

A

 Full Type System of SEDEL

 $\boxed{\Gamma \vdash A}$

(Well formedness)

WF-TOP $\frac{}{\Gamma \vdash \top}$	WF-INT $\frac{}{\Gamma \vdash \text{Int}}$	WF-ARROW $\frac{\Gamma \vdash A \quad \Gamma \vdash B}{\Gamma \vdash A \rightarrow B}$	WF-RECORD $\frac{\Gamma \vdash A}{\Gamma \vdash \{l : A\}}$	WF-VAR $\frac{\alpha * A \in \Gamma}{\Gamma \vdash \alpha}$
WF-AND $\frac{\Gamma \vdash A \quad \Gamma \vdash B \quad \Gamma \vdash A * B}{\Gamma \vdash A \& B}$	WF-FORALL $\frac{\Gamma \vdash A \quad \Gamma, \alpha * A \vdash B}{\Gamma \vdash \forall (\alpha * A). B}$	WF-TRAIT $\frac{\Gamma \vdash A \quad \Gamma \vdash B}{\Gamma \vdash \mathbf{Trait}[A, B]}$		

 $\boxed{A <: B}$

(Subtyping)

SUB-TOP	SUB-ANDR	SUB-INT	SUB-ANDL1
$\frac{}{A <: \top}$	$\frac{A_1 <: A_2 \quad A_1 <: A_3}{A_1 <: A_2 \& A_3}$	$\frac{}{\text{Int} <: \text{Int}}$	$\frac{A_1 <: A_3 \quad A_3 \text{ ordinary}}{A_1 \& A_2 <: A_3}$
SUB-ANDL2	SUB-REC	SUB-VAR	SUB-ARR
$\frac{A_2 <: A_3 \quad A_3 \text{ ordinary}}{A_1 \& A_2 <: A_3}$	$\frac{A <: B}{\{l : A\} <: \{l : B\}}$	$\frac{}{\alpha <: \alpha}$	$\frac{B_1 <: A_1 \quad A_2 <: B_2}{A_1 \rightarrow A_2 <: B_1 \rightarrow B_2}$
SUB-FORALL	SUB-TRAIT		
$\frac{B_1 <: B_2 \quad A_2 <: A_1}{\forall (\alpha * A_1). B_1 <: \forall (\alpha * A_2). B_2}$	$\frac{B_1 <: A_1 \quad A_2 <: B_2}{\mathbf{Trait} [A_1, A_2] <: \mathbf{Trait} [B_1, B_2]}$		

 $\boxed{\Gamma \vdash A * B}$

(Disjointness)

$\frac{\text{D-TOP}}{\Gamma \vdash \top * A}$	$\frac{\text{D-TOPSYM}}{\Gamma \vdash A * \top}$	$\frac{\text{D-VAR} \quad \alpha * A \in \Gamma \quad A <: B}{\Gamma \vdash \alpha * B}$	$\frac{\text{D-VARSYM} \quad \alpha * A \in \Gamma \quad A <: B}{\Gamma \vdash B * \alpha}$
$\frac{\text{D-FORALL} \quad \Gamma, \alpha * A_1 \& A_2 \vdash B * C}{\Gamma \vdash \forall (\alpha * A_1). B * \forall (\alpha * A_2). C}$	$\frac{\text{D-REC} \quad \Gamma \vdash A * B}{\Gamma \vdash \{l : A\} * \{l : B\}}$	$\frac{\text{D-REC} \quad l_1 \neq l_2}{\Gamma \vdash \{l_1 : A\} * \{l_2 : B\}}$	
$\frac{\text{D-ARROW} \quad \Gamma \vdash A_2 * B_2}{\Gamma \vdash A_1 \rightarrow A_2 * B_1 \rightarrow B_2}$	$\frac{\text{D-ANDL} \quad \Gamma \vdash A_1 * B \quad \Gamma \vdash A_2 * B}{\Gamma \vdash A_1 \& A_2 * B}$	$\frac{\text{D-ANDR} \quad \Gamma \vdash A * B_1 \quad \Gamma \vdash A * B_2}{\Gamma \vdash A * B_1 \& B_2}$	
$\frac{\text{D-TRAIT} \quad \Gamma \vdash A_2 * B_2}{\Gamma \vdash \mathbf{Trait}[A_1, A_2] * \mathbf{Trait}[B_1, B_2]}$	$\frac{\text{D-TRAITARR1} \quad \Gamma \vdash A_2 * B_2}{\Gamma \vdash \mathbf{Trait}[A_1, A_2] * B_1 \rightarrow B_2}$		
$\frac{\text{D-TRAITARR2} \quad \Gamma \vdash A_2 * B_2}{\Gamma \vdash A_1 \rightarrow A_2 * \mathbf{Trait}[B_1, B_2]}$	$\frac{\text{D-AX} \quad A *_{ax} B}{\Gamma \vdash A * B}$		

