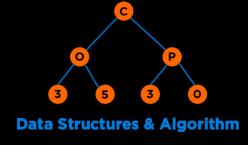
Lists, Stacks and Queues



Categories of Data Structures

Linear Ordered

Non-linear Ordered

Not Ordered

Lists

Trees

Sets

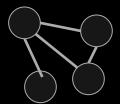
Stacks

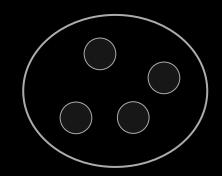
Graphs

Tables/Maps

Queues









Agenda

- Data Structures
- Abstract Data Types
- Lists
 - Array Implementation
 - Linked List Implementation and its types
 - Lists in C++
- Stacks
 - Array Based and Linked List Based
 - Stacks in C++
 - Use Cases
- Queues
 - Array Based and Linked List Based
 - Queue in C++
 - Use cases



Data Structures

Data Structures

A data structure is a way to store and organize data

- Mathematical or Logical models (Abstract Data Types)
 - List: Store, Read, Modify
- Implementation (Concrete)
 - Arrays, Linked List, Vector, ArrayList, List



Abstract Data Types



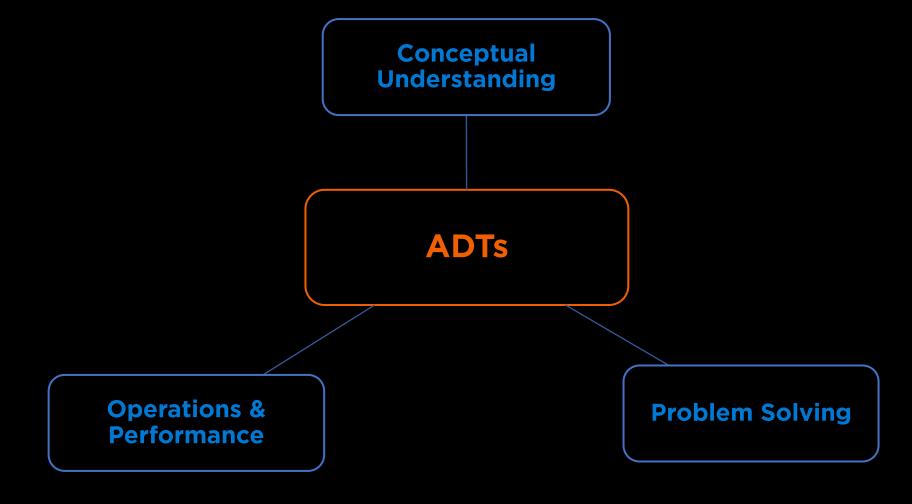
Abstract Data Types

- Class of objects whose logical behavior is defined by a set of values and a set of operations
 - Define data (properties) and operations (behavior)
 - Don't care about implementation





Abstract Data Types (ADTs)



Lists



List

- Ordered Collection of Data (Ordered = Position)
- Elements have some position
- Linear Structure
- Can have some size or grow/shrink
- No limit on nature of elements

List - Characteristics

Data

Operations

List - Characteristics

Data

- Items
- Number of Items (Size)
- Capacity

Operations

- Read/Write an element
- Add or remove an element
- Find an element
- Count
- Traverse the list (Printing)



| Characteristics | | |
|-----------------|--|--|
| Operations | | |
| Performance | | |
| | | |
| | | |
| | | |
| Benefits | | |
| Drawbacks | | |

| Characteristics | Fixed Size Stores Similar Elements Contiguous Indices Elements are stored contiguously in memory Allows Random Access | | |
|-----------------|---|--|--|
| Operations | | | |
| Performance | | | |
| | | | |
| | | | |
| | | | |
| Benefits | | | |
| Drawbacks | | | |

| Characteristics | Fixed Size Stores Similar Elements Contiguous Indices Elements are stored contiguously in memory Allows Random Access | | |
|-------------------|---|-----|--------|
| Operations | Adding - Beginning, Middle, EndRemoval - Beginning, Middle, End | | |
| Performance | | Add | Remove |
| | Beginning | | |
| | End | | |
| | Middle | | |
| Benefits | | | |
| Drawbacks | | | |

| Characteristics | Fixed Size Stores Similar Elements Contiguous Indices Elements are stored contiguously in memory Allows Random Access | | |
|-----------------|---|------|--------|
| Operations | Adding - Beginning, Middle, EndRemoval - Beginning, Middle, End | | |
| Performance | | Add | Remove |
| | Beginning | 0(n) | 0(n) |
| | End | 0(1) | 0(1) |
| | Middle | 0(n) | 0(n) |
| Benefits | | | |
| Drawbacks | | | |

| Characteristics | Fixed Size Stores Similar Elements Contiguous Indices Elements are stored contiguously in memory Allows Random Access | | |
|-----------------|---|--------------------|--------------------|
| Operations | Adding - Beginning, Middle, EndRemoval - Beginning, Middle, End | | |
| Performance | | Add | Remove |
| | Beginning | O(n) | 0(n) |
| | End | 0(1) | 0(1) |
| | Middle | 0(n) | 0(n) |
| Benefits | ■ Constant Access | Time, arr[i] = arr | + (i*sizeOf(type)) |
| Drawbacks | | | |

