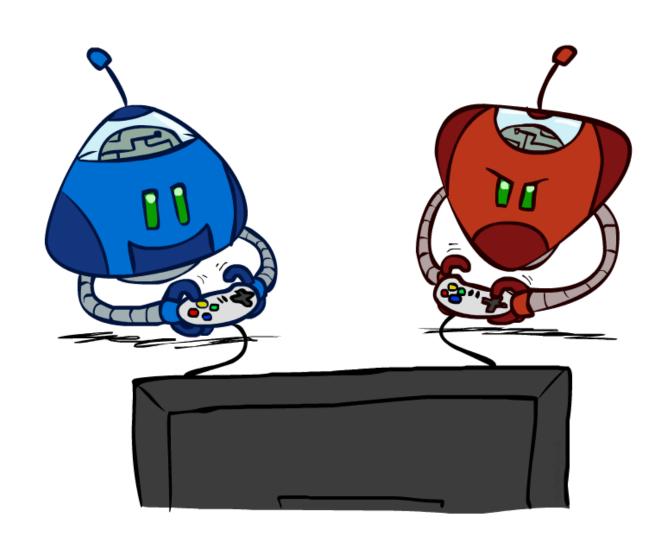




Om Saha Naav[au]-Avatu Saha Nau Bhunaktu Saha Viiryam Karavaavahai Tejasvi Naav[au]-Adhiitam-Astu Maa Vidvissaavahai Om Shaantih Shaantih Shaantih

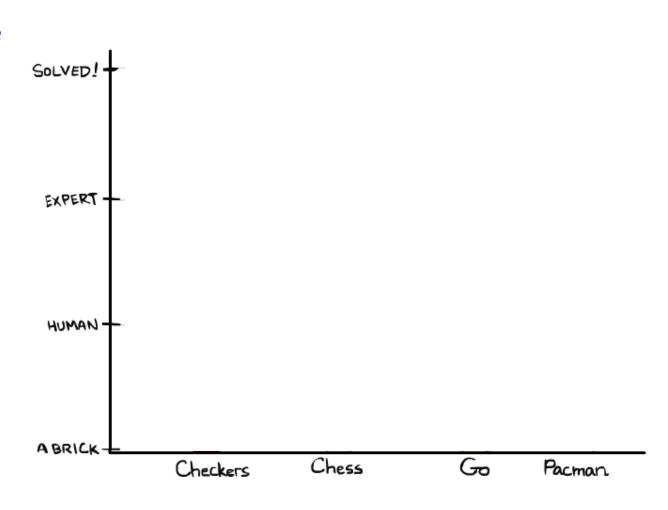
Om, May we all be protected
May we all be nourished
May we work together with great energy
May our intelect be sharpened (may our study be effective)
Let there be no Animosity amongst us
Om, peace (in me), peace (in nature), peace (in divine forces)

### **Adversarial Search and Game Trees**



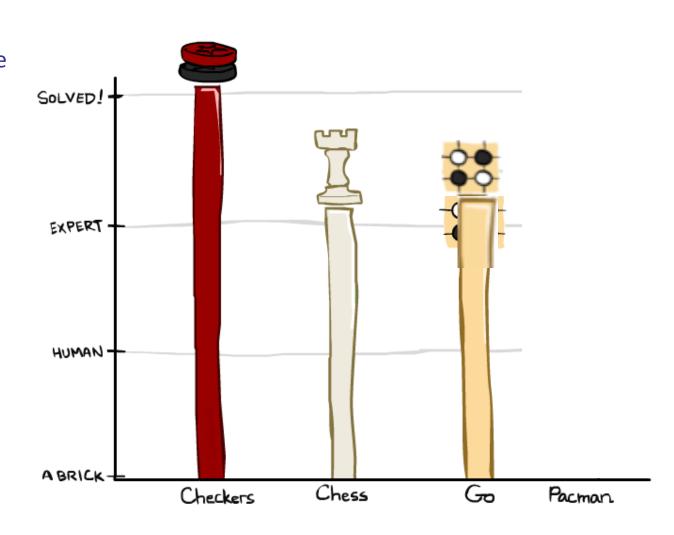
## Game Playing State-of-the-Art

- Checkers: 1950: First computer player. 1994: First computer champion: Chinook ended 40-year-reign of human champion Marion Tinsley using complete 8-piece endgame. 2007: Checkers solved!
- Chess: 1997: Deep Blue defeats human champion Gary Kasparov in a six-game match. Deep Blue examined 200M positions per second, used very sophisticated evaluation and undisclosed methods for extending some lines of search up to 40 ply. Current programs are even better, if less historic.
- Go: Human champions are now starting to be challenged by machines. In go, b > 300! Classic programs use pattern knowledge bases, but big recent advances use Monte Carlo (randomized) expansion methods.

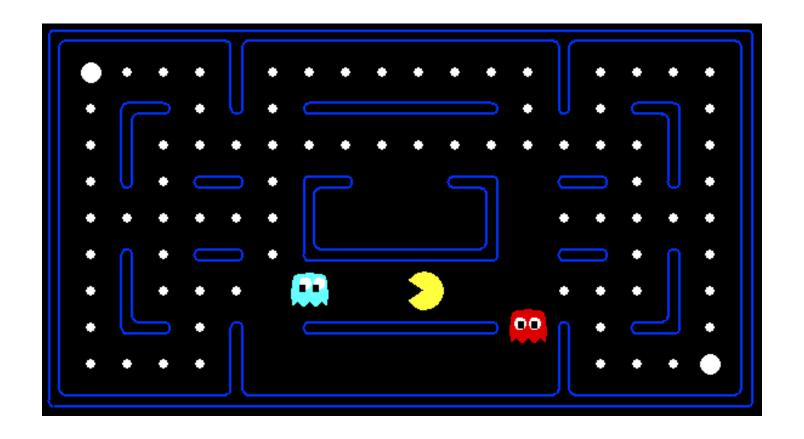


## Game Playing State-of-the-Art

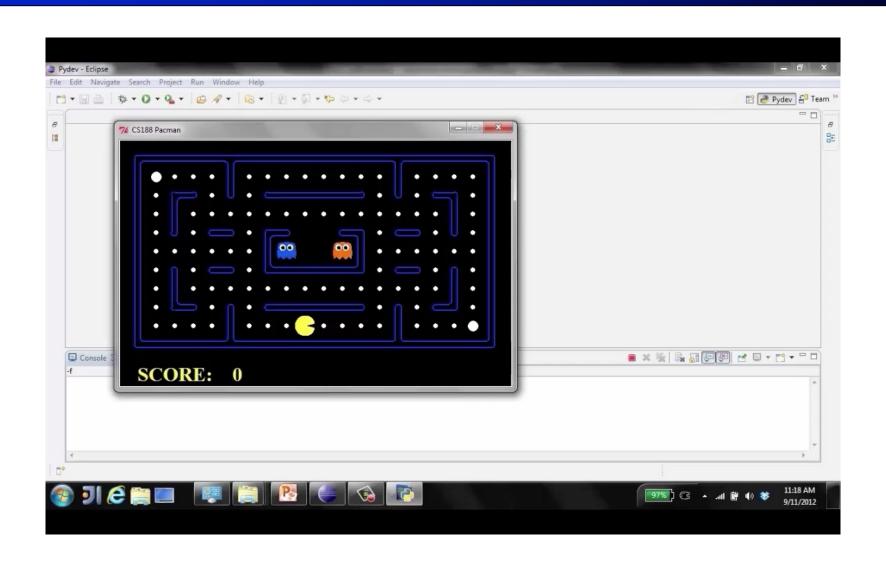
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- Go: AlphaGo have surpassed human champions using DNN, RL, and Monte Carlo Tree Search (MCTS). AlphaGo defeated world champion Lee Sedol in 2016, and its successor, AlphaZero, achieved superhuman performance by learning from scratch without human data
- Pacman



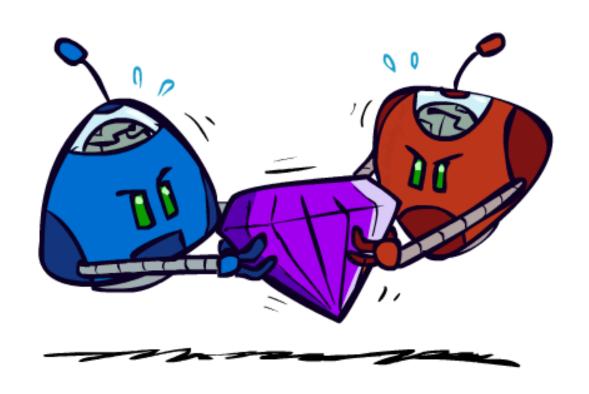
# **Behavior from Computation**



## Video of Demo Mystery Pacman



### **Adversarial Games**



### Types of Games

Many different kinds of games!

### Axes:

- Deterministic or stochastic?
- One, two, or more players?
- Zero sum?
- Perfect information (can you see the state)?