 $\boxed{A *_{ax} B}$

(Disjointness axiom)

DAX-SYM $\frac{B *_{ax} A}{A *_{ax} B}$	DAX-INTARR $\frac{}{\text{Int} *_{ax} A_1 \rightarrow A_2}$	DAX-INTREC $\frac{}{\text{Int} *_{ax} \{l : A\}}$	DAX-INTFORALL $\frac{}{\text{Int} *_{ax} \forall (\alpha * B_1). B_2}$
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$\frac{}{\text{Int} *_{ax} \mathbf{Trait} [A_1, A_2]}$	$\frac{}{A_1 \rightarrow A_2 *_{ax} \forall (\alpha * B_1). B_2}$	$\frac{}{A_1 \rightarrow A_2 *_{ax} \{l : B\}}$
DAX-TRAITFORALL	DAX-TRAITREC	DAX-ARRREC
$\frac{}{\mathbf{Trait} [A_1, A_2] *_{ax} \forall (\alpha * B_1). B_2}$	$\frac{}{\mathbf{Trait} [A_1, A_2] *_{ax} \{l : B\}}$	$\frac{}{\forall (\alpha * A_1). A_2 *_{ax} \{l : B\}}$
<div style="border: 1px solid black; display: inline-block; padding: 2px 5px;"> $\Gamma \vdash E \Rightarrow A \rightsquigarrow e$ </div> <div style="margin-left: 20px;">(Infer)</div>		
$\frac{}{\Gamma \vdash \top \Rightarrow \top \rightsquigarrow \top}$	$\frac{}{\Gamma \vdash i \Rightarrow \text{Int} \rightsquigarrow i}$	$\frac{}{\Gamma \vdash x \Rightarrow A \rightsquigarrow x}$
INF-TOP	INF-INT	INF-VAR
		$\frac{x : A \in \Gamma}{\Gamma \vdash x \Rightarrow A \rightsquigarrow x}$
		$\frac{}{\Gamma \vdash E \Leftarrow A \rightsquigarrow e}$
		$\frac{}{\Gamma \vdash E : A \Rightarrow A \rightsquigarrow e : A }$
$\frac{}{\Gamma \vdash E_1 \Rightarrow A_1 \rightarrow A_2 \rightsquigarrow e_1 \quad \Gamma \vdash E_2 \Leftarrow A_1 \rightsquigarrow e_2}$		
$\frac{}{\Gamma \vdash E_1 E_2 \Rightarrow A_2 \rightsquigarrow e_1 e_2}$		
$\frac{}{\Gamma \vdash E \Rightarrow \forall (\alpha * B). C \rightsquigarrow e \quad \Gamma \vdash A \quad \Gamma \vdash A * B}$		
$\frac{}{\Gamma \vdash E A \Rightarrow [A/\alpha]C \rightsquigarrow e A }$		
$\frac{}{\Gamma \vdash E_1 \Rightarrow A \rightsquigarrow e_1 \quad \Gamma \vdash E_2 \Rightarrow B \rightsquigarrow e_2 \quad \Gamma \vdash A * B}$		
$\frac{}{\Gamma \vdash E_1, , E_2 \Rightarrow A \& B \rightsquigarrow e_1, , e_2}$		
$\frac{}{\Gamma \vdash E \Rightarrow \{l : A\} \rightsquigarrow e}$		
$\frac{}{\Gamma \vdash \Lambda (\alpha * A). E \Rightarrow \forall (\alpha * A). B \rightsquigarrow \Lambda (\alpha * A). e}$		
$\frac{}{\Gamma \vdash E \Rightarrow A \rightsquigarrow e}$		
$\frac{}{\Gamma \vdash \{l = E\} \Rightarrow \{l : A\} \rightsquigarrow \{l = e\}}$		
$\frac{}{\Gamma \vdash E.l \Rightarrow A \rightsquigarrow e.l}$		
$\frac{}{\Gamma \vdash \Lambda (\alpha * A). E \Rightarrow \forall (\alpha * A). B \rightsquigarrow \Lambda (\alpha * A). e}$		