| Characteristics | Fixed Size Stores Similar Elements Contiguous Indices Elements are stored contiguously in memory Allows Random Access | | |
|-------------------|---|------|--------|
| Operations | Adding - Beginning, Middle, EndRemoval - Beginning, Middle, End | | |
| Performance | | Add | Remove |
| | Beginning | 0(n) | 0(n) |
| | End | 0(1) | 0(1) |
| | Middle | 0(n) | 0(n) |
| Benefits | ■ Constant Access Time, arr[i] = arr + (i*sizeOf(type)) | | |
| Drawbacks | Expensive for adding/removing elements from front | | |

```
class Node
{
    public:
        Node *next;
        int data;
};
```

```
int main ()
  Node* obj = new Node;
  obj -> data = 10;
  obj -> next = nullptr;
  obj -> next = new Node;
  obj -> next -> data = 20;
  obj -> next -> next = nullptr;
  obj -> next -> next = new Node;
  obj -> next -> next -> data = 30;
  obj -> next -> next -> next = nullptr;
  return 0;
```

```
class Node
{
    public:
        Node *next;
        int data;
};
```

```
int main ()
  Node* obj = new Node;
  obj -> data = 10;
  obj -> next = nullptr;
  obj -> next = new Node;
  obj -> next -> data = 20;
  obj -> next -> next = nullptr;
  obj -> next -> next = new Node;
  obj -> next -> next -> data = 30;
  obj -> next -> next -> next = nullptr;
  return 0;
```

Problem:



```
class Node
                                          int main ()
    public:
                                              Node* obj = new Node(10, nullptr);
        Node *next;
                                              obj = new Node(20, obj);
                                              obj = new Node(30, obj);
        int data;
        Node(int d, Node* n);
                                              cout << "Size = "<< obj -> size();
                                              return 0;
    private:
        int size();
};
Node::Node(int d, Node* n)
    this -> data = d;
    this \rightarrow next = n;
int Node::size()
    if(this == nullptr)
        return 0;
    else
        return 1 + (this->next)->size();
```

```
class Node
    public:
        Node *next;
        int data;
        Node(int d, Node* n);
    private:
        int size();
};
Node::Node(int d, Node* n)
    this -> data = d;
    this -> next = n;
int Node::size()
    if(this == nullptr)
        return 0;
    else
        return 1 + (this->next)->size();
```

```
int main ()
{
    Node* obj = new Node(10, nullptr);
    obj = new Node(20, obj);
    obj = new Node(30, obj);
    cout << "Size = "<< obj -> size();
    return 0;
}
```

Problem:

Too much details



Encapsulate Node: Better Design

```
class List
struct Node
                                     public:
 int data;
                                        int size;
 Node *next;
                                        Node* head;
};
                                        Node* tail;
or
                                     public:
                                        List();
class Node
                                        void push_front(int);
    public:
        Node *next;
        int data;
};
```

Encapsulate Node: Better Design

```
struct Node
{
   int data;
   Node *next;
};

or

class Node
{
   public:
        Node *next;
        int data;
};
```

```
class List
     public:
        int size;
        Node* head;
        Node* tail;
     public:
        List();
        void push front(int);
};
List::List()
    size = 0;
    head = new Node;
    tail = new Node;
    head -> next = tail;
```

```
void List::push front(int x)
    Node* temp = new Node;
    temp -> data = x;
    temp -> next = head->next;
    head -> next = temp;
    size++;
int main()
    List obj;
    obj.push_front(15);
    obj.push front(20);
    obj.push_front(15);
    return 0;
```



Do not expose internal details to the client

```
int main()
{
    List myList;
    myList.push_front(15);
    myList.push_front(20);
    myList.push_front(15);
    return 0;
}
```

| Characteristics | | |
|-----------------|--|--|
| Operations | | |
| Performance | | |
| | | |
| | | |
| | | |
| Benefits | | |
| Drawbacks | | |
| | | |



| Characteristics | Consists of Nodes □ Data □ Pointer to Next Node Stores Similar Elements ■ Elements are linked in memory but stored non-contiguously ■ Does not allow Random Access |
|-----------------|---|
| Operations | |
| Performance | |
| Benefits | |
| Drawbacks | 2 |



| Characteristics | Consists of Nodes Data Pointer to Next Node Stores Similar Elements Elements are linked in memory but stored non-contiguously Does not allow Random Access |
|-----------------|---|
| Operations | Adding (Add) - PushFront(Key), PushBack(Key) Removal (Remove) - PopFront, PopBack Access (Get) - TopFront, TopBack Find(Key), Erase(Key), Empty() AddBefore(Node, Key), AddAfter(Node, Key) |
| Performance | |
| Benefits | |
| Drawbacks | |



| Characteristics | Consists of Nodes Data Pointer to Next Node Stores Similar Elements Elements are linked in memory but stored non-contiguously Does not allow Random Access | | |
|-----------------|---|----------|-----------|
| Operations | Adding (Add) - PushFront(Key), PushBack(Key) Removal (Remove) - PopFront, PopBack Access (Get) - TopFront, TopBack Find(Key), Erase(Key), Empty() AddBefore(Node, Key), AddAfter(Node, Key) | | |
| Performance | PushFront | PushBack | AddBefore |
| | PopFront | PopBack | AddAfter |
| | TopFront | TopBack | |
| | Find | Erase | Empty |
| Benefits | | | |
| Drawbacks | | | |