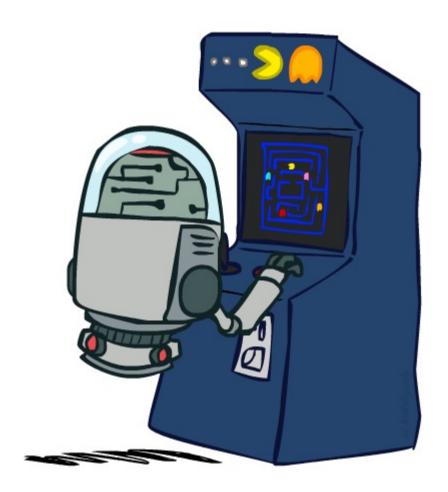


 Want algorithms for calculating a strategy (policy) which recommends a move from each state

### **Deterministic Games**

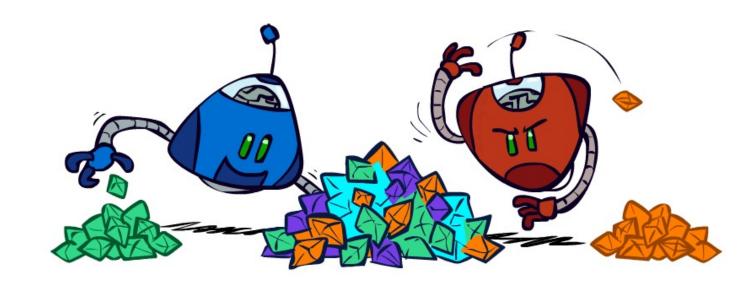
- Many possible formalizations, one is:
  - States: S (start at s<sub>0</sub>)
  - Players: P={1...N} (usually take turns)
  - Actions: A (may depend on player / state)
  - Transition Function:  $SxA \rightarrow S$
  - Terminal Test:  $S \rightarrow \{t,f\}$
  - Terminal Utilities:  $SxP \rightarrow R$

• Solution for a player is a policy:  $S \rightarrow A$ 



### Zero-Sum Games





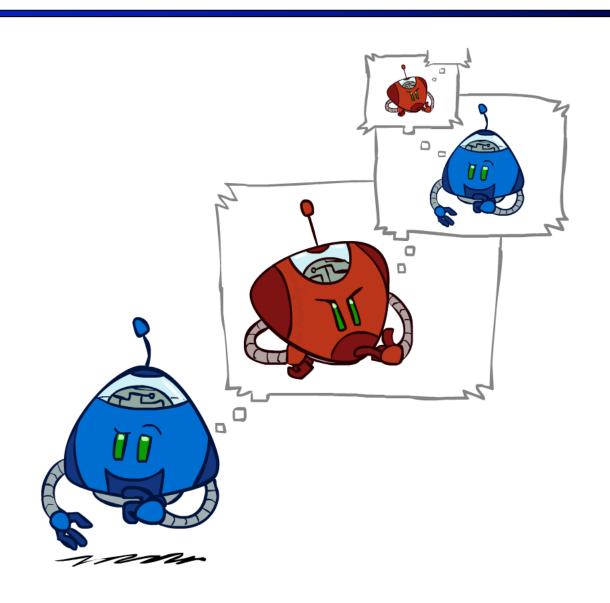
#### Zero-Sum Games

- Agents have opposite utilities (values on outcomes)
- Lets us think of a single value that one maximizes and the other minimizes
- Adversarial, pure competition

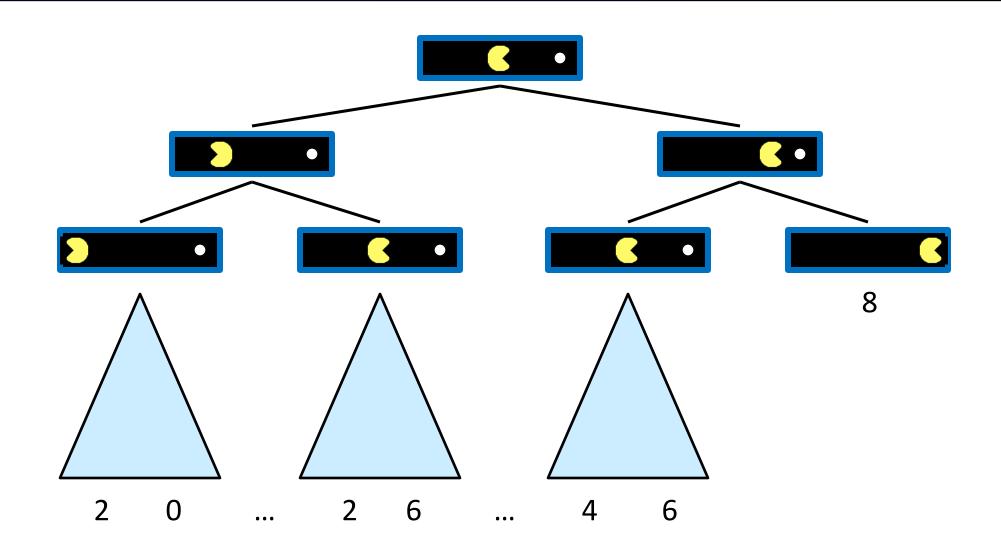
#### General Games

- Agents have independent utilities (values on outcomes)
- Cooperation, indifference, competition, and more are all possible
- More later on non-zero-sum games