$\frac{}{\Gamma, x : A \vdash E_1 \Leftarrow A \rightsquigarrow e_1 \quad \Gamma, x : A \vdash E_2 \Rightarrow B \rightsquigarrow e_2}$		
$\frac{}{\Gamma \vdash \mathbf{letrec} x : A = E_1 \mathbf{in} E_2 \Rightarrow B \rightsquigarrow \mathbf{letrec} x : A = e_1 \mathbf{in} e_2}$		
$\frac{}{\Gamma, \mathbf{self} : B \vdash E_i \Rightarrow \mathbf{Trait} [B_i, C_i] \rightsquigarrow e_i^{i \in 1..n} \quad \Gamma, \mathbf{self} : B \vdash \{\bar{l}_j = \bar{E}'_j^{j \in 1..m}\} \Rightarrow C \rightsquigarrow e}$		
$\frac{}{\Gamma \vdash \mathbf{trait} [\mathbf{self} : B] \mathbf{inherits} \bar{E}_i^{i \in 1..n} \{\bar{l}_j = \bar{E}'_j^{j \in 1..m}\} : A \Rightarrow \mathbf{Trait} [B, A] \rightsquigarrow \lambda (\mathbf{self} : B). ((\bar{e}_i \mathbf{self})^{i \in 1..n}), , e}$		
$\frac{}{\Gamma, \mathbf{self} : B \vdash E_i \Rightarrow \mathbf{Trait} [B_i, C_i] \rightsquigarrow e_i^{i \in 1..n} \quad \Gamma, \mathbf{self} : B, \mathbf{super} : C_1 \& .. \& C_n \vdash \{\bar{l}_j = \bar{E}'_j^{j \in 1..m}\} \Rightarrow C \rightsquigarrow e}$		
$\frac{}{\Gamma \vdash \mathbf{trait} [\mathbf{self} : B] \mathbf{inherits} \bar{E}_i^{i \in 1..n} \{\bar{l}_j = \bar{E}'_j^{j \in 1..m}\} : A \Rightarrow \mathbf{Trait} [B, A] \rightsquigarrow \lambda (\mathbf{self} : B). \mathbf{let} \mathbf{super} = (\bar{e}_i \mathbf{self})^{i \in 1..n} \mathbf{in} \mathbf{super}, , e}$		
$\frac{}{\Gamma \vdash E_i \Rightarrow \mathbf{Trait} [A_i, B_i] \rightsquigarrow e_i^{i \in 1..n} \quad \bar{A} <: \bar{A}_i^{i \in 1..n} \quad \Gamma \vdash B_1 \& .. \& B_n \quad B_1 \& .. \& B_n <: A}$		
$\frac{}{\Gamma \vdash \mathbf{new} [A] (\bar{E}_i^{i \in 1..n}) \Rightarrow A \rightsquigarrow \mathbf{letrec} \mathbf{self} : A = (\bar{e}_i \mathbf{self})^{i \in 1..n} \mathbf{in} \mathbf{self}}$		
$\frac{}{\Gamma \vdash E_1 \Rightarrow \mathbf{Trait} [A, B] \rightsquigarrow e_1 \quad \Gamma \vdash E_2 \Leftarrow A \rightsquigarrow e_2}$		
$\frac{}{\Gamma \vdash E_1 \wedge E_2 \Rightarrow B \rightsquigarrow e_1 e_2}$		

$$\boxed{\Gamma \vdash E \Leftarrow A \rightsquigarrow e}$$

(Check)

$$\frac{\text{CHK-ABS} \quad \Gamma \vdash A \quad \Gamma, x : A \vdash E \Leftarrow B \rightsquigarrow e}{\Gamma \vdash \lambda x. E \Leftarrow A \rightarrow B \rightsquigarrow \lambda x. e}$$

$$\frac{\text{CHK-SUB} \quad \Gamma \vdash E \Rightarrow A \rightsquigarrow e \quad A <: B \quad \Gamma \vdash B}{\Gamma \vdash E \Leftarrow B \rightsquigarrow e}$$

