97 c
             }
b5e
             fix_path(p);
7c3
             delete tmp:
510
        }
15b
        void pop_front() {
             if (!s.size()) return;
abe
342
             s.pop_back();
436
             rem_suf(root, s.size());
c6e
             tag.pop_back();
        }
987
530
        int query(node* x, 11 1, 11 r, 11 a, 11 b) {
             if (!x \text{ or } tag[x->sa] == -1 \text{ or } r < a \text{ or } b < 1) \text{ return}
e51
   s.size();
ef5
             if (a <= l and r <= b) return x->mi;
8eb
             int ans = s.size();
             if (a \le tag[x->sa]  and tag[x->sa] \le b) ans = min(ans,
e1f
   x - > lcp):
             ans = min(ans, query(x->1, 1, tag[x->sa]-1, a, b));
d99
261
             ans = min(ans, query(x->r, tag[x->sa]+1, r, a, b));
ba7
             return ans:
4c8
        int query(int i, int j) { // lcp(s[i..], s[j..])
588
             if (i == j) return s.size() - i:
209
             11 a = tag[mirror(i)], b = tag[mirror(j)];
29e
             int ret = query(root, 0, 1e18, min(a, b)+1, max(a, b));
710
             return ret;
edf
84e
        // optional: get rank[i], sa[i] and lcp[i]
044
        int rank(int i) {
396
             i = mirror(i):
             node* x = root;
52f
7c9
             int ret = 0:
f4c
             while (x) {
33e
                 if (tag[x->sa] < tag[i]) {</pre>
                     ret += size(x->1)+1;
f9d
                     x = x - > r:
a91
```

```
6dc
                 } else x = x - > 1;
a 19
edf
             return ret;
153
        }
649
         pair<int, int> operator[](int i) {
52f
             node* x = root;
31e
             while (1) {
d4d
                 if (i < size(x->1)) x = x->1;
4e6
                 else {
85f
                      i \rightarrow size(x\rightarrow 1):
e03
                      if (!i) return {mirror(x->sa), x->lcp};
040
                      i--, x = x->r:
b9b
                 }
7a2
             }
        }
90c
4c2 };
```

#### 7.11 Suffix Automaton

```
// Automato que aceita os sufixos de uma string
// Todas as funcoes sao lineares
16e namespace sam {
c1a
        int cur, sz, len[2*MAX], link[2*MAX], acc[2*MAX];
0b8
        int nxt[2*MAX][26];
        void add(int c) {
e6a
17a
            int at = cur:
9a6
            len[sz] = len[cur]+1, cur = sz++;
500
            while (at != -1 and !nxt[at][c]) nxt[at][c] = cur, at =
   link[at]:
7ea
            if (at == -1) { link[cur] = 0; return; }
654
            int q = nxt[at][c];
            if (len[q] == len[at]+1) { link[cur] = q; return; }
fd9
31f
            int qq = sz++;
2c3
            len[qq] = len[at]+1, link[qq] = link[q];
9a9
            for (int i = 0; i < 26; i++) nxt[qq][i] = nxt[q][i];
            while (at !=-1 and nxt[at][c] == q) nxt[at][c] = qq, at =
e76
   link[at]:
            link[cur] = link[q] = qq;
8b8
61a
        }
94 e
        void build(string& s) {
889
            cur = 0, sz = 0, len[0] = 0, link[0] = -1, sz++;
9fe
            for (auto i : s) add(i-'a');
17a
            int at = cur;
121
            while (at) acc[at] = 1, at = link[at];
```