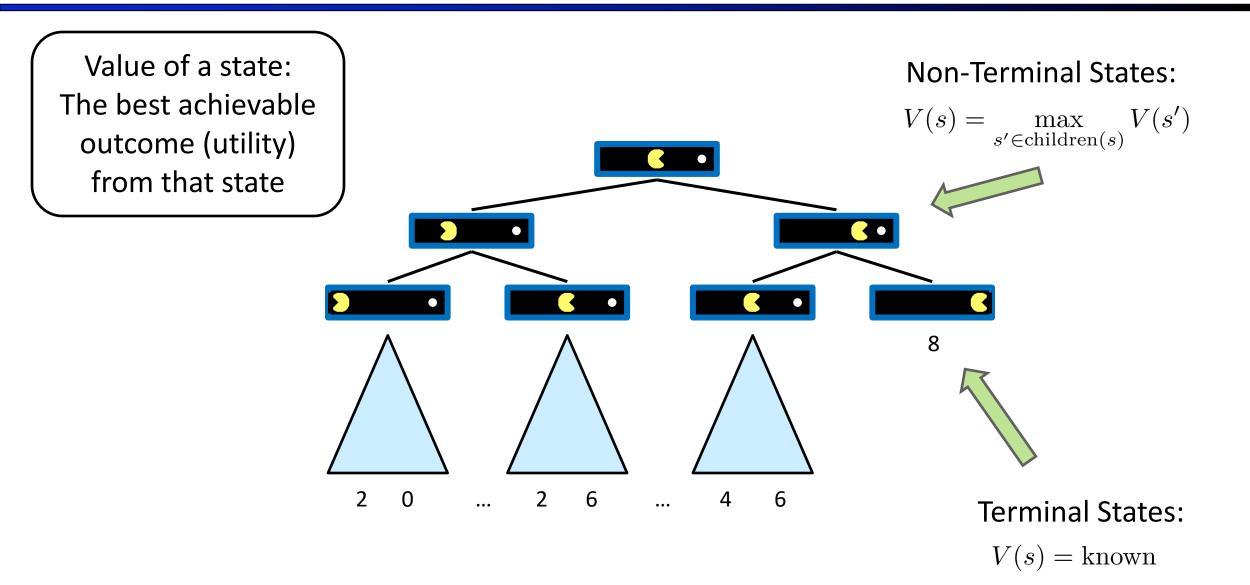
## Adversarial Search



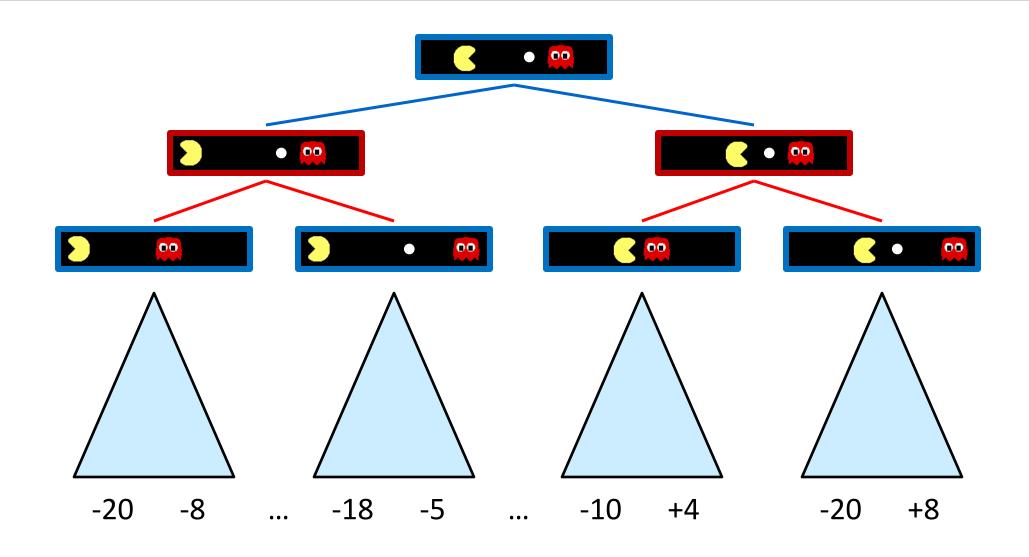
# Single-Agent Trees



### Value of a State

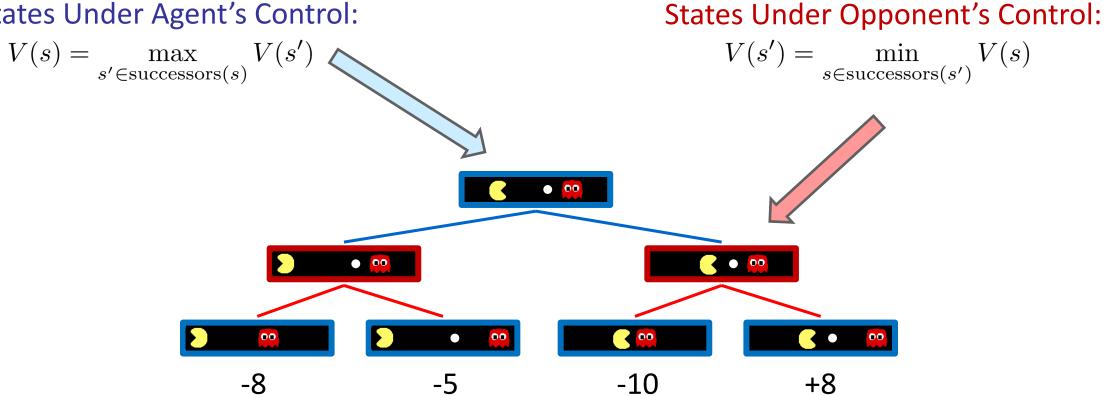


### Adversarial Game Trees



### Minimax Values

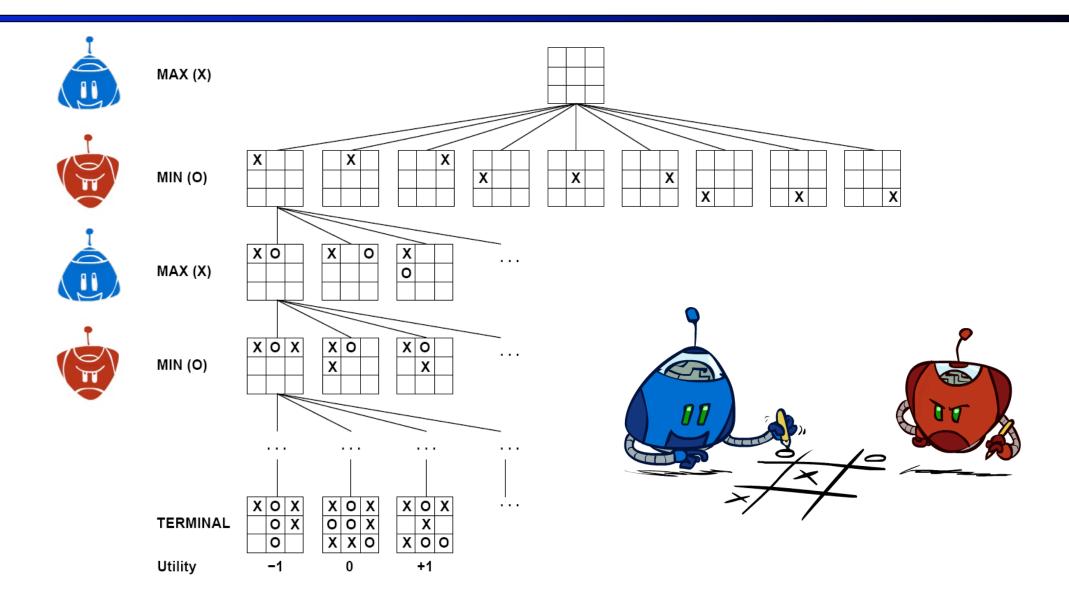
### States Under Agent's Control:



### **Terminal States:**

$$V(s) = \text{known}$$

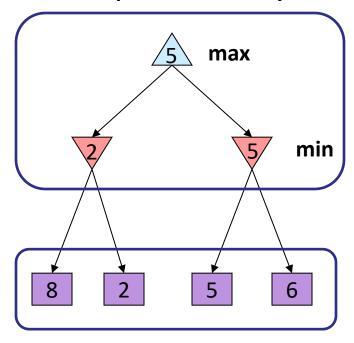
### Tic-Tac-Toe Game Tree



## Adversarial Search (Minimax)

- Deterministic, zero-sum games:
  - Tic-tac-toe, chess, checkers
  - One player maximizes result
  - The other minimizes result
- Minimax search:
  - A state-space search tree
  - Players alternate turns
  - Compute each node's minimax value: the best achievable utility against a rational (optimal) adversary

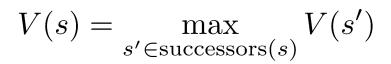
Minimax values: computed recursively



Terminal values: part of the game

### Minimax Implementation

# def max-value(state): initialize v = -∞ for each successor of state: v = max(v, min-value(successor)) return v





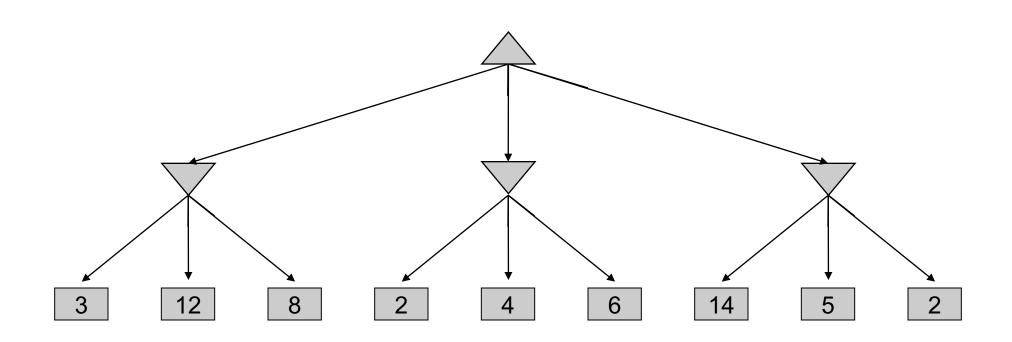
# def min-value(state): initialize v = +∞ for each successor of state: v = min(v, max-value(successor)) return v

$$V(s') = \min_{s \in \text{successors}(s')} V(s)$$

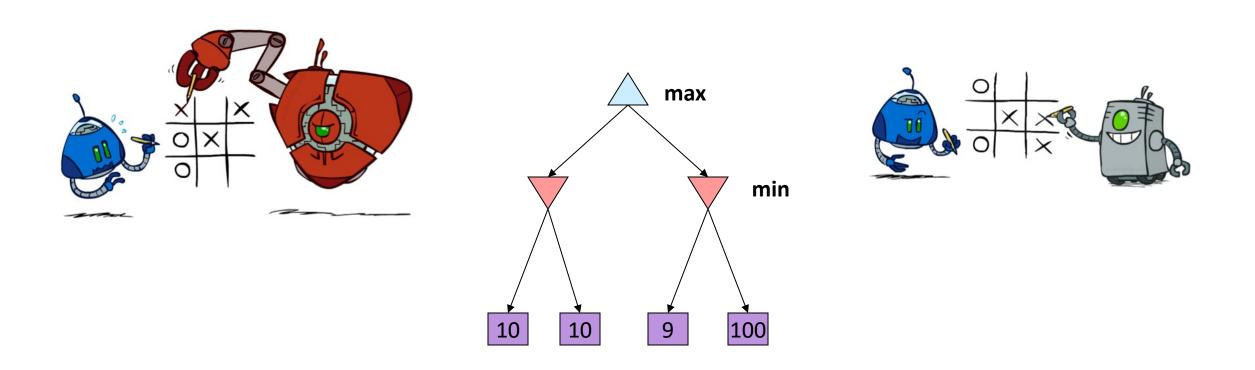
## Minimax Implementation (Dispatch)

```
def value(state):
                      if the state is a terminal state: return the state's utility
                      if the next agent is MAX: return max-value(state)
                      if the next agent is MIN: return min-value(state)
def max-value(state):
                                                             def min-value(state):
    initialize v = -\infty
                                                                 initialize v = +\infty
   for each successor of state:
                                                                 for each successor of state:
       v = max(v, value(successor))
                                                                     v = min(v, value(successor))
                                                                 return v
    return v
```

# Minimax Example



### **Minimax Properties**

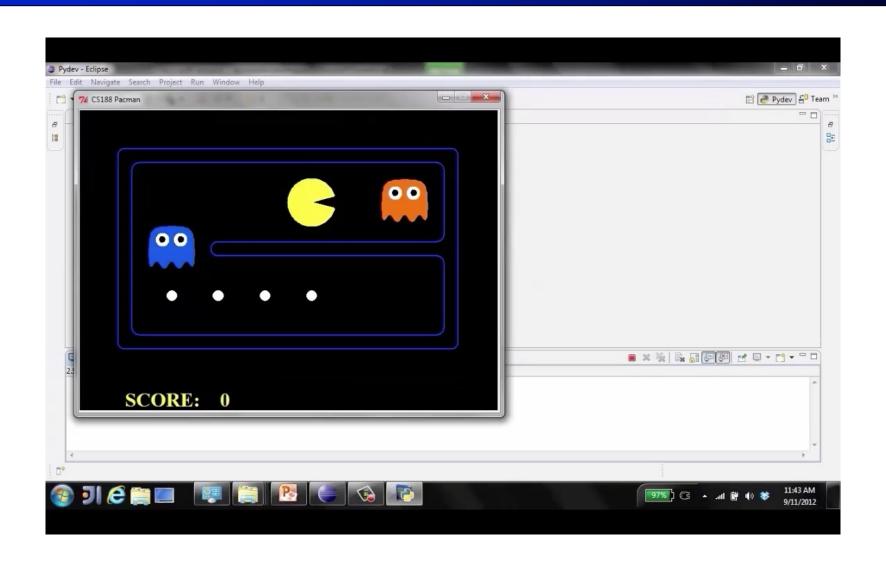


Optimal against a perfect player. Otherwise?