| Characteristics | Consists of Nodes Data Pointer to Next Node Stores Similar Element Elements are linked in Does not allow Random | nts In memory but stored r | non-contiguously |
|-----------------|---|-------------------------------|------------------|
| Operations | Adding (Add) - PushFront(Key), PushBack(Key) Removal (Remove) - PopFront, PopBack Access (Get) - TopFront, TopBack Find(Key), Erase(Key), Empty() AddBefore(Node, Key), AddAfter(Node, Key) | | |
| Performance | PushFront - 0(1) | PushBack - O(n) | AddBefore - O(n) |
| | PopFront - 0(1) | PopBack - O(n) | AddAfter - O(1) |
| | TopFront - 0(1) | TopBack - O(n) | |
| | Find - O(n) | Erase - O(n) | Empty - 0(1) |
| Benefits | | | |
| Drawbacks | | | |



| Characteristics | Consists of Nodes Data Pointer to Next Node Stores Similar Element Elements are linked in Does not allow Random | nts In memory but stored r | non-contiguously |
|-----------------|---|-------------------------------|------------------|
| Operations | Adding (Add) - PushFront(Key), PushBack(Key) Removal (Remove) - PopFront, PopBack Access (Get) - TopFront, TopBack Find(Key), Erase(Key), Empty() AddBefore(Node, Key), AddAfter(Node, Key) | | |
| Performance | PushFront - 0(1) | PushBack - O(n) | AddBefore - O(n) |
| | PopFront - 0(1) | PopBack - O(n) | AddAfter - O(1) |
| | TopFront - 0(1) | TopBack - O(n) | |
| | Find - O(n) | Erase - O(n) | Empty - 0(1) |
| Benefits | Adding/Removing in front is faster, O(1) | | |
| Drawbacks | | | , and a second |



| Characteristics | Consists of Nodes Data Pointer to Next Node Stores Similar Elements Elements are linked in memory but stored non-contiguously Does not allow Random Access | | | |
|-----------------|---|-----------------|------------------|--|
| Operations | Adding (Add) - PushFront(Key), PushBack(Key) Removal (Remove) - PopFront, PopBack Access (Get) - TopFront, TopBack Find(Key), Erase(Key), Empty() AddBefore(Node, Key), AddAfter(Node, Key) | | | |
| Performance | PushFront - 0(1) | PushBack - O(n) | AddBefore - O(n) | |
| | PopFront - 0(1) | PopBack - O(n) | AddAfter - 0(1) | |
| | TopFront - 0(1) | TopBack - O(n) | | |
| | Find - O(n) | Erase - O(n) | Empty - 0(1) | |
| Benefits | Adding/Removing in front is faster, O(1) | | | |
| Drawbacks | Expensive for random access, O(n)TopBack, PushBack, PopBack and AddBefore are expensive | | | |



List: Single Linked List with Tail

| Characteristics | Consists of Nodes and has a Tail Data Pointer to Next Node Stores Similar Elements Elements are linked in memory but stored non-contiguously Does not allow Random Access | | |
|-----------------|--|-----------------|------------------|
| Operations | Adding - PushFront(Key), PushBack(Key) Removal - PopFront, PopBack Access - TopFront, TopBack Find(Key), Erase(Key), Empty() AddBefore(Node, Key), AddAfter(Node, Key) | | |
| Performance | PushFront - 0(1) | PushBack - 0(1) | AddBefore - O(n) |
| | PopFront - 0(1) | PopBack - O(n) | AddAfter - O(1) |
| | TopFront - 0(1) | TopBack - 0(1) | |
| | Find - O(n) | Erase - O(n) | Empty - O(1) |
| Benefits | Solves the issue of pushing back and getting last element | | |
| Drawbacks | ■ Expensive for PopBack and AddBefore | | |



List: Doubly Linked List with Tail

| Characteristics | Consists of Nodes Data Pointer to Next and Previous Node Stores Similar Elements Elements are linked in memory but stored non-contiguously Does not allow Random Access | | |
|-----------------|--|-----------------|------------------|
| Operations | Adding - PushFront(Key), PushBack(Key) Removal - PopFront, PopBack Access - TopFront, TopBack Find(Key), Erase(Key), Empty() AddBefore(Node, Key), AddAfter(Node, Key) | | |
| Performance | PushFront - 0(1) | PushBack - 0(1) | AddBefore - O(1) |
| | PopFront - 0(1) | PopBack - O(1) | AddAfter - O(1) |
| | TopFront - 0(1) | TopBack - 0(1) | |
| | Find - O(n) | Erase - O(n) | Empty - O(1) |
| Benefits | Solves the issue of PopBack and AddBefore | | |
| Drawbacks | ■ Extra memory | | |