► **Definition 7** (Type translation).

$$\begin{aligned} |\alpha| &= \alpha \\ |\text{Int}| &= \text{Int} \\ |\top| &= \top \\ |A \rightarrow B| &= |A| \rightarrow |B| \\ |A \& B| &= |A| \& |B| \\ |\{l : A\}| &= \{l : |A|\} \\ |\forall (\alpha * A). B| &= \forall (\alpha * |A|). |B| \\ |\mathbf{Trait}[A, B]| &= |A| \rightarrow |B| \end{aligned}$$

B Metatheory

► **Lemma 8.** If $\Gamma \vdash A$ then $|\Gamma| \vdash |A|$.

Proof. By simple induction on the derivation of the judgement. ◀

► **Lemma 9.** If $A *_{ax} B$ then $|A| *_{ax} |B|$.

Proof. Note that $|\mathbf{Trait}[A, B]| = |A| \rightarrow |B|$, the rest are immediate. ◀

► **Lemma 10.** If $A <: B$ then $|A| <: |B|$.

Proof. Most of them are just F_i subtyping. We only show rule SUB-TRAIT.

■

$$\frac{\text{SUB-TRAIT} \quad B_1 <: A_1 \quad A_2 <: B_2}{\mathbf{Trait}[A_1, A_2] <: \mathbf{Trait}[B_1, B_2]}$$

$$\begin{array}{ll} |B_1| <: |A_1| & \text{By i.h.} \\ |A_2| <: |B_2| & \text{By i.h.} \\ |A_1| \rightarrow |A_2| <: |B_1| \rightarrow |B_2| & \text{By } S \rightarrow \end{array}$$

◀

► **Lemma 11.** If $\Gamma \vdash A * B$ then $|\Gamma| \vdash |A| * |B|$.

Proof. By induction on the derivation of the judgement.

■ Rules D-TOP, D-TOPSYM, and D-RECN are immediate.

■

$$\frac{\text{D-VAR} \quad \alpha * A \in \Gamma \quad A <: B}{\Gamma \vdash \alpha * B}$$

$|A| <: |B|$ By Lemma 10
 $\alpha * A \in \Gamma$ Given
 $\alpha * |A| \in |\Gamma|$ Above
 $|\Gamma| \vdash \alpha * |B|$ By $D\alpha$

■

$$\frac{\text{D-VARSYM} \quad \alpha * A \in \Gamma \quad A <: B}{\Gamma \vdash B * \alpha}$$

$|A| <: |B|$ By Lemma 10
 $\alpha * A \in \Gamma$ Given
 $\alpha * |A| \in |\Gamma|$ Above
 $|\Gamma| \vdash |B| * \alpha$ By $D\alpha\text{SYM}$

■

$$\frac{\text{D-FORALL} \quad \Gamma, \alpha * A_1 \ \& \ A_2 \vdash B * C}{\Gamma \vdash \forall (\alpha * A_1). B * \forall (\alpha * A_2). C}$$

$|\Gamma|, \alpha * |A_1| \ \& \ |A_2| \vdash |B| * |C|$ By i.h.
 $|\Gamma| \vdash \forall (\alpha * |A_1|). B * \forall (\alpha * |A_2|). C$ By $D\forall$

■

$$\frac{\text{D-REC} \quad \Gamma \vdash A * B}{\Gamma \vdash \{l : A\} * \{l : B\}}$$

$|\Gamma| \vdash |A| * |B|$ By i.h.
 $|\Gamma| \vdash \{l : |A|\} * \{l : |B|\}$ By $D\text{REC}=_$

■

$$\frac{\text{D-ARROW} \quad \Gamma \vdash A_2 * B_2}{\Gamma \vdash A_1 \rightarrow A_2 * B_1 \rightarrow B_2}$$

$|\Gamma| \vdash |A_2| * |B_2|$ By i.h.
 $|\Gamma| \vdash |A_1| \rightarrow |A_2| * |B_1| \rightarrow |B_2|$ By $D\rightarrow$

■

$$\frac{\text{D-ANDL} \quad \Gamma \vdash A_1 * B \quad \Gamma \vdash A_2 * B}{\Gamma \vdash A_1 \ \& \ A_2 * B}$$

