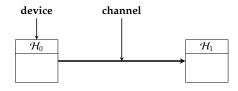
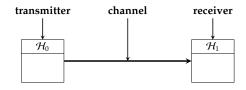
▶ We can *already* form a simple **point-to-point** communication channel



using TIA-232-F.

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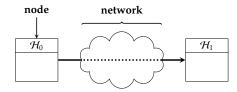
▶ We can *already* form a simple **point-to-point** communication channel



### using TIA-232-F, st.

- ► Good:
  - modular wrt. communication medium and protocol,
  - uses standardised components,
  - · ...
- ► Bad:
  - well defined, but quite limited functionality,
  - the organisation of components is fixed,
  - there are limits wrt. physical locality,
  - •

► We can *already* form a simple **point-to-point** communication channel



using TIA-232-F.

► Challenge: expand our remit to use of a **computer network** ...

COMS20001 lecture: week #20

▶ ... how?



COMS20001 lecture: week #20

- ▶ ... how?
  - 1. Requirements: what do we expect from a network?
    - supports a high degree of connectivity,
    - allows inter-operation between heterogeneous components,
    - uses appropriate level of abstraction to provide useful functionality,
    - satisfies any relevant quality metrics (e.g., efficiency, reliability),
    - can be (dynamically) scaled wrt. components and usage.

COMS20001 lecture: week #20

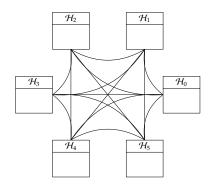
▶ ... how?

- 2. Architecture: how should we approach the design of a network?
  - ▶ (physical or logical) **topology**, i.e., the structure of components in the network, plus
  - control i.e., how do we use those components to communicate, and
  - standardisation, i.e., how do we ensure components which communicate can inter-operate.

▶ Idea #1: *fully*-connected, st.  $\mathcal{H}_i$  is connected to  $\mathcal{H}_j$  for  $j \in S = \{0, 1, ..., n-1\} \setminus \{i\}$ .

- Note that access to a given channel may be
  - 1. dedicated, or
  - shared, implying a need to control access but also the posibility to broadcast to multiple receiving nodes.

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- Example: a mesh topology, i.e.,



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- Example: a bus topology, i.e.,

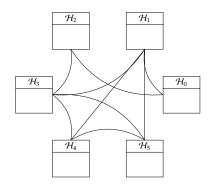


- ▶ Note that access to a given channel may be
  - 1. dedicated, or
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▶ Idea #2: *partially*-connected, st.  $\mathcal{H}_i$  is connected to  $\mathcal{H}_j$  for  $j \in S \subset \{0, 1, ..., n-1\} \setminus \{i\}$ .

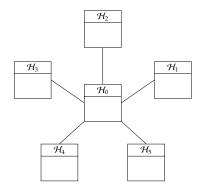
- Note that connectivity may now be either
  - 1. direct, or
  - indirect, implying a need for intermediate hops, e.g., as realised via store-and-forward by (intermediate) switching nodes.

- ▶ Idea #2: partially-connected, st.  $\mathcal{H}_i$  is connected to  $\mathcal{H}_j$  for  $j \in S \subset \{0, 1, ..., n-1\} \setminus \{i\}$ .
- **Example:** a **mesh topology**, i.e.,



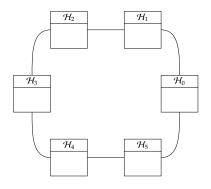
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- Example: a star topology, i.e.,



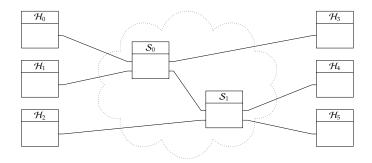
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- ▶ Idea #2: *partially*-connected, st.  $\mathcal{H}_i$  is connected to  $\mathcal{H}_j$  for  $j \in S \subset \{0, 1, ..., n-1\} \setminus \{i\}$ .
- Example: a ring topology, i.e.,