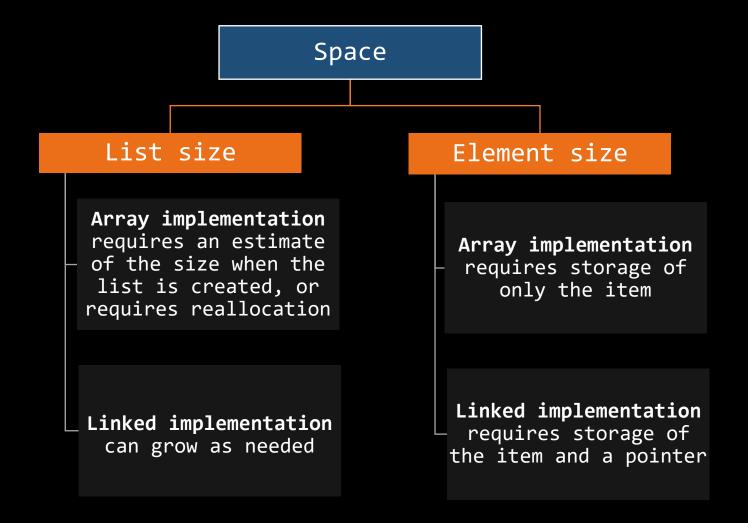


List - Example: Circular Linked List

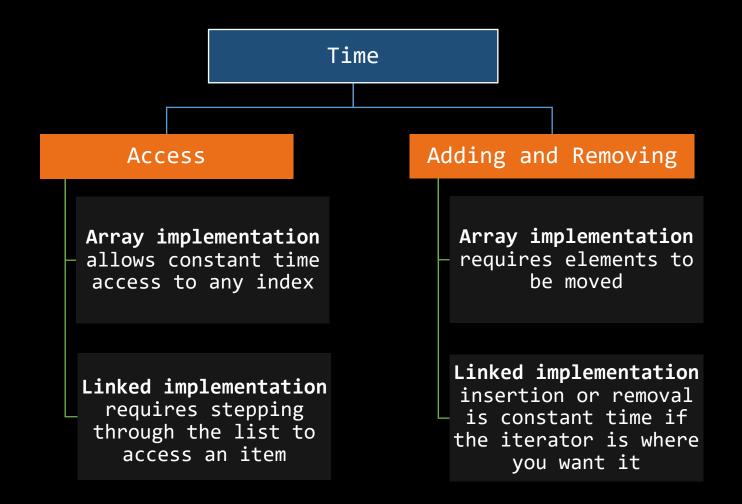
Single Circular

Doubly Circular

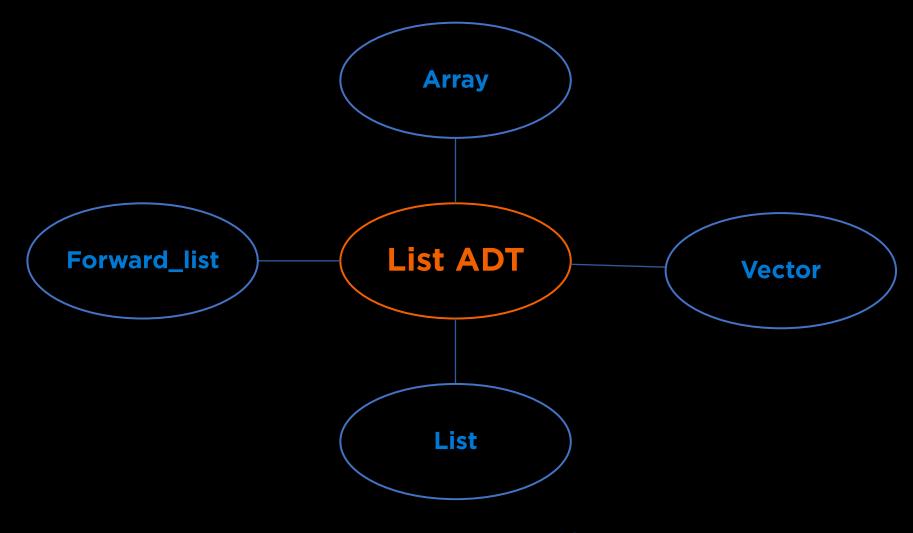
Array vs Linked List: Space



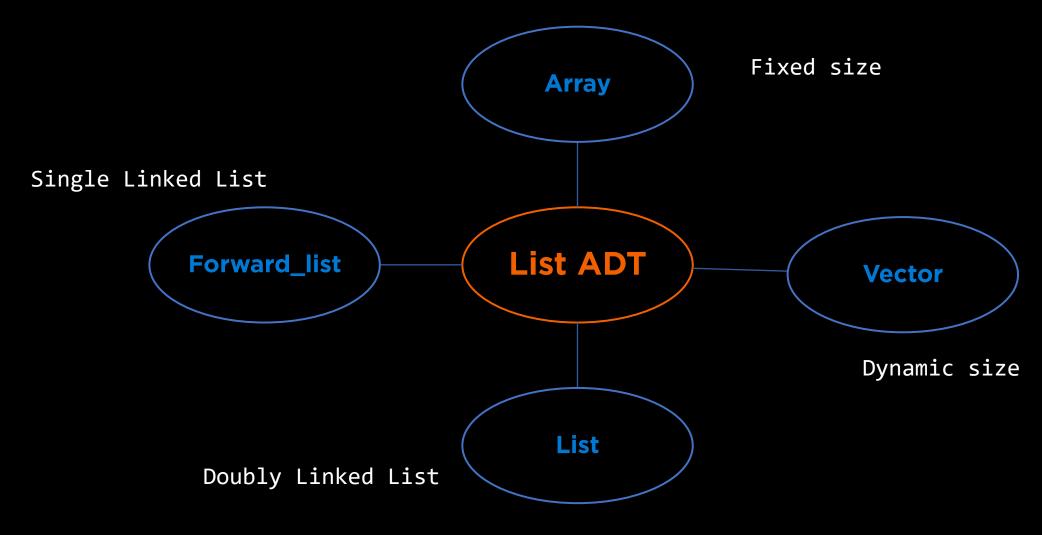
Array vs Linked List: Time



Lists in C++ Standard Template Library



Lists in C++ Standard Template Library



Lists in C++ STL: List

List containers are implemented as doubly-linked lists

| Function | Operation |
|----------------------------------|---|
| Constructor(n, val) | Constructs a container with n elements. Each element is a copy of val (if provided). |
| <pre>void push_front (val)</pre> | Inserts a new element at the beginning of the list, right before its current first element. |

```
std::list<int> mylist (4,100); [100, 100, 100, 100]
mylist.push_front (200); [200, 100, 100, 100]
```



Lists in C++ STL: List

Printing a List

```
01 list<int> mylist (4,100);  // [100, 100, 100, 100]
02 mylist.push_front (200);  // [200, 100, 100, 100, 100]
03 for(list<int>::iterator it = mylist.begin(); it != mylist.end(); ++it)
04 cout << ' ' << *it;</pre>
```

Prints: 200 100 100 100 100

All containers have Iterators!



Iterators

- Variables to keep track of where we are in a data set
- Iterator Class
 - Operators to advance to next data (++)
 - Dereference Operator (*) to access data
 - Operators to compare two iterators (!=)
 - Assignment operator (=)
- Container or the main data structure you are implementing Class
 - begin() methods
 - end() methods



Iterators

| | category | | | properties | valid expressions |
|------------------|---------------|---------|--|--|------------------------------------|
| all categories | | | copy-constructible, copy-assignable and destructible | X b(a); b = a; | |
| | | | Can be incremented | ++a a++ | |
| | | | T | Supports equality/inequality comparisons | a == b a != b |
| | | Forward | Input | Can be dereferenced as an <i>rvalue</i> | *a a->m |
| | | | Output | Can be dereferenced as an <i>Ivalue</i> (only for <i>mutable iterator types</i>) | *a = t *a++ = t |
| Random Access | Bidirectional | | | default-constructible | X a; X() |
| | | | | Multi-pass: neither dereferencing nor incrementing affects dereferenceability | { b=a; *a++; *b; } |
| | | | | Can be decremented | a a *a |
| | | | | Supports arithmetic operators + and - | a + n n + a a - n a - b |
| | | | | Supports inequality comparisons (<, >, <= and >=) between iterators | a < b a > b a <= b a >= b |