# Min. vs Exp. Adversary (Min.)



# Min. vs Exp. Adversary (Exp.)



## Minimax Efficiency

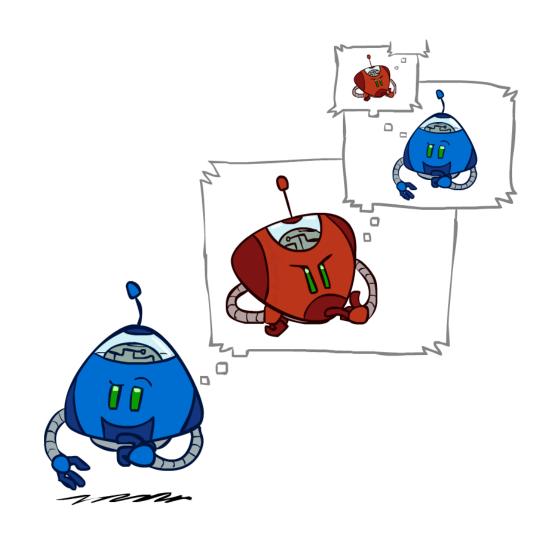
### How efficient is minimax?

Just like (exhaustive) DFS

■ Time: O(b<sup>m</sup>)

Space: O(bm)

- Example: For chess,  $b \approx 35$ ,  $m \approx 100$ 
  - Exact solution is completely infeasible
  - But, do we need to explore the whole tree?

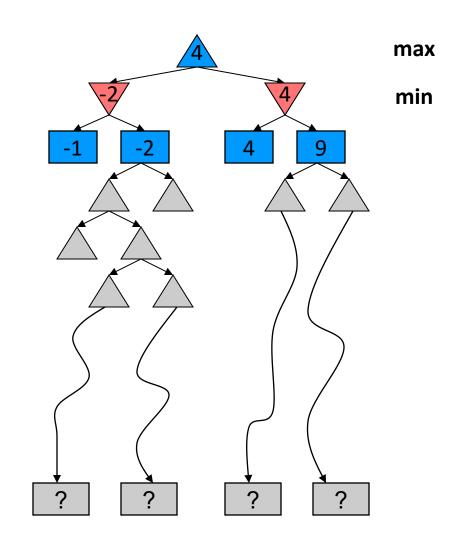


### **Resource Limits**



### **Resource Limits**

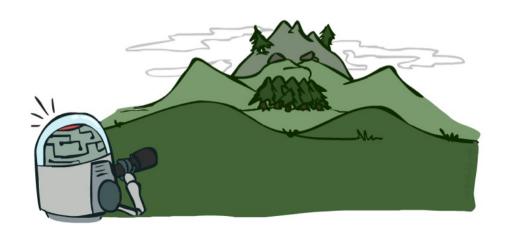
- Problem: In realistic games, cannot search to leaves!
- Solution: Depth-limited search
  - Instead, search only to a limited depth in the tree
  - Replace terminal utilities with an evaluation function for non-terminal positions
- Example:
  - Suppose we have 100 seconds, can explore 10K nodes / sec
  - So can check 1M nodes per move
  - $\alpha$ - $\beta$  reaches about depth 8 decent chess program
- Guarantee of optimal play is gone
- Use iterative deepening for an anytime algorithm



### **Depth Matters**

- Evaluation functions are always imperfect
- The deeper in the tree the evaluation function is buried, the less the quality of the evaluation function matters
- An important example of the tradeoff between complexity of features and complexity of computation





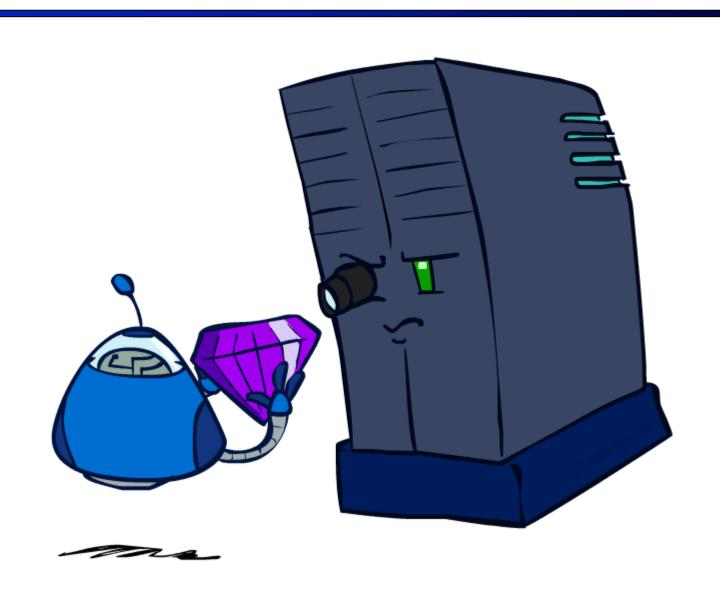
# Video of Demo Limited Depth (2)



# Video of Demo Limited Depth (10)

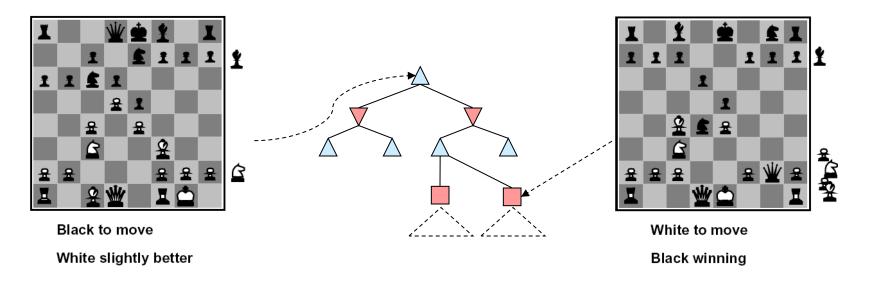


### **Evaluation Functions**



### **Evaluation Functions**

Evaluation functions score non-terminals in depth-limited search



- Ideal function: returns the actual minimax value of the position
- In practice: typically weighted linear sum of features:

$$Eval(s) = w_1 f_1(s) + w_2 f_2(s) + \dots + w_n f_n(s)$$

• e.g.  $f_1(s)$  = (num white queens – num black queens), etc.

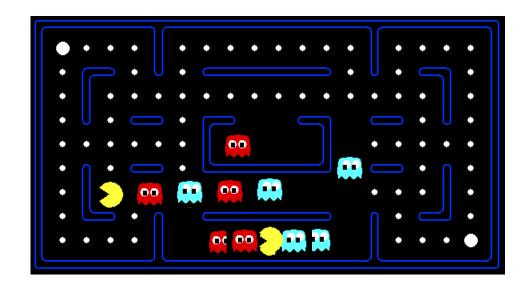
### Stockfish Evaluation

```
function main evaluation(pos) {
 var mg = middle game evaluation(pos);
 var eg = end game evaluation(pos);
 var p = phase(pos), t = tempo(pos);
 eg = eg * scale factor(pos, eg) / 64;
 return ((((mg * p + eg * (128 - p)) << 0) / 128) << 0) + t;
function middle game evaluation(pos) {
 var v = 0;
 v += piece value mg(pos) - piece value mg(colorflip(pos));
 v += psqt mg(pos) - psqt_mg(colorflip(pos));
 v += imbalance total(pos);
 v += pawns mg(pos) - pawns_mg(colorflip(pos));
 v += pieces mg(pos) - pieces mg(colorflip(pos));
 v += mobility mg(pos) - mobility mg(colorflip(pos));
 v += threats mg(pos) - threats mg(colorflip(pos));
 v += passed mg(pos) - passed mg(colorflip(pos));
 v += space(pos) - space(colorflip(pos));
 v += king mg(pos) - king mg(colorflip(pos));
 return v;
```