```
0e7
        }
                                                                             // T.insert(s) - 0(|s|*sigma)
        // coisas que da pra fazer:
                                                                             // T.erase(s) - O(|s|)
                                                                             // T.find(s) retorna a posicao, 0 se nao achar - O(|s|)
28c
        11 distinct_substrings() {
                                                                             // T.count_pref(s) numero de strings que possuem s como prefixo -
04b
            11 \text{ ans} = 0:
                                                                                O(|s|)
a1e
            for (int i = 1; i < sz; i++) ans += len[i] - len[link[i]];</pre>
                                                                             //
ba7
0d7
                                                                             // Nao funciona para string vazia
        string longest_common_substring(string& S, string& T) {
a6c
419
                                                                             ab5 struct trie {
            build(S):
            int at = 0, 1 = 0, ans = 0, pos = -1;
                                                                                     vector < vector < int >> to;
111
                                                                             e1a
d59
            for (int i = 0; i < T.size(); i++) {</pre>
                                                                             450
                                                                                     vector<int> end, pref;
f2c
                while (at and !nxt[at][T[i]-'a']) at = link[at]. 1 =
                                                                                     int sigma: char norm:
                                                                             af0
   len[at];
                                                                             bb1
                                                                                     trie(int sigma_=26, char norm_='a') : sigma(sigma_),
efa
                if (nxt[at][T[i]-'a']) at = nxt[at][T[i]-'a'], l++;
                                                                                norm(norm) {
                else at = 0, 1 = 0;
749
                                                                             58a
                                                                                         to = {vector < int > (sigma)};
                if (1 > ans) ans = 1, pos = i;
                                                                                         end = \{0\}, pref = \{0\};
a1a
                                                                             86e
2b3
                                                                                     }
                                                                             fe1
20f
            return T.substr(pos-ans+1, ans);
                                                                             64e
                                                                                     void insert(string s) {
930
        }
                                                                             c67
                                                                                         int x = 0:
46e
        11 dp [2*MAX];
                                                                             7e7
                                                                                         for(auto c : s) {
        11 paths(int i) {
                                                                             800
                                                                                             int &nxt = to[x][c-norm];
455
2a8
            auto& x = dp[i];
                                                                             dd7
                                                                                             if(!nxt) {
            if (x) return x;
                                                                                                  nxt = to.size();
dee
                                                                             0aa
483
                                                                             526
                                                                                                  to.push_back(vector<int>(sigma));
71c
            for (int j = 0; j < 26; j++) if (nxt[i][j]) x +=
                                                                             770
                                                                                                  end.push_back(0), pref.push_back(0);
   paths(nxt[i][j]);
                                                                             933
                                                                                             }
ea5
            return x;
                                                                             827
                                                                                             x = nxt, pref[x]++;
d88
                                                                             34 c
                                                                                         }
        void kth_substring(int k, int at=0) { // k=1 : menor substring
                                                                                         end[x]++;
105
                                                                             e4e
   lexicog.
                                                                             e6b
9d2
            for (int i = 0; i < 26; i++) if (k and nxt[at][i]) {
                                                                             6b2
                                                                                     void erase(string s) {
                if (paths(nxt[at][i]) >= k) {
d58
                                                                             c67
                                                                                         int x = 0:
d02
                     cout << char('a'+i):
                                                                             b4f
                                                                                         for(char c : s) {
                     kth_substring(k-1, nxt[at][i]);
                                                                                             int &nxt = to[x][c-norm];
c43
                                                                             800
505
                     return;
                                                                             10c
                                                                                             x = nxt, pref[x] --;
                                                                                             if(!pref[x]) nxt = 0;
69a
                }
                                                                             d8e
5f4
                k -= paths(nxt[at][i]);
                                                                             885
                                                                                         }
ef6
            }
                                                                                         end[x]--;
                                                                             bf0
a13
                                                                             ddd
                                                                                     }
c37 };
                                                                             aee
                                                                                     int find(string s) {
                                                                                         int x = 0;
                                                                             c67
                                                                             7e7
                                                                                         for(auto c : s) {
                                                                             2ec
                                                                                             x = to[x][c-norm];
7.12
      Trie
                                                                                             if(!x) return 0;
                                                                             a66
                                                                                         }
                                                                             e12
// trie T() constroi uma trie para o alfabeto das letras minusculas
                                                                             ea5
                                                                                         return x;
// trie T(tamanho do alfabeto, menor caracter) tambem pode ser usado
```

```
e77
839
        int count_pref(string s) {
e2f
            return pref[find(s)];
        }
f40
979 };
7.13 Z
// z[i] = lcp(s, s[i..n))
// Complexidades:
//z - 0(|s|)
// \text{ match - } O(|s| + |p|)
a19 vector < int > get_z(string s) {
163
        int n = s.size();
2b1
        vector < int > z(n, 0);
       int 1 = 0, r = 0;
fae
        for (int i = 1; i < n; i++) {</pre>
6f5
0af
            if (i \le r) z[i] = min(r - i + 1, z[i - 1]);
            while (i + z[i] < n \text{ and } s[z[i]] == s[i + z[i]]) z[i]++;
457
            if (i + z[i] - 1 > r) l = i, r = i + z[i] - 1:
65e
5cd
070
        return z;
74a }
```

## Extra

## 8.1 hash.sh

