

Aggregate Method:

- Let's say c_i be the cost of i^{th} operation, It is given that, $c_i = i$ if i is an exact power of 2 and $c_i = 1$ otherwise. This implies we will have a cost like this.

Operation : i	1	2	3	4	5	6	7	8	9
Cost : c_i	1	2	1	4	1	1	1	8	1

$$\sum_{i=1}^n c_i = \sum_{j=0}^{\log n} 2^j + \sum_{j=1}^n 1 - \sum_{j=1}^{\log n - 1} 1$$

$$= \sum_{j=1}^{\log n} 2^j + n - (\log n - 1)$$

$$\leq n + \sum_{j=1}^{\log n} 2^j \text{ for all } n > 0.$$

$$\leq n + (2n - 1)$$

$$< 3n \Rightarrow O(3n)$$

Average Cost of Operation = Total Cost / No of Operations = $3n / n = 3$.

So, the amortized cost per operation = $O(3) = O(1)$.

Accounting Analysis Method:

Operation : i	1	2	3	4	5	6	7	8	9
Cost : c_i	1	2	1	4	1	1	1	8	1
Amortized Cost	3	3	3	3	3	3	3	3	3

Let's assume amortized cost be 3 and we need to prove that 3 is actually amortized value.

$$\text{This implies amortized Cost} = \sum_{i=1}^n c_{i(\text{amortized})} = 3n$$

The summation of cost for each operation of i from time to time cannot be greater than the value of the sum of 3 as the value of i increases.

$$\sum_{i=1}^n c_{i(\text{amortized})} \geq \sum_{i=1}^n c_i$$

This implies the amortized cost the operation become $O(3) = O(1)$.