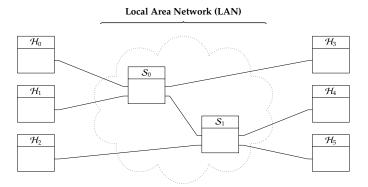


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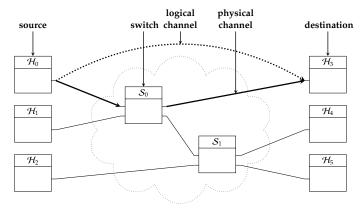
► We can *generalise* partially connected network topologies



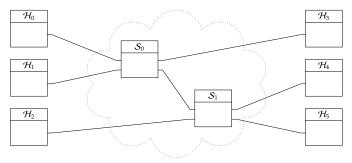
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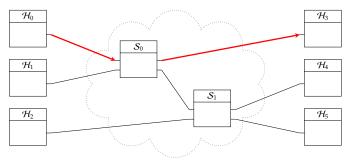


We can generalise partially connected network topologies



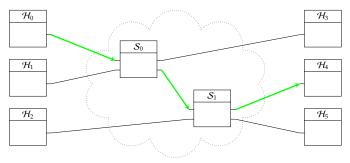
- circuit switching and
- packet switching (to support connection-based or connection-less channels).

We can generalise partially connected network topologies



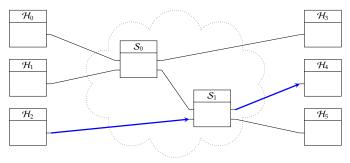
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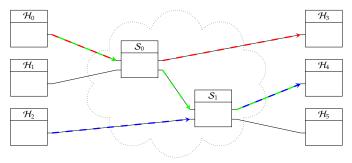
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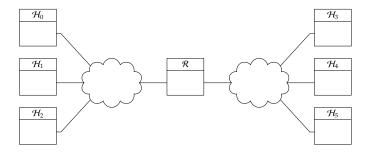
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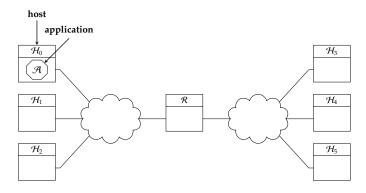


#### then, finally, noting

- a network provides connectivity between nodes, while
- ▶ an *inter*-network connects networks themselves together.



We can generalise partially connected network topologies

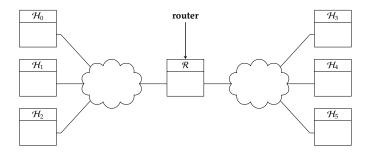


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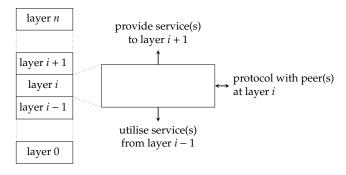
- a network provides connectivity between nodes, while
- an *inter*-network connects networks themselves together.



- An inter-network architecture seems a good approach!
  - ► Good:
    - scalable by virtue of a flexible and modular design, and
    - offers resilience against failure.
  - Bad: we need to cope with the fact
    - b ideally we support a multiplicity of communication mediums and protocols, and share access to both,
    - each host needs a globally unique address, and
    - router(s) need to know or discover how to communicate (or route) data from one host to another

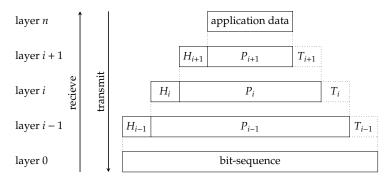
the difficulty of which is enhanced by a need to avoid impact on scalability (e.g., avoid centralised solutions).

Step #1: rather than use one monolithic protocol, decompose control, e.g.,



and hence provide various advantages stemming from modularity ...