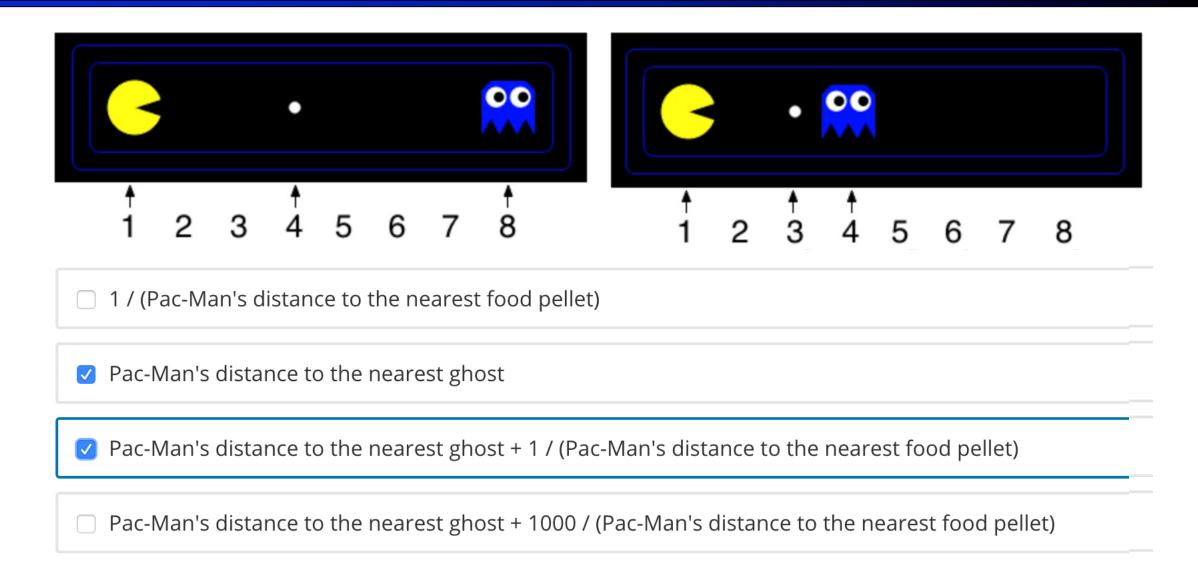


https://hxim.github.io/Stockfish-Evaluation-Guide/

### **Evaluation for Pacman**



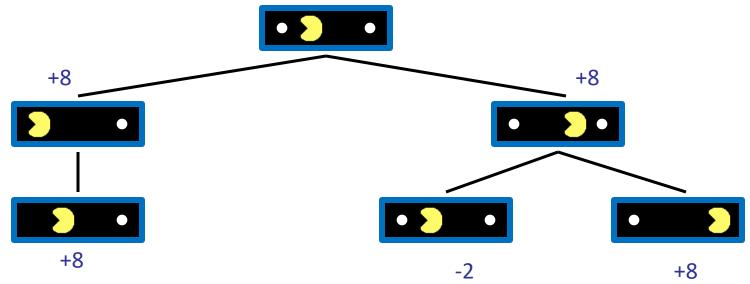
## Target Evaluation function: left > right



# Video of Demo Thrashing (d=2)



## Why Pacman Starves (horizon, d=2)



- A danger of replanning agents!
  - He knows his score will go up by eating the dot now (west, east)
  - He knows his score will go up just as much by eating the dot later (east, east)
  - There are no point-scoring opportunities after eating the dot (within the horizon, two here)
  - Therefore going to east is as good as west. And once it goes to west, it's as good as coming back to east.
  - Therefore, waiting seems just as good as eating: he may go east, then back west in the next round of replanning!

## Video of Demo Thrashing -- Fixed (d=2)



#### Co-ordination and Collaboration

This happens naturally when multiple min-agent acts again a max agent.

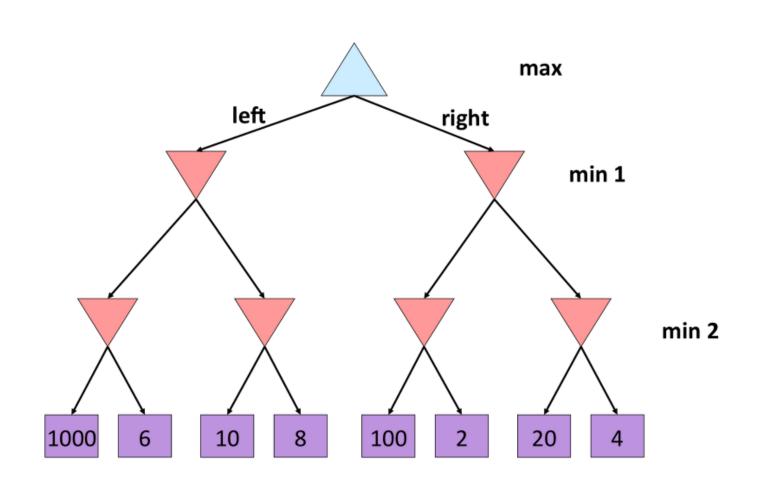
## Video of Demo Smart Ghosts (Coordination)



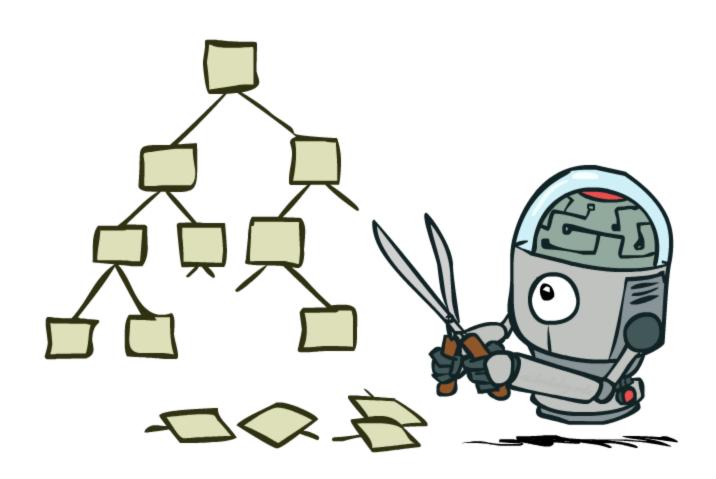
## Video of Demo Smart Ghosts (Coordination) – Zoomed In



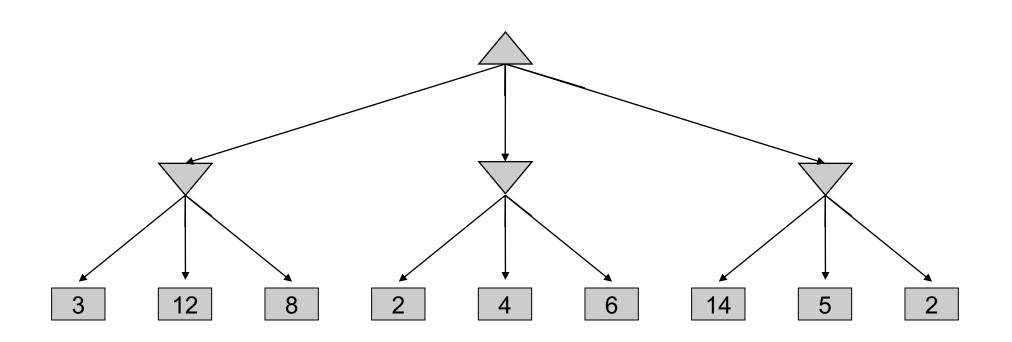
## Quiz: What is the minimax value of this game tree?



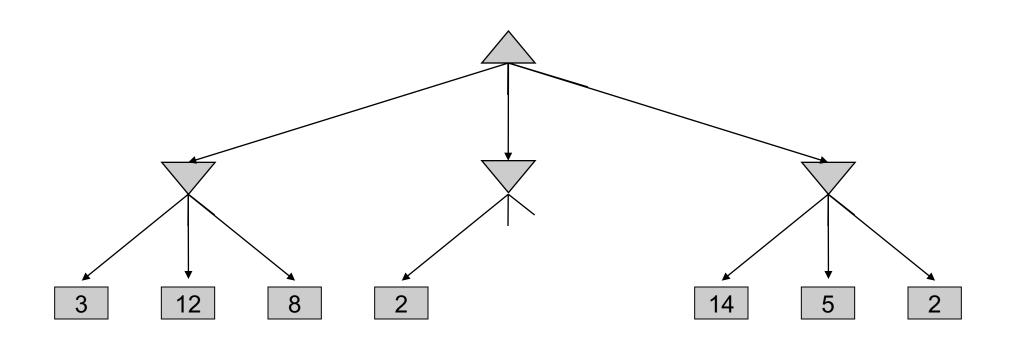
## Game Tree Pruning



## Minimax Example



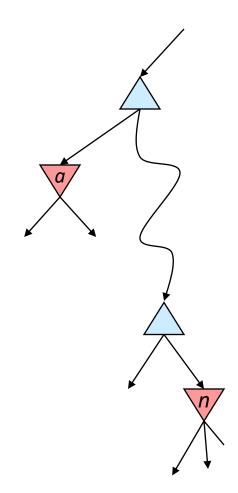
# Minimax Pruning



## Alpha-Beta Pruning

- General configuration (MIN version)
  - We're computing the MIN-VALUE at some node n
  - We're looping over *n*'s children
  - n's estimate of the childrens' min is dropping
  - Who cares about n's value? MAX
  - Let a be the best value that MAX can get at any choice point along the current path from the root
  - If *n* becomes worse than *a*, MAX will avoid it, so we can stop considering *n*'s other children (it's already bad enough that it won't be played)

MAX MIN MAX MIN



MAX version is symmetric

## Alpha-Beta Implementation

```
def value(state, \alpha, \beta):
    if the state is a terminal state: return the state's utility if the next agent is MAX: return max-value(state, \alpha, \beta) if the next agent is MIN: return min-value(state, \alpha, \beta)
```