$|\Gamma| \vdash |A_1| * |B|$ By i.h.
 $|\Gamma| \vdash |A_2| * |B|$ By i.h.
 $|\Gamma| \vdash |A_1| \& |A_2| * |B|$ By D&L

■

$$\frac{\text{D-ANDR} \quad \Gamma \vdash A * B_1 \quad \Gamma \vdash A * B_2}{\Gamma \vdash A * B_1 \& B_2}$$

$|\Gamma| \vdash |A| * |B_1|$ By i.h.
 $|\Gamma| \vdash |A| * |B_2|$ By i.h.
 $|\Gamma| \vdash |A| * |B_1| \& |B_2|$ By D&R

■

$$\frac{\text{D-TRAIT} \quad \Gamma \vdash A_2 * B_2}{\Gamma \vdash \mathbf{Trait}[A_1, A_2] * \mathbf{Trait}[B_1, B_2]}$$

$|\Gamma| \vdash |A_2| * |B_2|$ By i.h.
 $|\Gamma| \vdash |A_1| \rightarrow |A_2| * |B_1| \rightarrow |B_2|$ By D \rightarrow

■

$$\frac{\text{D-TRAITARR1} \quad \Gamma \vdash A_2 * B_2}{\Gamma \vdash \mathbf{Trait}[A_1, A_2] * B_1 \rightarrow B_2}$$

$|\Gamma| \vdash |A_2| * |B_2|$ By i.h.
 $|\Gamma| \vdash |A_1| \rightarrow |A_2| * |B_1| \rightarrow |B_2|$ By D \rightarrow

■

$$\frac{\text{D-TRAITARR2} \quad \Gamma \vdash A_2 * B_2}{\Gamma \vdash A_1 \rightarrow A_2 * \mathbf{Trait}[B_1, B_2]}$$

$|\Gamma| \vdash |A_2| * |B_2|$ By i.h.
 $|\Gamma| \vdash |A_1| \rightarrow |A_2| * |B_1| \rightarrow |B_2|$ By D \rightarrow

■

$$\frac{\text{D-AX} \quad A *_{ax} B}{\Gamma \vdash A * B}$$

$|A| *_{ax} |B|$ By Lemma 9
 $|\Gamma| \vdash |A| * |B|$ By DAX



► **Theorem 12** (Type-safe translation). *We have that:*

- If $\Gamma \vdash E \Rightarrow A \rightsquigarrow e$ then $|\Gamma| \vdash e \Rightarrow |A|$.
- If $\Gamma \vdash E \Leftarrow A \rightsquigarrow e$ then $|\Gamma| \vdash e \Leftarrow |A|$.

Proof. By induction on the typing judgement.

- Rules INF-TOP, INF-INT, and INF-VAR are immediate.

■

$$\frac{\text{INF-ANNO} \quad \Gamma \vdash E \Leftarrow A \rightsquigarrow e}{\Gamma \vdash E : A \Rightarrow A \rightsquigarrow e : |A|}$$

$$\begin{array}{ll} |\Gamma| \vdash e \Leftarrow |A| & \text{By i.h.} \\ |\Gamma| \vdash e : |A| \Rightarrow |A| & \text{By TI-ANNO} \end{array}$$

■

$$\frac{\text{INF-APP} \quad \Gamma \vdash E_1 \Rightarrow A_1 \rightarrow A_2 \rightsquigarrow e_1 \quad \Gamma \vdash E_2 \Leftarrow A_1 \rightsquigarrow e_2}{\Gamma \vdash E_1 E_2 \Rightarrow A_2 \rightsquigarrow e_1 e_2}$$

$$\begin{array}{ll} |\Gamma| \vdash e_1 \Rightarrow |A_1| \rightarrow |A_2| & \text{By i.h.} \\ |\Gamma| \vdash e_2 \Leftarrow |A_2| & \text{By i.h.} \\ |\Gamma| \vdash e_1 e_2 \Rightarrow |A_2| & \text{By TI-APP} \end{array}$$

■

$$\frac{\text{INF-TAPP} \quad \Gamma \vdash E \Rightarrow \forall (\alpha * B). C \rightsquigarrow e \quad \Gamma \vdash A \quad \Gamma \vdash A * B}{\Gamma \vdash E A \Rightarrow [A/\alpha]C \rightsquigarrow e |A|}$$

$$\begin{array}{ll} |\Gamma| \vdash e \Rightarrow \forall (\alpha * |B|). |C| & \text{By i.h.} \\ |\Gamma| \vdash |A| & \text{By Lemma 8} \\ |\Gamma| \vdash |A| * |B| & \text{By Lemma 11} \\ |\Gamma| \vdash e |A| \Rightarrow [\alpha \mapsto |A|]|C| & \text{By TI-TAPP} \end{array}$$

