

Student Information

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Answer 1

a)

(i) $D = A \cap (B \cup C)$

(ii) $E = (A \cap B) \cup C$

(iii) $F = (A - B) \cup (A \cap C)$

b)

Let $A = \{e,n,t\}$, $B = \{t,e,r\}$, and $C = \{s,e,r\}$

(i) $(A \times B) \times C = \{(e,t),(e,e),(e,r),(n,t),(n,e),(n,r),(t,t),(t,e),(t,r)\} \times \{s,e,r\} = \{(e,t,s),(e,e,s),(e,r,s),(n,t,s),(n,e,s),(n,r,s),(t,t,s),(t,e,s),(t,r,s),(e,t,e),(e,e,e),(e,r,e),(n,t,e),(n,e,e),(n,r,e),(t,t,e),(t,e,e),(t,r,e),(e,t,r),(e,e,r),(e,r,r),(n,t,r),(n,e,r),(n,r,r),(t,t,r),(t,e,r),(t,r,r)\} = \{e,n,t\} \times \{(t,s),(t,e),(t,r),(e,s),(e,e),(e,r),(r,s),(r,e),(r,s)\} = A \times (B \times C)$

(ii) $(A \cap B) \cap C = \{e,t\} \cap \{s,e,r\} = \{e\} = \{e,n,t\} \cap \{e,r\} = A \cap (B \cap C)$

(iii) $(A \oplus B) \oplus C = \{n,r\} \oplus \{s,e,r\} = \{s,e,n\} = \{e,n,t\} \oplus \{t,s\} = \{e,n,s\} = A \oplus (B \oplus C)$

Answer 2

a)

Take any $x \in A_0$, then we have $f(x) \in f(A_0)$. By definition, $f^{-1}(f(A_0)) = \{a \mid f(a) \in f(A_0)\}$. Since $f(x) \in f(A_0)$, x satisfies the condition(1): $f(a) \in f(A_0)$; we have $x \in f^{-1}(f(A_0))$. We have proved $x \in A_0 \rightarrow x \in f^{-1}(f(A_0))$. That means $A_0 \subseteq f^{-1}(f(A_0))$. Now, suppose f injective. We want to prove that $A_0 = f^{-1}(f(A_0))$. Suppose not; that means we can find an $x \in f^{-1}(f(A_0))$ such that $x \notin A_0$. Because we already know $A_0 \subseteq f^{-1}(f(A_0))$, $x \in f^{-1}(f(A_0)) \rightarrow x$ satisfies the condition(1) then $f(x) = u \in f(A_0)$. $u \in f(A_0)$ means, there exists a $y \in A_0$ with $f(y) = u$. Since f is injective and $f(x) = u = f(y)$, we must have $x=y$. But $x \notin A_0$ and $y \in A_0$; contradiction.

b)

Take any $y \in f^{-1}(f(B_0)) \rightarrow y = f(x)$ for some $x \in f^{-1}(B_0)$. $x \in f^{-1}(B_0) = \{a \mid f(a) \in B_0\}$ means x satisfies the condition(1); $f(x) \in B_0 \rightarrow y \in B_0$. So, $f(f^{-1}(B_0)) \subseteq B_0$. Suppose f is surjective, but $f(f^{-1}(B_0)) \neq B_0$. That means we can find a $y \in B_0$ such that $y \notin f(f^{-1}(B_0))$. By surjectivity, there exists an x such that $f(x) = y \rightarrow x$ satisfies the condition(1); $f(x) \in B_0 \rightarrow x \in f^{-1}(B_0) \rightarrow f(x) \in f(f^{-1}(B_0)) \rightarrow y \in f(f^{-1}(B_0))$; hence contradiction.

Answer 3

Let A be a non-empty set.

(i) A is countable

(ii) There is a surjective function $f : \mathbb{Z}^+ \rightarrow A$

(iii) There is an injective function $g : A \rightarrow \mathbb{Z}^+$

(i) \rightarrow (ii): If A is countably infinite, then there exists a bijection $f : \mathbb{Z}^+ \rightarrow A$ and then (ii) follows. If A is finite, then there is bijection $h : \{1, 2, \dots, n\} \rightarrow A$ for some n . Then the function $f : \mathbb{Z}^+ \rightarrow A$ defined by;

$$f(t) = \begin{cases} h(t) & 1 \leq t \leq n, \\ h(n) & t > n \end{cases} \quad \text{is a surjection.}$$

(ii) \rightarrow (iii): Assume that $f : \mathbb{Z}^+ \rightarrow A$ is a surjection. We claim that there is an injection $g : A \rightarrow \mathbb{Z}^+$. To define g note that if $a \in A$, then $f^{-1}(\{a\}) \neq \emptyset$. Hence we set $g(a) = \min f^{-1}(\{a\})$. Now note that if $a \neq a'$, then the sets $f^{-1}(\{a\}) \cap f^{-1}(\{a'\}) = \emptyset$ which implies $\min^{-1}(\{a\}) \neq \min^{-1}(\{a'\})$. Hence $g^{-1}(a) \neq g^{-1}(a')$ and $g : A \rightarrow \mathbb{Z}^+$ is an injective function.