### Alpha-Beta Implementation

α: MAX's best option on path to root

β: MIN's best option on path to root

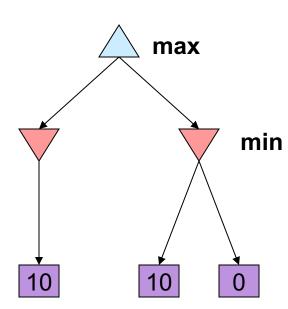
Initially:  $\alpha = -\infty$ ,  $\beta = +\infty$ Always check for if  $\alpha \ge \beta$ 

```
def max-value(state, \alpha, \beta):
    initialize v = -\infty
    for each successor of state:
        succ-value = value(successor, \alpha, \beta)
    v = \max(v, \text{succ-value})
    \alpha = \max(\alpha, v)
    if \alpha \ge \beta return v
return v
```

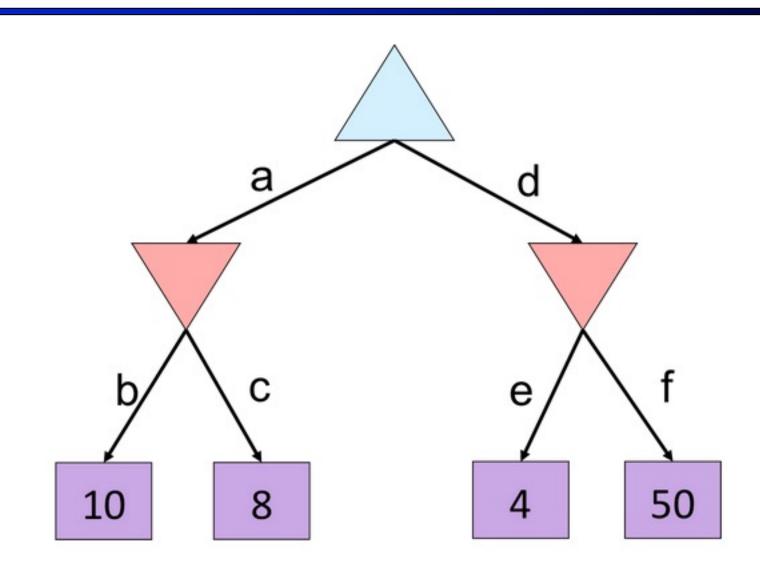
```
def min-value(state , \alpha, \beta):
    initialize v = +\infty
    for each successor of state:
        succ-value = value(successor, \alpha, \beta)
        v = \min(v, \text{ succ-value})
        \beta = \min(\beta, v)
        if \alpha \ge \beta return v
```

## Alpha-Beta Pruning Properties

- This pruning has no effect on minimax value computed for the root!
- Values of intermediate nodes might be wrong
  - Important: children of the root may have the wrong value
  - So the most naïve version won't let you do action selection
- Good child ordering improves effectiveness of pruning
- With "perfect ordering":
  - Time complexity drops to O(b<sup>m/2</sup>)
  - Doubles solvable depth!
  - Full search of, e.g. chess, is still hopeless...



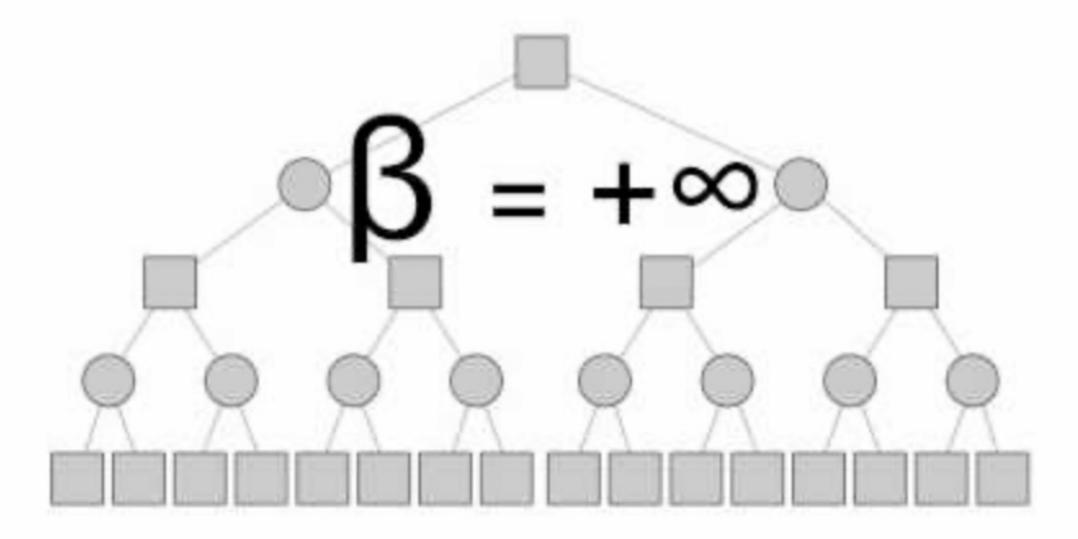
# Alpha-Beta Quiz



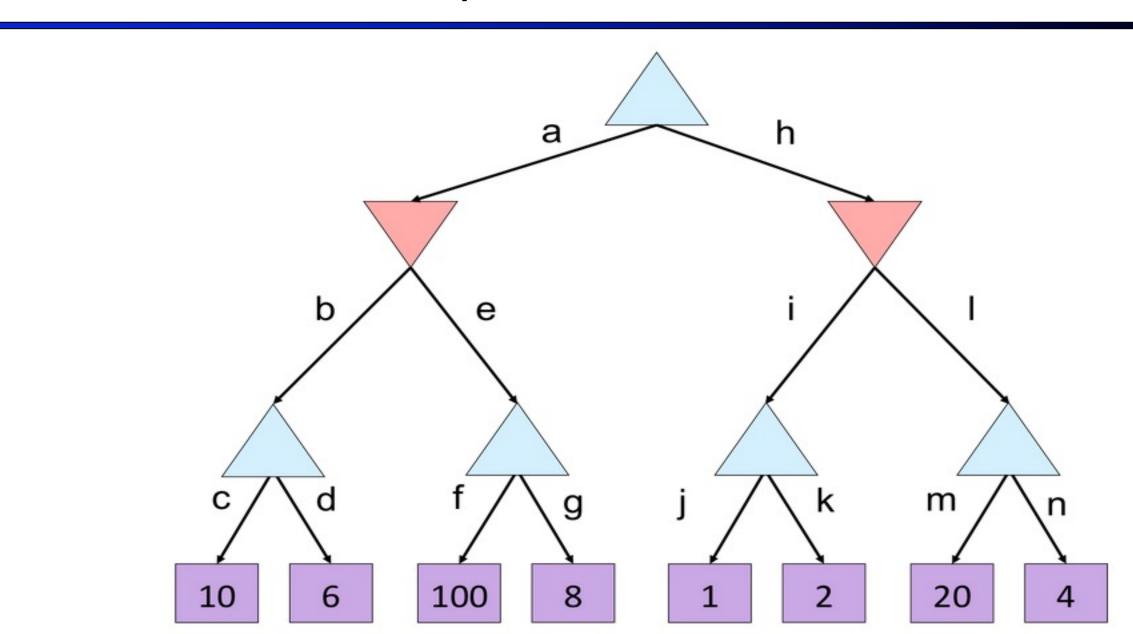
#### YouTube Link

https://youtu.be/J1GoI5WHBto?t=1579

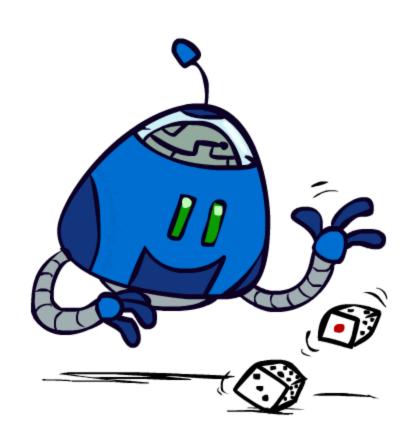
Initially:  $\alpha = -\infty$ ,  $\beta = +\infty$ Always check for if  $\alpha \ge \beta$ 



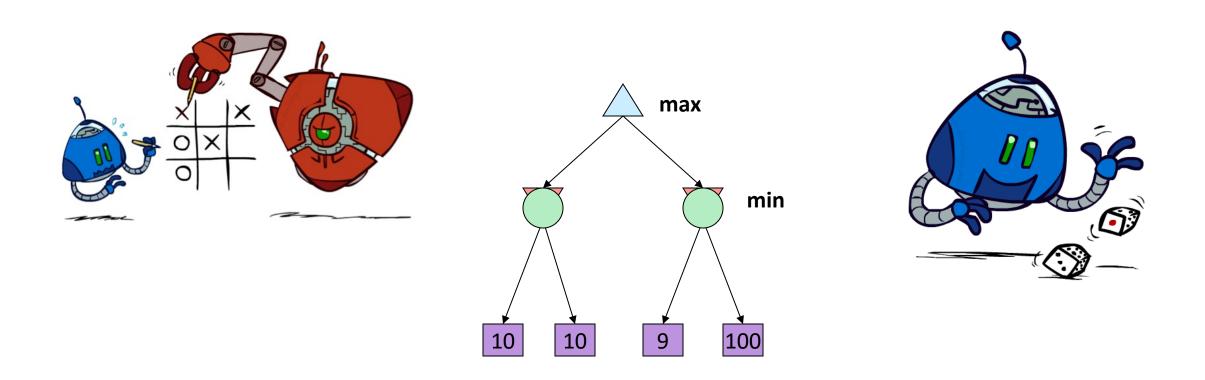
## Alpha-Beta Quiz 2



### **Uncertain Outcomes**



## Worst-Case vs. Average Case



Idea: Uncertain outcomes controlled by chance, not an adversary!

