

I haven't thought of a title yet

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# What is the first

This is my first time to write notes all in English.

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## 1 Preliminaries

### 1.1 Set and Collections

**Definition 1.1.1** *Let*  $\mathcal{D}$  *be a collection of sets. The union of*  $\mathcal{C}$ *, denoted by*  $\bigcup \mathcal{C}$ *, is the set defined by* 

$$\bigcup \mathcal{C} = \{x | x \in S \text{ for some } S \in \mathcal{C}\}.$$

If C is a non-empty collection, its intersection is the set  $\bigcap C$  given by

$$\bigcap \mathcal{C} = \{x | x \in S \text{ for every } S \in \mathcal{C}\}.$$

Here, the generalized union and generalized intersection of sets simply involve performing the union or intersection operation on each element within the set.

Note that if  $\mathcal{C}$  and  $\mathcal{D}$  are two collections such that  $\mathcal{C} \subseteq \mathcal{D}$ , then

$$\bigcup \mathcal{C} \subseteq \bigcup \mathcal{D} \text{ and } \bigcap \mathcal{D} \subseteq \bigcap \mathcal{C}$$

Within the framework of collections of subsets of a given set S, we extend the previous definition by taking  $\bigcap \emptyset = S$  for the empty collection of subsets of S.

This is consistent with the fact that  $\emptyset \subseteq \mathcal{C}$  implies  $\bigcap \mathcal{C} \subseteq \mathcal{S}$ .

**Definition 1.1.2** The symmetric difference of sets denoted by  $\oplus$  is defined by  $U \oplus V = (U - V) \cup (V - U)$  for all sets U, V.

The symmetric difference operation is easily shown to satisfy symmetry and associativity, and also  $U \oplus U = \emptyset$ . Next, we will prove the associativity.

#### **Proof 1.1.1**

$$LEFT = ((U - V) \cup (V - U)) \oplus T$$

$$= (((U - V) \cup (V - U)) - T) \cup (T - ((U - V) \cup (V - U)))$$

$$= ((U\bar{V} \cup V\bar{U}) \cap \bar{T}) \cup (T \cap (\bar{U}\bar{V} \cup UV))$$

$$= \sum_{i=1,2,4,7} m_i$$

RIGHT = LEFT

However, the above method is rather cumbersome. We can instead adopt the characteristic function approach for the proof. Here, we first supplement the relevant concepts of characteristic functions.

**Definition 1.1.3** For any set A, the function

$$\chi_A(x) = \begin{cases} 1, & x \in A, \\ 0, & x \notin A \end{cases}$$

is called the characteristic function of the set A.

The characteristic function has the following properties:

- 1. The necessary and sufficient condition for A=X is  $\chi_A(x)\equiv 1$ , and the necessary and sufficient condition for  $A=\varnothing$  is  $\chi_A(x)\equiv 0$ ;
- 2. The necessary and sufficient condition for  $A \subset B$  is

$$\chi_A(x) \leqslant \chi_B(x), (\forall x \in X);$$

- 3.  $\chi_{\bar{P}}(x) = 1 \chi_{P}(x)$ .
- 4.  $\chi_{A \cup B}(x) = \chi_A(x) + \chi_B(x) \chi_{A \cap B}(x)$ .
- 5.  $\chi_{A \cap B}(x) = \chi_A(x) \cdot \chi_B(x)$ ;
- 6.

$$\chi_{\bigcup_{\alpha \in \Lambda} A_{\alpha}}(x) = \max_{\alpha \in \Lambda} \chi_{A_{\alpha}}(x)$$

$$\chi_{\bigcap_{\alpha \in \Lambda} A_{\alpha}}(x) = \min_{\alpha \in \Lambda} \chi_{A_{\alpha}}(x)$$

7. Let  $\{A_k\}$  be an arbitrary sequence of sets, then

$$\chi_{\overline{\lim}_{k\to\infty}A_k}(x) = \overline{\lim}_{k\to\infty}\chi_{A_k}(x)$$

$$\chi_{\underline{\lim}_{k\to\infty}A_k}(x) = \underline{\lim}_{k\to\infty}\chi_{A_k}(x);$$

8. The necessary and sufficient condition for  $\lim_{k\to\infty} A_k$  to exist is that  $\lim_{k\to\infty} \chi_{A_k}(x)$  exists  $(\forall x\in X)$ , and when the limit exists, we have

$$\chi_{\lim_{k\to\infty} A_k}(x) = \lim_{k\to\infty} \chi_{A_k}(x) \quad (x\in X).$$

$$\overline{\lim}_{k \to \infty} A_k = \bigcap_{n=1}^{\infty} \bigcup_{k=n}^{\infty} A_k$$