| | | | | Supports compound assignment operations += and -= | a += n a -= n |
| | | | | Supports offset dereference operator ([]) | a[n] |

Iterators: Forward

```
1.  // Forward Lists - support forward iterators: https://www.cplusplus.com/reference/forward_list/forward_list/
2.
3.  std::forward_list<int> flist_container {1, 2, 3, 4, 5};
4.
5.  for(auto it = flist_container.begin(); it != flist_container.end(); it++)
6.  {
7.    std::cout << *it << std::endl;
8.    std::cout << *(--it); //This line will throw an error as iterator does not support going backward
9.    std::cout << *(it + 2); //This line will throw an error as iterator does not support random access
10.    break;
11. }</pre>
```

Iterators: Bidirectional

```
1.  // Lists - support bidirectional iterators: https://www.cplusplus.com/reference/list/list/
2.  std::list<int> list_container {1, 2, 3, 4, 5};
3.
4.  for(auto it = list_container.begin(); it != list_container.end(); it++)
5.  {
6.    std::cout << *it << " ";
7.    std::cout << *(--it) << std::endl; //This line is undefined behavior
8.    //std::cout << *(it + 2); //This line will throw an error as iterator does not support random access break;
10. }</pre>
```

Iterators: Random Access

```
// Vectors - support random access iterators: https://cplusplus.com/reference/vector/
2.
3.
      std::vector<int> vector_container {1, 2, 3, 4, 5};
      for(auto it = vector_container.begin(); it != vector_container.end(); it++)
5.
6.
          std::cout << *it << " ";</pre>
          std::cout << *(--it) << " "; //This line is undefined behavior</pre>
8.
          std::cout << *(it + 2);</pre>
9.
10.
          break;
11.
12.
```

Merge Two Sorted Linked Lists of Integers



Merge Two Sorted Linked Lists of Integers

```
Pseudocode:
given ListA, ListB both sorted
create mergedList
place iterA at the head of list A, iterB at the head of list B
itemA is item iterA is pointing to and itemB is item iterB is pointing to
while iterA is not at end of ListA and iterB is not at end of ListB
         if itemA == itemB
                  make a copy of itemA and add it to mergedList
                  move iterA and iterB forward
         else if itemA < itemB
                  make a copy of itemA and add it to mergedList
                   move iterA forward
         else
                  make a copy of itemB and add it to mergedList
                  move iterB forward
```