(iii) \rightarrow (i): Assume that $g : A \rightarrow \mathbb{Z}^+$ is an injection. We want to show that A is countable. Since $g : A \rightarrow g(A)$ is a bijection and $g(A) \subset \mathbb{Z}^+$, proposition "Any subset of a countable set is countable." implies that A is countable.

Answer 4

a)

Let \mathbf{F} be the set containing finite binary strings and ϵ denote the empty sequence (the sequence with no terms). Then, the sequence $\mathbf{F} = \{\epsilon, 0, 1, 00, 01, 10, 11, 000, 001, 010, 011, \dots\}$ in which the binary sequences of length 0 are listed, then the binary sequences of length 1 are listed in increasing numeric order, then the binary sequences of length 2 are listed in increasing numeric order, and so on, contains every finite binary sequence exactly once.

b)

Let \mathbf{I} be the set of all infinite sequences of 0s and 1s. We use Cantor's diagonal argument. So we assume (toward a contradiction) that we have an enumeration of the elements of \mathbf{I} , say as $\mathbf{I} = \{i_1 i_2 i_3 \dots\}$ where each s_n is an infinite sequence of 0s and 1s. We will write $s_1 = s_{1,1} s_{1,2} s_{1,3} \dots$, $s_2 = s_{2,1} s_{2,2} s_{2,3} \dots$, and so on; so $s_n = s_{n,1} s_{n,2} s_{n,3} \dots$. So we denote the m th element of s_n by $s_{n,m}$. Now we create a new sequence $t = t_1 t_2 t_3 t_4 \dots$ of 0s and 1s as follows: $t_n = s_{n,n} - 1$ (so $t_n = 1$ if $s_{n,n} = 0$ and $t_n = 0$ if the $s_{n,n}$ is 1). It is clear that t is an element of \mathbf{I} - it is an infinite sequence of 0s and 1s. However, we will now see that t is not in the list above. Suppose that $t = i_k$ for some value of k . Then $t_k = t_{k,k}$, but by the construction, $t_k \neq i_{k,k}$, so this is not possible. We conclude that \mathbf{I} is not countable, is uncountable.

Answer 5

a)

$f(n)$ is $\Theta(g(n))$, where $f(n)=n\log(n)$, $g(n)=\log(n!)$ and there exist constants c and k , if:

(i) $f(n)$ is $O(g(n)) \rightarrow f(n) \leq cg(n)$ whenever $n > k$

$$\begin{aligned}\log(n!) &= \log(1) + \log(2) + \dots + \log(n) \\ &\leq \log(n) + \log(n) + \dots + \log(n) \\ \log(n!) &\leq n\log(n); \text{ where } c = 1 \text{ and } n > 1\end{aligned}$$

(ii) $f(n)$ is $\Omega(g(n)) \rightarrow f(n) \geq cg(n)$

$$\begin{aligned}\log(n!) &= \sum_{i=1}^n \log(i) \\ &\geq \sum_{i=n/2}^n \log(i) \\ &\geq \sum_{i=n/2}^n \log(n/2) \\ &\geq \frac{n}{2} \log(n/2) \\ &= \frac{n}{2} (\log(n) - \log(2)); \text{ where } c = \frac{1}{2} \text{ and } n > 1 \\ &= \Omega(n\log(n))\end{aligned}$$

b)

$$\begin{aligned}2^n &= \underbrace{2 * 2 * \dots * 2 * 2}_{n \text{ times}} \\ n! &= \underbrace{1 * 2 * \dots * (n-1) * n}_{n \text{ elements}}\end{aligned}$$

$$\frac{2^n}{n!} = \frac{2}{1} * \frac{2}{2} * \frac{2}{3} * \dots * \frac{2}{n-1} * \frac{2}{n}$$

$$\frac{2^n}{n!} \leq \frac{2}{1} * \frac{2}{2} * 1 * \dots * 1 * \frac{2}{n} = \frac{4}{n}$$

Hence 2^n is $O(n!)$ such that $n!$ grows faster than 2^n ($|2^n| \leq c|n!|$); where $c = 1$ and $n \geq 4$