

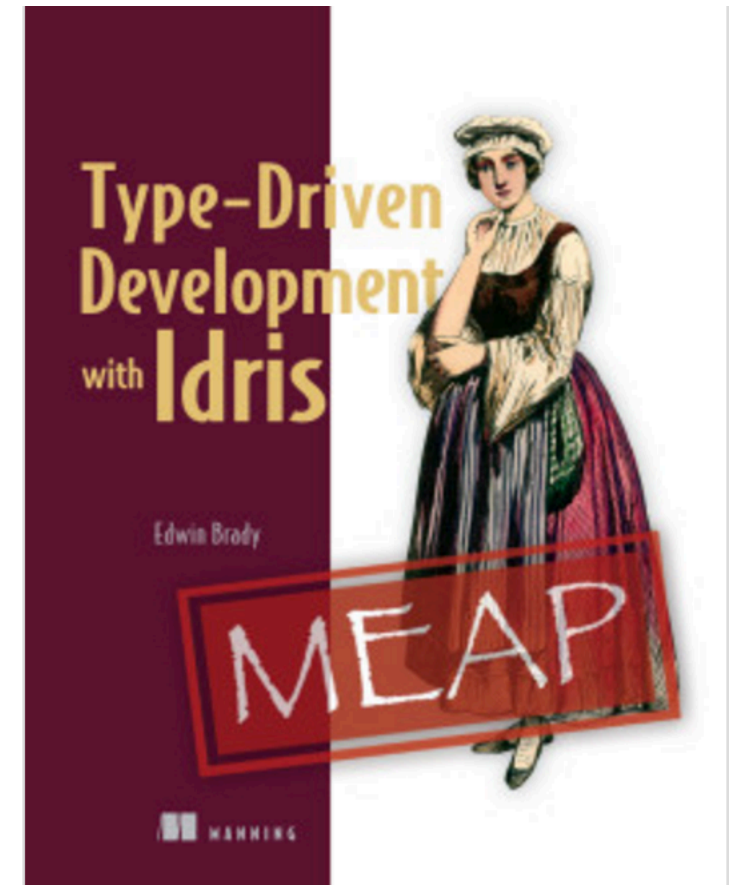
Introduction to Programming with Dependent Types in Idris

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Agenda

- Dependent types in a nutshell
- Getting started with Idris
- Interactive development with Types
- Programming with First Class Type

- This book (still work-in-progress) is written by the author of Idris
- Beginner-friendly with lots of examples and exercises (no particular background assumed)
- The talk will follow what has been written up to now



Functional Programming

- As opposed to imperative programming in C, C++, Java, etc.

```
quicksort [] = []
quicksort (p:xs) = (quicksort lesser) ++
                  [p] ++
                  (quicksort greater)
where
  lesser = filter (< p) xs
  greater = filter (>= p) xs
```

```
// To sort array a[] of size n: qsort(a,0,n-1)
void qsort(int a[], int lo, int hi)
{
    int h, l, p, t;

    if (lo < hi) {
        l = lo;
        h = hi;
        p = a[hi];

        do {
            while ((l < h) && (a[l] <= p))
                l = l+1;
            while ((h > l) && (a[h] >= p))
                h = h-1;
            if (l < h) {
                t = a[l];
                a[l] = a[h];
                a[h] = t;
            }
        } while (l < h);

        a[hi] = a[l];
        a[l] = p;

        qsort( a, lo, l-1 );
        qsort( a, l+1, hi );
    }
}
```

Types Matter (I)

- Types are a formal language that proves the absence of certain behaviors
- It has deep connections with varieties of logic, via the *Curry-Howard correspondence*

Types Matter (II)

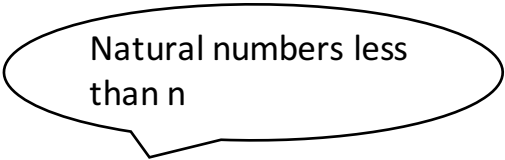
- The trend over the last 20 years goes toward more *expressive types* in modern programming languages
- They are capable of guaranteeing a large variety of program correctness properties:
 - Preventing resource allocation errors
 - Enforcing security in communication protocols
 - Preventing buffer overruns
 - Preventing data races and deadlocks in concurrent systems
 - ...

Dependent Types (I)

- Dependent types can depend on the values of terms
- Example: *Vect 3 Int*
- *appendV :: Vect n A -> Vect m A -> Vect (m + n) A*

Dependent Types (II)

- More broadly, dependent types allow programmers to put arbitrary specifications in types
- Example: avoid *out-of-bounds* error at compile time



Natural numbers less than n

lookup :: Vec Int n -> Fin n -> Int

lookup :: List Int -> Int -> Int

Getting Started with Idris

- Download <http://www.idris-lang.org/download/>
- Editors: Vim, Emacs, Atom, ...
- Fire up REPL 😊
 - :t
 - :doc

Example 1: Hello World

Types as First-class Constructs

- In most languages, there is syntactic separation between types and values (e.g., `x = int` is not allowed in Java)
- In Idris there are no such restrictions
 - Types are first-class
 - Any construct can appear as part of a type

Example 2: Calculating a Type

Interactive Programming with Types

- Dependent Types also give the compiler more information
- We follow the process of “Type, Define, Refine”
- Our editor provides interactive editing mode

Type, Define, Refine

An approach of type-driven development:

1. Write a *Type* to represent the system we are modelling
2. *Define* a function over that type
3. *Refine* the type and definition as necessary to capture any missing properties, or to help ensure totality

Interactive editing mode helps with the above process!

Example 3 : Type, Define, Refine

Type-directed Search

- Given enough information in the type, Idris can search for a valid expression satisfying that type
- This is extremely handy and shows how smart compiler can be

Example 4: Zipping two vectors

Parametricity

- Also known as “Free Theorem”
- You can know a lot by just looking at the type
- Example: $a \rightarrow a$, $a \rightarrow b \rightarrow a$
- A paper (“Theorem for free!”) written by Philip Wadler

Implicit Arguments

- Look again at the following type

appendV : Vect n A -> Vect m A -> Vect (m + n) A

- *n* and *A* are called type level variables
- They are not declared anywhere, and are thus referred as *implicit* arguments
- How can they be useful? How to use them in definition?

Example 5: Implicit arguments

Example 6: Sorting a vector

Example 7: Transposing a matrix

Programming with First Class Type

- In Idris, types can be manipulated just like any other language constructs.
- For example, we can give informative names to complex types (representing polygon)
- Dependent pattern matching (recall *valToString*)
- Differences with ordinary functions (runtime-representation, totality)

Example 8: A Type safe lookup function

- It should be possible to know at compile time whether the index is out of bounds when the program is run
- The *index* function is a type safe lookup function (from *Data.Vect* module):

$$\text{index} : \text{Fin } n \rightarrow \text{Vect } n \ a \rightarrow a$$

- *Fin n* denotes numbers that are strictly less than *n*
- But if the index is read from user input, the number isn't always going to be within the bound of the vector

Example 9: A Type safe *printf* Function

Want to Know More?

- Idris can also do theorem proving
- Idris Tutorial: <http://docs.idris-lang.org/en/latest/tutorial/>
- Excellent Talk by the author: https://youtu.be/X36ye-1x_HQ
- Questions?