Merge Two Sorted Linked Lists of Integers

```
given ListA, ListB both sorted
create mergedList
place iterA at the head of list A, iterB at the head of list B
itemA is item iterA is pointing to and itemB is item iterB is pointing to
while iterA is not at end of ListA and iterB is not at end of ListB
         if itemA == itemB
                  make a copy of itemA and add it to mergedList
                  move iterA and iterB forward
         else if itemA < itemB
                  make a copy of itemA and add it to mergedList
                   move iterA forward
         else
                  make a copy of itemB and add it to mergedList
                  move iterB forward
while iterA is not at end of list
         make a copy of itemA and add it to mergedList
         move iterA forward
while iterB is not at end of list
         make a copy of itemB and add it to mergedList
         move iterB forward
```

Pseudocode:

Merge Two Sorted Linked Lists of Integers (Union)

```
int arr1[] = {16,22,77,129};
     int arr2[] = {1,2,7,29,77,155,166};
03
     //Creating Lists from Arrays
     list<int> list1 (arr1, (arr1 + (sizeof(arr1) / sizeof(int))) );
     list<int> list2 (arr2, (arr2 + (sizeof(arr2) / sizeof(int))) );
07
     //Final Merged Sorted List
     list<int> list3;
10
     list<int>::iterator l1 = list1.begin();
12
     list<int>::iterator 12 = list2.begin();
13
     while(l1 != list1.end() && l2 != list2.end())
15
16
         if((*11) < (*12))
17
18
             list3.push_back(*l1);
19
             11++;
20
21
         else
22
         if((*12) < (*11))
23
24
             list3.push_back(*12);
25
             12++;
26
27
         else
28
29
             list3.push back(*11);
30
             11++;
31
             12++;
32
33
```

```
34
35 while(l1 != list1.end())
36 {
37     list3.push_back(*l1);
38     l1++;
39  }
40
41 while(l2 != list2.end())
42 {
43     list3.push_back(*l2);
44     l2++;
45 }
46
```



Median of Elements in a Linked List (3.3d)



Median of Elements in a Linked List (3.3d)

```
01
   //runner technique, tortoise and hare
   float median(Node* head)
02
03
04
       Node* fastptr = head;
05
       Node* slowptr = head;
06
       while (fastptr->next != NULL && fastptr->next->next != NULL)
07
98
           slowptr = slowptr->next;
09
           fastptr = fastptr->next->next;
10
11
       return fastptr->next == NULL ? slowptr->value : (slowptr->value + slowptr->next->value)/2.0;
12
```

Recommended Resources

- http://www.cplusplus.com/reference/stl/
- Linked List Questions on Stepik
- https://stackoverflow.com/questions/5384358/how-does-a-sentinel-node-offerbenefits-over-null
- OpenDSA Ch-9.1-9.7
- **Templated Singly Linked List:** https://cathyatseneca.gitbooks.io/data-structures-and-algorithms/lists/list_declaration.html
- https://www.cplusplus.com/reference/iterator/
- https://www.geeksforgeeks.org/input-iterators-in-cpp/
- https://onlinegdb.com/Nk9XgDjG-

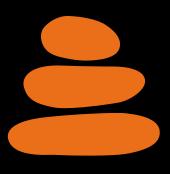
Stacks

Stack

Last in First Out (LIFO)







Stack ADT

Data

- o **Items**
- Number of Items
- o Top

Operations

o push(item): inserts an element

o pop(): removes and returns the last inserted element

o peek(): returns the last inserted element without removing it

o size(): returns the number of elements stored

o isEmpty(): indicates whether no elements are stored



| Characteristics | |
|-----------------|--|
| Operations | |
| Performance | |
| Benefits | |
| Drawbacks | |

| Characteristics | Fixed SizeStores Similar ElementsAccess through top index |
|-----------------|---|
| Operations | |
| Performance | |
| Benefits | |
| Drawbacks | |

| Characteristics | Fixed Size Stores Similar Elements Access through top index |
|-----------------|---|
| Operations | <pre>push(item) pop() peek() size() isEmpty()</pre> |
| Performance | |
| Benefits | |
| Drawbacks | |

| Characteristics | Fixed Size Stores Similar Elements Access through top index |
|-----------------|---|
| Operations | <pre>push(item) pop() peek() size() isEmpty()</pre> |
| Performance | <pre>push(item) pop() peek() size() isEmpty()</pre> |
| Benefits | |
| Drawbacks | |

| Characteristics | Fixed Size Stores Similar Elements Access through top index |
|-----------------|---|
| Operations | <pre>push(item) pop() peek() size() isEmpty()</pre> |
| Performance | <pre>push(item) - 0(1) pop() - 0(1) peek() - 0(1) size() - 0(1) isEmpty() - 0(1)</pre> |
| Benefits | |
| Drawbacks | |

| Characteristics | Fixed Size Stores Similar Elements Access through top index |
|-----------------|---|
| Operations | <pre>push(item) pop() peek() size() isEmpty()</pre> |
| Performance | push(item) - O(1) pop() - O(1) peek() - O(1) size() - O(1) isEmpty() - O(1) |
| Benefits | Constant time to add and remove elements |
| Drawbacks | |

| Characteristics | Fixed Size Stores Similar Elements Access through top index |
|-----------------|---|
| Operations | <pre>push(item) pop() peek() size() isEmpty()</pre> |
| Performance | push(item) - O(1) pop() - O(1) peek() - O(1) size() - O(1) isEmpty() - O(1) |
| Benefits | Constant time to add and remove elements |
| Drawbacks | ■ Fixed size and no random access |