■

$$\frac{\text{INF-MERGE} \quad \Gamma \vdash E_1 \Rightarrow A \rightsquigarrow e_1 \quad \Gamma \vdash E_2 \Rightarrow B \rightsquigarrow e_2 \quad \Gamma \vdash A * B}{\Gamma \vdash E_1, E_2 \Rightarrow A \& B \rightsquigarrow e_1, e_2}$$

$$\begin{array}{ll} |\Gamma| \vdash e_1 \Rightarrow |A| & \text{By i.h.} \\ |\Gamma| \vdash e_2 \Rightarrow |B| & \text{By i.h.} \\ |\Gamma| \vdash |A| * |B| & \text{By Lemma 11} \\ |\Gamma| \vdash e_1, e_2 \Rightarrow |A| \& |B| & \text{By TI-MERGE} \end{array}$$

■

$$\frac{\text{INF-REC} \quad \Gamma \vdash E \Rightarrow A \rightsquigarrow e}{\Gamma \vdash \{l = E\} \Rightarrow \{l : A\} \rightsquigarrow \{l = e\}}$$

$$\begin{array}{ll} |\Gamma| \vdash e \Rightarrow |A| & \text{By i.h.} \\ |\Gamma| \vdash \{l = e\} \Rightarrow \{l : |A|\} & \text{By TI-REC} \end{array}$$

■

$$\frac{\text{INF-PROJ} \quad \Gamma \vdash E \Rightarrow \{l : A\} \rightsquigarrow e}{\Gamma \vdash E.l \Rightarrow A \rightsquigarrow e.l}$$

$$\begin{array}{ll} |\Gamma| \vdash e \Rightarrow \{l : |A|\} & \text{By i.h.} \\ |\Gamma| \vdash e.l \Rightarrow |A| & \text{By TI-PROJ} \end{array}$$

■

$$\frac{\text{INF-BLAM} \quad \Gamma \vdash A \quad \Gamma, \alpha * A \vdash E \Rightarrow B \rightsquigarrow e}{\Gamma \vdash \Lambda(\alpha * A).E \Rightarrow \forall(\alpha * A).B \rightsquigarrow \Lambda(\alpha * |A|).e}$$

$$\begin{array}{ll} |\Gamma| \vdash |A| & \text{By Lemma 8} \\ |\Gamma|, \alpha * |A| \vdash e \Rightarrow |B| & \text{By i.h.} \\ |\Gamma| \vdash \Lambda(\alpha * |A|).e \Rightarrow \forall(\alpha * |A|).|B| & \text{By TI-BLAM} \end{array}$$

■

$$\frac{\text{INF-LETE} \quad \Gamma, x : A \vdash E_1 \Leftarrow A \rightsquigarrow e_1 \quad \Gamma, x : A \vdash E_2 \Rightarrow B \rightsquigarrow e_2}{\Gamma \vdash \mathbf{letrec} \, x : A = E_1 \mathbf{in} \, E_2 \Rightarrow B \rightsquigarrow \mathbf{letrec} \, x : |A| = e_1 \mathbf{in} \, e_2}$$

$$\begin{array}{ll} |\Gamma|, x : |A| \vdash e_1 \Leftarrow |A| & \text{By i.h.} \\ |\Gamma|, x : |A| \vdash e_2 \Rightarrow |B| & \text{By i.h.} \\ |\Gamma| \vdash \mathbf{letrec} \, x : |A| = e_1 \mathbf{in} \, e_2 \Rightarrow |B| & \text{By TI-LETE} \end{array}$$