## Video of Demo Min vs. Exp (Min)

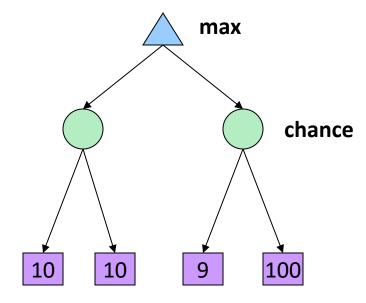


## Video of Demo Min vs. Exp (Exp)



### **Expectimax Search**

- Why wouldn't we know what the result of an action will be?
  - Explicit randomness: rolling dice
  - Unpredictable opponents: the ghosts respond randomly
  - Actions can fail: when moving a robot, wheels might slip
- Values should now reflect average-case (expectimax) outcomes, not worst-case (minimax) outcomes
- Expectimax search: compute the average score under optimal play
  - Max nodes as in minimax search
  - Chance nodes are like min nodes but the outcome is uncertain
  - Calculate their expected utilities
  - I.e. take weighted average (expectation) of children
- Later, we'll learn how to formalize the underlying uncertainresult problems as Markov Decision Processes



### Expectimax Pseudocode

```
def value(state):
    if the state is a terminal state: return the state's utility
    if the next agent is MAX: return max-value(state)
    if the next agent is EXP: return exp-value(state)
```

#### def max-value(state):

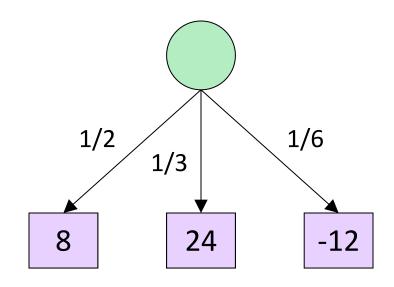
initialize v = -∞
for each successor of state:
 v = max(v, value(successor))
return v

def exp-value(state):

initialize v = 0
for each successor of state:
 p = probability(successor)
 v += p \* value(successor)
return v

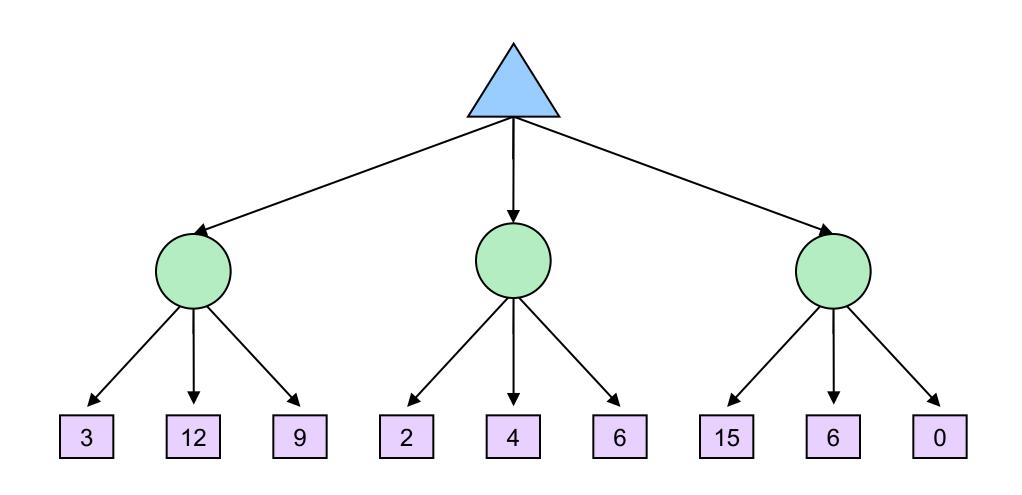
#### **Expectimax Pseudocode**

```
def exp-value(state):
    initialize v = 0
    for each successor of state:
        p = probability(successor)
        v += p * value(successor)
    return v
```

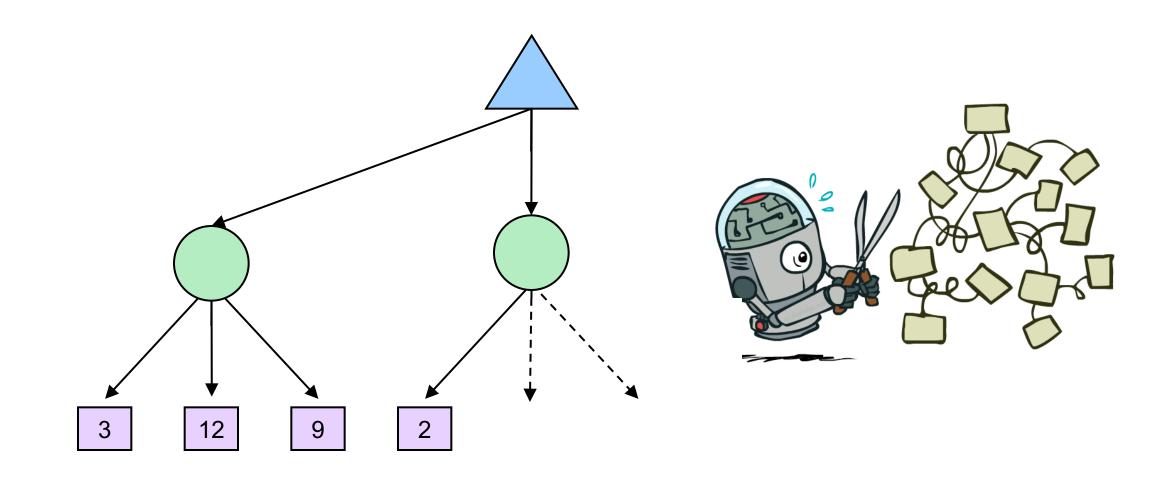


$$v = (1/2)(8) + (1/3)(24) + (1/6)(-12) = 10$$

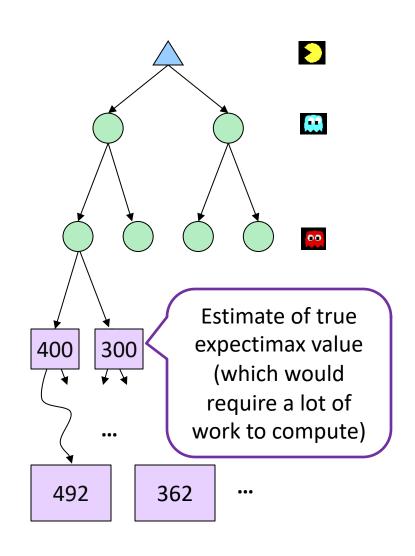
## **Expectimax Example**



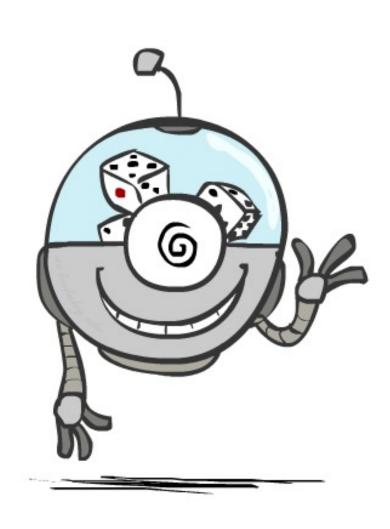
## **Expectimax Pruning?**



## **Depth-Limited Expectimax**



## **Probabilities**

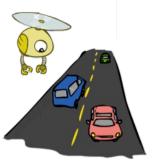


#### Reminder: Probabilities

- A random variable represents an event whose outcome is unknown.
- A probability distribution is an assignment of weights to outcomes
- Example: Traffic on freeway
  - Random variable: T = whether there's traffic
  - Outcomes: T in {none, light, heavy}
  - Distribution: P(T=none) = 0.25, P(T=light) = 0.50, P(T=heavy) = 0.25
- Some laws of probability:
  - Probabilities are always non-negative
  - Probabilities over all possible outcomes sum to one



0.25



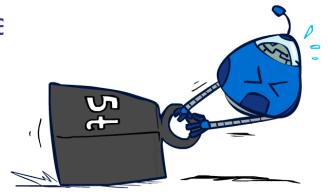
0.50



0.25

### Reminder: Expectations

■ The expected value of a function of a random variable is the average, weighted by the probability distribution over outcomes



• Example: How long to get to the airport?