| Characteristics | |
|-----------------|--|
| Operations | |
| Performance | |
| Benefits | |
| Drawbacks | |



| Characteristics | Flexible SizeStores Similar ElementsAccess through top pointer |
|-----------------|--|
| Operations | |
| Performance | |
| Benefits | |
| Drawbacks | |

| Characteristics | Flexible SizeStores Similar ElementsAccess through top pointer |
|-----------------|--|
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| Performance | |
| Benefits | |
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Stack Implementation - Linked List

| Characteristics | Flexible Size Stores Similar Elements Access through top pointer |
|-----------------|---|
| Operations | push(item) pop() peek() size() isEmpty() |
| Performance | push(item) - O(1) pop() - O(1) peek() - O(1) size() - O(1) isEmpty() - O(1) |
| Benefits | Constant time to add and remove elements; Variable size |
| Drawbacks | ■ More memory |

Stack in C++ STL

| Operations | C++ STL |
|-----------------------|---|
| <pre>push(item)</pre> | <pre>push(g) - Adds the element 'g' at the top of the stack</pre> |
| <pre>pop()</pre> | pop() – Deletes the topmost element of the stack |
| <pre>peek()</pre> | top() - Returns a reference to the topmost element of the stack |
| <pre>size()</pre> | <pre>size() - Returns the size of the stack</pre> |
| <pre>isEmpty()</pre> | empty() - Returns whether the stack is empty |

Stack in C++ STL: Check Palindrome

Stack in C++ STL: Check Palindrome

```
bool checkPalindrome(string s)
02
03
        stack<char> stk;
         int midIndex = s.length()/2;
                                                                                      Base case to return true if
05
                                                                                      string has one letter or is
06
        if(s.length() < 2)
                                                                                      an empty string
07
            return true;
08
                                                                                      Pushes first half of string
        for(int i = 0; i < s.length()/2; i++)
                                                                                       characters on stack
10
             stk.push(s.at(i));
11
         if(s.length() % 2 != 0)
12
                                                                                      Ignores central element if
            midIndex += 1;
13
                                                                                      the size of string is odd
14
         for(int i = midIndex; i < s.length(); i++)</pre>
15
16
                                                                                      Compare each element in the
            if(stk.top() != s.at(i))
                                                                                      second half of the string
17
                                                                                      with elements pushed in the
                return false;
                                                                                      stack
            stk.pop();
19
20
21
        return true;
22
                                                                         https://onlinegdb.com/Hy2vC3D\
```

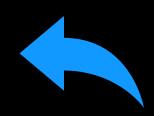
Stack Use Cases

Real World

- Plates
- Books
- Pringle Chips







Computers

- Function Call Stack
- Evaluate Expressions
- Backtracking
- Balanced Parenthesis
- Undo (CTRL + Z)
- Back button



```
main()
  int i=5;
  foo(i);
     foo(int j)
        int k = j+1;
        bar(k);
           bar(int m)
```



```
main()
  int i=5;
  foo(i);
     foo(int j)
        int k = j+1;
        bar(k);
           bar(int m)
```

main PC = 2i = 5

```
main()
  int i=5;
  foo(i);
     foo(int j)
        int k = j+1;
        bar(k);
           bar(int m)
```

foo PC = 3 j = 5 k = 6

main PC = 2 i = 5



```
main()
  int i=5;
  foo(i);
     foo(int j)
        int k = j+1;
        bar(k);
           bar(int m)
```



bar

$$PC = 1$$

 $m = 6$

foo

main

$$PC = 2$$

$$i = 5$$



Stack Use Cases: Balanced Parenthesis

Algorithm for method isBalanced

- 1. Create an empty stack of characters
- 2. Assume that the expression is balanced (balanced is true).
- 3. Set index to 0
- 4. while balanced is **true** and index < the expression's length
 - 5. Get the next character in the data string
 - 6. if the next character is an opening parenthesis
 - 7. Push it onto the stack.
 - 8. else if the next character is a closing parenthesis
 - 9. Pop the top of the stack
 - 10. if stack was empty or its top does not match the closing parenthesis 11. Set balanced to **false**
 - 12. Increment index
- 13. Return true if balanced is true and the stack is empty



Stack Use Cases: Expression Evaluation

- Dijkstra's two-stack algorithm
 - Value: push on the value stack
 - Operator: push on the operator stack
 - Left parenthesis, (: ignore
 - Right parenthesis,): pop operator and two values; push the result of applying that operator to the values on the value stack

Stack Use Cases: Expression Evaluation

Postfix Expression: Removes Parenthesis

| Postfix Expression | Infix Expression | Value |
|--------------------|-------------------|-------|
| 4 7 * | 4 * 7 | 28 |
| 4 7 2 + * | 4 * (7 + 2) | 36 |
| 4 7 * 20 - | (4 * 7) - 20 | 8 |
| 3 4 7 * 2 / + | 3 + ((4 * 7) / 2) | 17 |

Stack Use Cases: Postfix Evaluation

```
1. create an empty stack of integers
 2. while there are more tokens
     get the next token
     if the first character of the token is a digit
         push the token on the stack
5.
      else if the token is an operator
         pop the right operand off the stack
7.
         pop the left operand off the stack
8.
         evaluate the operation
9.
         push the result onto the stack
10.
11. pop the stack and return the result
```