■

$$\frac{\text{INF-NEW} \quad \frac{\Gamma \vdash E_i \Rightarrow \mathbf{Trait}[A_i, B_i] \rightsquigarrow e_i}{A <: A_i}^{i \in 1..n} \quad \Gamma \vdash B_1 \& \dots \& B_n \quad B_1 \& \dots \& B_n <: A}{\Gamma \vdash \mathbf{new} \, [A](\overline{E_i}^{i \in 1..n}) \Rightarrow A \rightsquigarrow \mathbf{letrec} \, \mathbf{self} : |A| = (\overline{e_i \mathbf{self}})^{i \in 1..n} \mathbf{in} \, \mathbf{self}}$$

$$\begin{array}{ll} |\Gamma| \vdash e_i \Rightarrow |A_i| \rightarrow |B_i| & \text{By i.h.} \\ |A| <: |A_i| & \text{By Lemma 10} \\ |\Gamma| \vdash |B_1| \& \dots \& |B_n| & \text{By Lemma 11} \end{array}$$

$ B_1 \& \dots \& B_n <: A $	By Lemma 10
$ \Gamma , \text{self} : A \vdash \text{self} \Rightarrow A $	By TI-VAR
$ \Gamma , \text{self} : A \vdash \text{self} \Leftarrow A_i $	By TC-SUB
$ \Gamma , \text{self} : A \vdash e_i \text{ self} \Rightarrow B_i $	By TI-APP
$ \Gamma , \text{self} : A \vdash (e_1 \text{ self}), \dots, (e_n \text{ self}) \Rightarrow B_1 \& \dots \& B_n $	By TI-MERGE and above
$ \Gamma , \text{self} : A \vdash (e_1 \text{ self}), \dots, (e_n \text{ self}) \Leftarrow A $	By TC-SUB
$ \Gamma \vdash \text{letrec self} : A = (e_1 \text{ self}), \dots, (e_n \text{ self}) \text{ in self} \Rightarrow A $	By TI-LETE

INF-TRAIT

$$\frac{\frac{\frac{\Gamma, \text{self} : B \vdash E_i \Rightarrow \mathbf{Trait}[B_i, C_i] \rightsquigarrow e_i^{i \in 1..n}}{\Gamma, \text{self} : B \vdash \{ \overline{l_j = E_j^{j \in 1..m}}} \Rightarrow C \rightsquigarrow e}}{B <: \overline{B_i^{i \in 1..n}} \quad \Gamma \vdash C_1 \& \dots \& C_n \& C \quad C_1 \& \dots \& C_n \& C <: A}}{\Gamma \vdash \mathbf{trait}[\text{self} : B] \mathbf{inherits} \overline{E_i^{i \in 1..n}} \{ \overline{l_j = E_j^{j \in 1..m}}} : A \Rightarrow \mathbf{Trait}[B, A] \rightsquigarrow \lambda(\text{self} : |B|). ((e_i \text{ self})^{i \in 1..n}), e}$$

$ \Gamma , \text{self} : B \vdash e_i \Rightarrow B_i \rightarrow C_i $	By i.h.
$ \Gamma , \text{self} : B \vdash e \Rightarrow C $	By i.h.
$ B <: B_i $	By Lemma 10
$ \Gamma \vdash C_1 \& \dots \& C_n \& C $	By Lemma 8
$ C_1 \& \dots \& C_n \& C <: A $	By Lemma 10
$ \Gamma , \text{self} : B \vdash \text{self} \Rightarrow B $	By TI-VAR
$ \Gamma , \text{self} : B \vdash \text{self} \Leftarrow B_i $	By TC-SUB
$ \Gamma , \text{self} : B \vdash e_i \text{ self} \Rightarrow C_i $	By TI-APP
$ \Gamma , \text{self} : B \vdash (e_1 \text{ self}), \dots, (e_n \text{ self}), e \Rightarrow C_1 \& \dots \& C_n \& C $	By TI-MERGE
$ \Gamma , \text{self} : B \vdash (e_1 \text{ self}), \dots, (e_n \text{ self}), e \Leftarrow A $	By TC-SUB
$ \Gamma \vdash \lambda(\text{self} : B). (e_1 \text{ self}), \dots, (e_n \text{ self}), e \Rightarrow B \rightarrow A $	By TI-ABS (annotated lambda typing)