The limit superior of a sequence of sets reflects the collection of elements that "repeatedly appear" in the sequence of sets.

$$\overline{\lim}_{k\to\infty}A_k=\bigcup_{n=1}^\infty\bigcap_{k=n}^\infty A_k$$

The limit inferior of a sequence of sets reflects the collection of elements that exhibit "stable membership" within the sequence of sets.

$$\varlimsup_{k\to\infty}\chi_{A_k}(x)=\lim_{n\to\infty}(\sup_{k\geqslant n}\chi_{A_k(x)})$$

$$\underline{\lim}_{k \to \infty} \chi_{A_k}(x) = \lim_{n \to \infty} (\inf_{k \geqslant n} \chi_{A_k(x)})$$

Next, we will employ the method of characteristic functions to complete the proof of the associativity of the symmetric difference operation.

#### **Proof 1.1.2**

$$\chi_{A \oplus B}(x) = \chi_{(A-B) \cup (B-A)}(x)$$

$$= \chi_{A-B}(x) + \chi_{B-A}(x)$$

$$= \chi_{A}(x) \cdot \chi_{\bar{B}}(x) + \chi_{\bar{A}}(x) \cdot \chi_{B}(x)$$

$$= \chi_{A}(x) + \chi_{B}(x) - 2\chi_{A \cap B}(x)$$

$$LEFT = \chi_{(U \oplus V) \oplus T}(x)$$

$$= \chi_{U}(x) + \chi_{V}(x) + \chi_{T}(x) - 2\chi_{U \cap V}(x) - 2\chi_{V \cap T}(x) - 2\chi_{U \cap T}(x) + 4\chi_{U \cap V \cap T}(x)$$

$$= RIGHT$$

### **Definition 1.1.4** Below are several concepts pertaining to sets.

- 1. An ordered pair is a collection of sets  $\{\{x,y\},\{x\}\}$ . It can be readily verified that x and y are determined uniquely.
- 2. Let  $\{\{x,y\},\{x\}\}$  be an ordered pair. Then x is the first component of p and y is the second component of p.
- 3. Let X, Y be two sets. Their product is the set  $X \times Y$  that consists of all pairs of the form (x, y) where  $x \in X$  and  $y \in Y$ .
- 4. Let C and D be two collections of sets such that  $\bigcup C = \bigcup D$ . D is a **refinement** of C if, for every  $D \in D$ , there exists  $C \in C$  such that  $D \subseteq C$ . This is denoted by  $C \subseteq D$ .
- 5. A collection of sets C is **hereditary** if  $U \in C$  and  $W \subseteq U$  implies  $W \in C$ .
- 6. The set of subsets of S that contain k elements is denoted by  $\mathcal{P}_k(S)$ . Clearly, for every set S, we have  $\mathcal{P}_0(S) = \{\emptyset\}$ . The set of all finite subsets of a set S is denoted by  $\mathcal{P}_{fin}(S) = \bigcup_{k \in \mathcal{N}} \mathcal{P}_k(S)$ .
- 7. Let C be a collection of sets and let U be a set. The **trace** of the collection C on the set U is the collection  $C_U = \{U \cap C | C \in C\}$ .

#### **Definition 1.1.5** *Let* C *and* D *be two collections of sets.*

- 1.  $C \vee D = \{C \cup D | C \in C \text{ and } D \in D\},\$
- 2.  $C \wedge D = \{C \cap D | C \in C \text{ and } D \in D\},\$
- 3.  $C D = \{C D | C \in C \text{ and } D \in D\}.$

Attention, unlike  $\cup$  and cap, the operations  $\vee$  and  $\wedge$  between collections of sets are not idempotent. Indeed, we have, for example,

$$\mathcal{D} \vee \mathcal{D} = \{\{y\}, \{x, y\}, \{u, y, z\}, \{u, x, y, z\}\} \neq \mathcal{D}.$$

**Definition 1.1.6** A partition of a non-empty subsets of S that are pairwise disjoint and whose union equals S. The members of  $\pi$  are referred to as the blocks of the partition  $\pi$ . The collection of partitions of a set S is denoted by PART(S). A partition is finite if it has a finite number of blocks. The set of finite partitions of S is denoted by  $PART_{fin}(S)$ .

If  $\pi \in PART(S)$  then a subset T of S is  $\pi$ -saturated if it is a union of blocks of  $\pi$ .