Mentimeter

Menti.com

8369 7506

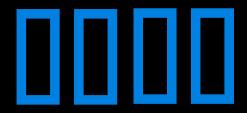




Queues

Queue

First in First Out (FIFO)













Queue ADT

Data

- o **Items**
- Number of Items
- Front and Back

Operations

o enqueue(item): inserts an element to the back of the queue

o dequeue(): removes the element from the front

o size(): returns the number of elements stored

o isEmpty(): indicates whether no elements are stored

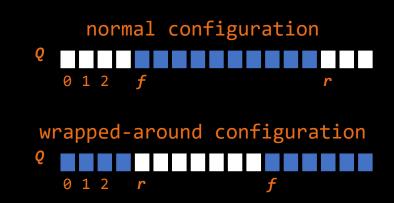


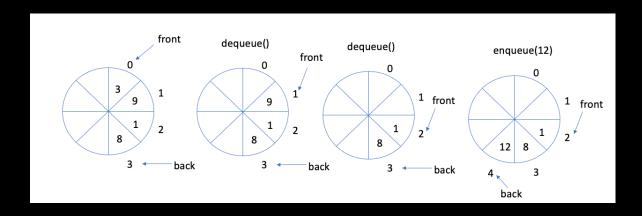
Queue Implementation - Array

| Characteristics | Fixed Size Stores Similar Elements Access through front index |
|-----------------|---|
| Operations | enqueue(item)dequeue()size()isEmpty() |
| Performance | enqueue(item) - O(1) dequeue() - O(1) size() - O(1) isEmpty() - O(1) |
| Benefits | ■ Constant time to add and remove elements |
| Drawbacks | Limited spaceRightward drift problem |

Queue Implementation - Circular Array

```
Use an array of size N in a circular fashion
Two variables keep track of the front and rear
f index of the front element
r index immediately past the rear element
Array location r is kept empty
```





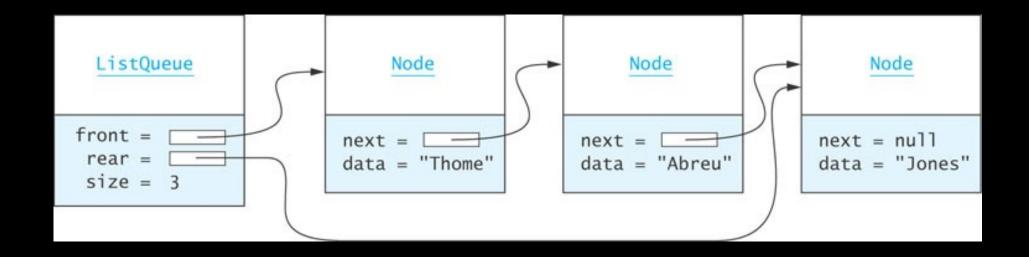
https://stepik.org/lesson/354738/step/1?unit=338782



Queue Implementation - Circular Array

| Characteristics | Fixed SizeStores Similar ElementsAccess through front index |
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| Operations | enqueue(item)dequeue()size()isEmpty() |
| Performance | enqueue(item) - 0(1) dequeue() - 0(1) size() - 0(1) isEmpty() - 0(1) |
| Benefits | ■ Constant time to add and remove elements |
| Drawbacks | ■ Limited space |

Queue Implementation - Linked List



Queue Implementation - Linked List

| Characteristics | Flexible Size Stores Similar Elements Access through front pointer |
|-----------------|---|
| Operations | enqueue(item)dequeue()size()isEmpty() |
| Performance | enqueue(item) - O(1) dequeue() - O(1) size() - O(1) isEmpty() - O(1) |
| Benefits | Constant time to add and remove elements; Variable size |
| Drawbacks | ■ More memory |

Queue in C++ STL

| Operations | C++ STL |
|--|--|
| enqueue(item)dequeue()size()isEmpty() | push(g) - Adds the element 'g' at the end of the queue pop() - Deletes the first element of the queue size() - Returns the size of the queue empty() - Returns whether the queue is empty front() - Returns a reference to the first element of the queue back() - Returns a reference to the last element of the queue |

Queue Use Cases

- Real World
 - Buying Tickets
 - Drive thru at fast food chains
 - Appointments



- Print Queue
- Task Scheduling by OS
- Packet Forwarding by Routers



| ₡ HP LaserJet 4050 Series PS → Use P | rinter Offline | | | | _1 | UN |
|---|----------------|---------------|-------|---------|----------------------|-----|
| Printer Document View Help | | | | | | |
| Document Name | Status | Owner | Pages | Size | Submitted | Pi |
| Microsoft Word - Queues_Paul_1007.doc | | Paul Wolfgang | 52 | 9.75 MB | 1:53:18 PM 10/7/2003 | |
| Microsoft Word - Stacks.doc | | Paul Wolfgang | 46 | 9.05 MB | 1:53:57 PM 10/7/2003 | |
| Microsoft Word - Trees2.doc | | Paul Wolfgang | 54 | 38.4 MB | 1:54:41 PM 10/7/2003 | |
| | | | | | | Þ |
| 3 document(s) in queue | | | | | | //. |



Recommended Resources

- http://www.cplusplus.com/reference/stl/
- Stacks and Queue Questions on Stepik
- OpenDSA Ch-9.8-9.14