• ... then support inter-layer (peer or otherwise) communication using **encapsulation**:



► Step #2: specify (abstract) layers we need, e.g.,

|         |         |          |   | OSI model    | Internet model |  |
|---------|---------|----------|---|--------------|----------------|--|
| layer 7 | 1       |          |   | Application  |                |  |
| layer 6 |         |          |   | Presentation | Application    |  |
| layer 5 | e       | uit      |   | Session      |                |  |
| layer 4 | recieve | transmit |   | Transport    | Transport      |  |
| layer 3 | re      | tr       |   | Network      | Internet       |  |
| layer 2 |         |          |   | Data link    | Link           |  |
| layer 1 |         |          | , | Physical     | <br>LIIK       |  |

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|                                  |                |         |          |         | OSI model    |          | Internet model |
|----------------------------------|----------------|---------|----------|---------|--------------|----------|----------------|
| PID {                            | layer 7        | e.      |          |         | Application  |          | Application    |
|                                  | layer 6        |         |          |         | Presentation |          |                |
|                                  | layer 5        |         | į.<br>į  |         | Session      |          |                |
| port number $\left\{ \right.$    | layer 4        | recieve | transmit |         | Transport    |          | Transport      |
| global address $\left\{ \right.$ | ress { layer 3 | Ĕ       |          | Network |              | Internet |                |
| local address {                  | layer 2        |         |          |         | Data link    | I in     | Link           |
|                                  | layer 1        |         |          | ļ       | Physical     |          | LIIK           |

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|--------------------|---------|---------|----------|---|--------------|--|----------------|
|                    | layer 7 |         |          |   | Application  |  |                |
| message (          | layer 6 |         |          |   | Presentation |  | Application    |
|                    | layer 5 | е       | uit      |   | Session      |  |                |
| segment, datagram  | layer 4 | recieve | transmit |   | Transport    |  | Transport      |
| packet {           | layer 3 | re      | tr       |   | Network      |  | Internet       |
| frame              | layer 2 |         |          |   | Data link    |  | Link           |
| symbol (e.g., bit) | layer 1 |         |          | , | Physical     |  | LIIK           |

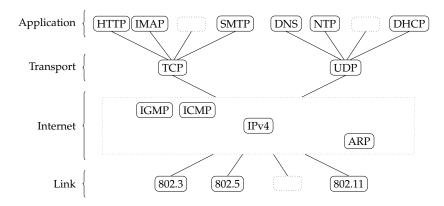
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|                    |         |         |          |  | OSI model    | Internet model |
|--------------------|---------|---------|----------|--|--------------|----------------|
|                    | layer 7 | ,       | 1        |  | Application  |                |
| haat               | layer 6 | recieve |          |  | Presentation | Application    |
| host               | layer 5 |         | li;      |  | Session      |                |
|                    | layer 4 |         | transmit |  | Transport    | Transport      |
| router             | layer 3 | re      | Ħ        |  | Network      | Internet       |
| switch, bridge     | layer 2 |         |          |  | Data link    | Link           |
| NIC, repeater, hub | layer 1 |         |          |  | Physical     | LIIK           |

► Step #2: specify (abstract) layers we need, e.g.,

|            |         |         |          |  | OSI model    | Internet model |
|------------|---------|---------|----------|--|--------------|----------------|
| user {     | layer 7 | 1       |          |  | Application  |                |
|            | layer 6 |         |          |  | Presentation | Application    |
|            | layer 5 | a       | uit      |  | Session      |                |
| kernel {   | layer 4 | recieve | transmit |  | Transport    | Transport      |
|            | layer 3 |         | ţį       |  | Network      | Internet       |
| hardware { | layer 2 |         |          |  | Data link    | Link           |
|            | layer 1 |         |          |  | Physical     | LIIK           |

Step #3: populate layers with (concrete) protocols, e.g.,

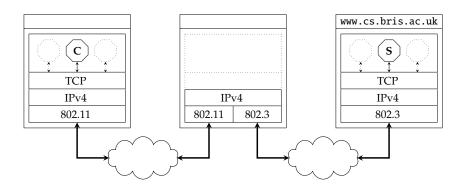


#### noting that

- a protocol graph is an abstract description of how protocols fit into the layered model, whereas
- ▶ a **protocol stack** is a concrete implementation of *one* top-to-bottom combination of protocols.