### 1.2 Relations and Functions

**Definition 1.2.1** Let X, Y be two sets. A relation on X, Y is a subset  $\rho$  of the set product  $X \times Y$ . If X = Y = S, we refer to  $\rho$  as a relation on S. The relation  $\rho$  on S is:

- reflexive if  $(x, x) \in \rho$  for every  $x \in S$ ;
- irreflexive if  $(x, x) \notin \rho$  for every  $x \in S$ ;

- symmetric if $(x, y) \in \rho$  implies  $(y, x) \in \rho$  for all  $x, y \in S$ ;
- antisymmetric if  $(x, y) \in \rho$  and  $(y, x) \in \rho$  imply x = y for all  $x, y \in S$ ;
- transitive if  $(x, y) \in \rho$  and  $(y, z) \in \rho$  imply  $(x, z) \in \rho$  for all  $x, y, z \in S$ .

A **partial order** on S is a relation  $\rho$  that belongs to  $\text{REFL}(S) \cap \text{ANTISYMM}(S) \cap \text{TRAN}(S)$ , that is, a relation that is reflexive, symmetric and transitive.

In current mathematical practice, we often write  $x\rho y$  instead on  $(x,y) \in \rho$ , where  $\rho$  is a relation of S and  $x,y \in S$ . This alternative way to denote the fact that (x,y) belongs to  $\rho$  is known as the infix notation.

**Definition 1.2.2** Let X, Y be two sets. A function(or a mapping) from X to Y is a relation f on X, Y such that  $(x,y),(x,y') \in f$  implies y=y'.

Let X, Y be two sets and let  $f: X \to Y$ .

The domain of f is the set

$$Dom(f) = \{x \in X | y = f(x) \text{ for some } y \in Y\}.$$

The range of f is the set

$$Ran(f) = \{ y \in Y | y = f(x) \text{ for some } x \in X \}.$$

**Definition 1.2.3** Let X be a set,  $Y = \{0,1\}$  and let L be a subset of S. The characteristic function is discussed above. The indicator function of L is the function  $I_L: S \to \mathcal{R} \cup \infty$  defined by

$$I_L(x) = \begin{cases} 1 & \text{if } x \in L. \\ \infty & \text{otherwise} \end{cases}$$

for  $x \in S$ .

**Definition 1.2.4** A function  $f: X \to Y$  is:

- 1. injective or one-to-one if  $f(x_1) = f(x_2)$  implies  $x_1 = x_2$  for  $x_1, x_2 \in Dom(f)$ ;
- 2. surjective or onto if Ran(f) = Y;
- 3. total if Dom(f) = X.

*If f both injective and surjective, then it is a bijective function.* 

**Theorem 1.2.1** A function  $f: X \to Y$  is injective if and only if there exists a function  $g: Y \to X$  such that g(f(x)) = x for every  $x \in \mathbf{Dom}(f)$ .

A function  $f: X \to Y$  is surjective if and only if there exists a function  $h: X \to Y$  such that f(h(y)) = y for every  $y \in Y$ .

Here we provide the proof for the latter theorem.

**Proof 1.2.1** Suppose that f is a surjective function. The collection  $\{f^{-1}(y)|y\in Y\}$  indexed by Y consists of non-empty sets. By the Axiom of Choice there exists a choice function for this collection, that is a function  $h:Y\to\bigcup_{y\in Y}f^{-1}(y)$  such that  $h(y)\in f^{-1}(y)$ , or f(h(y))=y for  $y\in Y$ .

Conversely, suppose that there exists a function  $h: X \to Y$  such that f(h(y)) = y for every  $y \in Y$ . Then, f(x) = y for y = h(y), which shows that f is surjective.

The Axiom of Choice: Let  $C = \{C_i | i \in I\}$  be a collection of non-empty sets indexed by a set I. There exists a function  $\phi : I \to \bigcup C$  (known as a choice function) such that  $\phi(i) \in C_i$  for each  $i \in I$ .

**Theorem 1.2.2** There is a bijection  $\Psi: \mathcal{P}(S) \to (S \to \{0,1\})$  between the set of subsets of S and the set of characteristic functions defined on S.

**Definition 1.2.5** A set S is indexed be a set I if there exists a surjection  $f: I \to S$ . In this case we refer to I as an index set.

If S is indexed by the function  $f: I \to S$  we write the element f(i) just as  $s_i$ , if there is no risk of confusion.

**Definition 1.2.6** A sequence of length n on a set X is a function  $x : \{0, 1, \dots, n-1\} \to X$ .