```
# Para usar (hash das linhas [11, 12]):
# bash hash.sh arquivo.cpp 11 12
sed -n $2','$3' p' $1 | sed '/^#w/d' | cpp -dD -P -fpreprocessed | tr
   -d '[:space:]' | md5sum | cut -c-6
8.2 vimrc
d79 set ts=4 sw=4 mouse=a nu ai si undofile
3b3 function Hash(1)
        return system("sed '/^#/d' | cpp -dD -P -fpreprocessed | tr -d
   '[:space:]' | md5sum", a:1)
Obe endfunction
db5 function PrintHash() range
        for i in range(a:firstline, a:lastline)
fb0
            let k = i
ccc
            let l = getline(i)
82d
            for j in range(len(1))
ca3
                if 1[i] == "{}"[1]
860
                    call cursor(i, j+1)
                    let k = searchpair('{', '', '}', 'b')
016
e96
                endif
bf9
            endfor
269
            echo Hash(join("\n", getline(k, i)))[0:2] 1
            call cursor(i, len(1))
139
bf9
        endfor
Obe endfunction
6b0 vmap <C-H> :call PrintHash() <CR>
3b5 "" }
8.3 template.cpp
```

```
#include <bits/stdc++.h>
using namespace std;
#define _ ios_base::sync_with_stdio(0);cin.tie(0);
#define endl '\n'
```

```
typedef long long 11;
const int INF = 0x3f3f3f3f;
const 11 LINF = 0x3f3f3f3f3f3f3f3f11;
int main() { _
    exit(0);
8.4 pragma.cpp
// Otimizacoes agressivas, pode deixar mais rapido ou mais devagar
#pragma GCC optimize("Ofast")
// Auto explicativo
#pragma GCC optimize("unroll-loops")
// Vetorizacao
#pragma GCC target("avx2")
// Para operacoes com bits
#pragma GCC target("bmi,bmi2,popcnt,lzcnt")
8.5 stress.sh
P=a
make ${P} ${P}2 gen || exit 1
for ((i = 1; ; i++)) do
    ./gen $i > in
    ./${P} < in > out
    ./${P}2 < in > out2
    if (! cmp -s out out2) then
        echo "--> entrada:"
        cat in
        echo "--> saida1:"
        cat out
        echo "--> saida2:"
        cat out2
        break;
    fi
    echo $i
done
```

## 8.6 makefile

```
CXX = g++
CXXFLAGS = -fsanitize=address, undefined -fno-omit-frame-pointer -g
    -Wall -Wshadow -std=c++17 -Wno-unused-result -Wno-sign-compare
   -Wno-char-subscripts #-fuse-ld=gold
8.7 fastIO.cpp
int read_int() {
    bool minus = false;
    int result = 0;
    char ch;
    ch = getchar();
    while (1) {
        if (ch == '-') break;
        if (ch >= '0' && ch <= '9') break;
        ch = getchar();
    if (ch == '-') minus = true;
    else result = ch-'0';
    while (1) {
        ch = getchar();
        if (ch < '0' || ch > '9') break:
        result = result *10 + (ch - '0');
    if (minus) return -result;
    else return result;
}
8.8 rand.cpp
mt19937 rng((int)
   chrono::steady_clock::now().time_since_epoch().count());
int uniform(int 1, int r){
    uniform_int_distribution < int > uid(1, r);
    return uid(rng);
}
8.9 timer.cpp
```

```
// timer T; T() -> retorna o tempo em ms desde que declarou
using namespace chrono;
struct timer : high_resolution_clock {
    const time_point start;
    timer(): start(now()) {}
    int operator()() {
        return duration_cast < milliseconds > (now() - start).count();
    }
};
8.10 debug.cpp
void debug_out(string s, int line) { cerr << endl; }</pre>
template < typename H, typename ... T>
void debug_out(string s, int line, H h, T... t) {
    if (s[0] != ',') cerr << "Line(" << line << ") ";</pre>
    do { cerr << s[0]; s = s.substr(1);</pre>
    } while (s.size() and s[0] != ',');
    cerr << " = " << h;
    debug_out(s, line, t...);
#ifdef DEBUG
#define debug(...) debug_out(#__VA_ARGS__, __LINE__, __VA_ARGS__)
#else
#define debug(...) 42
```

#endif