INF-TRAITSUPER

$$\frac{\frac{\frac{\Gamma, \text{self} : B \vdash E_i \Rightarrow \mathbf{Trait}[B_i, C_i] \rightsquigarrow e_i^{i \in 1..n}}{\Gamma, \text{self} : B, \text{super} : C_1 \& \dots \& C_n \vdash \{ \overline{l_j = E_j^{j \in 1..m}}} \Rightarrow C \rightsquigarrow e}}{B <: \overline{B_i^{i \in 1..n}} \quad \Gamma \vdash C_1 \& \dots \& C_n \& C \quad C_1 \& \dots \& C_n \& C <: A}}{\Gamma \vdash \mathbf{trait}[\text{self} : B] \mathbf{inherits} \overline{E_i^{i \in 1..n}} \{ \overline{l_j = E_j^{j \in 1..m}}} : A \Rightarrow \mathbf{Trait}[B, A] \rightsquigarrow \lambda(\text{self} : |B|). \text{let super} = (\overline{e_i \text{ self}})^{i \in 1..n} \text{ in super}, e}$$

$ \Gamma , \text{self} : B \vdash e_i \Rightarrow B_i \rightarrow C_i $	By i.h.
$ \Gamma , \text{self} : B , \text{super} : C_1 \& \dots \& C_n \vdash e \Rightarrow C $	By i.h.
$ B <: B_i $	By Lemma 10
$ \Gamma \vdash C_1 \& \dots \& C_n \& C $	By Lemma 8
$ C_1 \& \dots \& C_n \& C <: A $	By Lemma 10
$ \Gamma , \text{self} : B \vdash \text{self} \Rightarrow B $	By TI-VAR
$ \Gamma , \text{self} : B \vdash \text{self} \Leftarrow B_i $	By TC-SUB
$ \Gamma , \text{self} : B \vdash e_i \text{ self} \Rightarrow C_i $	By TI-APP
$ \Gamma , \text{self} : B \vdash (e_1 \text{ self}), \dots, (e_n \text{ self}) \Rightarrow C_1 \& \dots \& C_n $	By TI-MERGE
$ \Gamma , \text{self} : B , \text{super} : C_1 \& \dots \& C_n \vdash \text{super}, e \Rightarrow C_1 \& \dots \& C_n \& C $	By TI-MERGE

$ \Gamma , \text{self} : B \vdash \text{let super} = \overline{(e_i \text{ self})}^{i \in 1..n} \text{ in super}, , e \Rightarrow C_1 \& \dots \& C_n \& C $	By TI-LET (non-recursive let)
$ \Gamma , \text{self} : B \vdash \text{let super} = \overline{(e_i \text{ self})}^{i \in 1..n} \text{ in super}, , e \Leftarrow A $	By TC-SUB
$ \Gamma \vdash \lambda(\text{self} : B). \text{let super} = \overline{(e_i \text{ self})}^{i \in 1..n} \text{ in super}, , e \Rightarrow B \rightarrow A $	By TI-ABS (annotated lambda typing)

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$$\frac{\text{INF-FORWARD} \quad \Gamma \vdash E_1 \Rightarrow \mathbf{Trait} [A, B] \rightsquigarrow e_1 \quad \Gamma \vdash E_2 \Leftarrow A \rightsquigarrow e_2}{\Gamma \vdash E_1 \wedge E_2 \Rightarrow B \rightsquigarrow e_1 e_2}$$

$ \Gamma \vdash e_1 \Rightarrow A \rightarrow B $	By i.h.
$ \Gamma \vdash e_2 \Leftarrow A $	By i.h.
$ \Gamma \vdash e_1 e_2 \Rightarrow B $	By TI-APP

$$\frac{\text{CHK-ABS} \quad \Gamma \vdash A \quad \Gamma, x : A \vdash E \Leftarrow B \rightsquigarrow e}{\Gamma \vdash \lambda x. E \Leftarrow A \rightarrow B \rightsquigarrow \lambda x. e}$$

■

$ \Gamma \vdash A $	By Lemma 8
$ \Gamma , x : A \vdash e \Leftarrow B $	By i.h.
$ \Gamma \vdash \lambda x. e \Leftarrow A \rightarrow B $	By TC-ABS

$$\frac{\text{CHK-SUB} \quad \Gamma \vdash E \Rightarrow A \rightsquigarrow e \quad A <: B \quad \Gamma \vdash B}{\Gamma \vdash E \Leftarrow B \rightsquigarrow e}$$

■

$ \Gamma \vdash e \Rightarrow A $	By i.h.
$ A <: B $	By Lemma 10
$ \Gamma \vdash B $	By Lemma 8
$ \Gamma \vdash e \Leftarrow B $	By TC-SUB

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