Time: 20 min

Probability:

X

0.25

+

30 min

0.50

+

60 min

X

0.25



35 min







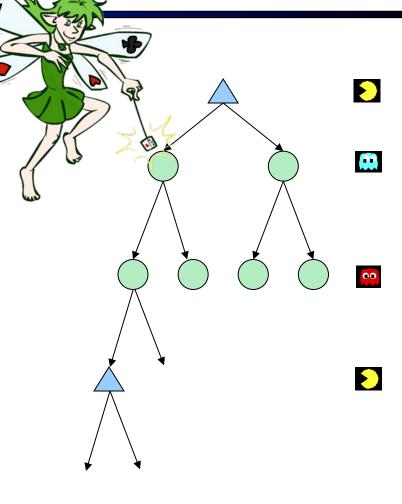
## What Probabilities to Use?

 In expectimax search, we have a probabilistic model of how the opponent (or environment) will behave in any state

Model could be a simple uniform distribution (roll a die)

 Model could be sophisticated and require a great deal of computation

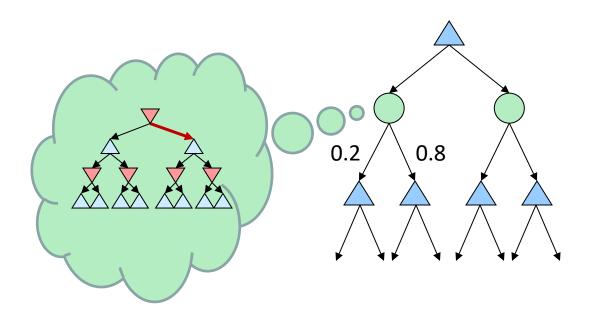
- We have a chance node for any outcome out of our control: opponent or environment
- The model might say that adversarial actions are likely!
- For now, assume each chance node magically comes along with probabilities that specify the distribution over its outcomes



Having a probabilistic belief about another agent's action does not mean that the agent is flipping any coins!

#### Quiz: Informed Probabilities

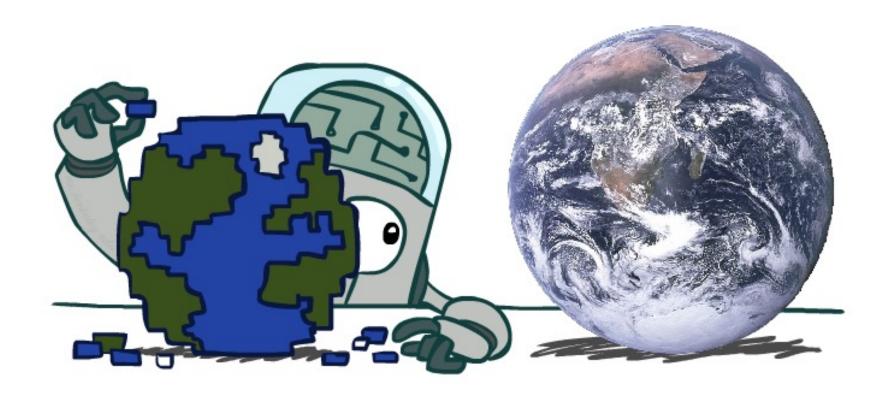
- Let's say you know that your opponent is actually running a depth 2 minimax, using the result 80% of the time, and moving randomly otherwise
- Question: What tree search should you use?



#### Answer: Expectimax!

- To figure out EACH chance node's probabilities,
   you have to run a simulation of your opponent
- This kind of thing gets very slow very quickly
- Even worse if you have to simulate your opponent simulating you...
- ... except for minimax, which has the nice property that it all collapses into one game tree

## **Modeling Assumptions**



## The Dangers of Optimism and Pessimism

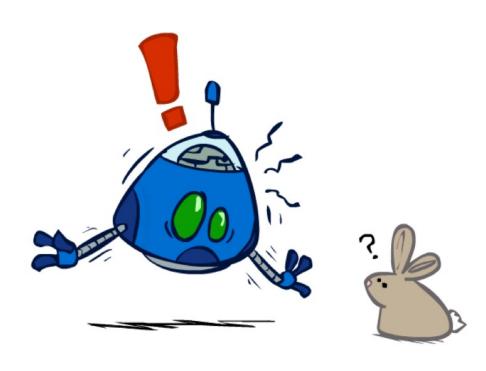
#### Dangerous Optimism

Assuming chance when the world is adversarial

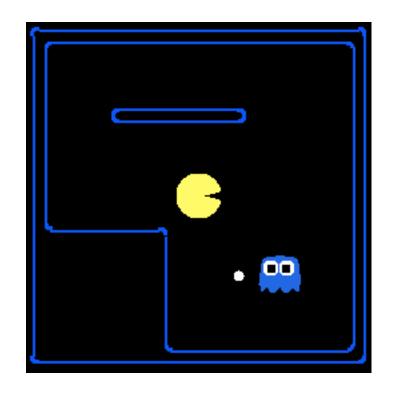


#### **Dangerous Pessimism**

Assuming the worst case when it's not likely



## Assumptions vs. Reality

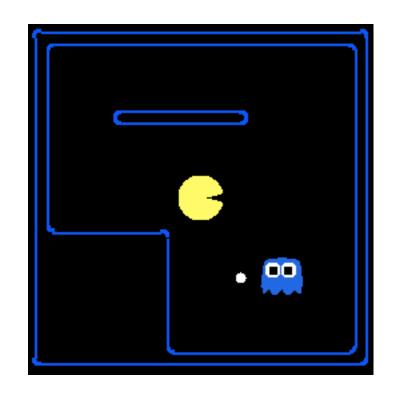


	Adversarial Ghost	Random Ghost
Minimax Pacman		
Expectimax Pacman		

Results from playing 5 games

Pacman used depth 4 search with an eval function that avoids trouble Ghost used depth 2 search with an eval function that seeks Pacman

## Assumptions vs. Reality



	Adversarial Ghost	Random Ghost
Minimax	Won 5/5	Won 5/5
Pacman	Avg. Score: 483	Avg. Score: 493
Expectimax	Won 1/5	Won 5/5
Pacman	Avg. Score: -303	Avg. Score: 503

Results from playing 5 games

Pacman used depth 4 search with an eval function that avoids trouble Ghost used depth 2 search with an eval function that seeks Pacman

[Demos: world assumptions (L7D3,4,5,6)]

## Video of Demo World Assumptions Random Ghost – Expectimax Pacman



## Video of Demo World Assumptions Adversarial Ghost – Minimax Pacman



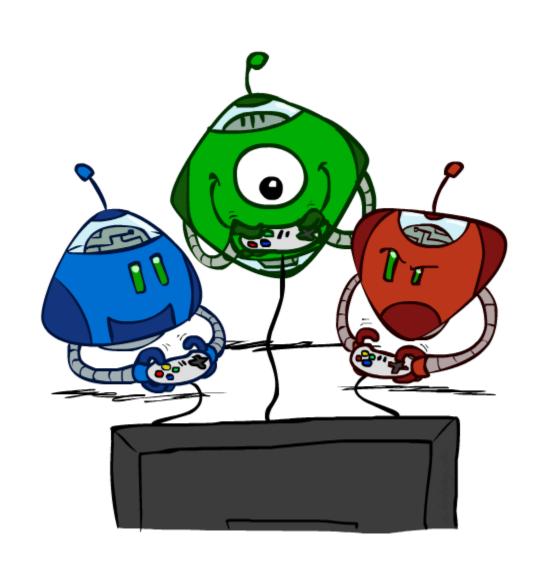
## Video of Demo World Assumptions Adversarial Ghost – Expectimax Pacman



### Video of Demo World Assumptions Random Ghost – Minimax Pacman

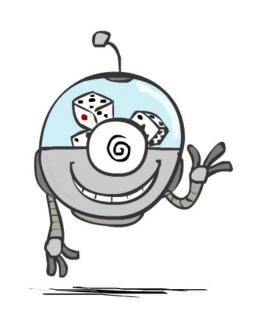


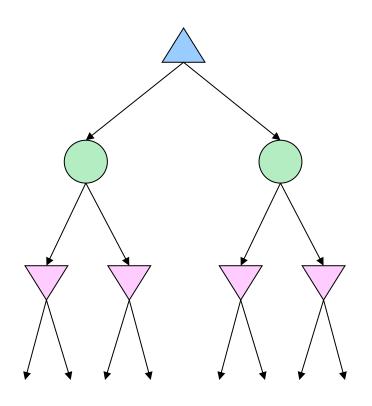
# Other Game Types



## Mixed Layer Types

- E.g. Backgammon
- Expectiminimax
  - Environment is an extra "random agent" player that moves after each min/max agent
  - Each node
     computes the
     appropriate
     combination of its
     children











## Example: Backgammon

- Dice rolls increase *b*: 21 possible rolls with 2 dice
  - Backgammon ≈ 20 legal moves
  - Depth  $2 = 20 \times (21 \times 20)^3 = 1.2 \times 10^9$
- As depth increases, probability of reaching a given search node shrinks
  - So usefulness of search is diminished
  - So limiting depth is less damaging
  - But pruning is trickier...
- Historic AI: TDGammon uses depth-2 search + very good evaluation function + reinforcement learning: world-champion level play
- 1st AI world champion in any game!





What if the game is not zero-sum, or has multiple players?

#### Generalization of minimax:

- Terminals have utility tuples
- Node values are also utility tuples
- Each player maximizes its own component
- Can give rise to cooperation and competition dynamically...