### Conclusions

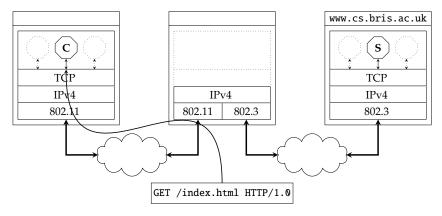
► As a final step, we can add some detail to our running example, e.g.,



with the rest of the unit aiming to further explain each layer.

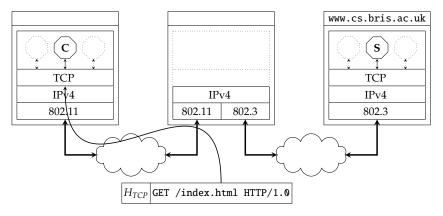
### Conclusions

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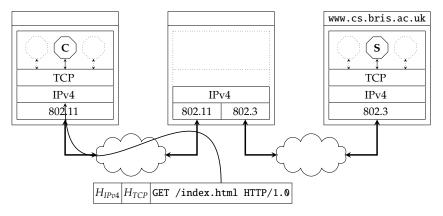


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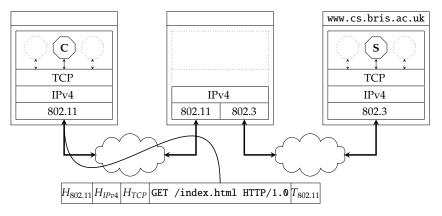
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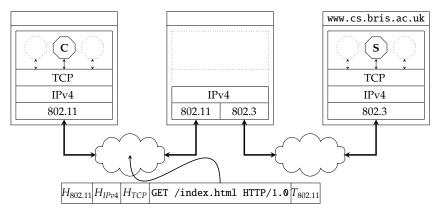
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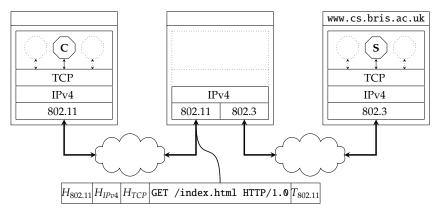
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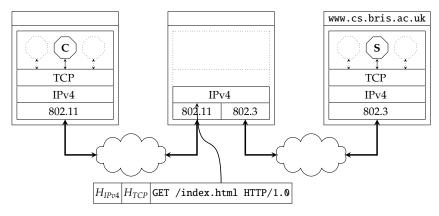
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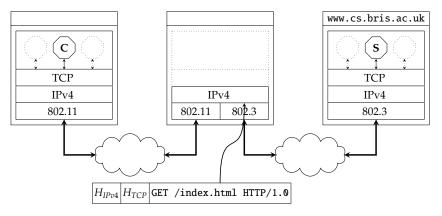
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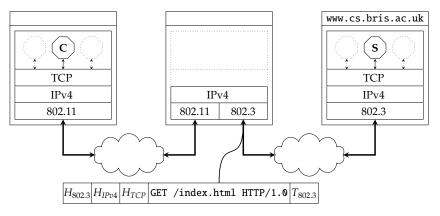
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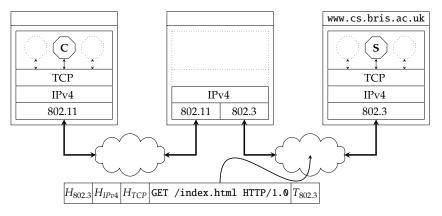
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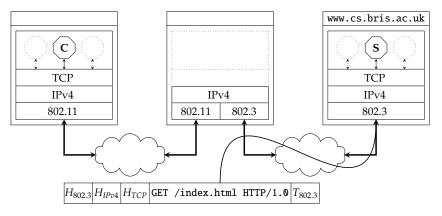
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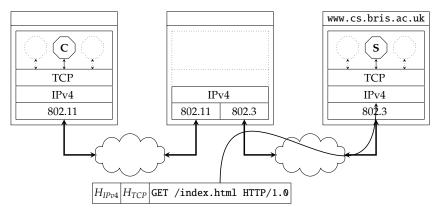
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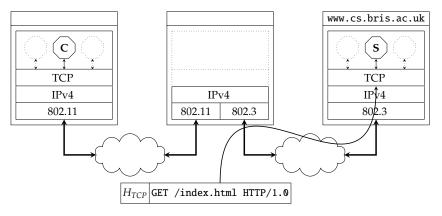
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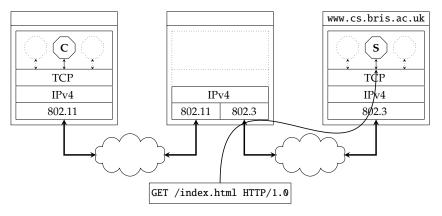
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#### Take away points:

- Although a lot of this content might seem very abstract, the point is that we now have a sensible high-level design.
- We'll fill in missing detail using a bottom-up approach by covering the
  - 1. link layer,
  - 2. internet layer,
  - 3. transport layer, and
  - 4. application layer

while emphasising the *general* concepts as applied in specific technologies ...

- ... keeping in mind the running, motivating example of HTTP.
- A central challenge throughout is design of solutions that work efficiently in the general-case, but guarentee correctness in (many) special- or corner-cases.

#### Additional Reading

- Wikipedia: Circuit switching. URL: http://en.wikipedia.org/wiki/Circuit\_switching.
- Wikipedia: Packet switching. url: http://en.wikipedia.org/wiki/Packet\_switching.
- ▶ Wikipedia: OSI model. URL: http://en.wikipedia.org/wiki/OSI\_model.
- Wikipedia: Internet protocol suite. URL: http://en.wikipedia.org/wiki/Internet\_protocol\_suite.
- W. Stallings. "Chapter 2: Data communications, data networks and the Internet". In: Data and Computer Communications. 9th ed. Pearson, 2010.
- W. Stallings. "Chapter 3: Protocol architecture, TCP/IP, and Internet-based applications". In: Data and Computer Communications.
  9th ed. Pearson, 2010.
- W. Stallings. "Chapter 11: Circuit switching and packet switching". In: Data and Computer Communications. 9th ed. Pearson, 2010.

#### References

- [1] Wikipedia: Circuit switching. url: http://en.wikipedia.org/wiki/Circuit\_switching (see p. 51).
- [2] Wikipedia: Internet protocol suite. URL: http://en.wikipedia.org/wiki/Internet\_protocol\_suite (see p. 51).
- [3] Wikipedia: OSI model. url: http://en.wikipedia.org/wiki/OSI\_model (see p. 51).
- [4] Wikipedia: Packet switching. url: http://en.wikipedia.org/wiki/Packet\_switching (see p. 51).
- [5] W. Stallings. "Chapter 11: Circuit switching and packet switching". In: Data and Computer Communications. 9th ed. Pearson, 2010 (see p. 51).
- [6] W. Stallings. "Chapter 2: Data communications, data networks and the Internet". In: Data and Computer Communications. 9th ed. Pearson, 2010 (see p. 51).
- [7] W. Stallings. "Chapter 3: Protocol architecture, TCP/IP, and Internet-based applications". In: Data and Computer Communications. 9th ed. Pearson, 2010 (see p. 51).
- [8] J. Postel. Transmision Control Protocol. Internet Engineering Task Force (IETF) Request for Comments (RFC) 793. 1981. URL: http://tools.ietf.org/html/rfc793.
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- [10] D. Cohen. "On Holy Wars and a Plea for Peace". In: IEEE Computer 14.10 (1981), pp. 48-54.

